

# Midterm 1

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|--|------------------------------|---------------------|
| <b>Due</b> No due date                   | <b>Points</b> 24             | <b>Questions</b> 19 |
| <b>Available</b> until Feb 16 at 11:20pm | <b>Time Limit</b> 62 Minutes |                     |

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## Instructions

- You have **60 minutes** to complete this exam.
- The exam will be stopped at **8pm**.
- This exam is closed book/closed notes; **you may not communicate with anyone** other than the instructor and/or TAs and LAs during the exam.
- Any question about the exam can be posted in campuswire; make sure it's **visible only to TAs and instructors**. We will also make broad announcements and clarifications in campuswire. You can join campuswire using <https://campuswire.com/p/GB6D343A6> (<https://campuswire.com/p/GB6D343A6>) and sign-up code of 8502.
- Once you begin the exam, you must complete it within the time limit. If you logout of Canvas or close your browser after you enter the exam, the **countdown will not stop**.
- Any kind of **cheating may result in failing the course**.

This quiz was locked Feb 16 at 11:20pm.

## Attempt History

|        | Attempt                   | Time       | Score          |
|--------|---------------------------|------------|----------------|
| LATEST | <a href="#">Attempt 1</a> | 62 minutes | 14.5 out of 24 |

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Score for this quiz: **14.5** out of 24

Submitted Feb 16 at 8:02pm

This attempt took 62 minutes.

### Question 1

0.5 / 0.5 pts

For an unambiguous grammar, there is exactly one parse tree for any terminal string that belongs to the language defined by the grammar.

Correct!

☒ True



☐ False

## Question 2

0.5 / 0.5 pts

Static binding means that the association between a program entity and a property is decided at compile time.

Correct!

☒ True

☐ False

## Question 3

0.5 / 0.5 pts

Dynamic binding is also called late binding.

Correct!

☒ True

☐ False

## Question 4

0.5 / 0.5 pts

Any language that can be defined by a regular expression can also be defined by a context-free grammar.

Correct!

☒ True



☐ False

### Question 5

0.5 / 0.5 pts

BNF grammars do not allow left-recursive rules.

☐ True

☒ False

Correct!

### Question 6

0.5 / 0.5 pts

In C, the scope of a static local variable is the entire program.

☐ True

☒ False

Correct!

### Question 7

0 / 0.5 pts

An operator's precedence determines whether it associates to the left or right.

☒ True

☐ False

Not Answered

Correct Answer



### Question 8

0 / 0.5 pts

For a deterministic FSA, a state may have multiple outgoing edges labeled with the same input character.

Not Answered

☒ True

Correct Answer

☐ False

### Question 9

1 / 1 pts

Suppose  $r$  stands for the regular expression  $(a|b)^*c^*$ . Which one of the following strings is NOT a sentence of the language generated by  $r$ .

☐ aaabbbccc

☒ aaacccbbb

☐ bbbaaa

☐ aaa

Correct!

### Question 10

1 / 1 pts

Which one of the following is true of E-BNF grammars (extended BNF)?



Correct!

- ☐ E-BNF grammars allow one to describe the syntax that can not be described by BNF grammars
- ☒ Any E-BNF grammar can also be expressed in a BNF grammar
- ☐ E-BNF grammars allow one to describe the semantics of a programming language
- ☐ Any E-BNF grammar can also be expressed in a regular expression

### Question 11

0 / 1 pts

Which one of the following is true of l-values and r-values?

Correct Answer

- ☐ An l-value refers to a variable's memory location while an r-value to its current value

You Answered

- ☒ l-values are always to the left of r-values
- ☐ l-values are local and r-values are relative.
- ☐ An l-value is a logical value, and an r-value is a real value.

### Question 12

0 / 1 pts

Which following scheme of variable scoping is used by most modern programming languages?

You Answered

- ☒ Both static and dynamic scoping



Incorrect Answer

☐ Dynamic scoping

☐ Static scoping

### Question 13

1 / 1 pts

The following BNF grammar defines the grammar of the scientific notation for non-negative floating point numbers.

$\langle \text{SNFloat} \rangle \rightarrow \langle \text{Float} \rangle \mid \langle \text{Float} \rangle E \langle \text{Exponent} \rangle$

$\langle \text{Float} \rangle \rightarrow \langle \text{NonZeroDigit} \rangle \mid \langle \text{NonZeroDigit} \rangle . \langle \text{Num} \rangle$

$\langle \text{Exponent} \rangle \rightarrow \langle \text{Num} \rangle \mid + \langle \text{Num} \rangle \mid - \langle \text{Num} \rangle$

$\langle \text{Num} \rangle \rightarrow \langle \text{Digit} \rangle \mid \langle \text{Digit} \rangle \langle \text{Num} \rangle$

$\langle \text{Digit} \rangle \rightarrow 0 \mid \langle \text{NonZeroDigit} \rangle$

$\langle \text{NonZeroDigit} \rangle \rightarrow 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

Show the left-most derivation for 5E+2; show every step.

Your Answer:

$\langle \text{SNFloat} \rangle \rightarrow \langle \text{Float} \rangle E \langle \text{Exponent} \rangle$

$\langle \text{Float} \rangle E \langle \text{Exponent} \rangle \rightarrow \langle \text{NonZeroDigit} \rangle E \langle \text{Exponent} \rangle$

$\langle \text{NonZeroDigit} \rangle E \langle \text{Exponent} \rangle \rightarrow 5 E \langle \text{Exponent} \rangle$

$5 E \langle \text{Exponent} \rangle \rightarrow 5 E + \langle \text{Num} \rangle$

$5 E + \langle \text{Num} \rangle \rightarrow 5 E + \langle \text{Digit} \rangle$

$5 E + \langle \text{Digit} \rangle \rightarrow 5 E + \langle \text{NonZeroDigit} \rangle$

$5 E + \langle \text{NonZeroDigit} \rangle \rightarrow 5 E + 2$

### Question 14

1 / 1 pts



The following BNF grammar defines the grammar of the scientific notation for non-negative floating point numbers.

$\langle \text{SNFloat} \rangle \rightarrow \langle \text{Float} \rangle \mid \langle \text{Float} \rangle \text{E} \langle \text{Exponent} \rangle$

$\langle \text{Float} \rangle \rightarrow \langle \text{NonZeroDigit} \rangle \mid \langle \text{NonZeroDigit} \rangle . \langle \text{Num} \rangle$

$\langle \text{Exponent} \rangle \rightarrow \langle \text{Num} \rangle \mid + \langle \text{Num} \rangle \mid - \langle \text{Num} \rangle$

$\langle \text{Num} \rangle \rightarrow \langle \text{Digit} \rangle \mid \langle \text{Digit} \rangle \langle \text{Num} \rangle$

$\langle \text{Digit} \rangle \rightarrow 0 \mid \langle \text{NonZeroDigit} \rangle$

$\langle \text{NonZeroDigit} \rangle \rightarrow 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

Explain why 100.4E+1 is not part of the language defined by the grammar.

Your Answer:

because the 100 in 100.4E+1 is not accessible by the grammar

for  $\langle \text{Float} \rangle \rightarrow \langle \text{NonZeroDigit} \rangle \mid \langle \text{NonZeroDigit} \rangle . \langle \text{Num} \rangle$ , you can only have NonZeroDigit before the decimal point.

### Question 15

1 / 1 pts

The following BNF grammar defines the grammar of the scientific notation for non-negative floating point numbers.

$\langle \text{SNFloat} \rangle \rightarrow \langle \text{Float} \rangle \mid \langle \text{Float} \rangle \text{E} \langle \text{Exponent} \rangle$

$\langle \text{Float} \rangle \rightarrow \langle \text{NonZeroDigit} \rangle \mid \langle \text{NonZeroDigit} \rangle . \langle \text{Num} \rangle$

$\langle \text{Exponent} \rangle \rightarrow \langle \text{Num} \rangle \mid + \langle \text{Num} \rangle \mid - \langle \text{Num} \rangle$

$\langle \text{Num} \rangle \rightarrow \langle \text{Digit} \rangle \mid \langle \text{Digit} \rangle \langle \text{Num} \rangle$

$\langle \text{Digit} \rangle \rightarrow 0 \mid \langle \text{NonZeroDigit} \rangle$

$\langle \text{NonZeroDigit} \rangle \rightarrow 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

Show the right-most derivation for 3.14E3; show every step.



Your Answer:

<SNFloat> -> <Float>E<Exponent>  
<Float>E<Exponent>-><Float>E<Num>  
<Float>E<Num>-><Float>E<Digit>  
<Float>E<Digit>-><Float>E<NonZeroDigit>  
<Float>E<NonZeroDigit>-><Float>E3  
<Float>E3-><NonZeroDigit>.<Num>E3  
<NonZeroDigit>.<Num>E3-><NonZeroDigit>.<Digit><Num>E3  
<NonZeroDigit>.<Digit><Num>E3-><NonZeroDigit>.<Digit><Digit>E3  
<NonZeroDigit>.<Digit><Digit>E3-><NonZeroDigit>.<Digit>  
<NonZeroDigit>E3  
<NonZeroDigit>.<Digit><NonZeroDigit>E3-><NonZeroDigit>.  
<Digit>4E3  
<NonZeroDigit>.<Digit>4E3-><NonZeroDigit>.<NonZeroDigit>4E3  
<NonZeroDigit>.<NonZeroDigit>4E3-><NonZeroDigit>.14E3  
<NonZeroDigit>.14E3->3.14E3

## Question 16

3 / 3 pts

For the following C program, you will be asked to determine which variables are visible in a number of different situations. In each case, identify each variable by its name and the line number of its declaration. Please clearly label the answer for each subquestion.

```
1 int a,b;  
2 void foo(int a) {  
3 ...  
4 }  
  
5 void bar () {  
6 int a;  
7 ...  
8 }
```





```
9 void main() {  
10 int b;  
11 ...  
12 }
```

(a) (1 point) C uses static scoping. Say which variables are visible in the bodies of each of the functions: main, foo, bar.

(b) (1 point) If C used dynamic scoping and the calling sequence is that main calls bar, say which variables would be visible in bar.

(c) (1 point) If C used dynamic scoping and the calling sequence is that main calls foo, and foo calls bar, say which variables would be visible in bar.

Your Answer:

A)

main - <b,10><a,1>

bar - <a,6><b,1>

foo - <a,2><b,1>

B)

<a,6><b,10>

C)

<a,6><b,10>

## Question 17

1 / 4 pts

Answer this question based on the following BNF grammar with the start non-terminal being <foo>.

<foo> -> <bar> | <bar> \$ <foo>

<bar> -> <baz> | <bar> # <baz>

<baz> -> x | y | ( <foo> )

- Which operator has higher precedence? \$

- The associativity of the \$ operator is [ Select ]





• The associativity of the # operator is [ Select ]

• Is the grammar ambiguous or unambiguous?

[ Select ]

**Answer 1:**

Incorrect Answer

#

Not Answered

\$

**Answer 2:**

Correct!

right

**Answer 3:**

Incorrect Answer

left

Not Answered

right

**Answer 4:**

Not Answered

ambiguous

Incorrect Answer

unambiguous

## Question 18

1 / 3 pts

Answer the following questions based on the following grammar:

<clause> -> <clause> and <phrase> | <phrase>

<phrase> -> ...

(a) (1 point) Explain in 1 or 2 sentences why we cannot directly use recursive-descent parsing on the grammar.

(b) (1 point) Transform the grammar into an equivalent grammar on which recursive-descent parsing can be applied.



(c) (1 point) Write some pseudo code for the parsing method for nonterminal `<clause>` in the transformed grammar, assuming there is already a parsing method for nonterminal `<phrase>`.

Please clearly label your answer for each subquestion.

Your Answer:

a)

because the recursive does not guarantee to stop.

B)

```
<clause> -> <clause> and <phrase> | <phrase>
<phrase> -> terminate
```

C)

While input != class:

```
<clause>;
```

```
if <phrase>:
```

```
    break;
```

(-0.5) (a) The answer is partially correct: does not mention the reason. (-0.5) (b) The answer does not remove left recursion. (-1.0) (c) The parsing is incorrect.

## Question 19

1.5 / 3 pts

The following E-BNF grammar defines the grammar of the scientific notation for non-negative floating point numbers. Note that "[", "]", "(", ")", "|", "{", and "}" are meta-symbols of E-BNF.

```
<SNFloat> -> <Float> [E<Exponent>]
```

```
<Float> -> <NonZeroDigit> [.<Num>]
```

```
<Exponent> -> [(+|-)]<Num>
```



```
<Num> -> <Digit>{<Digit>}  
<Digit> -> 0 | <NonZeroDigit>  
<NonZeroDigit> -> 1 | 2 | ... | 9
```

The rule for Exponent allows numbers such as 0023; change the rule for <Exponent> so that numbers with leading zeros are illegal in exponents (the number 0 itself should still be legal though). You can use either EBNF or BNF in your answer.

Your Answer:

```
<Exponent> -> (+|-)+<NonZeroDigit>+<Num>*|(+|-)+<Digit>
```

- "(+|-)+" not correct. - <NonZeroDigit>+ not correct - shouldn't use the regexp syntax.

Quiz Score: **14.5** out of 24