CMPSC 465
Data Structures and Algorithms
Spring 2022

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Linear Programming (Textbook, Section 7.1)

Duality of LP (I)

Consider

$$\begin{aligned} \text{maximize} \quad & x_1 + 2x_2 \\ \text{subject to} \quad & x_1 \leq 20 \\ & x_2 \leq 30 \\ & x_1 + x_2 \leq 40 \\ & x_1, x_2 \geq 0 \end{aligned}$$

Can we show the optimal solution is at least 60? Check (0, 30)

Can we show that optimal solution is at most 90? Use linear combinations constraints

Duality of LP (II)

Define a variable for each constraint

$$\begin{aligned} \text{maximize} \quad & x_1 + 2x_2 & \\ \text{subject to} \quad & x_1 \le 20 & \color{red}{y_1} \\ & x_2 \le 30 & \color{red}{y_2} \\ & x_1 + x_2 \le 40 & \color{red}{y_3} \\ & x_1, x_2 \ge 0 & \end{aligned}$$

Adding them together:

$(y_1 + y_3)x_1 + (y_2 + y_3)x_2 \le 20y_1 + 30y_2 + 40y_3$

We let $y_1 + y_3 \ge 1$ and $y_2 + y_3 \ge 2$ to get an upper bound on $x_1 + 2x_2$:

$x_1 + 2x_2 \le (y_1 + y_3)x_1 + (y_2 + y_3)x_2 \le 20y_1 + 30y_2 + 40y_3$

Duality of LP (III)

Primal LP

$$\begin{aligned} \text{maximize} \quad & x_1 + 2x_2 \\ \text{subject to} \quad & x_1 \leq 20 \\ & x_2 \leq 30 \\ & x_1 + x_2 \leq 40 \\ & x_1, x_2 \geq 0 \end{aligned}$$

Optimal solution: $(x_1, x_2) = (10, 30) \implies x_1 + 2x_2 = 70$

Dual LP

$$\begin{aligned} \text{minimize} \quad & 20y_1 + 30y_2 + 40y_3 \\ \text{subject to} \quad & y_1 + y_3 \geq 1 \\ & y_2 + y_3 \geq 2 \\ & y_1, y_2, y_3 \geq 0 \end{aligned}$$

Optimal solution: $(y_1, y_2, y_3) = (0, 1, 1) \implies 20y_1 + 30y_2 + 40y_3 = 70$

Duality of LP(IV)

More generally

Primal LP

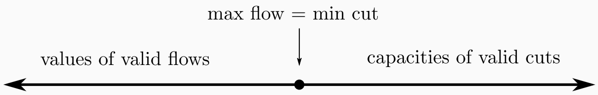
$$
\begin{aligned}
\max \quad & c_1x_1 + c_2x_2 + \cdots + c_nx_n \\
\text{s.t.} \quad & a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n \le b_1 \\
& a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n \le b_2 \\
& \quad \vdots \\
& a_{m1}x_1 + a_{m2}x_2 + \cdots + a_{mn}x_n \le b_m \\
& x_1, x_2, \ldots, x_n \ge 0
\end{aligned}
$$

Dual LP

$$
\begin{aligned}
\min \quad & b_1y_1 + b_2y_2 + \cdots + b_my_m \\
\text{s.t.} \quad & a_{11}y_1 + a_{21}y_2 + \cdots + a_{m1}y_m \ge c_1 \\
& a_{12}y_1 + a_{22}y_2 + \cdots + a_{m2}y_m \ge c_2 \\
& \quad \vdots \\
& a_{1n}y_1 + a_{2n}y_2 + \cdots + a_{mn}y_m \ge c_n \\
& y_1, y_2, \ldots, y_m \ge 0
\end{aligned}
$$

Duality of LP (V)

Duality of flow and cut



For LP we have:

Theorem (Weak Duality)

A feasible solution to the dual LP is an upper bound on any feasible solution to the primal LP

Theorem (Strong Duality)

The optimal solution to the dual LP is equal to the optimal solution to the primal LP