

P2: Exec and low-level I/O

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Pitfall



Anything might go wrong?

```
pid_t pid = fork();
if (pid < 0) {
    perror("fork()"); exit(1); // exit if fork() fails
} else if (pid == 0) {
    child_tasks();
} else {
    parent_tasks();
}
more_parent_tasks();</pre>
```

Pitfall



Anything might go wrong?

```
pid_t pid = fork();
if (pid < 0) {
   perror("fork()"); exit(1); // exit if fork() fails
} else if (pid == 0) {
   child_tasks();
   exit(0);
            // terminate the child process
} else {
   parent_tasks();
more_parent_tasks();
```

Process upgrades



- Usually....
 - A fresh clone wants to run different code
- This is done by
 - Loading another executable into the process address space
 - [picked up from the file system of course]
- Note
 - Opened files are NOT AFFECTED by the upgrade operation

The exec family



- The act of 'upgrading' is done by the child with a system call
 - Many variants. "man -S3 execl" for all details

- The path to the executable to load inside our own address space
- A list of arguments to be passed to the new executable
- A final NULL pointer to give the "end of argument list"
- If successful, execl () does not return! Started a new process





We will turn the child process into the following exectuable

```
#include <stdio.h>
#include <stdlib.h>

int main(int argc,char* argv[]) {
   int i, sum=0;
   for(i=1;i<argc;i++)
      sum += atoi(argv[i]);
   printf("sum is: %d\n",sum);
   return 0;
}</pre>
```

This is a simple "adder" program that computes the sum of its integer arguments



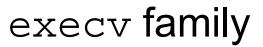
Parent Program

```
int main() { // complete code is in demo repo
  char *cmd1 = "./adder", *cmd2 = "expr";
 pid_t child = fork();
  if (child == 0) {
    printf("In child!\n");
    execl(cmd1,cmd1,"1","2","3","10",NULL);
    printf("Oops.... something went really wrong!\n");
    perror(cmd1);
    return -1;
  } else {
    printf("In parent!\n");
    execl(cmd2,cmd2,"100","+","300",NULL);
    printf("Oops.... something went really wrong!\n");
    perror(cmd2);
    return -1;
```





- Specify a path, like /bin/ls
- Specify a file, and the system searches in directories listed in PATH
 - echo \$PATH in bash to see directories separated by ':'





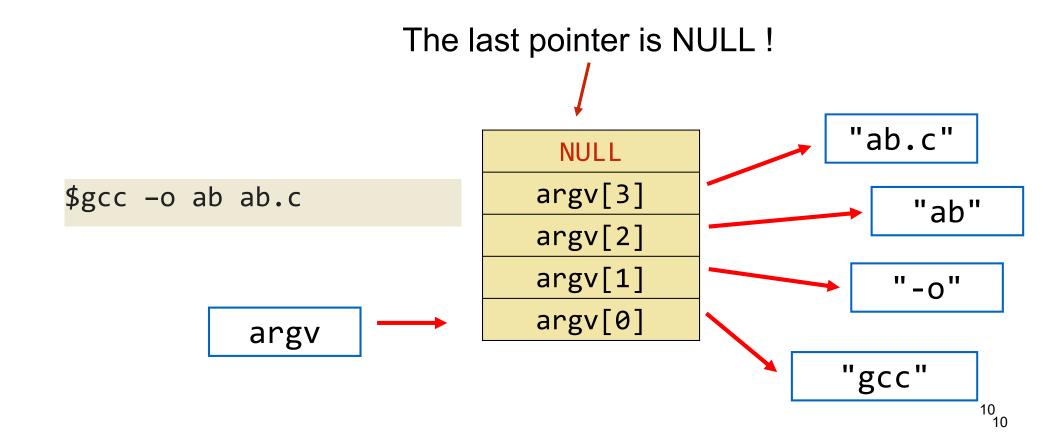
```
// If the number of arguments is unknown at compile time
#include <unistd.h>
int execv(const char *path, char *const argv[]);
int execvp(const char *file, char *const argv[]);
```

- The arguments in execl in are placed in an array
 - argv is the argv you see in the main function!
- execv needs a path while execvp can search file in PATH
- Start a new process if successful. Similar to execl

argv to execv and execvp



- Note the NULL pointer at the end
- Why?



Question



What might go wrong?

```
pid_t pid = fork();
if (pid < 0) {
   perror("fork()"); exit(1); // exit if fork() fails
else if (pid == 0) {
   // in child process
   execlp("genie", "genie", "clean the house", NULL);
// in parent process
online_shopping();
```

File APIs



- Remember the (C standard library) IO APIs
 - The "f" family (fopen, fclose, fread, fgetc, fscanf, fprintf,...)
 - All these use a FILE* abstraction to represent a file
 - Additional features: user-space buffering, line-ending translation, formatted I/O, etc.
- UNIX has lower-level APIs for file handling
 - Directly mapped to system calls
 - open, close, read, ...
 - Use file descriptors [which are just integers]
 - Deal with bytes only





• Read the man pages (man –s2 ...) for more functions

```
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
#include <unistd.h>
int open(const char *path, int oflag);
int open(const char *path, int oflag, int mode);
int close(int fd);
ssize t read(int fd, void *buf, size_t nbyte);
ssize t write(int fd, const void *buf, size_t nbyte);
off t lseek(int fd, off t offset, int whence);
```

Open a file



```
#include <fcntl.h>
#include <unistd.h>
int open(const char *path, int oflag);
```

- Parameters
 - path: the path to the file to be opened/created
 - oflag: read, write, or read and write, and more (on the next slide)
- The function returns a file descriptor, a small, nonnegative integer
 - Return -1 on error





- Must include one of the following:
 O_RDONLY (read only), O_WRONLY (write only), or O_RDWR (read and write)
- And or-ed (|) with many optional flags, for example,
 - O_TRUNC: Truncate the file (remove existing contents) if opening a file for write
 - O CREAT: Create a file if it does not exist.

Example:

```
// remember open() returns -1 on error
fd1 = open("a.txt", O_RDONLY); // open for read
fd1 = open("a.txt", O_RDWR); // open for read and write
fd1 = open("a.txt", O_RDWR|O_TRUNC); // read, write, truncate the file
```





```
// a mode must be provided if O_CREAT or O_TMPFILE is set
int open(const char *path, int oflag, int mode);
mode: specify permissions when a new, or temporary, file is created.
```

```
open("b.txt", O_WRONLY|O_TRUNC|O_CREAT, 0600);

// open b.txt for write. If the file exists, clear (truncate) the contents.
```

// if the file does not exist, create one, and set the permission so that the owner of the file can read and write, but other people cannot.



File descriptor vs stream

```
#include <stdio.h>
int fileno(FILE *stream);
// returns a file descriptor for a stream
```

| FD | FILE * |
|----|--------|
| 0 | stdin |
| 1 | stdout |
| 2 | stderr |



File descriptors after fork and exec

Opened files are NOT AFFECTED by the upgrade operation

```
pid_t pid = fork();
assert(pid >= 0);
if (pid == 0) {
   // Child process can access FDs 0, 1, and 2
   // if execl() is successful, gcc can access FDs 0, 1, and 2
   execlp("gcc", "gcc", "a.c", NULL);
   // If control gets here, execlp() failed.
   // Remember to terminate the child process!
   return 1;
```