## Exam1 Solution – Penn State CMPSC 461 Spring 2022, Prof. G. Tan.

Question about grammar derivations:

```
• Left-most derivation of 5E+2
```

- <SNFloat> -> <Float>E<Exponent>
- -> <NonZeroDigit>E<Exponent>
- -> 5E<Exponent>
- -> 5E+<Num>
- -> 5E+<Digit>
- -> 5E+<NonZeroDigit>
- -> 5E+2
- Right-most derivation of 3.14E3
  - <SNFloat> -> <Float>E<Exponent>
  - -> <Float>E<Digit>
  - -> <Float>E<NonZeroDigit>
  - -> <Float>E3
  - -> <NonZeroDigit>.<Num>E3
  - -> <NonZeroDigit>.<Digit><Digit>E3
  - -> <NonZeroDigit>.<Digit><NonZeroDigit>E3
  - -> <NonZeroDigit>.<Digit>4E3
  - -> <NonZeroDigit>.<NonZeroDigit>4E3
  - -> <NonZeroDigit>.14E3
  - -> 3.14E3
- Why 100.4E+1 is not part of the language defined by the grammar? There are multiple ways to answer this question.
  - (1) **Natural language** According the grammar, for 100.4E+1, 100.4 must be derived from <Float>, but <Float> can have only one <NonZeroDigit> before its "."; thus 100.4E+1 cannot be derived from the grammar.
  - (2) **Derivation** The derivation for the string will get stuck at the following step, with a left-most derivation.
    - <SNFloat> ->\* <NonZeroDigit>.<Num>E<Exponent>
  - (3) Parse tree A stuck parse tree would also be a valid answer.

## Question about scoping:

- For the version with first line being "int a b;"
  - (a) main: <b,10>, <a,1>
    - foo: <a,2>, <b,1>
    - bar: <a.6>. <b.1>
  - (b) < a,6 >, < b,10 >
  - (c) < a,6 >, < b,10 >

For the version with first line being "int a;"

```
(d) main: <b,10>, <a,1> foo: <b,2>, <a,3> bar: <a,5> (e) <a,5>, <b,10> (f) <a,5>, <b,2>
```

## Question about recursive descent parsing:

- For the version with grammar "<term> -> <term> \* <id> | <id>; <id> -> a | b | c"
  - (a) The grammar contains left recursion in the rules of <term>. A recursive-descent parser results in an infinite loop when parsing a string from such a grammar.

```
(b)
   <term> -> <id> {* <id>}
   <id> -> a | b | c
   <term> -> <id> * <term> | <id>
   <id> -> a | b | c
(c) void term() {
     id(); // parses an <id> and advance token to nextToken ()
    while ( token.type == timesOp ) {
      token = nextToken();
      id();
     }
   }
   or
   void term() {
    id(); // parses an <id> and advance token to nextToken ()
    if ( token.type == timesOp ) {
      token = nextToken();
       term();
     }
   }
```

- For the version with grammar "<clause> -> <clause> and <phrase> | <phrase>"
  - (a) The grammar contains left recursion in the rules of <clause>. A recursive-descent parser results in an infinite loop when parsing a string from such a grammar.

```
(b)
     <clause> -> <phrase> {and <phrase>}
     <phrase> -> ...
     or
     <clause> -> <phrase> and <clause> | <phrase>
     <phrase> -> ...
```

```
(c) void clause() {
    phrase(); // parses a <phrase> and advance token to nextToken ()
    while ( token.type == andOp ) {
        token = nextToken();
        phrase();
    }
    or
    void clause() {
        phrase(); // parses a <phrase> and advance token to nextToken ()
        if ( token.type == andOp ) {
            token = nextToken();
            clause();
        }
    }
}
```

Question about no nonleading zeros in <Exponent>

• One possible answer:

```
<Exponent> -> [(+|-)] <CNum> <CNum> -> <Digit> | <NonZeroDigit><Num>
```

• Another possible answer:

```
<Exponent> -> [(+|-)]<Digit> | [(+|-)]<NonZeroDigit> \{<Digit>\}
```