

# CMPSC 465 Assignment 04

Shi Qiu

TOTAL POINTS

**6.42 / 40**

## QUESTION 1

**P1** 15 pts

**1.1 P1.1** 4.8 / 6

- ✓ + **3 pts** proved/explained 1.a correctly.
- ✓ + **3 pts** proved/explained 1.b correctly.
- + **0 pts** No answer/Handwritten/More than 2 hours
- + **0.6 pts** 10% points for 6 points
- + **0 pts** Incorrect
- **1.2 Point adjustment**

**1.2 P1.2** 0.5 / 5

- + **5 pts** proved/ explained region correctly.
- + **3 pts** partially correct explanation of region.
- + **0 pts** No answer/Handwritten
- ✓ + **0.5 pts** 10% points
- + **0 pts** Incorrect

**1.3 P1.3** 0.32 / 4

- + **4 pts** proved that  $I^*$  is in the region  $s^*$
- + **2 pts** partially correct explanation of  $I^*$  is in the region  $s^*$ .
- ✓ + **0.4 pts** 10% points
- + **0 pts** No or wrong answer/Handwritten
- **0.08 Point adjustment**

## QUESTION 2

**P2** 15 pts

**2.1 P2.1** 0 / 3

- + **3 pts** Provided a correct explanation or proof.
- ✓ + **0.3 pts** I don't know how to answer this question.
- + **0 pts** Wrong Answer
- **0.6 Point adjustment**
- Late submission

**2.2 P2.2** 0 / 6

- + **6 pts** Provided a correct proof.
- + **3 pts** Provided an incomplete proof or explanation.
- ✓ + **0.6 pts** I don't know how to answer this question.
- + **0 pts** Wrong Answer
- **1.2 Point adjustment**
- Late submission

**2.3 P2.3** 0 / 6

- + **6 pts** Provided an algorithm in  $O(n^3 \log k)$
- + **4 pts** Provided an algorithm in  $O(n^{\log_2 7})$  ( $n^{\log_2 7}$  may come from the optimized matrix multiplication in the textbook)
- + **3 pts** Provided an algorithm in  $O(n^3 k)$
- + **3 pts** Provided an incomplete proof or explanation.
- ✓ + **0.6 pts** I don't know how to answer this question.
- + **0 pts** Wrong Answer
- **1.2 Point adjustment**
- Late submission

## QUESTION 3

**3 P3** 0.8 / 10

- + **3 pts** Described a correct way to find the uppermost and lowermost coordinates.
- + **4 pts** Described a correct way to find the leftmost and rightmost coordinates.
- + **3 pts** Provided a reasonable running time analysis for the algorithm, and it runs in  $O(\log(m+n))$  or  $O(\log(mn))$ .
- + **1 pts** I don't know how to answer this question.
- + **0 pts** No answer/Handwritten/More than 2 hours
- + **0.8 Point adjustment**
- Late submission

partial credit: +1, 20% off as late penalty

## Problem 2

Points:

a)

true:

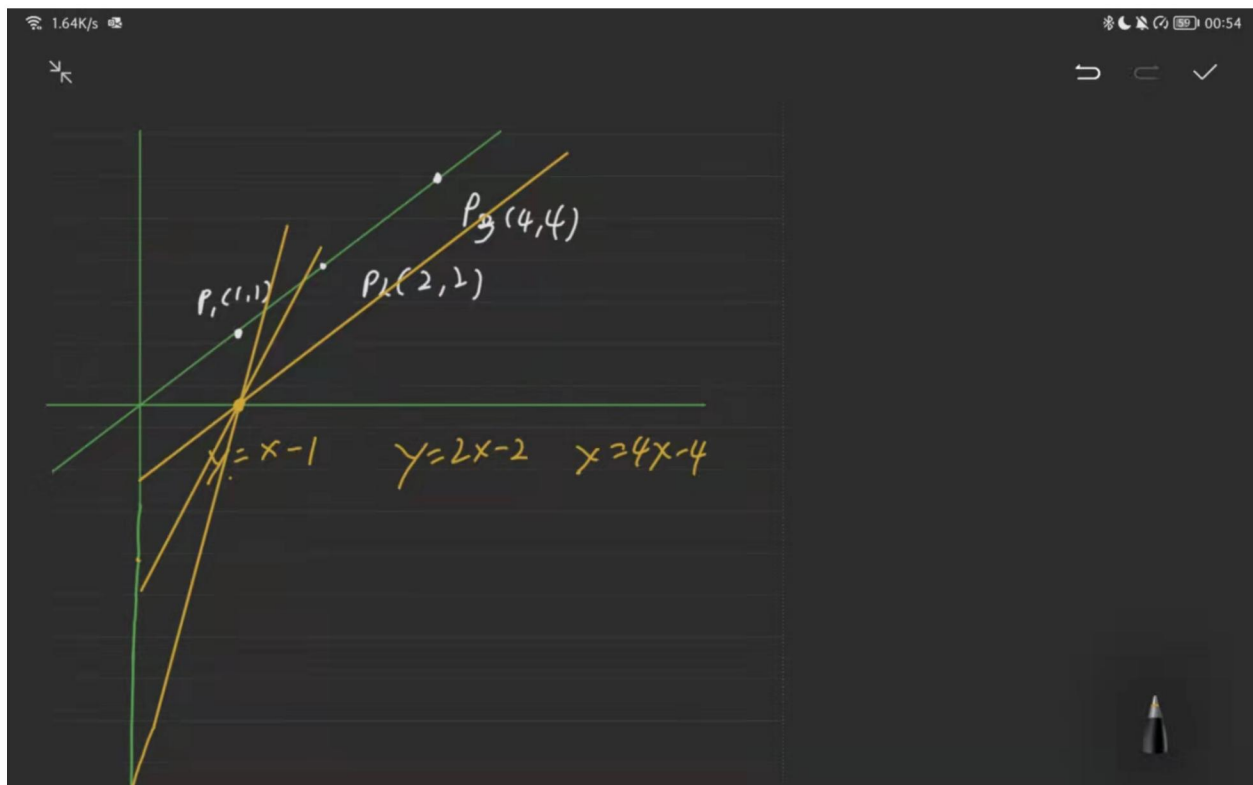
point  $p$  is on line  $l$  if and only if point  $l$  is on line  $p$ .

since  $p_1$  to  $p_n$  are on the common line, their line will pass a common point

For any point  $p$ , we have  $(p) = p$ . For any line  $l$ , we have  $(l) = l$ .

Point  $p$  is on line  $l$  if and only if point  $l$  is on line  $p$ .

Point  $p$  is above (resp. below) line  $l$  if and only if point  $l$  is above (resp. below) line  $p$ .



b)

visé versa, if all the lines go over a point, their point will be on the same line.

2.

I don't know how to answer this question

3.

I don't know how to answer this question

1.1 P1.1 4.8 / 6

- ✓ + 3 pts proved/explained 1.a correctly.
- ✓ + 3 pts proved/explained 1.b correctly.
  - + 0 pts No answer/Handwritten/More than 2 hours
  - + 0.6 pts 10% points for 6 points
  - + 0 pts Incorrect
- 1.2 Point adjustment

## Problem 2

Points:

a)

true:

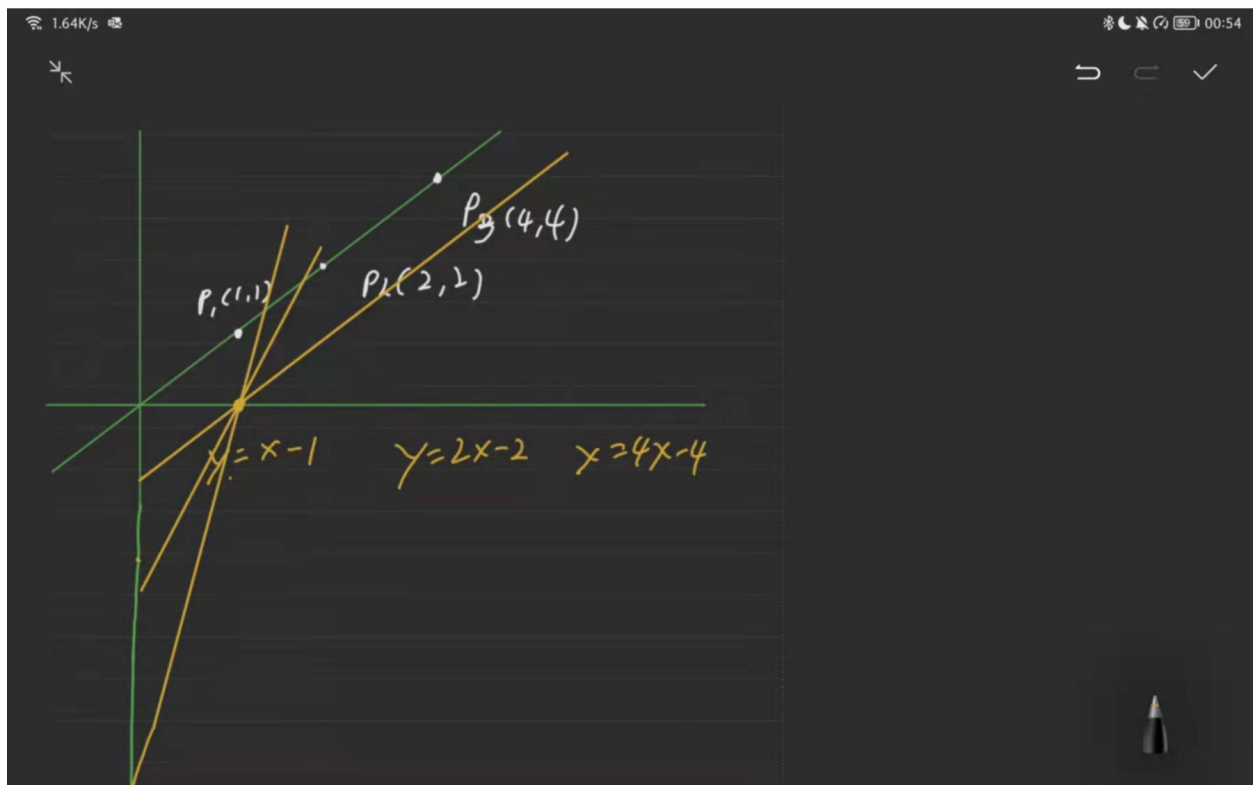
point  $p$  is on line  $l$  if and only if point  $l$  is on line  $p$ .

since  $p_1$  to  $p_n$  are on the common line, their line will pass a common point

For any point  $p$ , we have  $(p) = p$ . For any line  $l$ , we have  $(l) = l$ .

Point  $p$  is on line  $l$  if and only if point  $l$  is on line  $p$ .

Point  $p$  is above (resp. below) line  $l$  if and only if point  $l$  is above (resp. below) line  $p$ .



b)

visé versa, if all the lines go over a point, their point will be on the same line.

2.

I don't know how to answer this question

3.

I don't know how to answer this question

## 1.2 P1.2 0.5 / 5

- + **5 pts** proved/ explained region correctly.
- + **3 pts** partially correct explanation of region.
- + **0 pts** No answer/Handwritten
- ✓ + **0.5 pts** 10% points
- + **0 pts** Incorrect

## Problem 2

Points:

a)

true:

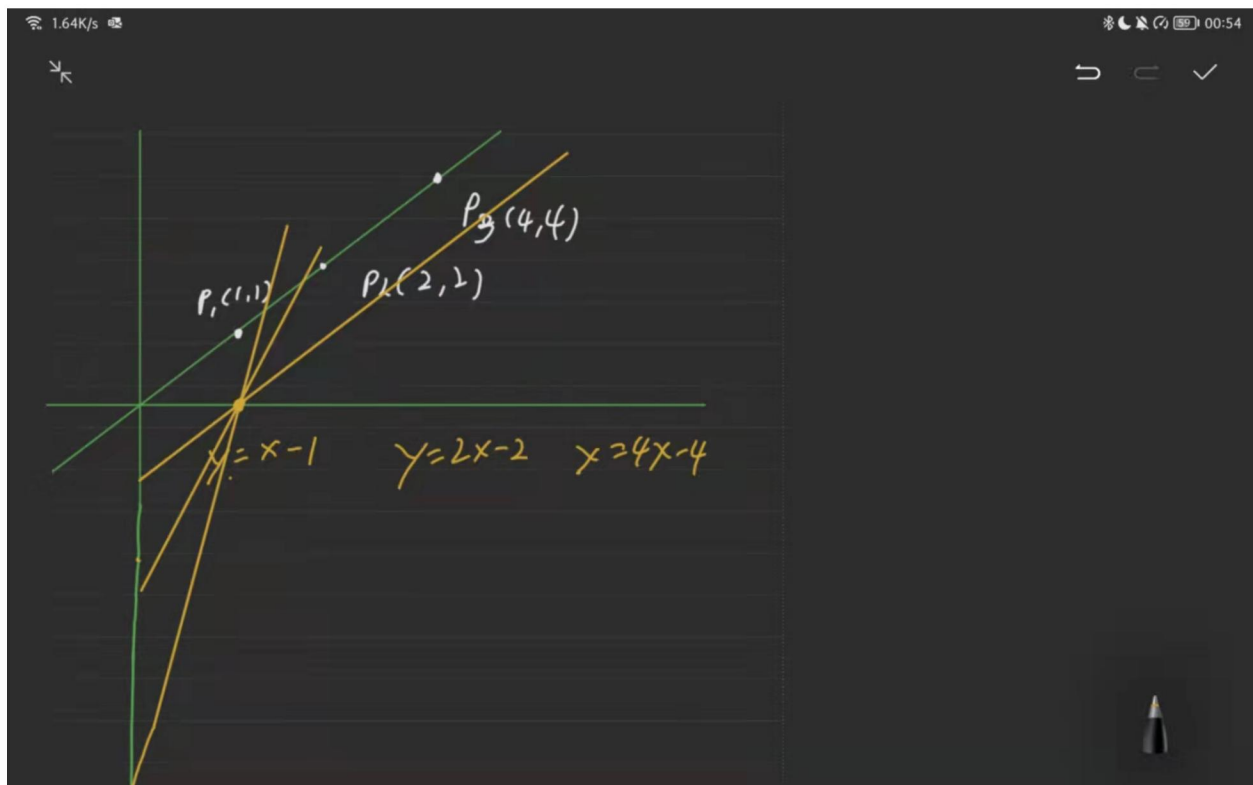
point  $p$  is on line  $l$  if and only if point  $l$  is on line  $p$ .

since  $p_1$  to  $p_n$  are on the common line, their line will pass a common point

For any point  $p$ , we have  $(p) = p$ . For any line  $l$ , we have  $(l) = l$ .

Point  $p$  is on line  $l$  if and only if point  $l$  is on line  $p$ .

Point  $p$  is above (resp. below) line  $l$  if and only if point  $l$  is above (resp. below) line  $p$ .



b)

visé versa, if all the lines go over a point, their point will be on the same line.

2.

I don't know how to answer this question

3.

I don't know how to answer this question

### 1.3 P1.3 0.32 / 4

+ 4 pts proved that  $\mathbb{S}^1 \times \mathbb{S}^1$  is in the region  $\mathbb{S}^2 \times \mathbb{S}^1$

+ 2 pts partially correct explanation of  $\mathbb{S}^1 \times \mathbb{S}^1$  is in the region  $\mathbb{S}^2 \times \mathbb{S}^1$ .

✓ + 0.4 pts 10% points

+ 0 pts No or wrong answer/Handwritten

- 0.08 Point adjustment



**Problem 3**

<b>Points:</b>
----------------

I don't know how to answer this question

2.1 P2.1 0 / 3

+ 3 pts Provided a correct explanation or proof.

✓ + 0.3 pts I don't know how to answer this question.

+ 0 pts Wrong Answer

- 0.6 Point adjustment

🗨 Late submission

**Problem 3**

<b>Points:</b>
----------------

I don't know how to answer this question

## 2.2 P2.2 0 / 6

+ 6 pts Provided a correct proof.

+ 3 pts Provided an incomplete proof or explanation.

✓ + 0.6 pts I don't know how to answer this question.

+ 0 pts Wrong Answer

- 1.2 Point adjustment

🗨 Late submission

**Problem 3**

<b>Points:</b>
----------------

I don't know how to answer this question

### 2.3 P2.3 0 / 6

- + 6 pts Provided an algorithm in  $O(n^3 \log k)$
- + 4 pts Provided an algorithm in  $O(n^{\log_2 7} k)$  ( $n^{\log_2 7}$  may come from the optimized matrix multiplication in the textbook)
- + 3 pts Provided an algorithm in  $O(n^3 k)$
- + 3 pts Provided an incomplete proof or explanation.
- ✓ + 0.6 pts I don't know how to answer this question.
- + 0 pts Wrong Answer
- 1.2 Point adjustment
- 🗨 Late submission

**Problem 4**

<b>Points:</b>
----------------

first set  $x_1$  be 0,  $y_1$  be 0,  $x_2$  be half of the  $n$  and  $y_2$  be  $m$

so we can check if it is true, and gradually minimize the area by test if half of the remaining area will return true.

Repeatedly divide the section that may contain the item in half until you have narrowed down the possible locations to just one.

### 3 P3 0.8 / 10

- + **3 pts** Described a correct way to find the uppermost and lowermost coordinates.
- + **4 pts** Described a correct way to find the leftmost and rightmost coordinates.
- + **3 pts** Provided a reasonable running time analysis for the algorithm, and it runs in  $O(\log(m+n))$  or  $O(\log(mn))$ .
- + **1 pts** I don't know how to answer this question.
- + **0 pts** No answer/Handwritten/More than 2 hours

#### + 0.8 Point adjustment

- 🗨 partial credit: +1, 20% off as late penalty