

README

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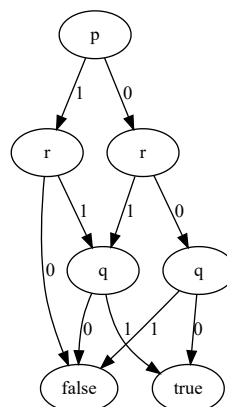
1 Introduction

This project is intended to parse the input propositional logic formula and generate the corresponding Reduced Ordered Binary Decision Diagram(ROBDD).

For instance, the input formula $(p \rightarrow r) \wedge (q \leftrightarrow (r \vee p))$ would lead to the following output

```
1  (p → r) & (q ↔ (r ∨ p))
2  digraph{
3      0 [label="false"]
4      1 [label="true"]
5      2 [label="p"]
6      3 [label="r"]
7      4 [label="r"]
8      2 → 3 [label="0"]
9      2 → 4 [label="1"]
10     5 [label="q"]
11     6 [label="q"]
12     3 → 5 [label="0"]
13     3 → 6 [label="1"]
14     5 → 1 [label="0"]
15     5 → 0 [label="1"]
16     6 → 0 [label="0"]
17     6 → 1 [label="1"]
18     4 → 0 [label="0"]
19     4 → 6 [label="1"]
20 }
```

which is [DOT](#) code. To visualize the diagram, one choice is to paste the generated code to [Viz.js](#) ([viz.js.com](#)). The example show above would generate the following graph



1.1 The Formula Grammar

1.1.1 Operators

The supported operations (in descending priority order):

1. `!` Not
2. `&` And
3. `|` Or
4. `->` Implication, `<->` Equivalence

Parentheses `(,)` can be used to alter the priority of the subexpressions.

1.1.2 Variables and Constants

The identifier of the variables should **consist of case-sensitive alphabetic characters and digits** (`[a-zA-Z0-9]+`).

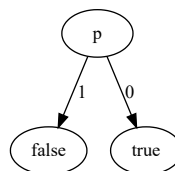
2 names are reserved: `T` for true and `F` for false.

For instance:

1. `T -> F`



2. `p -> F`



2 Build

The project is written completely in [Rust Programming Language](#). To build the project from source, follow the instructions on [Install Rust](#) to get the Rust compilation toolchain. The installation is supposed to be easy. If you do not want to rebuild the project, you can run the binary files located in directory `bin` (see [Run](#)).

With the tool chain ready, simply use `cargo` to build and run the program. One thing to notice is that the file `src/formula_parser/grammar.rs` is generated by [lalrpop/lalrpop: LR\(1\) parser generator for Rust \(github.com\)](#) from `src/formula_parser/grammar.lalrpop`. If you change the contents in `grammar.lalrpop`, a conversion is needed to update the `grammar.rs`. The following instructions should do the work:

```
1 | $ cargo install lalrpop          # only if you have not installed it
2 | $ lalrpop grammar.lalrpop
```

3 Run

The released `x86_linux-gnu` binary is located at `bin/command_line_interface`.

After executing the program, input the propositional logic formula and then hit `Enter`.

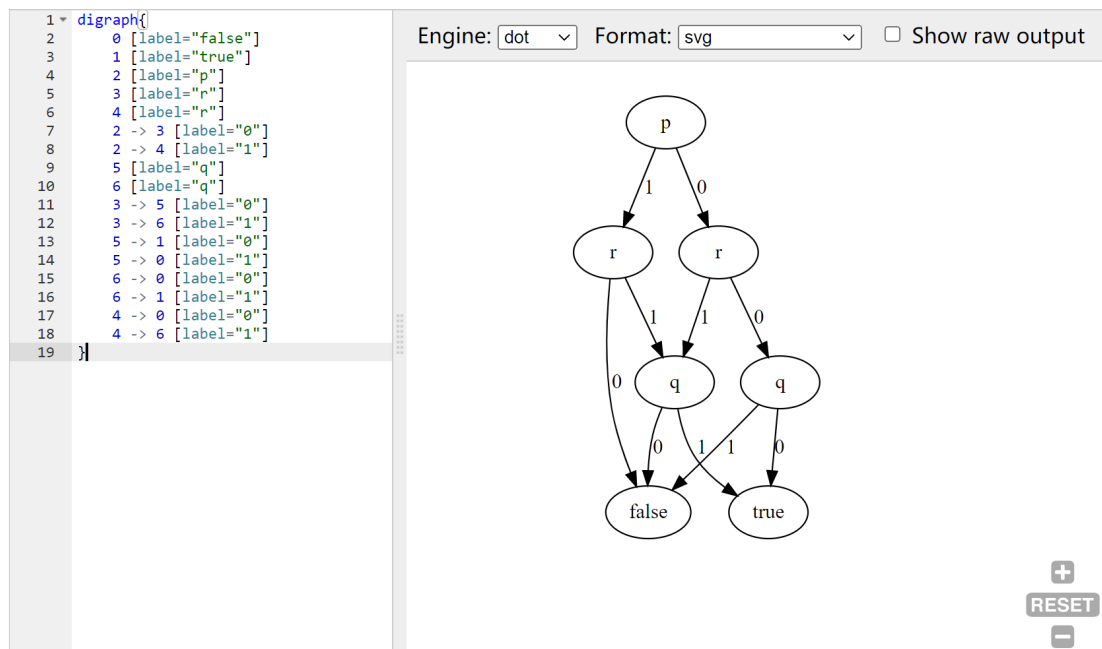
```

root reduced_ordered_binary_decision_diagram>./bin/command_line_interface
(p -> r) & ( q <-> (r | p))
digraph{
0 [label="false"]
1 [label="true"]
2 [label="p"]
3 [label="r"]
4 [label="r"]
2 -> 3 [label="0"]
2 -> 4 [label="1"]
5 [label="q"]
6 [label="q"]
3 -> 5 [label="0"]
3 -> 6 [label="1"]
5 -> 1 [label="0"]
5 -> 0 [label="1"]
6 -> 0 [label="0"]
6 -> 1 [label="1"]
4 -> 0 [label="0"]
4 -> 6 [label="1"]
}

To visualize the diagram, paste the output to http://viz-js.com/

```

To visualize the diagram, paste the output to [Viz.js \(viz-js.com\)](http://viz-js.com/):



4 Examples & Tests

1. F

```

root reduced_ordered_binary_decision_diagram>./bin/command_line_interface
F
digraph{
0 [label="false"]
}

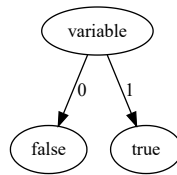
```

2. variable

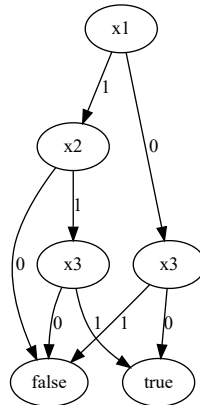
```

root reduced_ordered_binary_decision_diagram>./bin/command_line_interface
variable
digraph{
0 [label="false"]
1 [label="true"]
2 [label="variable"]
2 -> 0 [label="0"]
2 -> 1 [label="1"]
}

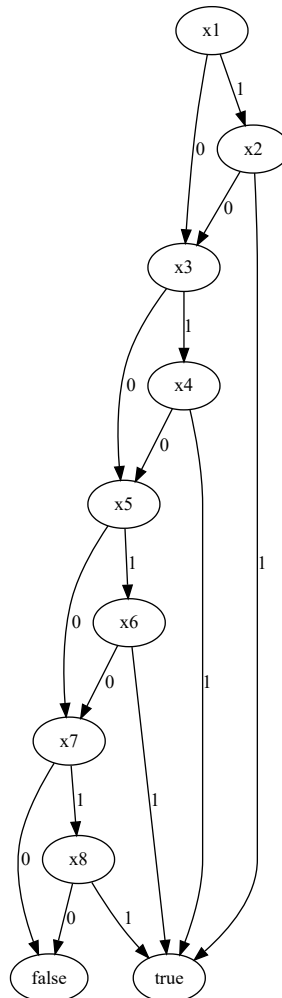
```



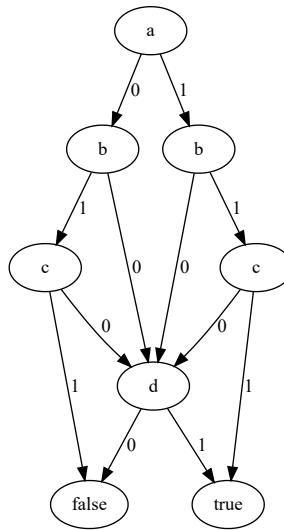
3. $(\neg x_1 \mid x_2) \ \& \ (x_1 \mid \neg x_3) \ \& \ (\neg x_1 \mid \neg x_2 \mid x_3)$



4. $x_1 \ \& \ x_2 \mid x_3 \ \& \ x_4 \mid x_5 \ \& \ x_6 \mid x_7 \ \& \ x_8$



5. $a \ \& \ b \ \& \ c \mid \neg b \ \& \ d \mid \neg c \ \& \ d$



5 Implementation

In this section, I would elucidate how I realized this ROBDD generator.

5.1 Data Structure: A Reduced Binary Tree

The structure of the primary data structure, that is, the one used to represent the BDD, is virtually the same as the diagrams shown above: each node has exactly 2 children and any number of parents; each node can represent a propositional logic formula.

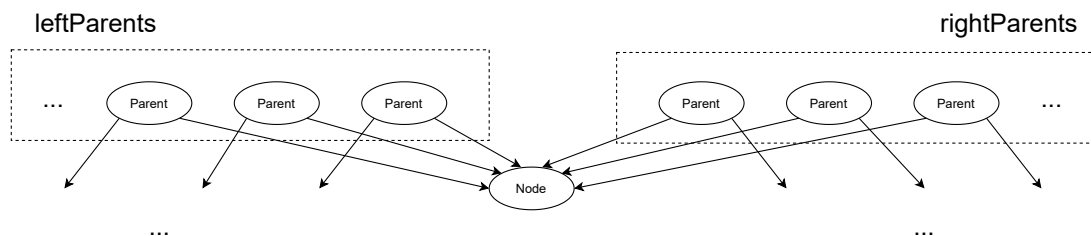
The nodes can never be duplicate since any addition of node to the graph would be check. This would be further discussed in [Merge and Elimination](#).

5.1.1 Inside a Node

The data structure maintains several field for each node:

1. The left and right child of the node
2. The value of the node, i.e, what the node represents
This can be either a variable or a Boolean constant.
3. The parents of this node

A node would maintain 2 sets of parents. One for the parents of which the node is the left child of, and one for the right. This field is primarily used for [merge and elimination](#).



5.1.2 Merge and Elimination

Every time a node is about to be added to the tree, duplication detection would be performed to eliminate duplicate/useless node.

Let $n.val$ denote the variable that the new node n represents, $n.left$ and $n.right$ denote the existing left and right child of the new node. The method `add_node_checked(n.value, n.left, n.right)` would either find an existing eligible node or add a new node to the tree and returns a pointer to it.

1. Eliminate redundant tests

If $n.\text{left} = n.\text{right}$, simply returns the child.

2. Merge equivalent leaves

There are only 2 leaves in the tree: one for `true` and one for `false`.

3. Merge isomorphic nodes

node x and node y are isomorphic if and only if

1. $x.\text{left} = y.\text{left}$, and
2. $x.\text{right} = y.\text{right}$, and
3. $x.\text{value} = y.\text{value}$

To efficiently determine if there is an existing node isomorphic to n , the program utilize the data field `parents` of the children, searches among $n.\text{left}.\text{rightParents} \cap n.\text{right}.\text{leftParents}$ for a node that has the same value as the new node.

5.2 Algorithm

The algorithm works in a recursive approach. It traverses the parse tree in depth-first order and employee function `apply()` on each operator node with its subtrees.

The pseudo code is as follows:

```
1  BDD_node construct ROBDD(parse_tree_node){
2      if (parse_tree_node is a constant) {                // `T` or `F`
3          return BDD::get_leaf(constant);
4      }else if (parse_tree_node is a variable){           // variable like
5          return BDD::add_variable(name);                 // a single-variable
6      }else if (parse_tree_node is a binary operator){    // `And`, `Or`,
7          BDD_node left_subtree = construct_ROBDD(parse_tree_node.left);
8          BDD_node right_subtree = construct_ROBDD(parse_tree_node.right);
9          return apply_binary(op, left_subtree, right_subtree);
10     }else{                                               // unary operator
11         BDD_node subtree = construct_ROBDD(parse_tree_node.child);
12         return apply_binary(op, subtree);
13     }
14 }
```

Line 2-5 handles the basic case where the node is a constant or variable. Line 6-12 invokes `apply` on the operator node depending on whether the operator is a binary one or a unary one.

5.2.1 `apply`

Function `apply()` constructs a new formula based on the given operator and existing operands.

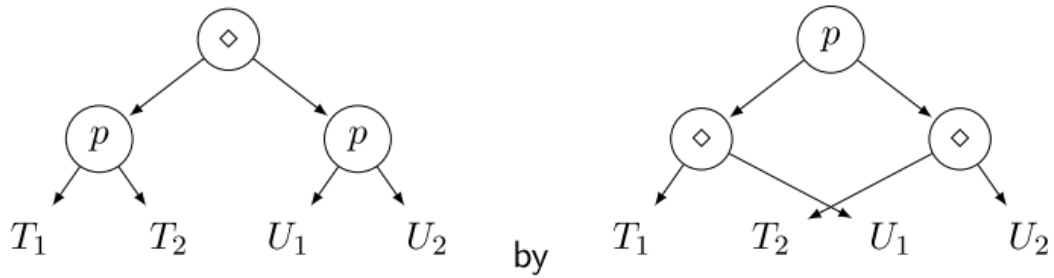
5.2.1.1 `apply_binary(op, left, right)`

1. Basic case

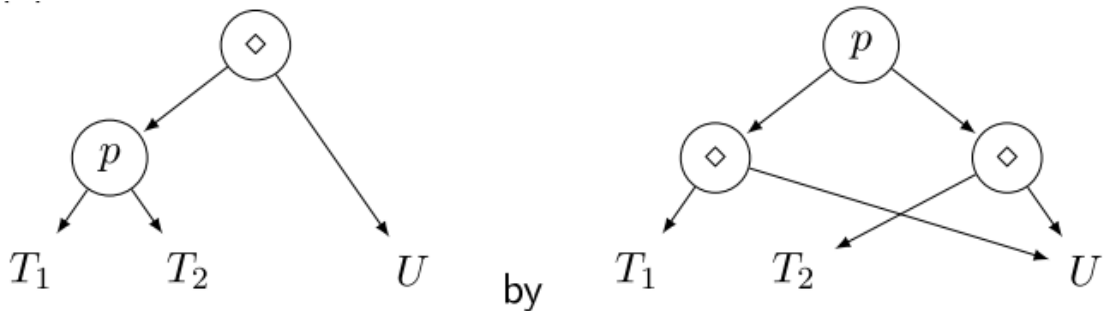
If either `left` or `right` is a constant, the new formula is a constant or determined by the other operand.

e.g. `apply_binary(And, false, right) = false`, `apply_binary(Or, false, right) = right`

2. If left and right have the same *smallest* variable, apply `op` on the grandchildren and add a new node with obtained children.



3. If the smallest variables appears on only one side, apply `op` on the grandchildren and children and add a new node with obtained children.

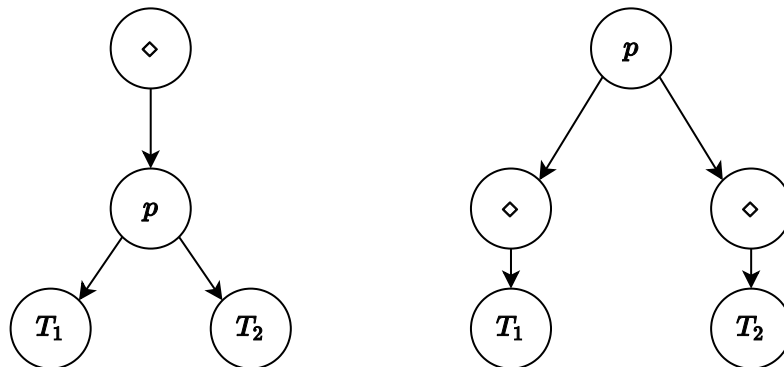


5.2.1.2 `apply_unary(op=Not, node)`

1. Basic case

If `node` is a constant, returns the negation of it.

2. otherwise, `node` is a variable. Obtain the negation of its children by `apply_unary(op, node.left)`, `apply_unary(op, node.right)` and add a new node with obtained children.



5.2.2 Variable Ordering

The order of variables is automatically determined by the program. It uses a simple strategy: the variable appearing earlier in the formula has a higher priority, and therefore appears closer to the root in the BDD.

6 Code Structure

All the source files are located in `src/`.

```
1  src/
2  |   bin
3  |   |   command_line_interface.rs
4  |   |   binary_decision_diagram
5  |   |   |   binary_index.rs
6  |   |   |   mod.rs
7  |   |   |   node_handler.rs
8  |   |   formula_parser
9  |   |   |   grammar.lalrpop
10 |   |   |   grammar.rs
11 |   |   |   mod.rs
12 |   |   lib.rs
13 |   |   utility.rs
```

1. `binary_decision_diagram`

Implementation of the data structure `BinaryDecisionDiagram`, which represent the reduced binary tree.

2. `lib.rs`

Implementation of the core algorithms, including `apply_binary`, `apply_unary`, mapping the variable names to numeric indices, constructing ROBDD from the parse tree.

3. `formula_parser`

Construct a parse tree from the input string.

7 Acknowledgement

[Course slide 5.2](#)