CS 461

Programming Language Concepts

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CH2 SYNTAX

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BNF Example

- ☐ A simple arithmetic expression language
 - Non-terminals: e, n, d; identify grammatical categories
 - Terminals: +, -, 0, 1, 2, 3, 4, 5, ..., 9; identify the basic alphabet
 - Production rules

<e> -> <n> | <e>+<e> | <e>-<e><
<n> -> <d> | <n><d>< | <n><d><n><d>< | <n>< | <n

- "|" indicates a choice
- the right hand side of a rule can be a sequence of terminal or nonterminal symbols
- Start symbol: e
- Note: the grammar has **recursion**
- ☐ Example numbers: 4, 27, 704
- ☐ Example expressions: 27 4 + 704

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Derivations

- ☐ Derivation: a sequence of replacement steps
 - starting from the start symbol
 - Replace a nonterminal by the rhs (right hand side) of a rule
 - Keep doing it until resulting in a string of terminals

<e> \rightarrow <e>+<e> \rightarrow <e>-<e>+<e> \rightarrow <n>-<n>+<n> \rightarrow <n>-<d>+<d> \rightarrow \rightarrow <d>-<d>+<d>+<d> \rightarrow ... \rightarrow 27 - 4 + 3

Grammar defines a language: any terminal string that can be

derived belongs to the language of the grammar

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Left vs. right-most derivations

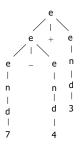
- □ Left-most derivation
 - At each step, always replace the leftmost nonterminal by one of its alternatives
- □ Right-most derivation
 - At each step, replace the rightmost nonterminal by one of its alternatives
- □ An example
 - 4 + 3
 - Both derivations correspond to the same parse tree

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Parse tree

□ Derivation represented by a tree

- Parse tree ignores the detail about the order of replacing nonterminals
- Two derivations may correspond to the same parse tree
- Tree shows parenthesization of expressions: 7-4+3 treated as (7-4) + 3



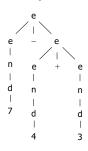
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Parse Tree and Ambiguity

• The expression 7 - 4 + 3 has two parse trees; another one



- Problem: $7 (4 + 3) \neq (7 4) + 3$
- A grammar is **ambiguous** if some terminal string has more than one parse tree. Otherwise it is unambiguous
 - Note: not two derivations; to show a grammar is ambiguous, need to find one terminal string with two parse

Ways to resolve operator ambiguity

□ Second way: design a new grammar that enforces associativity and precedence

<e> -> <n> | <e>+<n> | <e>-<n>

- 27 4 +3 has only one parse tree in this grammar; which one it is?
- Q: how many parse trees for 27 4 + 3

 The grammar makes the +/- operator left associative (by convention) By using left recursion
 - Q: what if we want to make + and right associative?
- □ However, what if we really want 27 (4 + 3)?
 <e> -> <t> | <e>+<t> | <e>-<t> <</p>
 <t> | <e>>
 - Ex: parse tree for 27 (4+3)

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Examples of enforcing associativity and precedence

- ☐ Four operators: +, -, *, /
 - they are left associative
 - *, / should have higher precedence than + and -

<exp> -> <term> | <exp> + <term> | <exp> - <term> <term> -> <factor> | <term> * <factor> | <term> / <factor>

<factor> -> <num> | (<exp>)

- Example: 7 + 4 * 3
 - Why can't it be parsed like (7 + 4) * 3?
 - The explicit parentheses allow the parsing of "(7+4)*3"
 - More examples: 7*4*(3 + 5)

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Ways to resolve operator ambiguity

☐ One way: add parentheses explicitly

- <e> -> <n> | (<e> + <e>) | (<e> <e>)
- 7 4 + 3 is not an expression

Associativity & Precedence

- □ Associativity
 - Parenthesize operators of equal precedence to the left (or the
 - + is left associative, parse 3 + 4 + 5 as (3 + 4) + 5
 - Parse 3-4+5 as (3-4)+5
 - Example : the exponentiation operator is right associative
- ☐ Precedence: an operator has a higher precedence than another operator if the first one binds tighter
 - Group * before +
 - Parse 2 + 3*4 as 2+ (3*4)

E-BNF (Extended BNF)

- □ Put alternative parts in parentheses and separate them with vertical bars
 - <exp> -> <exp> (+ | -) <exp>
 - Really just an abbreviation for two BNF rules
 - Note: "(", "|", and ")" are meta-symbols, while "+" and "-" are terminals
- □ Put repetition (0 or more) in curly braces { }
 - <num> -> <digit> {<digit>}
 - Example: 101
- ☐ Optional parts are placed in square brackets []
 - <if-stmt> -> if <test> then <stmt> [else <stmt>]
- □ Examples:
 - <Expr> -> <Term> { (+ | -) <Term>}
 - <proc_call> -> <ident> ([<expr_list>])

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E-BNF is no more powerful than BNF

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□ <exp> -> <exp> (+ | -) <exp>
   • Same as
   <exp> -> <exp> + <exp> |<exp> - <exp>
\square <num> -> <digit> {<digit>}
   • Same as
      <num> -> <digit> | <digit> <num>
□ <if-stmt> -> if <test> then <stmt> [else <stmt>]
      <if-stmt> -> if <test> then <stmt>
               | if <test> then <stmt> else <stmt>
                                                     13
```

Real Examples: Excerpt from Java Grammar

Variable declarations in Java

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→ <type-name> <declarator-list> ; <type-name> → boolean | byte | short | int | long | char | float | double <declarator-list> → <declarator> | <declarator> , <declarator-list>
> <variable-name> | <declarator> <variable-name> = <expr>

Assume <variable-name> and <expr> are defined elsewhere

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