

1. (a)

$\langle e \rangle \rightarrow \langle e \rangle / \langle e \rangle \rightarrow \langle d \rangle / \langle e \rangle \rightarrow 9 / \langle e \rangle \rightarrow 9 / \langle e \rangle * \langle e \rangle \rightarrow 9 / \langle d \rangle * \langle e \rangle \rightarrow$
 $9 / 3 * \langle e \rangle \rightarrow 9 / 3 * \langle d \rangle \rightarrow 9 / 3 * 3$

Another possible answer is: $\langle e \rangle \rightarrow \langle e \rangle * \langle e \rangle \rightarrow \dots$

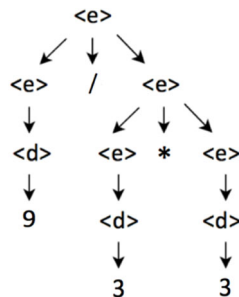
1. (b)

$\langle e \rangle \rightarrow \langle e \rangle * \langle e \rangle \rightarrow \langle e \rangle * \langle d \rangle \rightarrow \langle e \rangle * 3 \rightarrow \langle e \rangle / \langle e \rangle * 3 \rightarrow \langle e \rangle / \langle d \rangle * 3 \rightarrow$
 $\langle e \rangle / 3 * 3 \rightarrow \langle d \rangle / 3 * 3 \rightarrow 9 / 3 * 3$

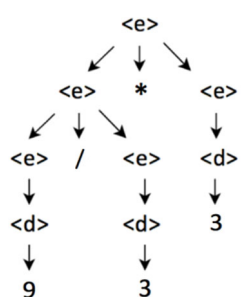
Another possible answer is: $\langle e \rangle \rightarrow \langle e \rangle / \langle e \rangle \rightarrow \dots$

1. (c)

For (a):



For (b):



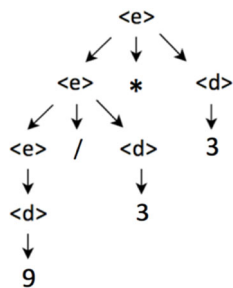
Note: depending on the answers for (b) and (c), it's possible that the parse trees for the two derivations are the same (one of the above).

1. (d)

$\langle e \rangle \rightarrow \langle d \rangle \mid \langle e \rangle / \langle d \rangle \mid \langle e \rangle * \langle d \rangle$

$\langle d \rangle \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

The parse tree for "9 / 3 * 3" based on the new grammar:



Argument for unambiguity: there are two choices to begin parsing 9/3*3:

(1) $\langle e \rangle \rightarrow \langle e \rangle / \langle d \rangle$ or (2) $\langle e \rangle \rightarrow \langle e \rangle * \langle d \rangle$

For option (1), it is impossible to get the expression 9/3*3 from $\langle e \rangle \rightarrow \langle e \rangle / \langle d \rangle \rightarrow \langle e \rangle / 3 \rightarrow \langle e \rangle * \langle d \rangle / 3$.

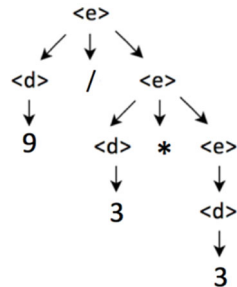
For option (2), it is shown above. Hence, (2) is the only choice.

1. (e)

$\langle e \rangle \rightarrow \langle d \rangle \mid \langle d \rangle / \langle e \rangle \mid \langle d \rangle * \langle e \rangle$

$\langle d \rangle \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

The parse tree for "9 / 3 * 3" based on the new grammar:

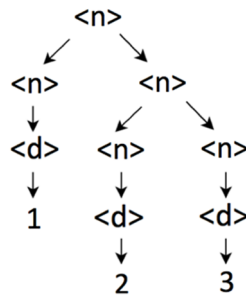


2.

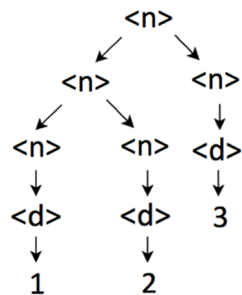
The grammar is ambiguous.

Example input: 123

Corresponding parse tree 1:



parse tree 2:



An equivalent unambiguous grammar:

$\langle n \rangle \rightarrow \langle n \rangle \langle d \rangle \mid \langle d \rangle$

$\langle d \rangle \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

3. $\langle S \rangle \rightarrow () \mid (\langle S \rangle)$

4.

$\langle e \rangle \rightarrow \langle u \rangle \mid \langle e \rangle + \langle u \rangle \mid \langle e \rangle - \langle u \rangle$

$\langle u \rangle \rightarrow \langle d \rangle \mid \sim \langle u \rangle$

$\langle d \rangle \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

5.

$\langle \text{email} \rangle \rightarrow \langle \text{account} \rangle @ \langle \text{subDomains} \rangle . \langle \text{topDomain} \rangle$

$\langle \text{account} \rangle \rightarrow \langle \text{letter} \rangle \mid \langle \text{account} \rangle \langle \text{letter} \rangle \mid \langle \text{account} \rangle \langle \text{digit} \rangle$

$\langle \text{subDomains} \rangle \rightarrow \langle \text{letterDigitSeq} \rangle \mid \langle \text{letterDigitSeq} \rangle . \langle \text{subDomains} \rangle$

$\langle \text{letterDigitSeq} \rangle \rightarrow \langle \text{letter} \rangle \mid \langle \text{digit} \rangle$

$\mid \langle \text{letter} \rangle \langle \text{letterDigitSeq} \rangle$

$\mid \langle \text{digit} \rangle \langle \text{letterDigitSeq} \rangle$

$\langle \text{topDomain} \rangle \rightarrow \text{com} \mid \text{org} \mid \text{gov} \mid \text{net}$

<letter> -> a | b | c | ... | z | A | B | C | ... | Z
<digit> -> 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9

6. BNF grammar:

<SNFloat> -> <Float> | <Float>E<Exponent>
<Float> -> <NonZeroDigit> | <NonZeroDigit>.<Num>
<Exponent> -> <Num> | +<Num> | -<Num>
<Num> -> <Digit> | <Digit><Num>
<Digit> -> 0 | <NonZeroDigit>
<NonZeroDigit> -> 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9