

Shannon's Expansion

Shannon's Expansion Theorem

- Any Boolean function $f(w_1, w_2, \dots, w_n)$ can be written in the form

$$f(w_1, w_2, \dots, w_n) = w_1' \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$

- Any of the n variables can be used.

Proof

It suffices to prove the theorem holds for all possible values of one of the variables.

Since this is a binary function each variable can only be 1 or 0.

Therefore we need only look at two cases.

Proof

Let $w_1 = 0$, then

$$\begin{aligned} f(w_1, w_2, \dots, w_n) &= 1 \cdot f(0, w_2, \dots, w_n) + 0 \cdot f(1, w_2, \dots, w_n) \\ &= f(0, w_2, \dots, w_n) \end{aligned}$$

Let $w_1 = 1$, then

$$\begin{aligned} f(w_1, w_2, \dots, w_n) &= 0 \cdot f(0, w_2, \dots, w_n) + 1 \cdot f(1, w_2, \dots, w_n) \\ &= f(1, w_2, \dots, w_n) \end{aligned}$$

The proof holds for any arbitrary x .

Example: 3 input majority circuit

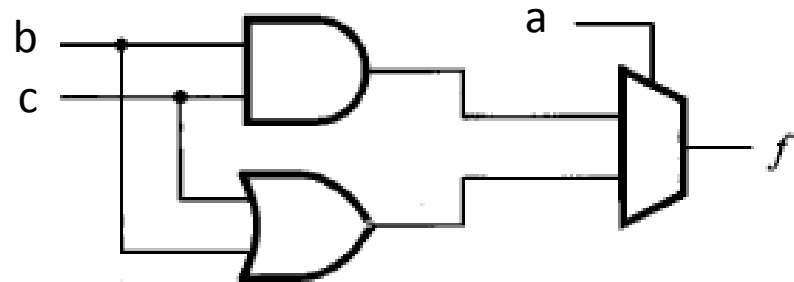
Given $f(a, b, c) = ab + ac + bc$

a	b	c	f
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

Grouping rows by 'a':

- Rows where $a = 0$ (rows 1-4) are grouped under a' .
- Rows where $a = 1$ (rows 5-8) are grouped under a .

a	f
0	bc
1	$b + c$



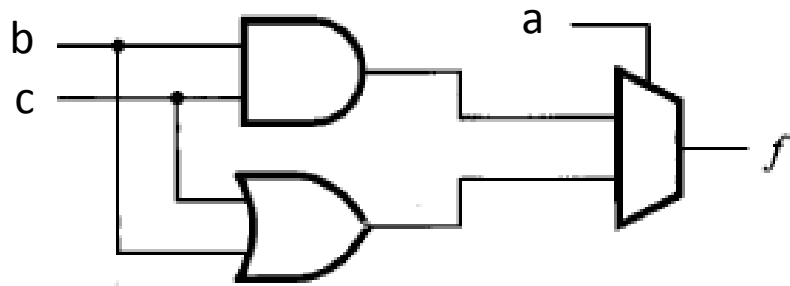
Shannon's Expansion Theorem

Given $f(a, b, c) = ab + ac + bc$

Expanding this function in terms of a gives

$$f = a'(0 \cdot b + 0 \cdot c + bc) + a(1 \cdot b + 1 \cdot c + bc)$$

$$f = a'(bc) + a(b + c)$$



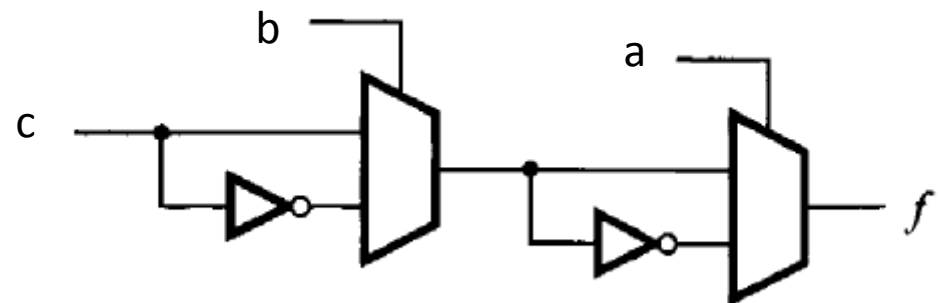
Example: 3 input XOR

a	b	c	f
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

a'

a

a	f
0	$b \text{ XOR } c$
1	$(b \text{ XOR } c)'$



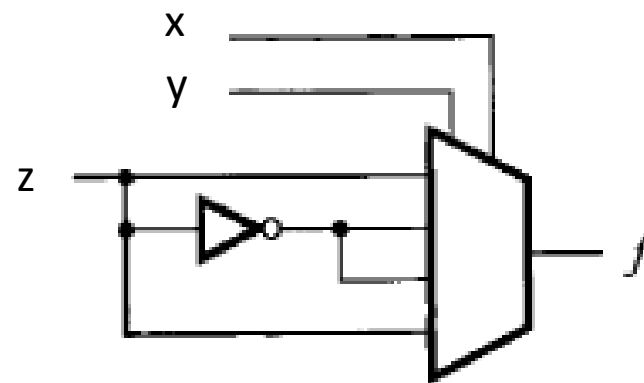
Example: 3 input XOR

x	y	z	f
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

Brackets on the right side of the table group the rows by the value of z :

- Rows 1 and 2 are grouped under z .
- Rows 3 and 4 are grouped under z' .
- Rows 5 and 6 are grouped under z' .
- Rows 7 and 8 are grouped under z .

xy	f
00	z
01	z'
10	z'
11	z



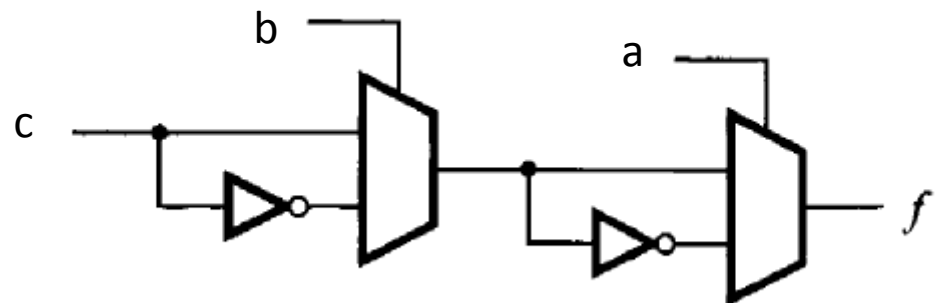
Shannon's Expansion Theorem

3 input XOR function

$$f = a \oplus b \oplus c$$

$$= a'(0 \oplus b \oplus c) + a(1 \oplus b \oplus c)$$

$$= a'(b \oplus c) + a(b \oplus c)'$$



Example

$$f = a'c + bc'$$

Expansion on a

$$f = a'(b + c) + a(bc')$$

Expansion on b

$$f = b'(a'c) + b(a' + c')$$

Expansion on c

$$f = c'(b) + c(a')$$

The point?

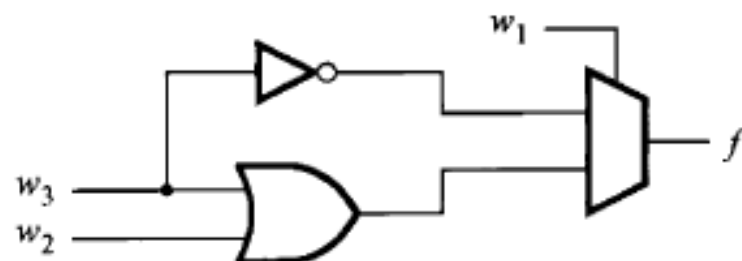
Some expansions are more efficient than others.

More problems

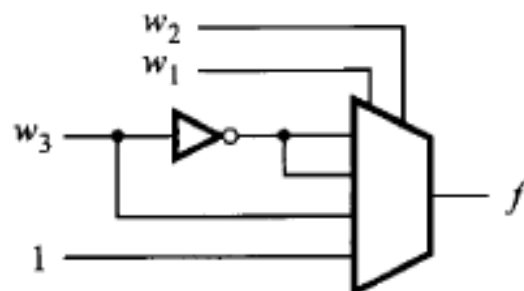
$$f = a'c' + ab + ac$$

Implement using a 2-to-1 multiplexer and any other necessary gates.

Implement the same function using a 4-to-1 multiplexer.



(a) Using a 2-to-1 multiplexer



(b) Using a 4-to-1 multiplexer