

# BAHUG 101 - Lecture 8

7th April 2016

# Outline of Today's Lecture

- ▶ Applicatives
- ▶ Parsing
- ▶ More monads!

# Motivation

# Employee Example

Consider the following **Employee** type:

```
type EName = String
```

```
data Employee = Employee { employee_name    :: EName  
                           , employee_phone  :: String }  
    deriving Show
```

# Employee Example

Consider the following **Employee** type:

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type EName = String
```

```
data Employee = Employee { employee_name    :: EName  
                           , employee_phone  :: String }  
    deriving Show
```

If we want to make an **Employee**, we of course can use the constructor.

```
Employee :: EName → String → Employee
```

# Maybe Employee

Can we convert a function of the following type

`EName → String → Employee`

to one that uses **Maybes**?

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```
maybeEmployee :: Maybe EName → Maybe String → Maybe Employee
maybeEmployee Nothing _          = Nothing
maybeEmployee _      Nothing     = Nothing
maybeEmployee (Just n) (Just p) = Just (Employee n p)
```

# Lists of Employees

Now, can we convert a function of the following type

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to one that uses `List`?

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`[EName] → [String] → [Employee]`

This could work in two ways

- ▶ Match each name with the associated phone in the second list.
- ▶ The Cartesian product of the name and phone lists.

# Using Reader with Employee?

Well what if we have the functions  $e \rightarrow \text{EName}$  and  $e \rightarrow \text{String}$  for some  $e$ . Can we make a function of

$\text{EName} \rightarrow \text{String} \rightarrow \text{Employee}$

into the following?

$(e \rightarrow \text{EName}) \rightarrow (e \rightarrow \text{String}) \rightarrow (e \rightarrow \text{Employee})$

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into the following?

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There is only one way to write this one.

# How do we generalize these examples?

What we would like to generalize our examples is a function like this

$$(a \rightarrow b \rightarrow c) \rightarrow (f\ a \rightarrow f\ b \rightarrow f\ c)$$

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$$(a \rightarrow b \rightarrow c) \rightarrow (f\ a \rightarrow f\ b \rightarrow f\ c)$$

Well that kinda looks like `fmap`.

$$\text{fmap} :: (a \rightarrow b) \rightarrow (f\ a \rightarrow f\ b)$$

But we need an `fmap` that can work for functions of two arguments

$$\text{fmap2} :: \text{Functor } f \Rightarrow (a \rightarrow b \rightarrow c) \rightarrow (f\ a \rightarrow f\ b \rightarrow f\ c)$$

`fmap2 h fa fb = undefined`

## But we can't use **Functor**!

Try hard as we might, however, **Functor** does not quite give us enough to implement **fmap2**. We have

```
h  :: a → (b → c)
```

```
fa :: f a
```

```
fb :: f b
```



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```
h  :: a → (b → c)
```

```
fa :: f a
```

```
fb :: f b
```

If we try to apply **fmap**, we get a function in our functor that we can't apply.

```
h      :: a → (b → c)
```

```
fmap h  :: f a → f (b → c)
```

```
fmap h fa :: f (b → c)
```

# Applicative

# Applicative Definition

Functors for which this sort of “contextual application” is possible are called **Applicative**.

```
class Functor f => Applicative f where  
  pure  :: a -> f a  
  (<*>) :: f (a -> b) -> f a -> f b — pronounced “ap”
```

On top of what a **Functor** offers, we have the following functions

- ▶ **pure**: promotes a value into the functor (ex.  $a \rightarrow [a]$ )
- ▶ **(<\*>)**: applies a function inside a functor

# Implementing `fmap2`

Now that we have `(<*>)`, we can implement `fmap2`, which in the standard library is actually called `liftA2`:

```
liftA2 :: Applicative f => (a -> b -> c) -> f a -> f b -> f c
liftA2 h fa fb = (h 'fmap' fa) <*> fb
```

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liftA2 h fa fb = (h 'fmap' fa) <*> fb
```

This pattern is common, so `(<$>)` as a synonym for `fmap`,

```
(<$>) :: Functor f => (a -> b) -> f a -> f b
(<$>) = fmap
```

so that we can write

```
liftA2 h fa fb = h <$> fa <*> fb
```

## What about `liftA3`?

Defining `liftA3` is as simple as it was to define `liftA2`.

```
liftA3 :: Applicative f => (a -> b -> c -> d) ->  
      f a -> f b -> f c -> f d  
liftA3 h fa fb fc = h <$> fa <*> fb <*> fc
```

We can use this pattern to create any `liftA#` without a new typeclass.

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```

We can use this pattern to create any `liftA#` without a new typeclass.

We can also use **Applicative** on regular values since we can put them in a **Functor** by using `pure`.

```
liftX :: Applicative f => (a -> b -> c -> d) ->
      f a -> b -> f c -> f d
liftX h fa b fc = h <$> fa <*> pure b <*> fc
```

## Defining `Applicative` for `Maybe`

```
instance Applicative Maybe where
  pure          = Just
  Nothing <*> _  = Nothing
  _ <*> Nothing = Nothing
  Just f <*> Just x = Just (f x)
```



## Applicative Maybe examples

```
m_name1, m_name2 :: Maybe EName  
m_name1 = Nothing  
m_name2 = Just "Brent"
```

```
m_phone1, m_phone2 :: Maybe String  
m_phone1 = Nothing  
m_phone2 = Just "555-1234"
```

```
ex01, ex02, ex03, ex04 :: Maybe Employee  
ex01 = Employee <$> m_name1 <*> m_phone1  
ex02 = Employee <$> m_name1 <*> m_phone2  
ex03 = Employee <$> m_name2 <*> m_phone1  
ex04 = Employee <$> m_name2 <*> m_phone2
```

# Parsing

# Parser Overview

A *parser* is an algorithm which takes unstructured data and produces structured data.

Examples include

- ▶ json parsers.
- ▶ the **ghci** REPL.
- ▶ any interpreter.

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For this lecture, we will be using the parsing package **Attoparsec**. It allows us to define simple parsers like so.

```
word :: Parser ByteString  
word = takeWhile $ inClass "a-zA-Z"
```

# Parsers for words

Now, let's write a parser for names. Names start with a capital letter and then contain lower case letters.

```
upper :: Parser Word8
upper = satisfy $ inClass "A-Z"

lword :: Parser ByteString
lword = takeWhile (inClass "a-z")
```

## Parsers for names

We can use these two primitive parsers, with `ByteString.cons`, to make the `name` parser. This is possible using the `Parser`'s `Applicative` definition.

```
name :: Parser ByteString  
name = BS.cons <$> upper <*> lword
```

Or, alternatively, we can *lift* the `cons` function into the `Parser` using `liftA2`:

```
name' :: Parser ByteString  
name' = liftA2 BS.cons upper lword
```

## Parsing full names

Suppose we want to parse full names (ie, first and last). Instead of returning a **ByteString**, we want to structure the output data. But first we need a way to skip spaces.

```
skipSpace :: Parser () — Don't care about captures characters.  
skipSpace = skipWhile isSpace_w8
```

Combining **skipSpace** with (**\*>**) allows us to ignore the parser's output. Now, let's write the full name parser:

```
firstLast :: Parser (ByteString, ByteString)  
firstLast = (,) <$> name <*> (skipSpace *> name)
```

# Full names with potential middle names

We want to target the following data type:

```
data Name = TwoName ByteString ByteString  
          | ThreeName ByteString ByteString ByteString  
          deriving Show
```

Unfortunately, there is no way to do this using applicative functors as our combinator because we can't make decisions with **Applicative**. We will have to use something stronger than applicative.



# Monads to the Rescue!

Recall that the **Monad** type class exposes two function:

```
class Monad m where
  return :: a → m a
  (>=)    :: m a → (a → m b) → m b
```

In general, monads can be used to sequence actions and make choices on prior results.

# First and last name parser

Let's write a parser for first and last name using monads instead of applicative functors like we did above.

```
firstLast' :: Parser (ByteString, ByteString)
firstLast' = do
  fname <- name
  lname <- skipSpace *> name
  return (fname, lname)
```

When possible, it is good practice to use **Applicative** instead of **Monad**.

# Full name parser

The full name parser can be written this way.

```
fullName :: Parser Name
fullName = do
  n1 <- name
  n2 <- skipSpace *> name
  mn <- skipSpace *> optional name
  case mn of
    Just n3 → return $ ThreeName n1 n2 n3
    Nothing → return $ TwoName n1 n2
```

Note that we used the `optional :: Parser a → Parser (Maybe a)` function above. This is a function defined by `Attoparsec` that allows a parser to fail without terminating the entire computation.

# Parsing an ambiguous strings

Consider the following string.

`"Haskell Brooks Curry Simon Peyton Jones".`

# Parsing an ambiguous strings

Consider the following string.

`"Haskell Brooks Curry Simon Peyton Jones".`

Should this string be parsed as the list

```
[TwoName "Haskell" "Brooks", TwoName "Curry" "Simon",  
 TwoName "Peyton" "Jones"]
```

or the list

```
[ThreeName "Haskell" "Brooks" "Curry",  
 ThreeName "Simon" "Peyton" "Jones"]
```

## Updated string

Since the prior string was ambiguous, we will parse the following.

**"[true, true] Haskell Brooks Curry Simon Peyton Jones"**

This signifies that there are two names in the sequence of words and both of them have middle names.

# Creating a bool list parser

We can create the bool list parser as follows.

```
bool :: Parser Bool
bool = do
  s <- word
  case s of
    "true"  → return True
    "false" → return False
    _       → fail $ show s ++ " is not a bool"
```

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list :: Parser a → Parser [a]
list p = char '[' *> sepBy p comma <*> char ']'
  where comma = skipSpace *> char ',' <*> skipSpace
```



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list :: Parser a → Parser [a]
list p = char '[' *> sepBy p comma <*> char ']'
  where comma = skipSpace *> char ',' <*> skipSpace

boolList :: Parser [Bool]
boolList = list bool
```

Note that the **sepBy** function creates a parser for a list of values that are separated by some other parser. In this case, the parser that separates the list elements is a comma surrounded by

# Complete full name parser

We can now use this list of **Bools** to figure out how to parse the names.

```
names :: Parser [Name]
names = boolList >=> mapM bToP
  where
    bToP :: Bool → Parser Name
    bToP True  = ThreeName <$> sn <*> sn <*> sn
    bToP False = TwoName  <$> sn <*> sn

    sn :: Parser ByteString
    sn = skipSpace *> name
```

This can be called using

```
result :: [Name]
result = fromRight $ parseOnly names
      "[true, true] Haskell Brooks Curry Simon Peyton Jones"
```