Compute the Spatial Distance Histogram of a Set of 2D Points using Parallel Computing

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Abstract — as parallel computing has evolved, all areas of computing are faced with rebuilding to get faster processing speed. This article focuses on one important issue in scientific simulation data analysis: the spatial distance histogram (SDH). When calculating the spatial distance from the histogram when the calculation is very huge. Usually these calculations should be kept for a short period of time, in the event of failure, but also to start again. We propose a more efficient algorithm that uses parallel computing and prove that this algorithm has a good fault tolerance. The core of our algorithm is to decompose the space of a particle into a quadtree-based structure and then obtain the number of particles in each node of the tree to determine the spatial distribution of the particle. Our research shows how we can implement these algorithms in a graphics processing unit (GPU) to achieve a certain level of fault tolerance. The efficiency of our proposed algorithm will have a wide range of effects on the data produced by the actual simulation study

 $\label{local_equation} \textbf{Index Terms} \ -- \ \textbf{spatial distance histogram, parallel computing, 2D Points}$

I. INTRODUCTION

year through computer simulation experiments will produce a large amount of data. Despite the use of a database management system to process this data, data management software is not sufficient to handle such a large amount of data [1-3]. Traditional database management system is established for business applications, not suitable for managing big data. So we need parallel computing to redesign the data management system. These programs usually require parallelism methods that use large numbers of GPUs for computations to significantly improve operational efficiency [4]. In this way, these scientific big data can be effectively used for scientific simulation experiments.

In order to analyze the data of 2D points, traditionally, scientists have to do a lot of complex calculations to create a spatial distance histogram [5-7]. In general, the queries used in this analysis are based on the location of each particle: the function of all m-tuple subsets involving data is called the m-body correlation function. One such analysis query discussed in this article is the socalled Spatial Distance Histogram (SDH). SDH is a histogram of the distances between all pairs of particles in the system, which represents a continuous probability distribution of distances (called the radial distribution function (RDF).) This type of query is very important in an MS database as one of the basic building blocks to describe a series of key quantities (such as total pressure and energy) required by physical systems.

With the development of GPU-based parallel computing in recent years, GPU parallel computing has become a fantastic choice when the CPU cannot handle large amounts of social networking data quickly[8, 9]. Through the GPU's large-scale processing power, can quickly construct SDH relationship between notes. In this study, we use parallel computing to construct SDH relationship in 2D points.

II. METHOD

The SDH problem can be formally described as follows: given the coordinates of N particles and a user-defined distance w, we need to compute the number of particle-to-particle distances falling into a series of ranges (named buckets) of width w: [0, w), [w, 2w), ..., [(l-1)w, lw]. Essentially, the SDH provides an ordered list of non-negative integers H = (h0, h1, ..., hl-1), where each $hi(0 \le i < l)$ is the number of distances falling into the bucket [iw, (i+1)w). We also use H[i] to denote hi in this paper.

Clearly, the bucket width w is the only parameter of this type of problem.

To capture the variations of system states over time, there is a need to compute SDH for a large number of consecutive frames. We denote the count in bucket i at frame j as Hj [i].

In this project, we use a computer to generate a random number of points. Then we find the two points farthest away to determine the maximum boundary and use it as the root node of the quad-tree

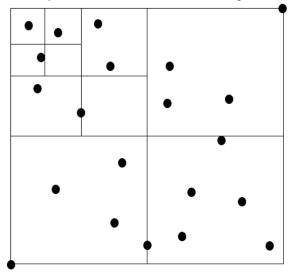


Fig 1. Data points distribution

Each node has 4 child nodes, down until it reaches a threshold, the last layer is the leaf node. The nodes above the leaf nodes need only store the number of internal points of the Cell (each square is a Cell). The leaf nodes not only store the number but also store the index of all the points in the cell, so as to directly calculate the distance between the points distance.

Down through the tree layer by layer until you find that level diagonal less than the width of the bucket to stop. This level is the work level (start level). The distances between points in each cell on this level will be less than the width of the bucket by adding all of them to the first histogram (histogram [0]). Then calculate the maximum distance and the minimum distance between the two cells with the cell of the same level, and if the maximum distance and the minimum distance are in the same bucket, then the mutual distance between the two cell middle points is in this bucket. If the maximum distance and

the minimum distance are not in the same bucket, compute their child node until it meets this condition or calculate the distance between the leaf nodes directly.

If you can not find the diagonal less than the width of the bucket level, the direct calculation of the distance between the leaves and the point.

III. RESULT

All the points we randomly generate falls between 0 and 23,000. First, we generate 10,000 points and build SDH with bucket width 500. The results for CPU (Figure 1) and GPU (Figure 2) are as below.

| | | t on SDH qu | | | | | | |
|------------|------|-------------|------|------|------|---------|--------|---------|
| | | distances | | | | | | |
| | | distances | | | | | | |
| | | distances | | | | | | |
| | | distances | | | | | | |
| | | | | | | | | |
| | | distances | | | | | | |
| | | distances | | | | | | |
| The number | | distances | fall | | | bucket: | 889374 | |
| The number | r of | distances | fall | | No.9 | bucket: | 974262 | |
| The number | | distances | fall | into | | bucket: | | 1054512 |
| | | distances | | | | | | |
| | | distances | | | | | | 1195985 |
| | | distances | | | | | | 1256512 |
| The number | | distances | | | | bucket: | | 1308940 |
| | | distances | | | | | | 1357247 |
| | | distances | | | | | | 1395996 |
| The number | | distances | fall | | | bucket: | | 1431645 |
| The number | r of | distances | fall | | No.1 | bucket: | | 1459501 |
| The number | r of | distances | fall | | | bucket: | | 1482999 |
| | | distances | | | | | | 1501710 |
| | | distances | | | | | | 1511927 |
| | | distances | | | | | | 1517963 |
| | | distances | | | | | | |
| | | | | | | | | 1511360 |
| The number | | distances | fall | | | bucket: | | 1499271 |
| The number | | | | | | | | 1482020 |

Fig 2. Time cost 4.267s for 10,000 points with bucket width 500 in CPU

| | | 344978 | 469453 | | |
|--|---------|--------|---------|---------|--|
| | 794132 | | 974262 | | |
| | 1195985 | | 1308940 | | |
| | | | | | |
| | 1517965 | | | | |
| | | | | 1374283 | |
| | | | 1196882 | | |
| | 1026491 | | | | |
| | | 601969 | | | |
| | | | | | |
| | | | | | |
| | | | | | |
| | | | | | |
| | | | | | |

Fig 3. Time cost 0.386s for 10,000 points with bucket width 500 in GPU

Second, we generate 100,000 points and build SDH with bucket width 1,000. The results for CPU (Figure 4) and GPU (Figure 5) are as below.

```
The time spent on SDM query: 249.630 (sec)

The number of distances fall into No.1 bucket: 28626442

The number of distances fall into No.2 bucket: 81637251

The number of distances fall into No.2 bucket: 181637251

The number of distances fall into No.5 bucket: 186233281

The number of distances fall into No.5 bucket: 186233281

The number of distances fall into No.5 bucket: 283717572

The number of distances fall into No.5 bucket: 233717572

The number of distances fall into No.7 bucket: 233717572

The number of distances fall into No.7 bucket: 27668701

The number of distances fall into No.9 bucket: 276188766

The number of distances fall into No.10 bucket: 309738996

The number of distances fall into No.10 bucket: 309736869

The number of distances fall into No.11 bucket: 300773689

The number of distances fall into No.12 bucket: 29678417

The number of distances fall into No.13 bucket: 29678417

The number of distances fall into No.14 bucket: 290709241

The number of distances fall into No.15 bucket: 278819768

The number of distances fall into No.18 bucket: 278819768

The number of distances fall into No.18 bucket: 278819768

The number of distances fall into No.18 bucket: 278819768

The number of distances fall into No.18 bucket: 278819768

The number of distances fall into No.18 bucket: 171635381

The number of distances fall into No.19 bucket: 171635381

The number of distances fall into No.19 bucket: 171635381

The number of distances fall into No.20 bucket: 171635381

The number of distances fall into No.21 bucket: 11129453

The number of distances fall into No.22 bucket: 11129453

The number of distances fall into No.23 bucket: 31011923

The number of distances fall into No.25 bucket: 11129453

The number of distances fall into No.25 bucket: 31011923

The number of distances fall into No.26 bucket: 31011923

The number of distances fall into No.27 bucket: 31011923

The number of distances fall into No.28 bucket: 4806075

The number of distances fall into No.28 bucket: 4806075

The number of distance
```

Fig 4. Time cost 249s for 100,000 points with bucket width 1,000 in CPU

| 72958 | 212698 | | 469453 | 585543 |
|---------|---------|--------|---------|---------|
| 695643 | 794132 | 889373 | 974262 | 1054512 |
| 1127704 | 1195985 | | 1308940 | |
| | 1431644 | | 1483001 | |
| | 1517965 | | | |
| | | | | 1374283 |
| | | | 1196882 | |
| | 1026491 | | | |
| | | 601969 | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |

Fig 5. Time cost 12.894s for 100,000 points with bucket width 1,000 in GPU

Third, we generate 100,000 points and build SDH with bucket width 500. The results for CPU (Figure 6) and GPU (Figure 7) are as below.

| | SDH query: 452.593(sec) | |
|--------------------|-----------------------------|--------------|
| The number of dist | ances fall into No.1 bucket | : 7286465 |
| The number of dist | ances fall into No.2 bucket | : 21339977 |
| The number of dist | ances fall into No.3 bucket | : 34600325 |
| The number of dist | ances fall into No.4 bucket | 47036926 |
| The number of dist | ances fall into No.5 bucket | : 58740267 |
| The number of dist | ances fall into No.6 bucket | : 69672872 |
| The number of dist | ances fall into No.7 bucket | 79873019 |
| The number of dist | ances fall into No.8 bucket | : 89360262 |
| The number of dist | ances fall into No.9 bucket | 98140915 |
| The number of dist | ances fall into No.10 bucke | t: 106175512 |
| The number of dist | ances fall into No.11 bucke | t: 113515718 |
| The number of dist | ances fall into No.12 bucke | t: 120201854 |
| The number of dist | ances fall into No.13 bucke | t: 126186358 |
| The number of dist | ances fall into No.14 bucke | t: 131480343 |
| The number of dist | ances fall into No.15 bucke | t: 136065561 |
| | ances fall into No.16 bucke | |
| The number of dist | ances fall into No.17 bucke | t: 143474421 |
| The number of dist | ances fall into No.18 bucke | t: 146251006 |
| The number of dist | ances fall into No.19 bucke | t: 148465625 |
| The number of dist | ances fall into No.20 bucke | t: 150073281 |
| The number of dist | ances fall into No.21 bucke | t: 151132839 |
| The number of dist | ances fall into No.22 bucke | t: 151630850 |
| The number of dist | ances fall into No.23 bucke | t: 151603751 |
| The number of dist | ances fall into No.24 bucke | t: 151109394 |
| The number of dist | ances fall into No.25 bucke | t: 150106703 |
| The number of dist | ances fall into No.26 bucke | t: 148571714 |
| The number of dist | ances fall into No.27 bucke | t: 146599109 |
| | ances fall into No.28 bucke | |
| The number of dist | ances fall into No.29 bucke | t: 141094534 |
| The number of dist | ances fall into No.30 bucke | t: 137725234 |
| The number of dist | ances fall into No.31 bucke | t: 133895236 |
| The number of dist | ances fall into No.32 bucke | t: 129651709 |

Fig 6. Time cost 452s for 100,000 points with bucket width 500 in CPU

| | 21339978 | | 47036929 | 58740276 |
|-----------|-----------|-----------|-----------|-----------|
| | | | | 106175534 |
| | 120201863 | | 131480343 | 136065577 |
| | 143474459 | | 148465659 | |
| 151132840 | | | | |
| | 146599098 | 144110130 | 141094545 | |
| | 129651676 | 125028450 | 119854480 | 114330019 |
| | 102497608 | 96049669 | 89325419 | 82309941 |
| | 67496725 | 59706485 | 51622927 | |
| 34814853 | | | 17485660 | 14026249 |
| | | 6873882 | | 3961925 |
| | | | | |
| | 136907 | | | |
| | | | | |

Fig 7. Time cost 10.969s for 100,000 points with bucket width 500 in GPU

Finally, we generate 100,000 points and build SDH with bucket width 5,000. The results for CPU (Figure 8) and GPU (Figure 9) are as below.

| | | ent on SDH q | | | | | | | |
|------|--------|---------------|-----------|---------|----------|----------|----------|----------|--|
| | | of distances | | | | | | | |
| | | of distances | | | | | | | |
| | mber o | | fall into | No.3 | bucket: | | 1712681 | | |
| e nu | mber o | f distances | fall into | No.4 | bucket: | | 1018716 | | |
| | mber c | f distances | | | | | 1185410 | | |
| e nu | mber c | f distances | fall into | No.6 | bucket: | | 462001 | | |
| e nu | mber o | f distances | fall into | | bucket: | | 4402 | | |
| | | mber of dis | | the en | | | 19995000 | | |
| | soluti | on $= 5000.0$ | 00 start | level | = 3 stop | levels : | | | |
| vel | 3 4 | 065.86 | Recursion | s (Pair | | 2016 | | Resolved | |
| vel | | | Recursion | s (Pair | | 32256 | | Resolved | |
| vel | 5 1 | 016.47 | Recursion | s (Pair | | | | Resolved | |
| | node d | listances re | | | | | | | |
| | | | | | | 779978 (| | | |

Fig 8. Time cost 1.634s for 100,000 points with bucket width 5,000 in CPU



Fig 9. Time cost 3.911s for 100,000 points with bucket width 5,000 in CPU

| Points; bucket | CPU | GPU |
|----------------|--------|---------|
| width | | |
| 10,000; 500 | 4.267s | 0.386s |
| 100,000; 1000 | 249s | 12.894s |
| 100,000; 500 | 452s | 10.969s |
| 10,000; 5,000 | 1.634s | 3.911s |

Table 1. Overall result

III. CONCLUSION

We found that for 10,000 dots, the GPU performed better when the width was small. When the width is large, the CPU performs better. For

100000 points, the GPU behaves better than the CPU, regardless of the width.

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