

## HW6 CS336

1. a.  $A = \text{Restaurant}$      $B = \text{Address}$      $C = \text{Cuisine}$   
 $D = \text{Rating}$      $E = \text{Rank}$      $F = \text{User-name}$   
 $G = \text{User-reviews}$      $H = \text{Date}$      $I = \text{Stars}$

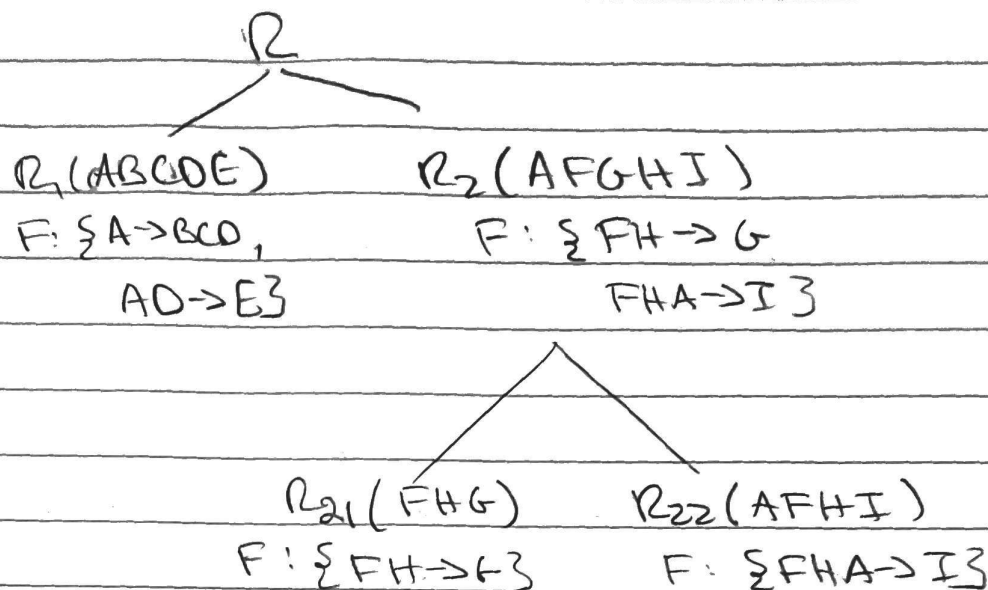
$$F = \{ A \rightarrow BCD, AD \rightarrow E, FH \rightarrow G, FHA \rightarrow I \}$$

- b.  $A^+ = \{ ABCDE \}$   
 $AD^+ = \{ ABCDE \}$   
 $FH^+ = \{ FHG \}$   
 $FHA^+ = \{ ABCDEFGHI \}$

FHA is a key for this scheme

- c.  $A \rightarrow BCD$   
 $AD \rightarrow ABCDE$   
 $FH \rightarrow G$   
 $FHA \rightarrow ABCDEFGHI$

d.



BCNF:  $R_1(ABCDE), R_{21}(FHG), R_{22}(AFHI)$

e.

	A	B	C	D	E	F	G	H	I
$R_1(ABCDE)$	$a_1$	$b_1$	$c_1$	$d_1$	$e_1$	$f_1$	$g_1$	$h_1$	$i_1$
$R_{21}(FHG)$	$a_2$	$b_2$	$c_2$	$d_2$	$e_2$	$f_2^F$	$g_2^G$	$h_2^H$	$i_2$
$R_{22}(AFHI)$	$a_3^A$	$b_3^B$	$c_3^C$	$d_3^D$	$e_3^E$	$f_3^F$	$g_3^G$	$h_3^H$	$i_3^I$

Tuple  $R_{22}(AFHI)$  has all of its placeholders replaced with actual values, therefore the BCNF is lossless

2.  $R = \{A, B, C, D, E\}$

$F = \{A \rightarrow B, B \rightarrow C, C \rightarrow D, CD \rightarrow E, D \rightarrow A\}$

Proc-processed:  $F = \{A \rightarrow BCDE, B \rightarrow COEA, C \rightarrow DAB, CD \rightarrow E, D \rightarrow AEC\}$

ACD  $F = \{A \rightarrow CD, C \rightarrow DA, D \rightarrow AC\}$

Yes, it is in BCNF. All keys are superkeys.

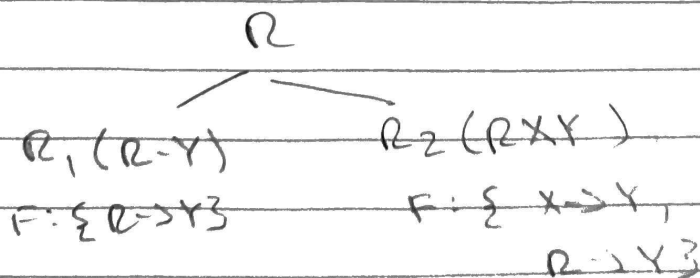
3. In general,  $N$  attributes with each being a candidate key means the number of possible superkeys are  $2^N - 1$ .

$R = \{A, B, C, D, E, F\}$  has 6 attributes.

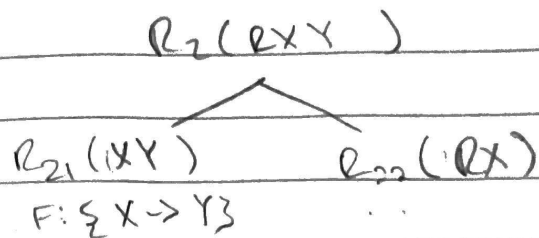
$$2^6 - 1 = 63.$$

There are 63 possible superkeys.

4.  $R = \{R, X, Y\}$



$X \rightarrow Y$  violates BCNF



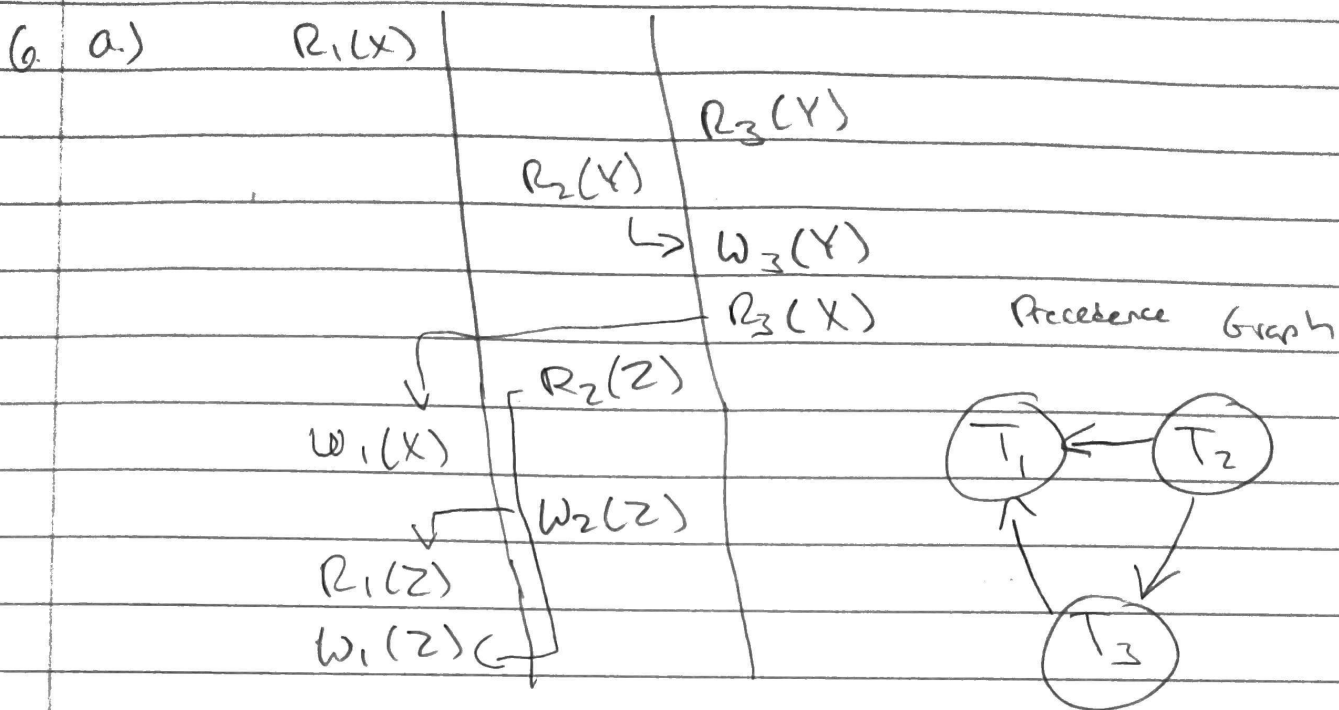
	R	X	Y
$R_1(RY)$	$r_1$	$x_1$	$y_1$
$R_{21}(XY)$	$r_2$	$x_2$	$y_2$
$R_{22}(RX)$	$r_3$	$x_3$	$y_3$

$R_{22}$  is tuple  $t$  that all placeholders are replaced with actual values, so lossless.

5 The consequences of deleting a beer from the table Beers will automatically delete all corresponding types in the Sells table.

Deleting a tuple from the Sells table is not much of an issue as corresponding data in other tables are not deleted.

Inserting a tuple into Sells can be an issue because we can only add tuples with pre-existing, corresponding data in Beers.



No Cycle between any vertices, so  
Conflict Serializable

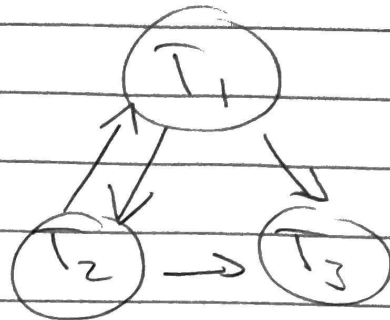
Serial Schedules:  $T_2, T_3, T_1$   
 $T_2, T_1, T_3$

b.

1

2

3

 $r_1(x)$  $r_2(y)$  $r_3(y)$  $w_3(y)$  $r_2(z)$  $w_1(x)$  $r_3(x)$  $r_1(z)$  $w_2(z)$  $w_1(z)$ 

Not Conflict

Serializable because

 $T_1$  &  $T_2$  are

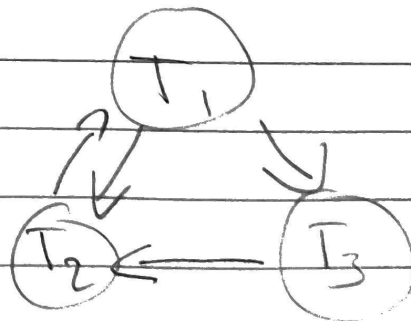
cyclic

c.

1

2

3

 $r_1(x)$  $w_1(x)$  $r_3(y)$  $r_1(z)$  $w_3(y)$  $r_2(y)$  $r_2(z)$  $r_3(x)$  $w_1(z)$  $w_2(z)$ 

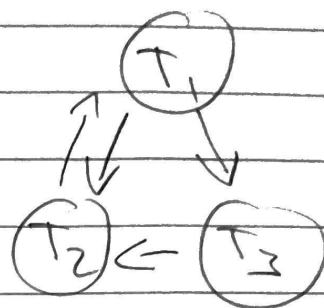
Not Conflict Serializable

because  $T_1$  &  $T_2$ 

are cyclic

d.

1	2	3
$w_1(x)$		
$r_1(z)$		$r_3(y)$
	$w_2(y)$	$w_3(y)$
	$r_2(y)$	
	$r_2(z)$	$r_3(x)$
	$w_2(z)$	
	$w_2(y)$	
$w_1(z)$		



$T_2$  and  $T_1$  are  
cyclic, so not  
Conflict serializable