Compilers Principles Lab5 Report Language Feature Extension

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Outline

- Type System Extension
 - Pointers and Arrays
 - Structures
 - Initialization
- Operator Extension
 - Context Free Grammar of Expressions
 - Operations on Structures
- Classes and Templates
 - Non-Static Member Functions
 - Operator Overriding
 - Class Templates



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Goals

```
Pointers 
int *ptr;
```

- Function int func(int);
- Arraysint array [42];

```
• Structures
    struct S {...};
    struct S s;
    struct S {...} s;
    struct {...} anonymous;
```

Declaration Grammar

Declaration grammar

 $declaration \rightarrow type-specifier declarator$;

Example

Variable declaration	Type specifier	Declarator
<pre>int *ptr;</pre>	int	*ptr
int array [42];	int	array [42]

Declaration Grammar

Declaration grammar

 $declaration \rightarrow type$ -specifier declarator;

For type specifiers, cminuf supports type int, float and void.

Type specifiers

type-specifier \rightarrow int |float|void

Declaration Grammar

Declaration grammar

 $declaration \rightarrow type$ -specifier declarator;

For declarators, cminuf supports pointer declarators, function declarators and array declarators.

- Pointer declarators
 - *ptr in int *ptr;
- Function declarators

Array declarators

```
array [42] in int array [42];
```



Crux: declarators can get mixed up

Trial 1: a naive solution

Trial 1: a naive solution

Example

Problem: ambiguity

Trial 2: eliminating ambiguity

Preced	Operation	
0	Parenthesis ()	
1	Call ()	
1	subscript []	
2	Dereference *	

Parentheses and identifiers

factor
$$\rightarrow$$
 ID \mid (decl)

• Function calls and array subscripting

$$\begin{array}{ccc} \mathsf{decl}\text{-}1 & \to & \mathsf{decl}\text{-}1(\mathsf{params}) \\ & | & \mathsf{decl}\text{-}1[\mathsf{int}] \\ & | & \mathsf{factor} \end{array}$$

Dereference

$$\det$$
 \rightarrow * \det \det \det \det \det \det \det \det

Trial 2: eliminating ambiguity

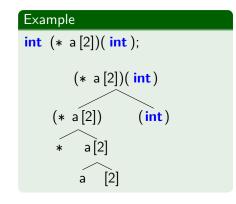
Parentheses and identifiers

$$\begin{array}{ccc} \mathsf{factor} & \to & \mathbf{ID} \\ & | & (\mathsf{decl}\) \end{array}$$

Function calls and array subscripting

$$\begin{array}{ccc} \mathsf{decl}\text{-}1 & \to & \mathsf{decl}\text{-}1(\mathsf{params}) \\ & | & \mathsf{decl}\text{-}1[\mathsf{int}] \\ & | & \mathsf{factor} \end{array}$$

Dereference



Goals

- Structure definition struct S {...};
- Variable definition struct S s;
- Structure and variable definition struct S {...} s;
- Anonymous structure struct {...} anonymous;

Extension of the declaration grammar

Declaration grammar

 $declaration \rightarrow type-specifier declarator$;

Declaration	Type specifier	Declarator
struct S {};	struct S {}	Ø
struct S s;	struct S	S
struct S {} s;	struct S {}	S
<pre>struct {} anonymous;</pre>	struct {}	anonymous

Extension of the declaration grammar

```
Declaration grammar
              declaration \rightarrow type-specifier declarator;
           type-specifier \rightarrow struct-definition
                                     struct ID
           struct-definition \rightarrow struct ID{definitions}
                                    struct{definitions}
           declarator
```

Self reference

Example

A naive implementation of the node of a forwarding linked list

```
struct list_node{
   int elem;
   struct list_node* next;
};
```

The identifier **struct** list_node is visible within the definition of the structure, although **struct** list_node is an incomplete type.

Self reference

Example

A naive implementation of the node of a forwarding linked list

```
struct list_node{
    int elem;
    struct list_node* next;
};
```

To render the identifier **struct** list_node available, construct an empty structure before parsing the body of the structure definition, and complete the type after the definition parsing is done.

Copy construction in declaration

Declaration grammar

 $declaration \rightarrow type$ -specifier declarator = expression

A declaration with initialization is transformed into a declaration with no initialization value and an Assignment.

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Supported operators

- Arithmatic operators
 +,-,*,/
- Relational operators>,<,>=,<=,==,!=
- Assignment operator
- Array subscript array [index]

- Pointer dereference
 - * ptr
- Address of & Ivalue
- Member access var.member
- Function call callable (params)

Precedence and CFG

Parenthesis ()

```
\begin{array}{cccc} \mathsf{expr-0} & \to & \textbf{ID} \\ & | & \textbf{Integer-literal} \\ & | & \textbf{Float-literal} \\ & | & (\mathsf{expr}\ ) \end{array}
```

Function call, array subscript, member access

Precedence and CFG

② Dereference, address of

expr-2
$$\rightarrow$$
 * expr-2
 | & expr-2
 | expr-1

Multiplication, division

expr-3
$$\rightarrow$$
 expr-3 **MulOp** expr-2 \mid expr-2

Addition subtraction

expr-4
$$\rightarrow$$
 expr-4 **AddOp** expr-3 \mid expr-3



Precedence and CFG

Relational operations

expr-5
$$\rightarrow$$
 expr-5 **RelOp** expr-4 \mid expr-4

Assignment

$$expr \rightarrow expr-5 = expr$$

| $expr-5$

Note

The assignment operator is right associated, so the production is $\exp r \rightarrow \exp r-5 = \exp r$ instead of $\exp r \rightarrow \exp r = \exp r-5$.

Problems with operations on structures

- Assignment to structuresstruct S s = t;
- Structures as parameters void func(struct S){...};
- Structures as return values struct S func (){...};

Assignment to structures

Source code

```
struct S
{
    int member1;
    float member2;
};
```

Transformed code

struct S

```
{
    int member1;
    float member2;
};
s.member1 = t.member1;
s.member2 = t.member2;
```

Structures as parameters

Source code

```
void func(struct S s)
{...}

int main()
{
    func(t);
}
```

Transformed code

```
void func(struct S *ptr)
{
    struct S s;
    s = *ptr;
}
int main()
{
    func(&t);
}
```

Structures as return values

Source code

Transformed code

```
void func(struct S* ret_ptr)
    *ret ptr = ret value;
int
    main()
    struct S ret value;
    func(&ret_value);
```

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Functions declared within a structure

Transformation of member functions

Source code

```
struct S
    void func(){}
int main()
    struct S s:
    s.func();
```

Transformed code

```
struct S
void S.func(struct S* this)
int main()
    struct S s;
    S.func(&s);
```

Override an operator

```
struct S {
    struct S operator+(struct S rhs)
    { . . . }
int main()
    struct S s:
    struct S t:
    s + t; // S::operator+ is called
```

Transform the overriding function

Source code

```
struct S {
    S operator + (S)
    \{\ldots\}
int main()
    struct S s:
    struct S t:
    s + t:
```

Transformed code

```
void S.operator+(
    struct S* ret ptr,
    struct S* this
    struct S* rhs ptr){
    struct S rhs;
    rhs = *rhs_ptr;
    *ret_ptr = ret_value;
int main(){
    struct S s: struct S t:
    struct S ret_value;
    S.operator+(&ret_value, &s, &t);
```

Some details about my implementation

Function tables for operators

The supported operators are arithmatic operators +,-,*,/. A function table for each of these operators is maintained, and is looked up every time the operation is performed.

Not operator Overloading

Function overloading is not implemented. Once the operator is overriden, the overriding function will be invoked every time the operation is performed.

Generic programing

Example

```
template < typename T>
struct stack{
   typename T top();
   void push(typename T);
};

int main()
{
   stack < int > a;
   stack < string > b;
   a.push(42);
   b.push("42");
}
```

Note

Due to the limitation of CFG, in the specification of a class template, the keyword *typename* is required before a template parameter.



CFG for class templates

Translation process of a class template

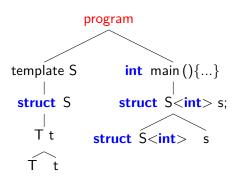
```
template < typename T >
struct S{
    typename T t;
};
int main(){
    struct S < int > s;
}
```

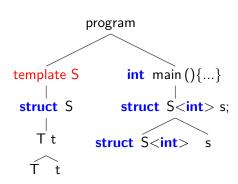
```
program

template S int main(){...}

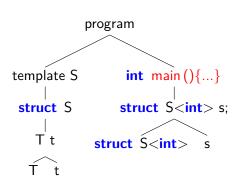
struct S struct S<int> s;

T t struct S<int> s
```

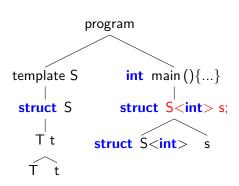




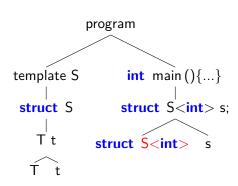
 Record the pointer of node template S in AST, do not analyse subtrees.



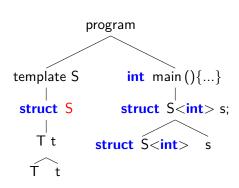
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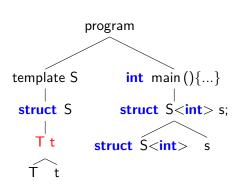
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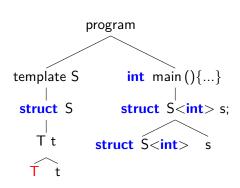
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- Map template T to int, go analyse class template



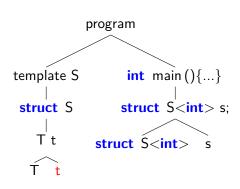
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- Map template T to int, go analyse class template
- Construct structure S<int>



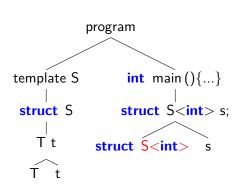
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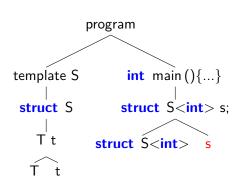
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- Construct structure S<int>
- Use mapping to replace T with int



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- Map template T to int, go analyse class template
- Construct structure S<int>
- Use mapping to replace T with int
- S<int> now is a complete type



- Record the pointer of node template S in AST, do not analyse subtrees.
- Map template T to int, go analyse class template
- Construct structure S<int>
- Use mapping to replace T with int
- S<int> now is a complete type
- Declare a variable of type S<int>



Non-Static Member Function Operator Overriding Class Templates

Thank You.