CS 461

Programming Language Concepts

Gary Tan Computer Science and Engineering Penn State University

Imperative Programming

☐ Oldest and most well-developed paradigm

• Mirrors computer architecture (von-neumann model)

□ Stateful computation

• A program's state: code, data; both in memory

• Memory/registers is the state: a map from addresses to

• Imperative constructs for modifying the state (memory) - E.g., x = y+1;

☐ Control flow: sequencing, loops, ...

□ Example Languages

· Fortran, Pascal

• C, Ada

2

4

6

Pure Functional Programming

□ Program defined as a set of functions

- Functions are defined in terms of the composition of other functions
- □ Stateless computation
 - Immutable values; no assignment statements
 - You just construct new values from old values
 - (define x 8); (define x (+ x 1))
 - (append I1 I2) constructs a new list out of old ones
 - Use GC to get rid of unused old values
- No construct can change the state
- ☐ Control flow: no sequencing of statements or loops; use recursion
- ☐ Examples: pure Racket, Core ML, Haskell

Imperative vs. Declarative Constructs

□ Imperative constructs

• int x = 1;

• x = x + 1; // increment x by one

☐ Declarative constructs (for declaring new entities)

• (define x 1)

• (define x (+ x 1))

• (define (f x) (+ x 1)) ☐ The distinction is between whether

• changing an existing value (change the state; side effects)

• or declaring a new value (purity)

Example:

(define x '(1 2 3)) (define y (cons 0 x))

(define x (cdr x))

what's the value of y at this point?

3

5

What Can Purity Give You? Referential Transparency

```
☐ (define (removeDup lst) ...)
```

(define lst1 (removeDup '(a b a a a)))

....; some complicated computations here (define lst2 (removeDup '(a b a a a)))

- $\hfill\square$ Observation: no need to recompute lst2
 - removeDup is given the same input, so the outputs should also be the same
- ☐ In general, in pure Racket, calling a function with the same input always produces the same result no matter where the fun call is
 - whenever you refer to it, it is always the same

No Referential Transparency in Imperative Languages such as C

 \square int y = 3;

int f (int x) { return x + y;}

int main () {

int z1 = f(3); ...; // some computation here

int $z^2 = f(3)$;

☐ Because the state can change, and the function's return value depends on the state;

Purity Enables Advanced Techniques

- □ Such as
 - Memoization; hash consing; automatic parallelization
- ☐ Memoization (caching): (fac n) -> (fac_mem n)
 - Initialize an empty association list at the very beginning
 - For (fac_mem n)
 - Check if n has a binding in the assoc list
 - If so, return the value in the binding
 - If not, call (fac n) and add (n, (fac n)) to the assoc list;
 - This is only possible because fac is a pure function

7

Control Flow: Iteration vs Recursion

```
□ Iteration/loop
x = 0; i = 1; j = 100;
while i < j do
   x = x + i*j; i = i + 1; j = j - 1
end while
return x
```

 \square becomes f(0,1,100), where f(x,i,j) is defined as if i < j then f(x+i*j, i+1, j-1) else x

9

Constructs with Side Effects in Racket

- □ I/O effects
 - (display "hello")
 - file operations
- ☐ Sequencing uses the special form begin
 - (begin (display "hello") (display "world"))
- □ Iteration
 - · do and foreach; see book
- ☐ Most of the time, we can program without using these constructs
 - after all, we have done a lot of programming in Racket already

Simulate "assignments" by Function Calls

☐ In general, you can get the effect of a series of assignments

> x = 0x = expr1x = expr2

from f3(f2(f1(0))), where each f expects the value of x as an argument, f1 returns expr1, and f2 returns expr2

Constructs with Side Effects in Racket

☐ set!: mutate the value associated with a name (define x 0)

(define (getCounter) x); a zero-argument function (define (inc) (set! x (+ x 1)))

(inc) (getCounter); returns 2

(getCounter); returns 3; note this loses referential transparency

- ☐ set-car!: mutate the car part of a list
- ☐ set-cdr!: mutate the cdr part of a list
 - (define abcde '(a b c d e))
 - (set-car! abcde `u) ; the value of abcde afterwards is (u b c d e)
 - (set-cdr! abcde '(d e)) ; the value of abcde afterwards is (u d e)

10

8

11