

Advanced Algorithms – COMS31900

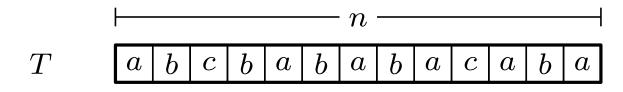
Pattern Matching part two

Suffix Arrays

Benjamin Sach

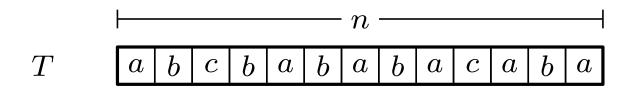


Preprocess a text string T (length n) to answer pattern matching queries. . .





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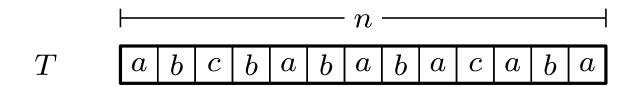


After preprocessing, a query is a pattern P (length m),

$$P \qquad \boxed{a \mid b \mid a}$$
$$\vdash m \dashv$$



Preprocess a text string T (length n) to answer pattern matching queries. . .

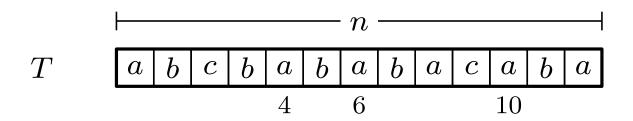


After preprocessing, a query is a pattern P (length m),

$$P \qquad \boxed{a \mid b \mid a} \\ \vdash m \dashv$$

the output is a list of all matches in T.

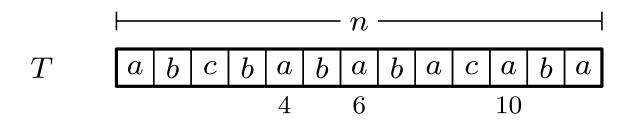
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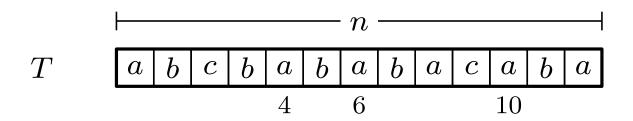
Preprocess a text string T (length n) to answer pattern matching queries...



After preprocessing, a query is a pattern P (length m),

• Last lecture we saw the that **text indexing** problem can be solved using a suffix tree which uses O(n) space (when it's stored compacted)

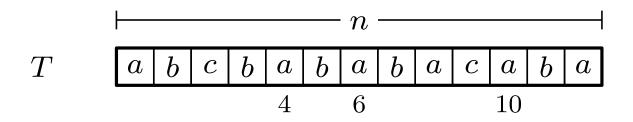
Preprocess a text string T (length n) to answer pattern matching queries...



After preprocessing, a query is a pattern P (length m),

- Last lecture we saw the that **text indexing** problem can be solved using a suffix tree which uses O(n) space (when it's stored compacted)
- Queries take $O(m + {\sf occ})$ time when the alphabet size is constant occ is the number of occurences (matches)

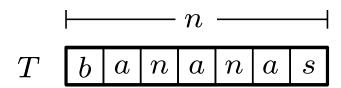
Preprocess a text string T (length n) to answer pattern matching queries...



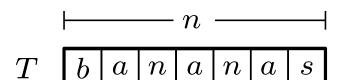
After preprocessing, a query is a pattern P (length m),

- Last lecture we saw the that **text indexing** problem can be solved using a suffix tree which uses O(n) space (when it's stored compacted)
- Queries take $O(m + {\sf occ})$ time when the alphabet size is constant ${\sf occ}$ is the number of occurences (matches)
- Suffix trees can be constructed in O(n) time (but we only saw how to achieve $O(n^2)$ time)



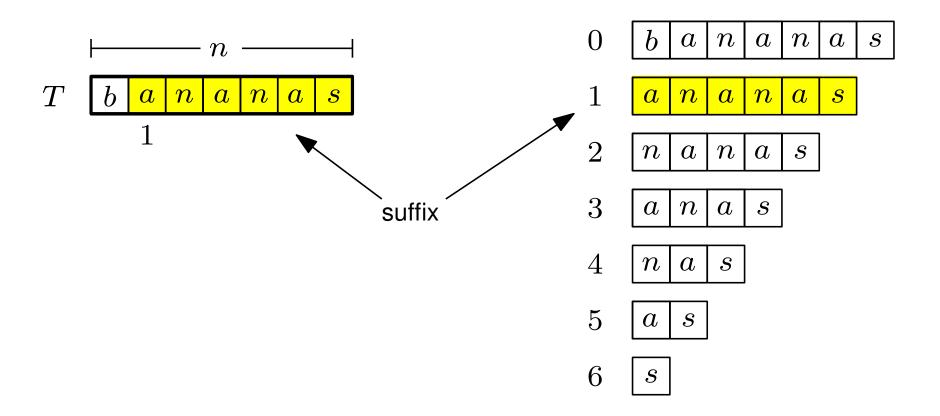




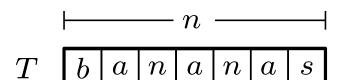


- $0 \quad \boxed{b \quad a \quad n \quad a \quad n \quad a \quad s}$
- $1 \quad \boxed{a \mid n \mid a \mid n \mid a \mid s}$
- $2 \quad \boxed{n \mid a \mid n \mid a \mid s}$
- $a \mid a \mid n \mid a \mid s$
- $4 \quad n \mid a \mid s$
- $5 \quad \boxed{a \mid s}$
- $6 \quad \boxed{s}$



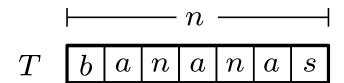






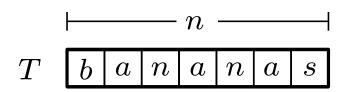
- $0 \quad \boxed{b \quad a \quad n \quad a \quad n \quad a \quad s}$
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- $a \mid a \mid n \mid a \mid s$
- $4 \quad n \mid a \mid s$
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- $a \mid a \mid n \mid a \mid s$
- $4 \quad \boxed{n \mid a \mid s}$
- $5 \quad \boxed{a \mid s}$
- $6 \quad \boxed{s}$





Sort the suffixes lexicographically

 The symbols themselves must have an order throughout we will use alphabetical order

0	b	\overline{a}	n	\overline{a}	n	a	s
---	---	----------------	---	----------------	---	---	---

$$1 \quad \boxed{a \mid n \mid a \mid n \mid a \mid s}$$

$$2 \quad \boxed{n \mid a \mid n \mid a \mid s}$$

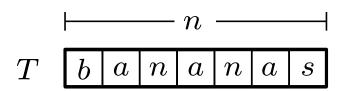
$$3 \quad a \mid n \mid a \mid s$$

$$4 \quad \boxed{n \mid a \mid s}$$

$$5 \quad \boxed{a \mid s}$$

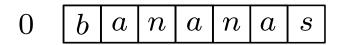
$$6 \quad \boxed{s}$$

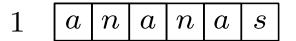




Sort the suffixes lexicographically

 The symbols themselves must have an order throughout we will use alphabetical order





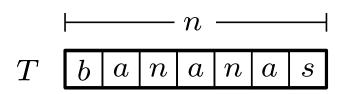
$$2 \quad \boxed{n \mid a \mid n \mid a \mid s}$$

$$a \mid a \mid n \mid a \mid s$$

$$4 \quad \boxed{n \mid a \mid s}$$

$$5 \quad \boxed{a \mid s}$$

$$6 \quad \boxed{s}$$



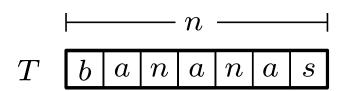
Sort the suffixes lexicographically

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- $0 \quad \boxed{b \mid a \mid n \mid a \mid n \mid a \mid s}$
- 1 $a \mid n \mid a \mid n \mid a \mid s$
- $2 \quad \boxed{n \mid a \mid n \mid a \mid s}$
- $3 \quad a \mid n \mid a \mid s$
- $4 \mid n \mid a \mid s$
- $5 \quad \boxed{a \mid s}$
- $6 \quad \boxed{s}$

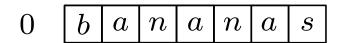
$$a \mid a \mid < b \mid a$$

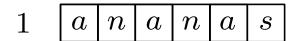




Sort the suffixes lexicographically

 The symbols themselves must have an order throughout we will use alphabetical order





$$2 \quad \boxed{n \mid a \mid n \mid a \mid s}$$

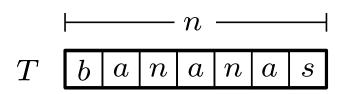
$$a \mid a \mid n \mid a \mid s$$

$$4 \mid n \mid a \mid s$$

$$5 \mid a \mid s$$

$$6 \quad \boxed{s}$$

$$\begin{bmatrix} a & a \end{bmatrix} < \begin{bmatrix} b & a \end{bmatrix}$$



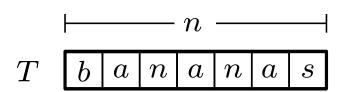
Sort the suffixes lexicographically

 The symbols themselves must have an order throughout we will use alphabetical order

- $0 \quad \boxed{b \mid a \mid n \mid a \mid n \mid a \mid s}$
- 1 $a \mid n \mid a \mid n \mid a \mid s$
- $2 \quad \boxed{n \mid a \mid n \mid a \mid s}$
- $3 \quad a \mid n \mid a \mid s$
- $4 \mid n \mid a \mid s$
- $5 \quad \boxed{a \mid s}$
- $6 \quad \boxed{s}$

$$a \mid a \mid < b \mid a$$



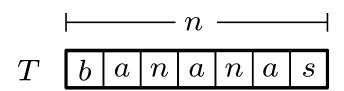


- Sort the suffixes lexicographically
- The symbols themselves must have an order throughout we will use alphabetical order

- $0 \quad \boxed{b \mid a \mid n \mid a \mid n \mid a \mid s}$
- $1 \quad \boxed{a \mid n \mid a \mid n \mid a \mid s}$
- $2 \quad \boxed{n \mid a \mid n \mid a \mid s}$
- $a \mid a \mid n \mid a \mid s$
- $4 \mid n \mid a \mid s$
- $5 \mid a \mid s$
- $6 \quad \boxed{s}$

$$a \mid a \mid c \mid c \mid c$$



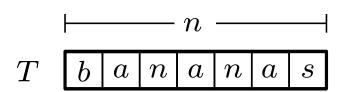


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- The symbols themselves must have an order throughout we will use alphabetical order

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- $a \mid a \mid n \mid a \mid s$
- $4 \mid n \mid a \mid s$
- $5 \mid a \mid s$
- $6 \quad \boxed{s}$

$$\begin{bmatrix} a & a \end{bmatrix} < \begin{bmatrix} b & a \end{bmatrix} < \begin{bmatrix} b & c \end{bmatrix}$$



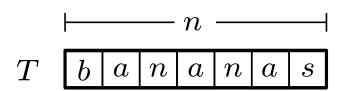


- Sort the suffixes lexicographically
- The symbols themselves must have an order throughout we will use alphabetical order

- $0 \quad \boxed{b \mid a \mid n \mid a \mid n \mid a \mid s}$
- $1 \quad \boxed{a \mid n \mid a \mid n \mid a \mid s}$
- $2 \quad \boxed{n \mid a \mid n \mid a \mid s}$
- $a \mid a \mid n \mid a \mid s$
- $4 \mid n \mid a \mid s$
- $5 \mid a \mid s$
- $6 \quad \boxed{s}$

$$a \mid a \mid c \mid c \mid c$$





Sort the suffixes

 $2 \quad \boxed{n \mid a \mid n \mid a \mid s}$

lexicographically

 $3 \quad \boxed{a \mid n \mid a \mid s}$

a

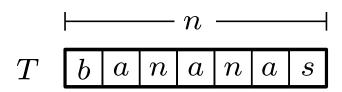
s

 $5 \quad \boxed{a \mid s}$

 The symbols themselves must have an order throughout we will use alphabetical order

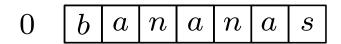
 $6 \quad \boxed{s}$

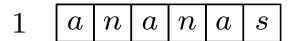




Sort the suffixes lexicographically

 The symbols themselves must have an order throughout we will use alphabetical order





$$2 \quad n \mid a \mid n \mid a \mid s$$

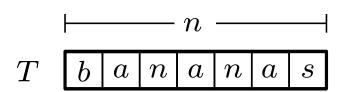
$$a \mid a \mid n \mid a \mid s$$

$$4 \mid n \mid a \mid s$$

$$5 \mid a \mid s$$

$$6 \quad \boxed{s}$$

The suffix array



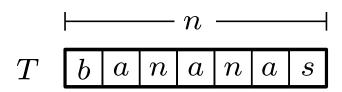
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- $4 \mid n \mid a \mid s$
- $5 \quad \boxed{a \mid s}$
- $6 \quad \boxed{s}$

In lexicographical ordering we sort strings based on the first symbol that differs:

The suffix array



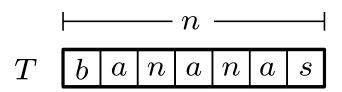
Sort the suffixes lexicographically

 The symbols themselves must have an order throughout we will use alphabetical order

- $0 \quad \boxed{b \mid a \mid n \mid a \mid n \mid a \mid s}$
- 1 $a \mid n \mid a \mid n \mid a \mid s$
- $2 \quad \boxed{n \mid a \mid n \mid a \mid s}$
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Sort the suffixes

lexicographically

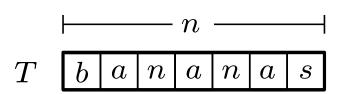
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- $1 \quad \boxed{a \mid n \mid a \mid n \mid a \mid s}$
- $a \mid a \mid n \mid a \mid s$
- $5 \mid a \mid s$
- $0 \quad b \mid a \mid n \mid a \mid n \mid a \mid s$
- $2 \mid n \mid a \mid n \mid a \mid s$
- $4 \quad \boxed{n \mid a \mid s}$
- $6 \quad \boxed{s}$

In lexicographical ordering we sort strings based on the first symbol that differs:

$$\begin{bmatrix} a & a \end{bmatrix} < \begin{bmatrix} b & a \end{bmatrix} < \begin{bmatrix} b & c \end{bmatrix} < \begin{bmatrix} b & c \end{bmatrix}$$

The suffix array



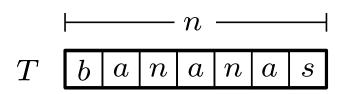
- $1 \quad \boxed{a \mid n \mid a \mid n \mid a \mid s}$
- $a \mid a \mid n \mid a \mid s$
- $5 \mid a \mid s$

Sort the suffixes lexicographically

- $2 \quad n \quad a \quad n \quad s$
- $4 \quad n \mid a \mid s$
- The symbols themselves must have an order throughout we will use alphabetical order
- 6 *s*

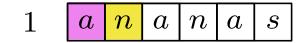
In lexicographical ordering we sort strings based on the first symbol that differs:

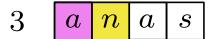
The suffix array



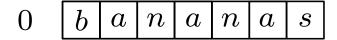
Sort the suffixes lexicographically

 The symbols themselves must have an order throughout we will use alphabetical order





$$5 \quad a \quad s$$



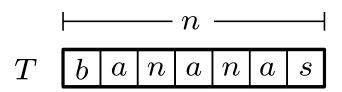
$$2 \quad \boxed{n \mid a \mid n \mid a \mid s}$$

$$4 \mid n \mid a \mid s$$

$$6 \quad \boxed{s}$$

In lexicographical ordering we sort strings based on the first symbol that differs:

The suffix array



Sort the suffixes

lexicographically

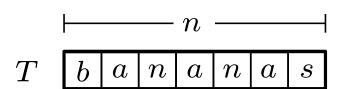
 The symbols themselves must have an order throughout we will use alphabetical order

- $1 \quad \boxed{a \mid n \mid a \mid n \mid a \mid s}$
- $a \mid a \mid n \mid a \mid s$
- $5 \mid a \mid s$
- $0 \quad b \mid a \mid n \mid a \mid n \mid a \mid s$
- $2 \mid n \mid a \mid n \mid a \mid s$
- $4 \quad \boxed{n \mid a \mid s}$
- $6 \quad \boxed{s}$

In lexicographical ordering we sort strings based on the first symbol that differs:

$$\begin{bmatrix} a & a \end{bmatrix} < \begin{bmatrix} b & a \end{bmatrix} < \begin{bmatrix} b & c \end{bmatrix} < \begin{bmatrix} b & c \end{bmatrix}$$

The suffix array



- $1 \quad \boxed{a \mid n \mid a \mid n \mid a \mid s}$
- $a \mid a \mid a \mid s$
- $5 \quad a \mid s$

Sort the suffixes lexicographically

- $0 \quad b \quad a \quad n \quad a \quad n \quad s$
- $2 \quad \boxed{n \mid a \mid n \mid a \mid s}$
- $4 \quad \boxed{n \mid a \mid s}$
- The symbols themselves must have an order throughout we will use alphabetical order

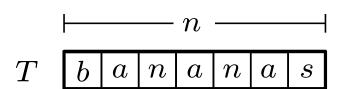
 $6 \quad \boxed{s}$

just a fancy name for the order the strings would appear in the dictionary

In lexicographical ordering we sort strings based on the first symbol that differs:

$$\begin{bmatrix} a & a \end{bmatrix} < \begin{bmatrix} b & a \end{bmatrix} < \begin{bmatrix} b & c \end{bmatrix} < \begin{bmatrix} b & c \end{bmatrix}$$

The suffix array



- $1 \quad \boxed{a \mid n \mid a \mid n \mid a \mid s}$
- a n a s
- $5 \quad a \mid s$

Sort the suffixes lexicographically

- $0 \quad b \quad a \quad n \quad a \quad n \quad s$
- $2 \quad \boxed{n \mid a \mid n \mid a \mid s}$
- $4 \quad \boxed{n \mid a \mid s}$
- The symbols themselves must have an order
 throughout we will use alphabetical order
- $6 \quad \boxed{s}$

j j

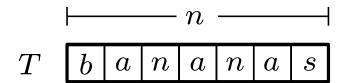
just a fancy name for the order the strings would appear in the dictionary

In lexicographical ordering we sort strings based on the first symbol that differs:

(in a 'tie', the shorter string is smaller)

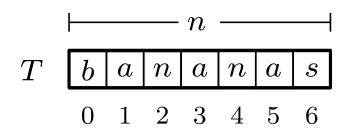
If the symbols don't have a natural order, we use their binary representation in memory





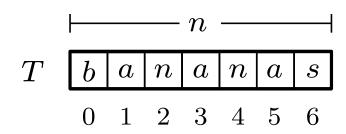
- 1 $a \mid n \mid a \mid n \mid a \mid s$
- $a \mid a \mid n \mid a \mid s$
- $5 \mid a \mid s$
- $0 \quad b \quad a \quad n \quad a \quad n \quad s$
- $2 \quad \boxed{n \mid a \mid n \mid a \mid s}$
- $4 \quad \boxed{n \mid a \mid s}$
- $6 \mid s$





- 1 $a \mid n \mid a \mid n \mid a \mid s$
- $a \mid a \mid n \mid a \mid s$
- $5 \mid a \mid s$
- $0 \mid b \mid a \mid n \mid a \mid n \mid a \mid s$
- $2 \mid n \mid a \mid n \mid a \mid s$
- $4 \quad \boxed{n \mid a \mid s}$
- $6 \mid s$





1 a n a n a s

$$a \mid a \mid n \mid a \mid s$$

$$5 \mid a \mid s$$

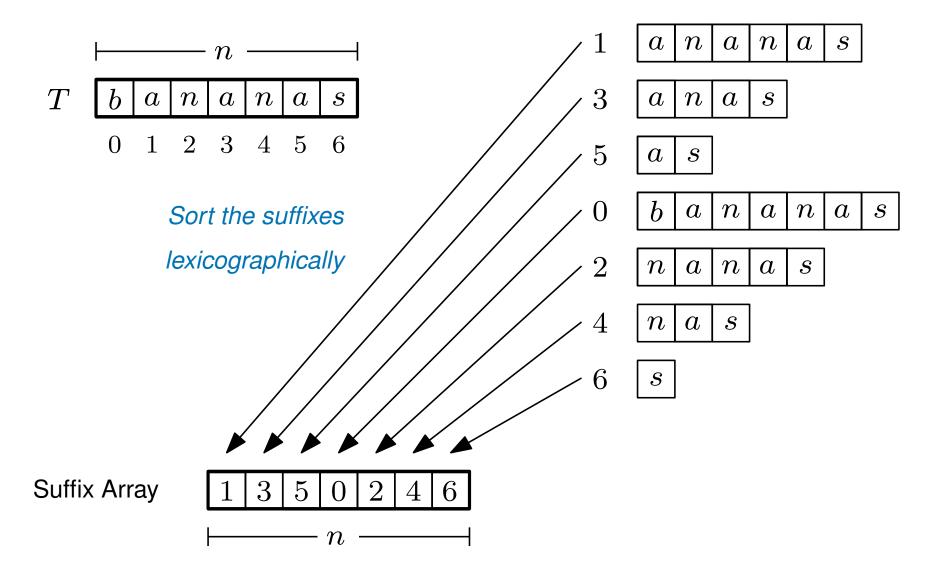
$$0 \quad b \mid a \mid n \mid a \mid n \mid a \mid s$$

$$2 \quad \boxed{n \mid a \mid n \mid a \mid s}$$

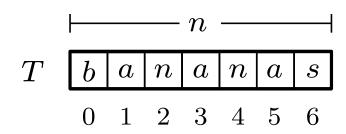
$$4 \quad \boxed{n \mid a \mid s}$$

$$6 \quad \boxed{s}$$









1 a n a n a s

$$a \mid a \mid n \mid a \mid s$$

$$5 \mid a \mid s$$

$$0 \quad b \mid a \mid n \mid a \mid n \mid a \mid s$$

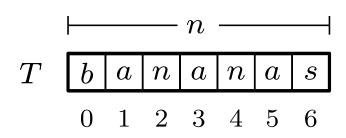
$$2 \quad \boxed{n \mid a \mid n \mid a \mid s}$$

$$4 \quad \boxed{n \mid a \mid s}$$

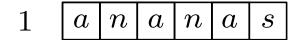
$$6 \quad \boxed{s}$$

University of BRISTOL

The suffix array



Sort the suffixes lexicographically



$$3 \quad a \mid n \mid a \mid s$$

$$5 \mid a \mid s$$

$$0 \quad b \mid a \mid n \mid a \mid n \mid a \mid s$$

$$2 \mid n \mid a \mid n \mid a \mid s$$

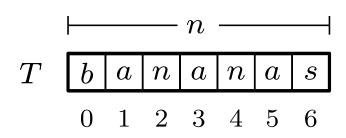
$$4 \quad \boxed{n \mid a \mid s}$$

$$6 \quad \boxed{s}$$

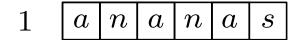
The suffix array is much smaller than the suffix tree (in terms of constants)

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The suffix array



Sort the suffixes lexicographically



$$3 \quad a \mid n \mid a \mid s$$

$$5 \mid a \mid s$$

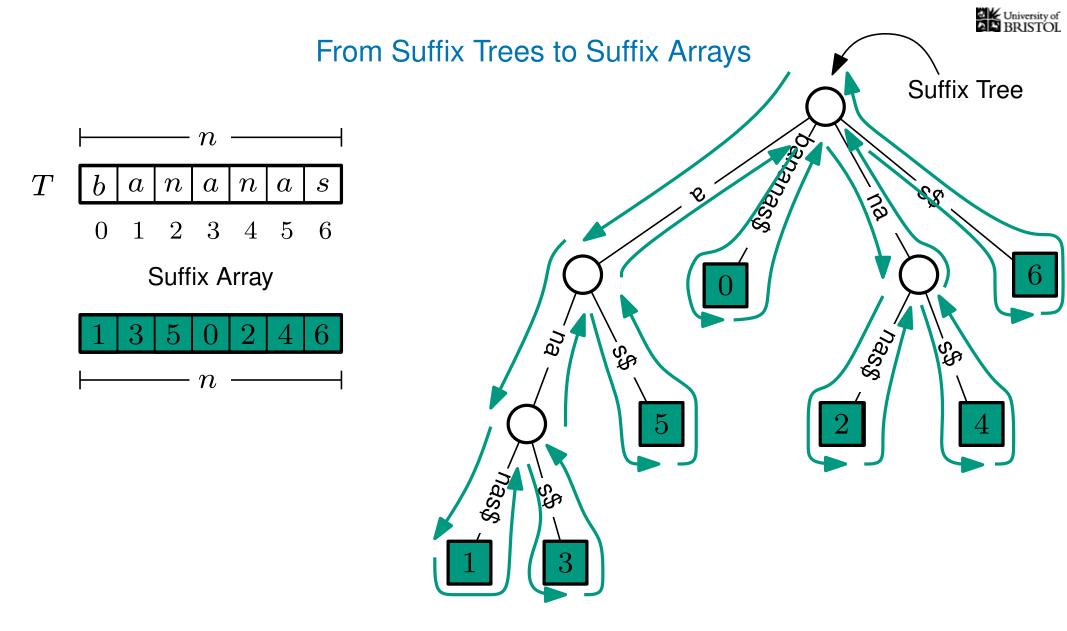
$$0 \quad b \mid a \mid n \mid a \mid n \mid a \mid s$$

$$2 \mid n \mid a \mid n \mid a \mid s$$

$$4 \quad \boxed{n \mid a \mid s}$$

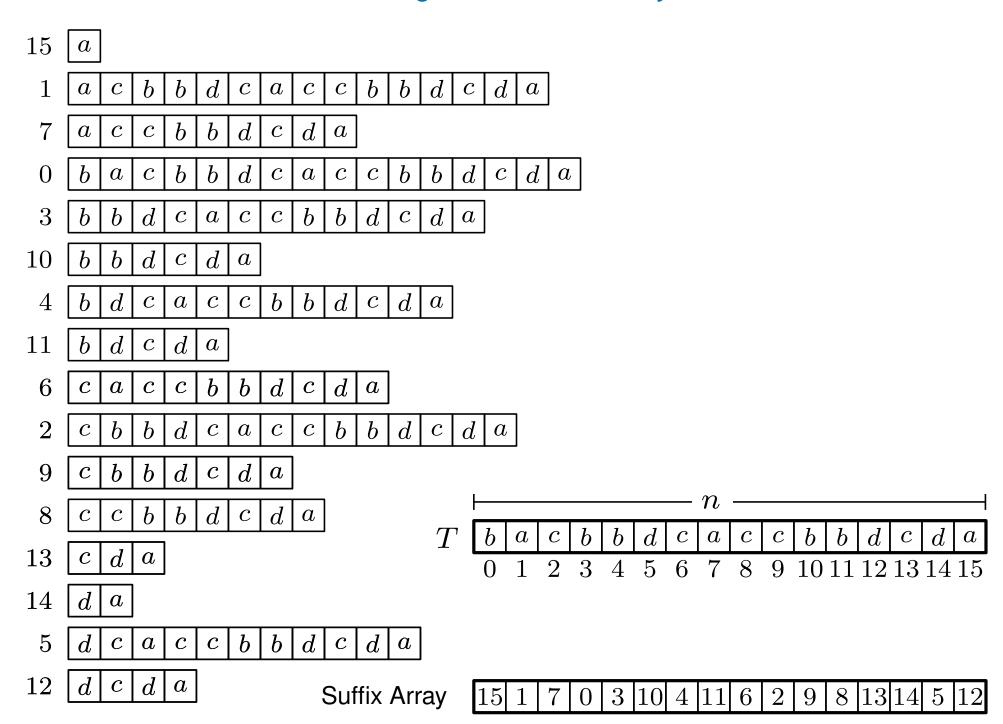
$$6 \quad \boxed{s}$$

The suffix array is much smaller than the suffix tree (in terms of constants)

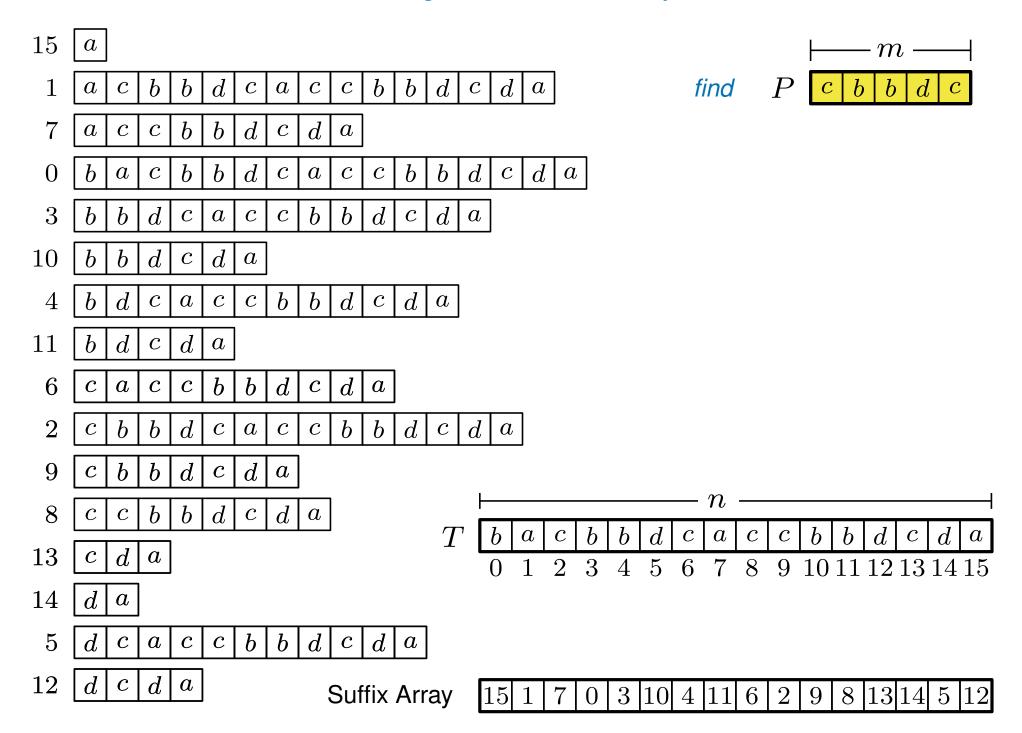


Recall that we can get get the Suffix Array from the Suffix Tree using depth-first search in O(n) time

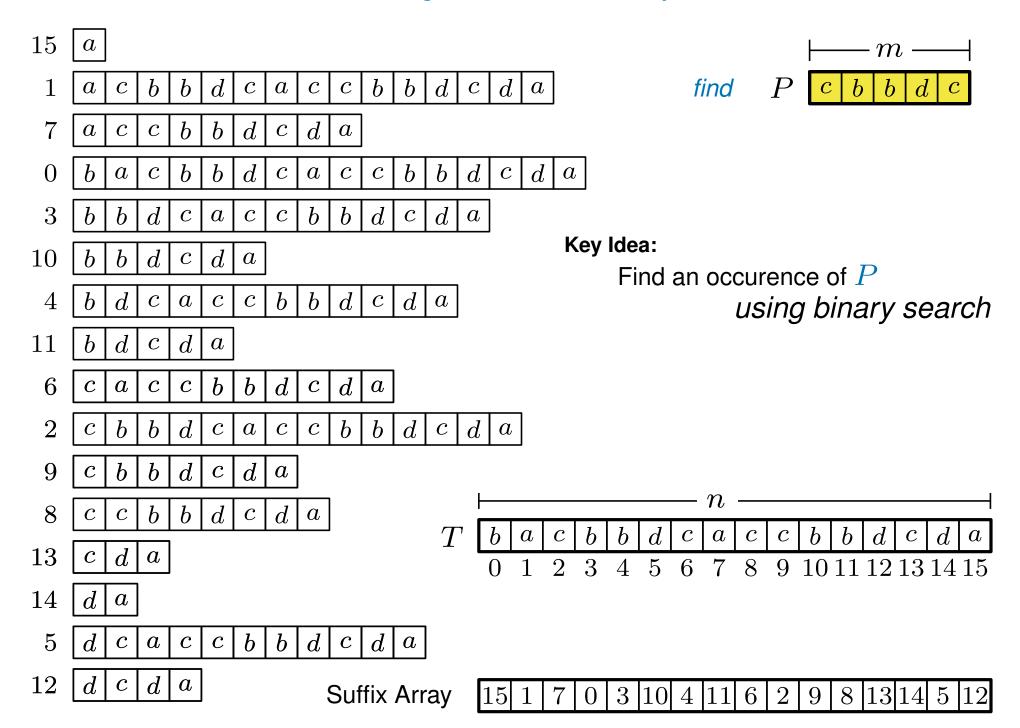




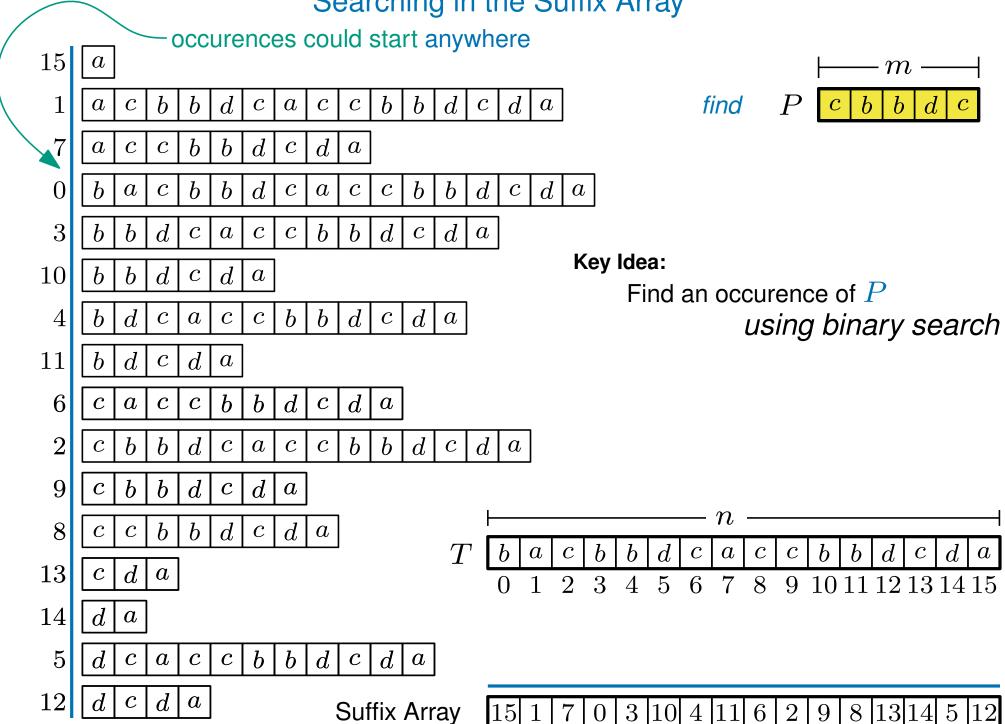






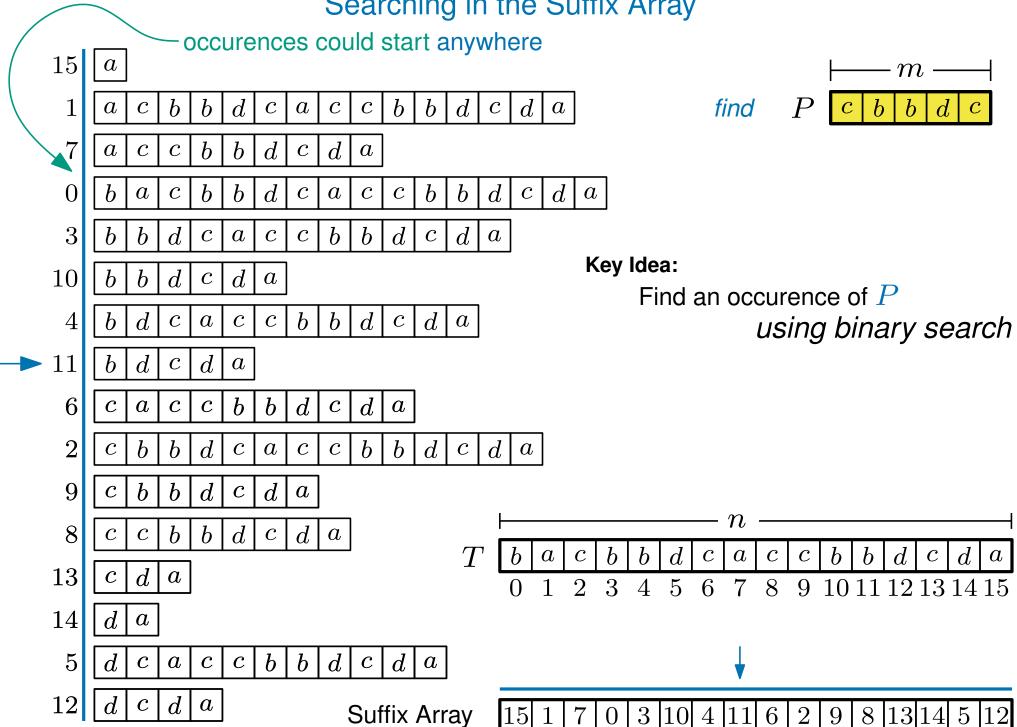




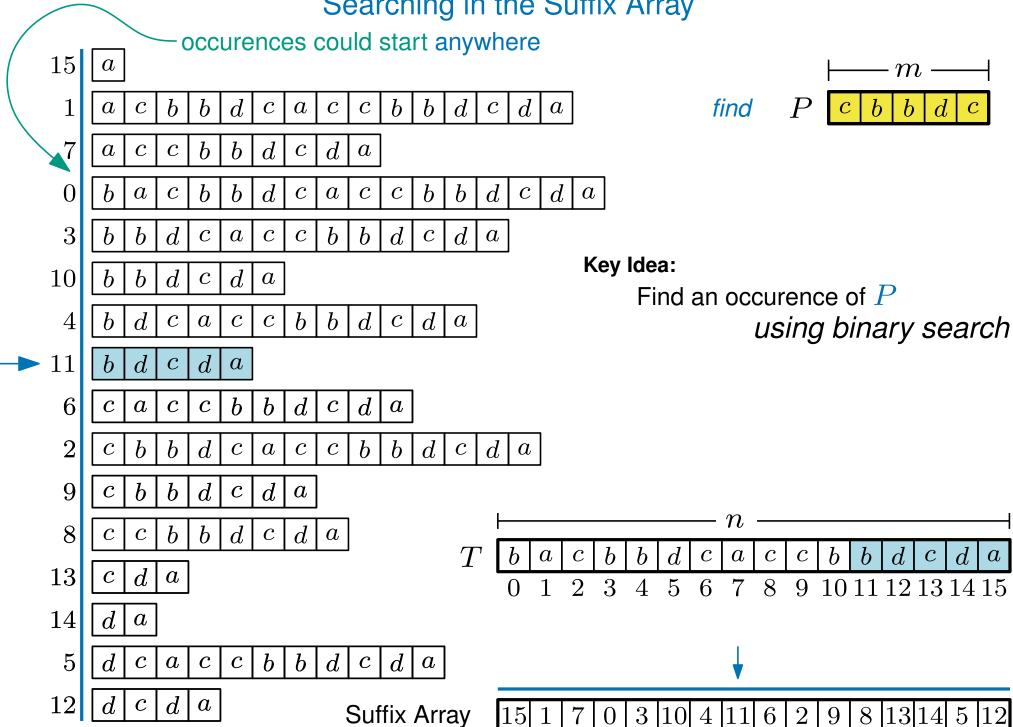




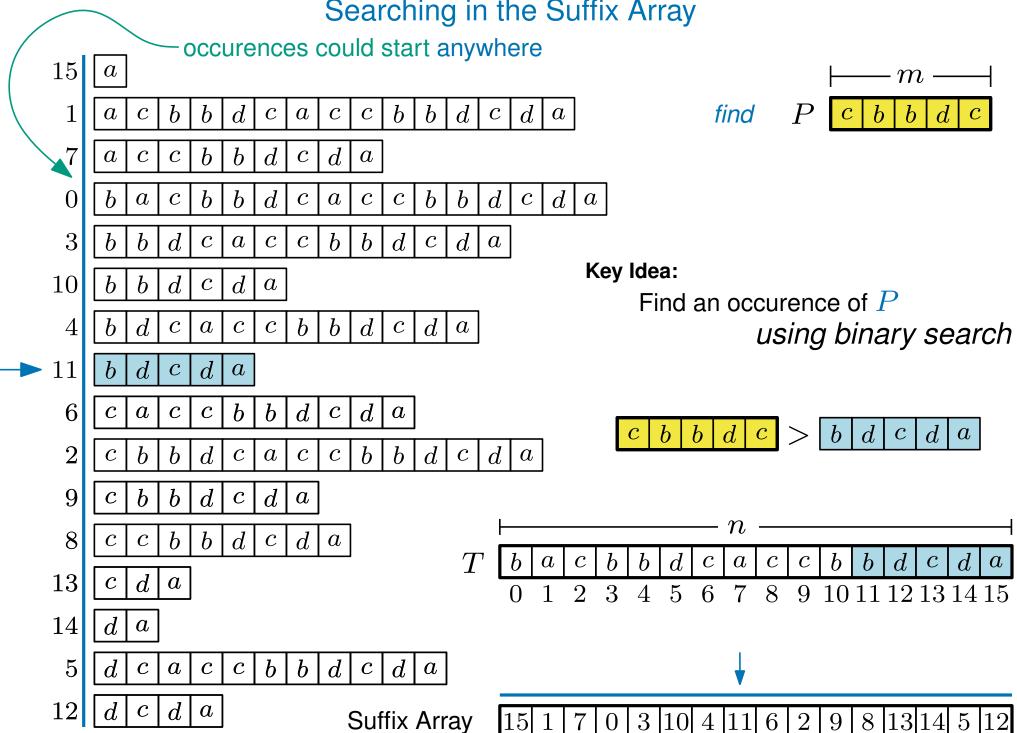




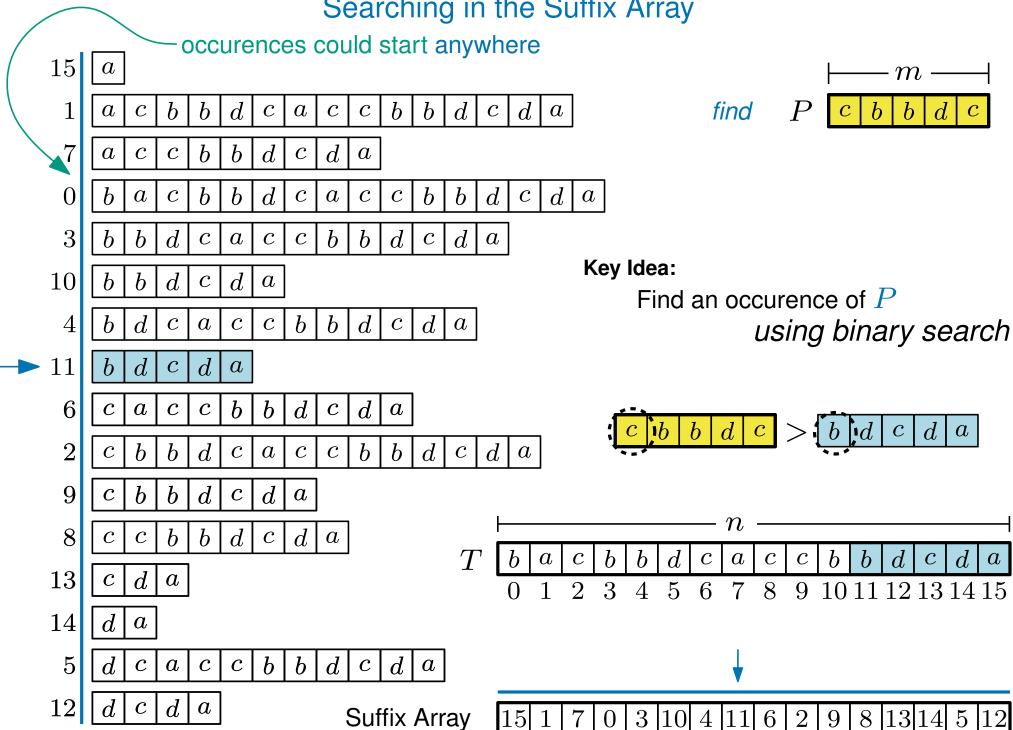




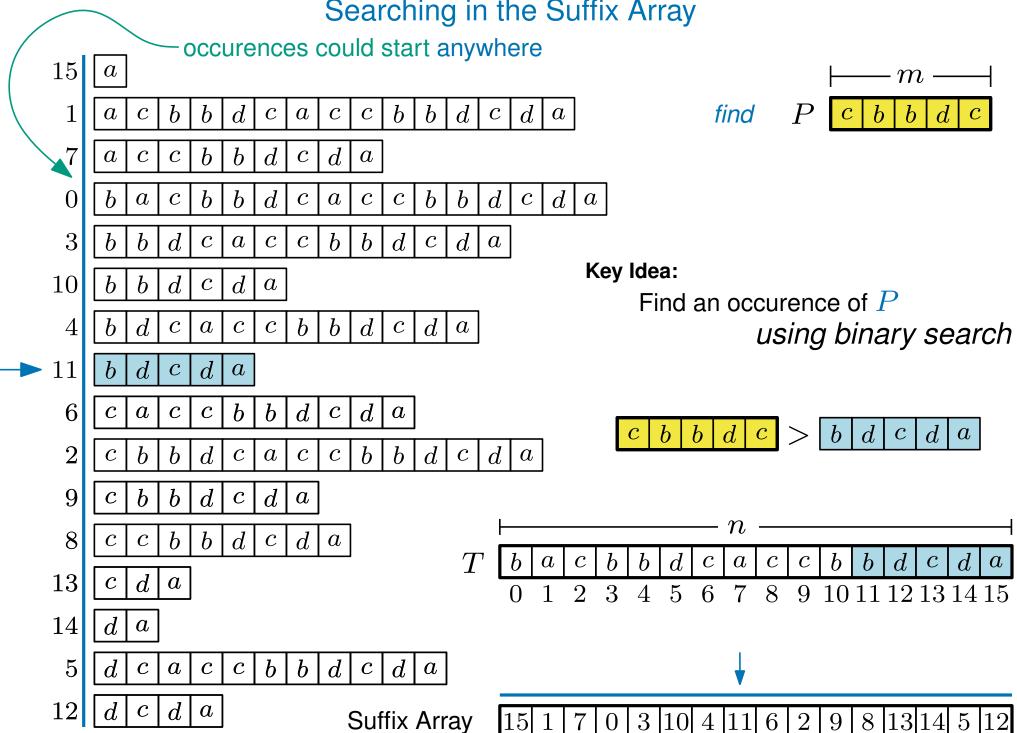




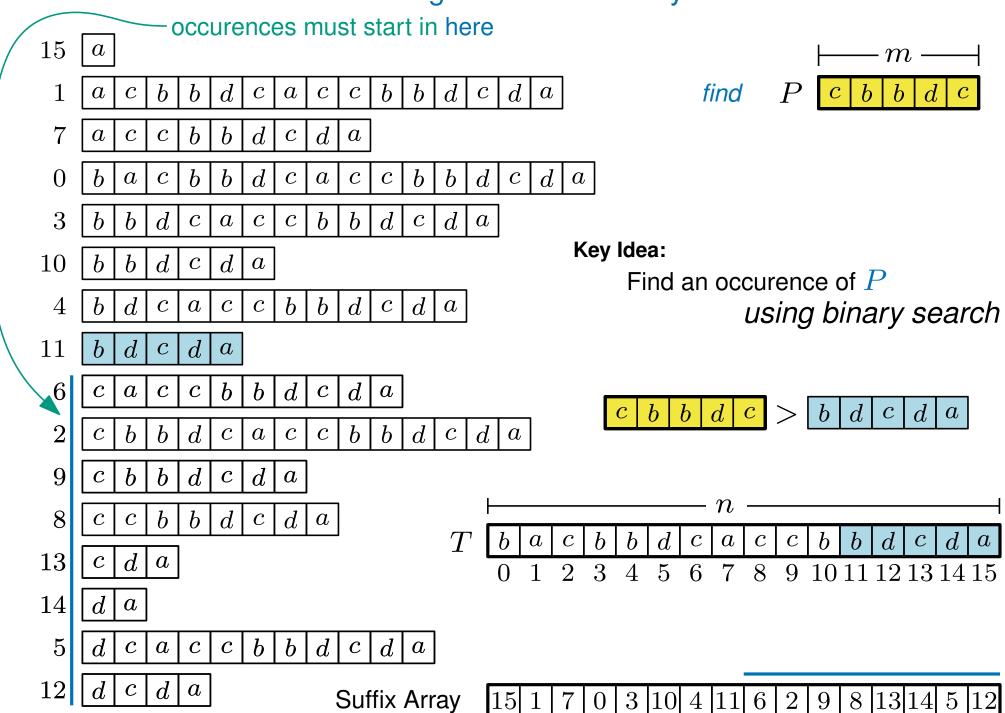




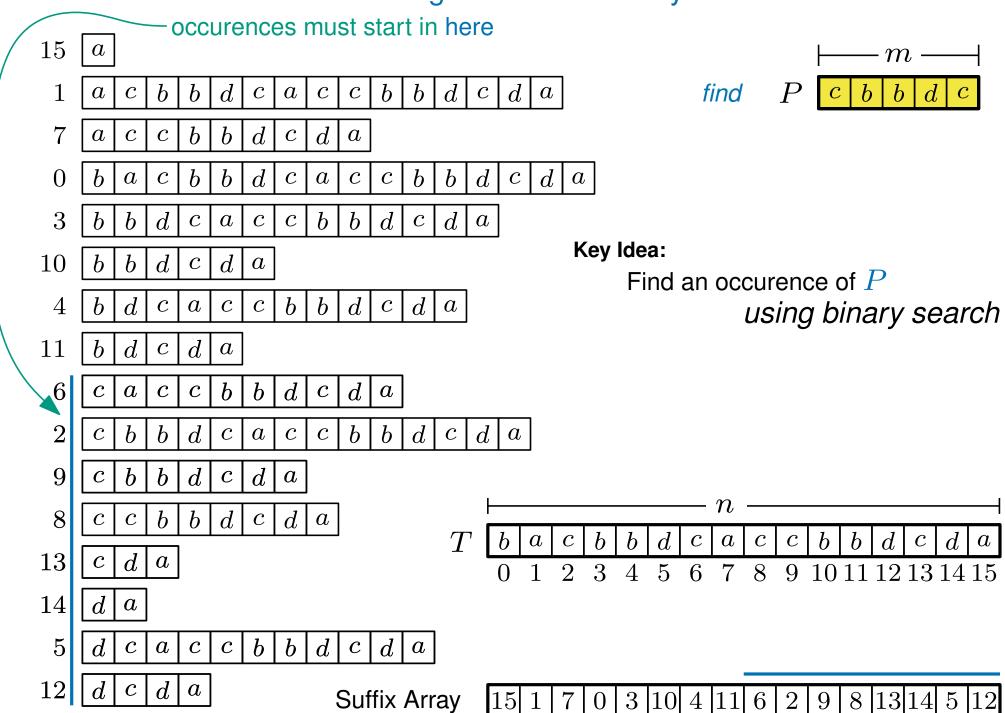




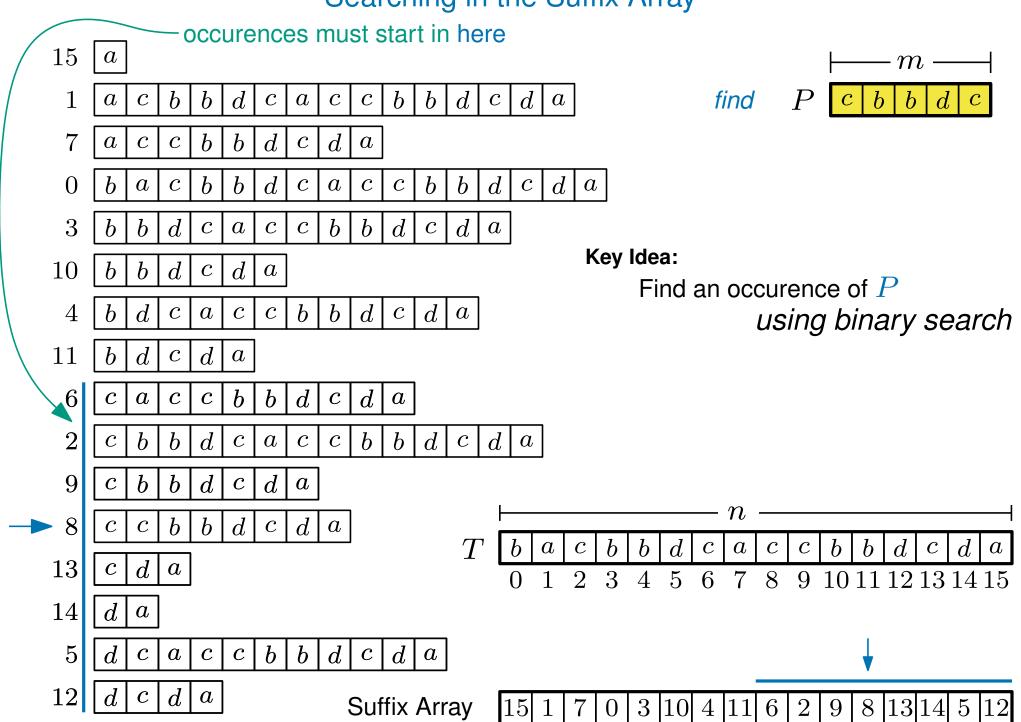




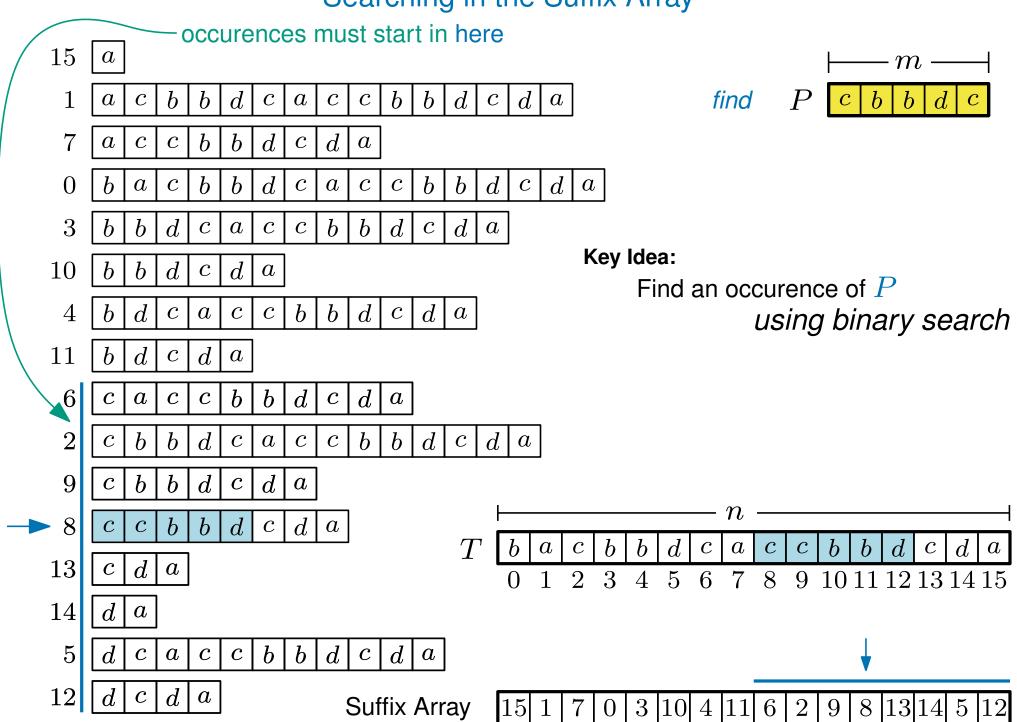




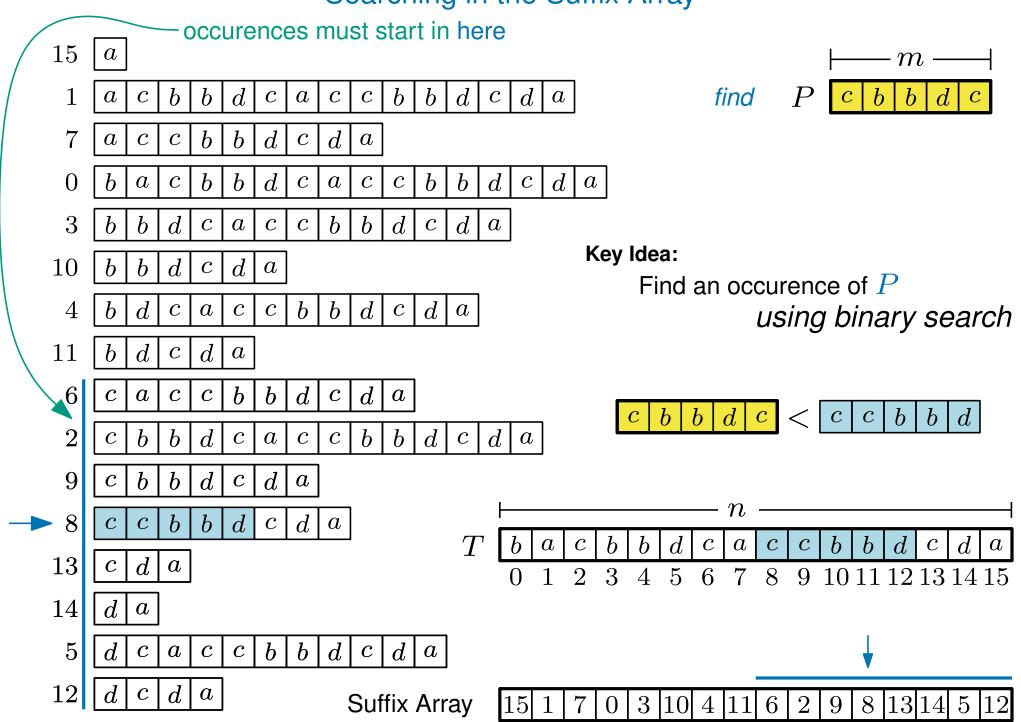




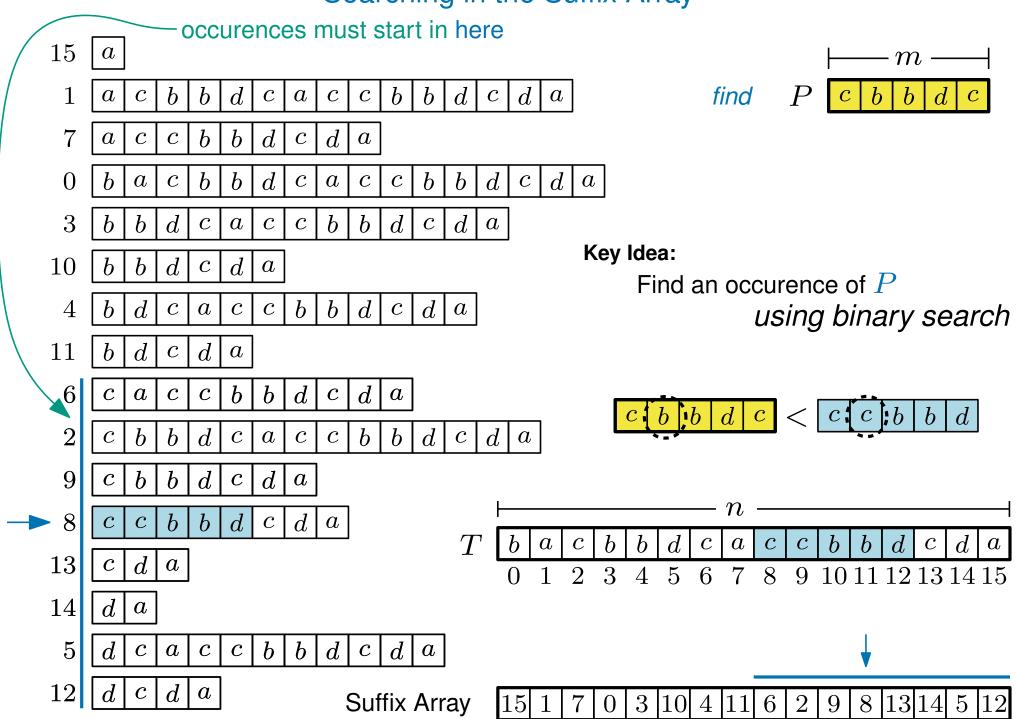




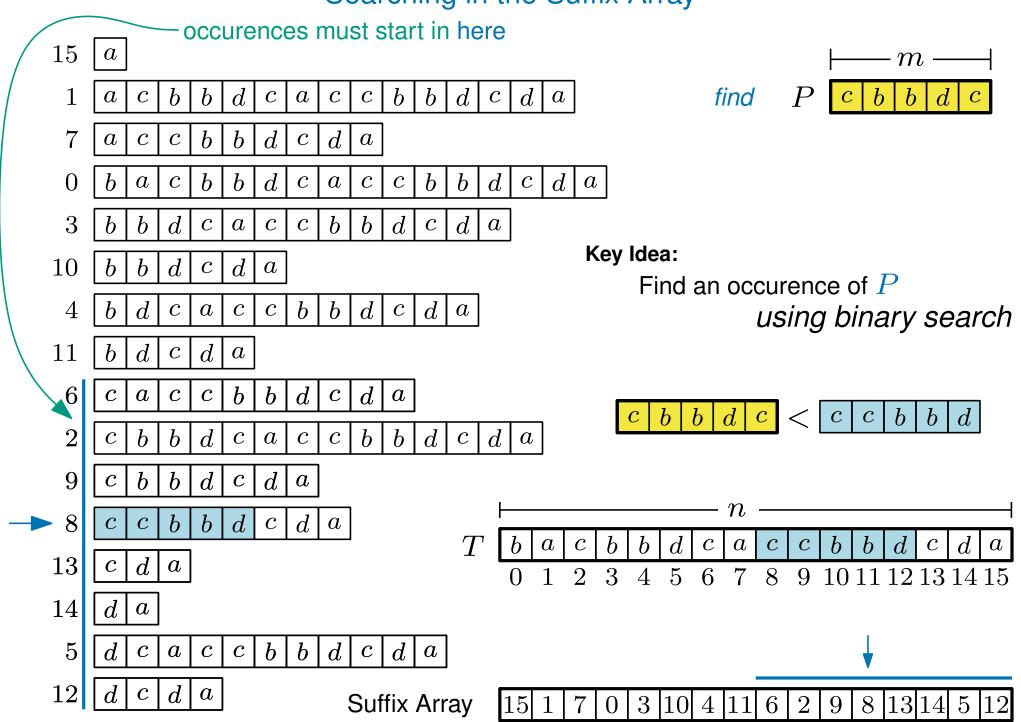




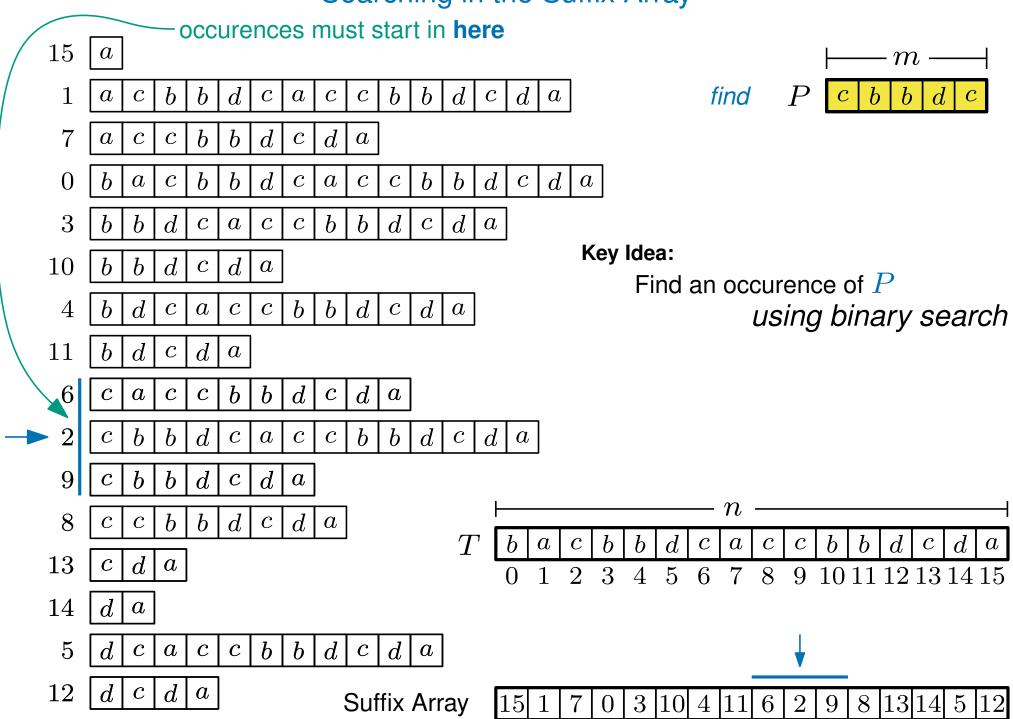




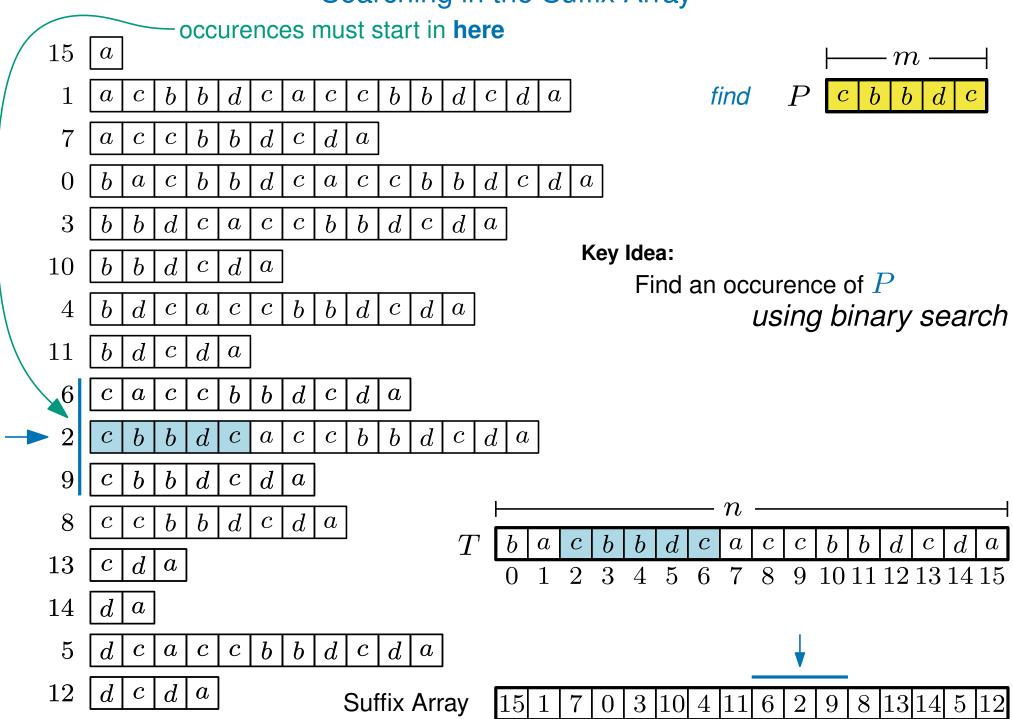




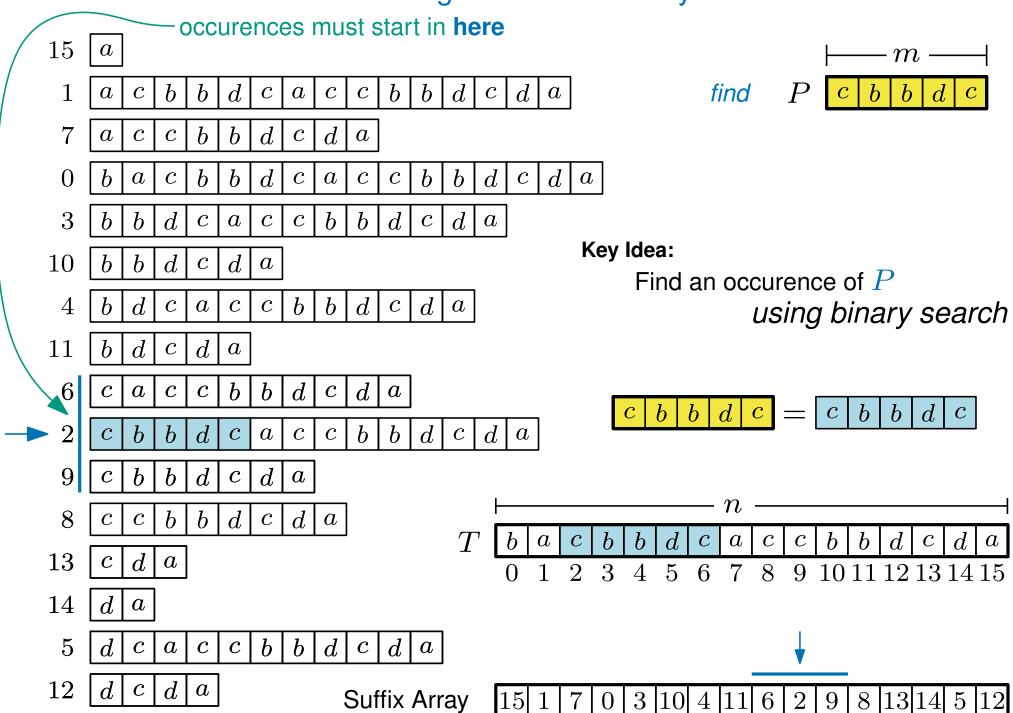




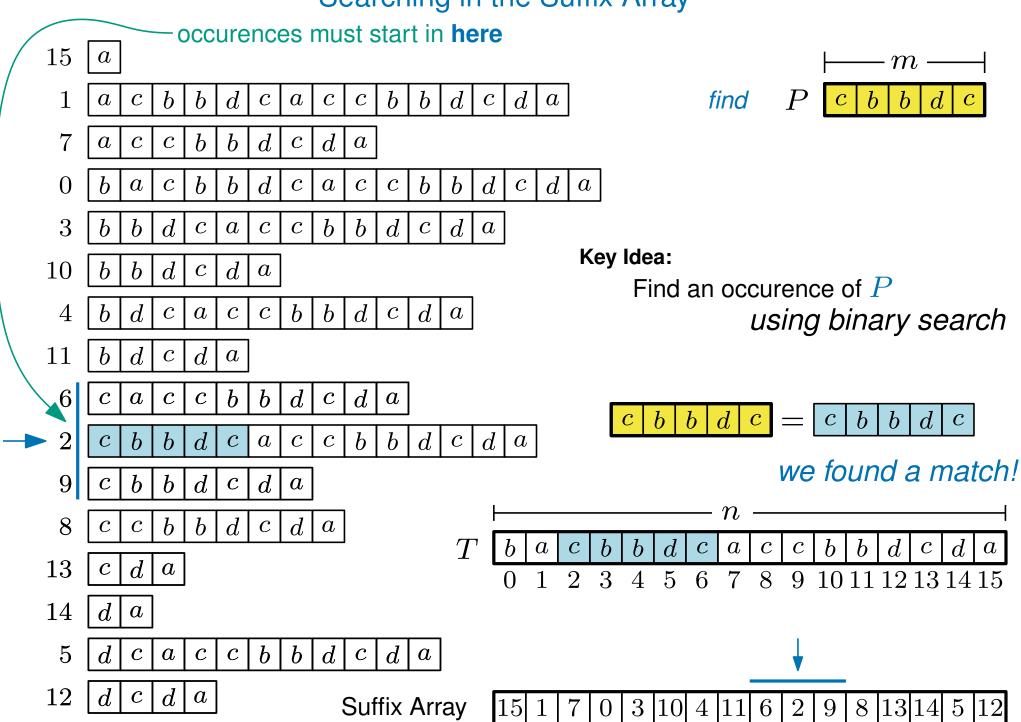




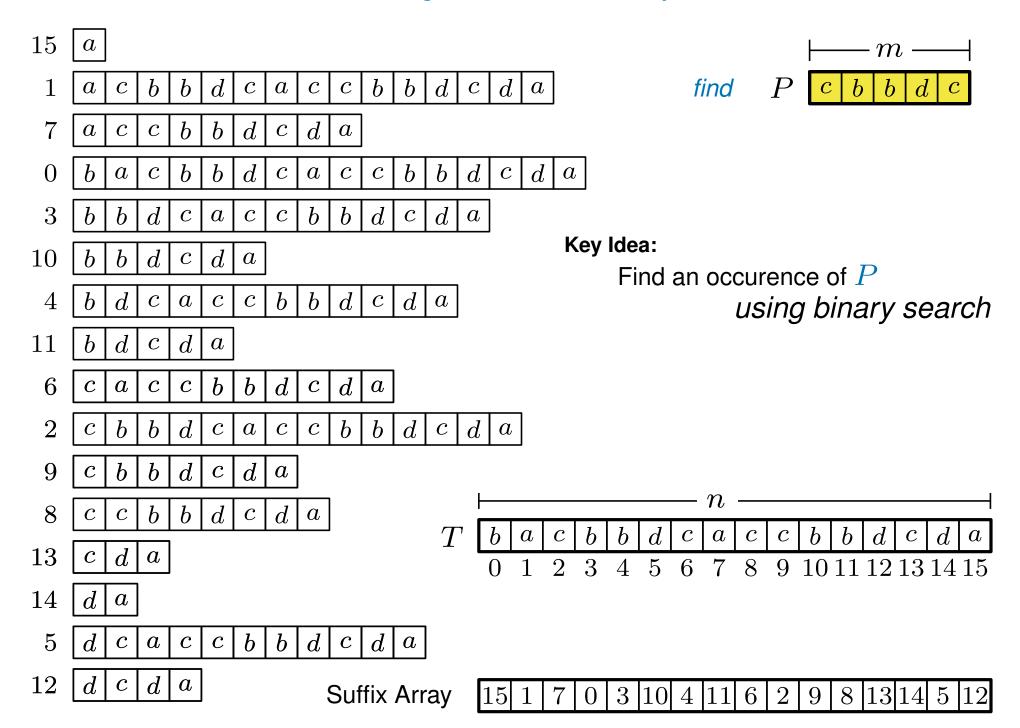




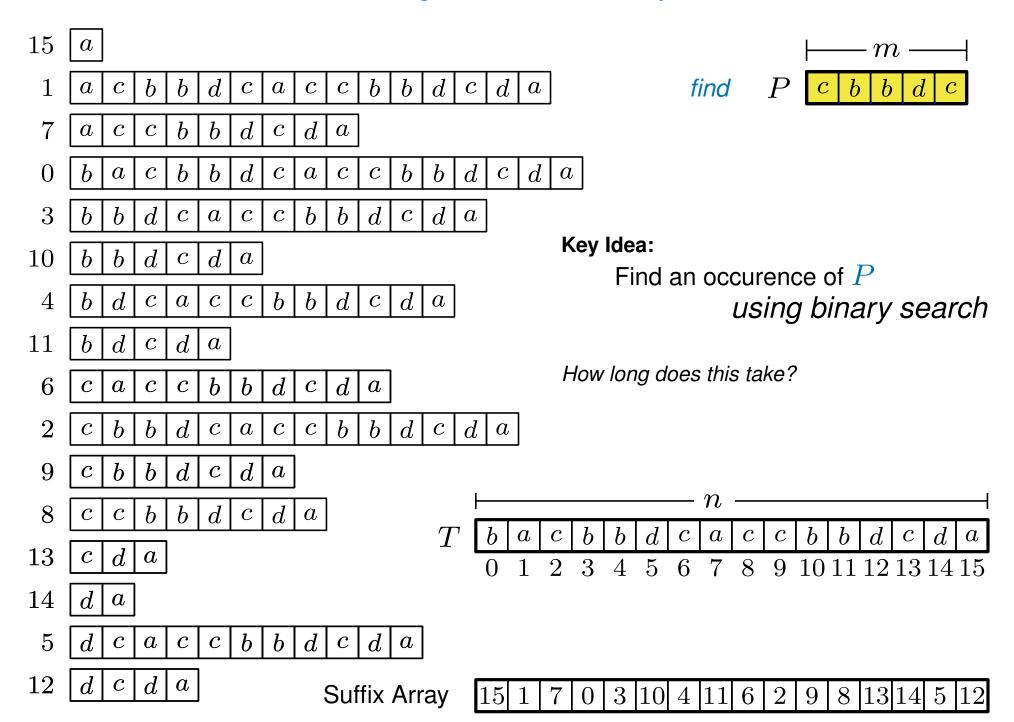




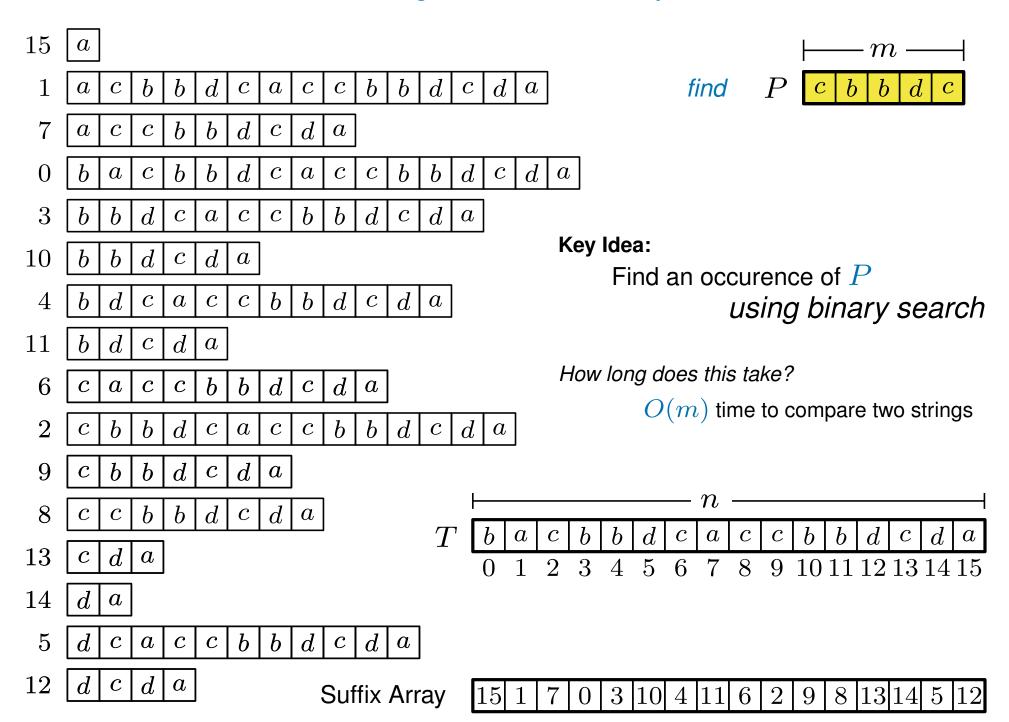




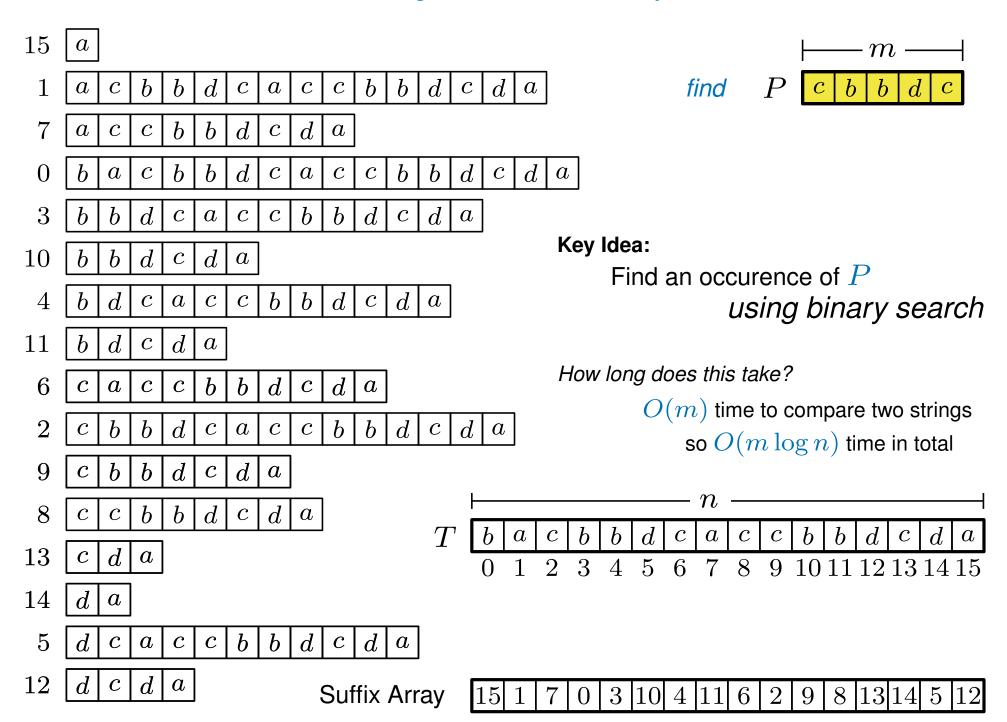




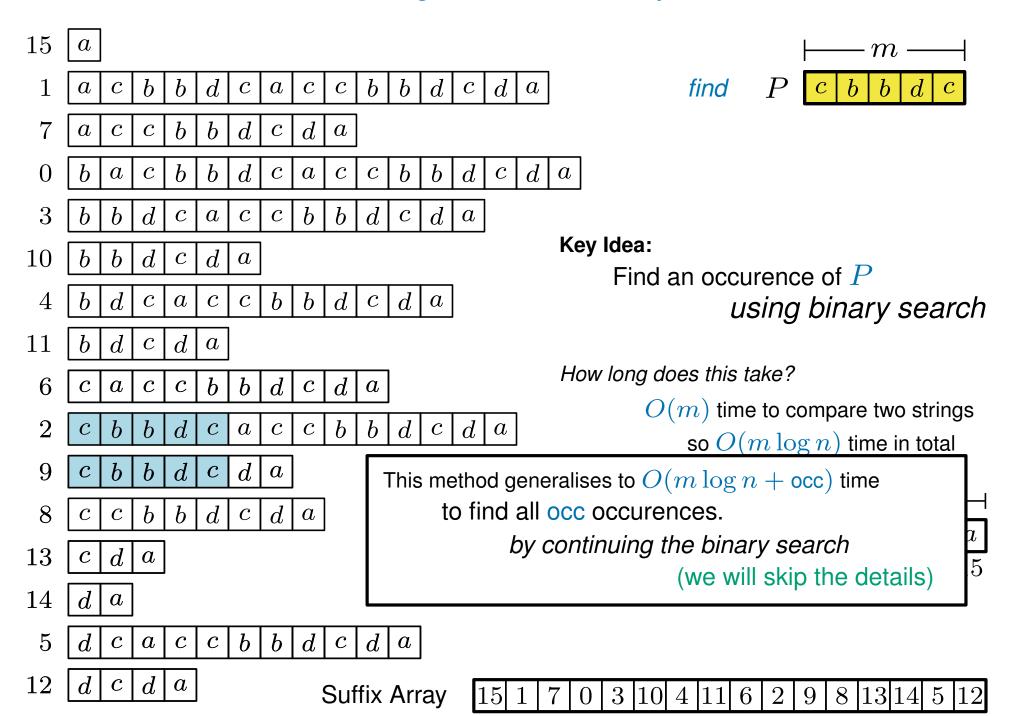






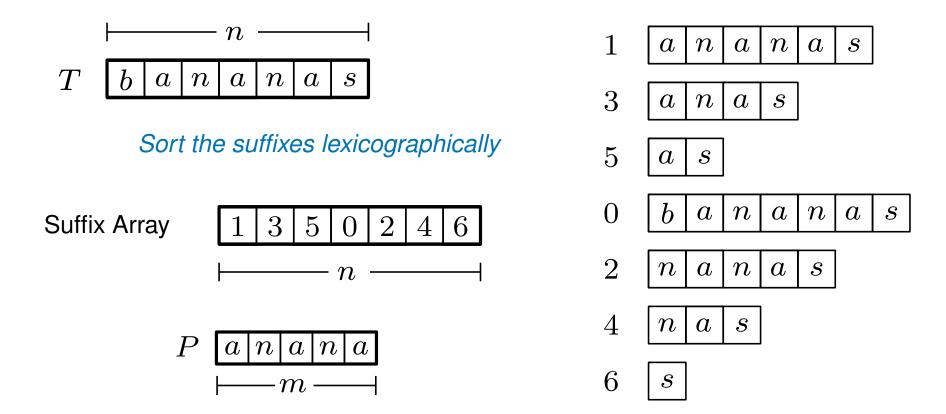








The suffix array



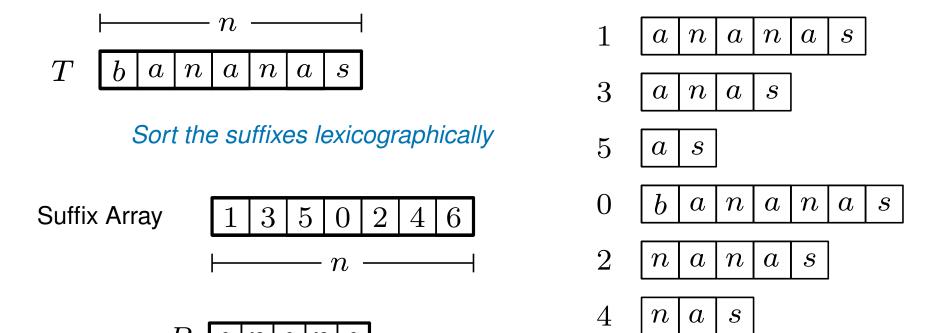
Finding an occurrence of a pattern (length m) takes $O(m \log n)$ time

Finding all occurrences takes $O(m \log n + occ)$ time

where occ is the number of occurences



The suffix array



Finding an occurrence of a pattern (length m) takes $O(m \log n)$ time

- $\,m$ -

Finding all occurrences takes $O(m \log n + occ)$ time

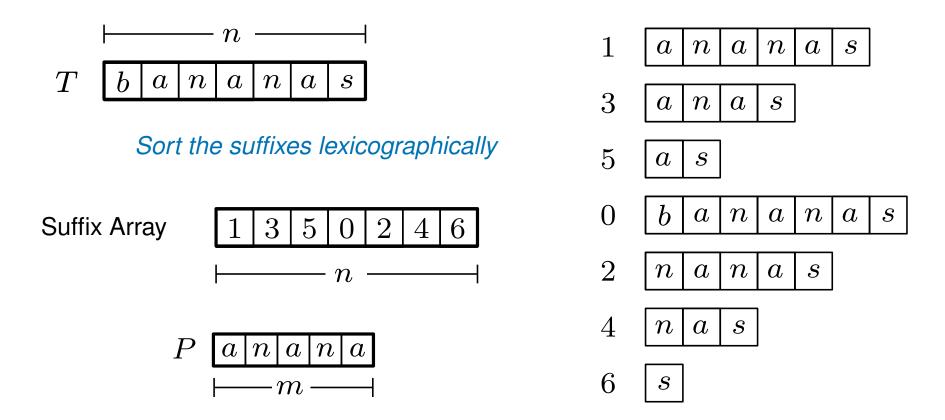
6

where occ is the number of occurences

This can be further improved to $O(m + \log n + \operatorname{occ})$ time (using LCP queries which we will see in a future lecture)



The suffix array



Finding an occurrence of a pattern (length m) takes $O(m \log n)$ time

Finding all occurrences takes $O(m \log n + occ)$ time

where occ is the number of occurences

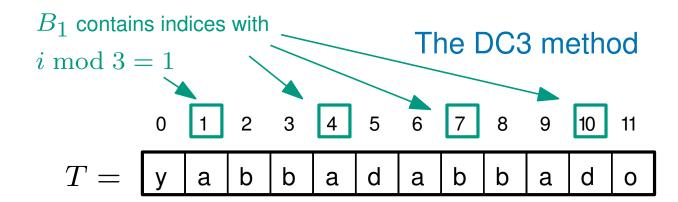
This can be further improved to $O(m + \log n + \operatorname{occ})$ time (using LCP queries which we will see in a future lecture)

Do we really need to build the suffix tree to construct the suffix array?



The DC3 method

			2									
T =	У	а	b	b	а	d	а	b	b	а	d	0





The DC3 method

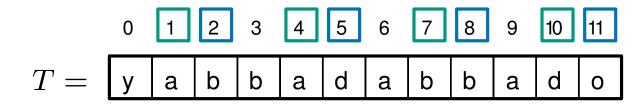
 $i \mod 3 = 1$



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 B_2 contains indices with

 $i \mod 3 = 2$





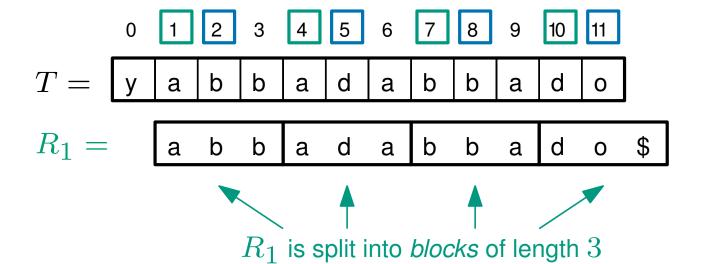
 B_2 contains indices with $i \bmod 3 = 2$



$$0$$
 1 2 3 4 5 6 7 8 9 10 11 $T=\begin{bmatrix}y&a&b&b&a&d&a&b&b&a&d&o\end{bmatrix}$

 B_2 contains indices with $i \bmod 3 = 2$





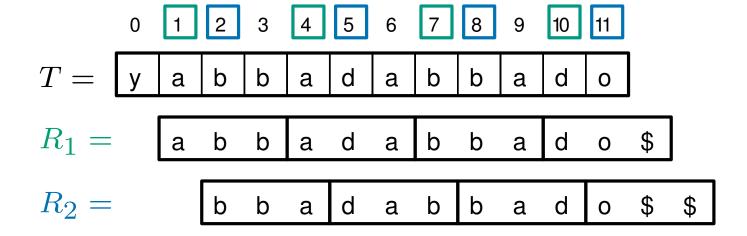
 B_2 contains indices with $i \bmod 3 = 2$



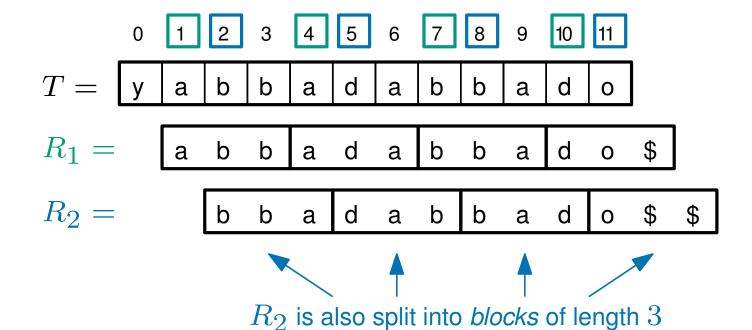
$$0$$
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 B_2 contains indices with $i \bmod 3 = 2$



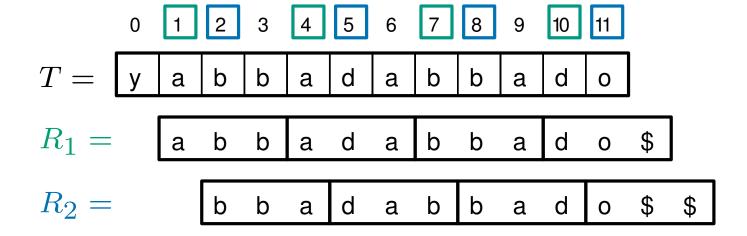






 B_2 contains indices with $i \bmod 3 = 2$





 $i \mod 3 = 2$

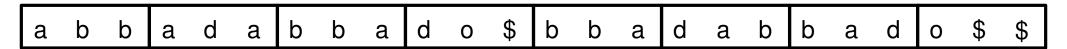


9 b b b b d d a 0 d b b d b b a a a 0 $R_2 =$ d b b d b a a a

Introduce a new

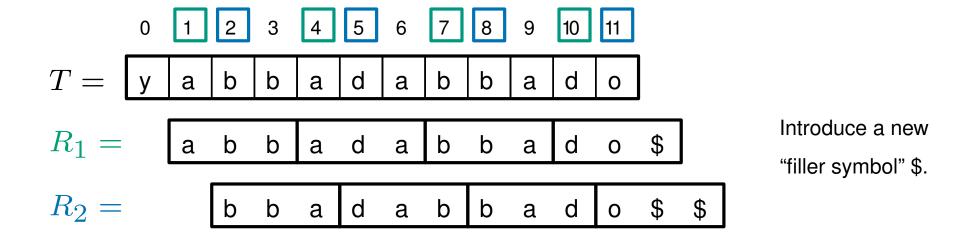
"filler symbol" \$.

Concatenate R_1 and R_2 to obtain R:

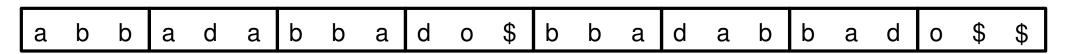


 B_2 contains indices with $i \mod 3 = 2$





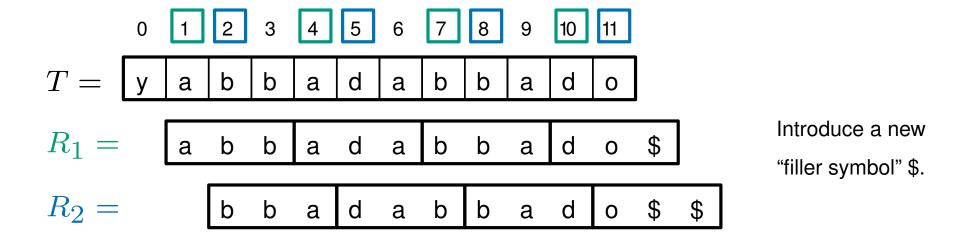
Concatenate R_1 and R_2 to obtain R:



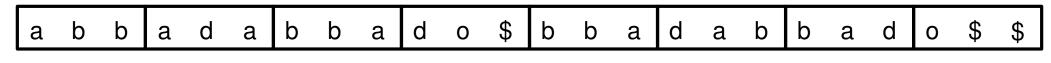
Number the **blocks** in lexicographical order

 B_2 contains indices with $i \mod 3 = 2$





Concatenate R_1 and R_2 to obtain R:



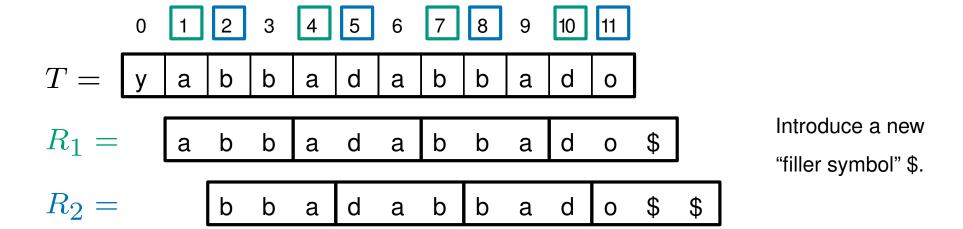
1

Number the **blocks** in lexicographical order

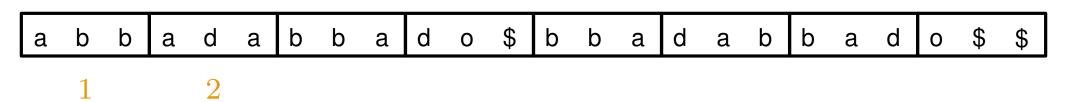
 B_2 contains indices with

 $i \mod 3 = 2$





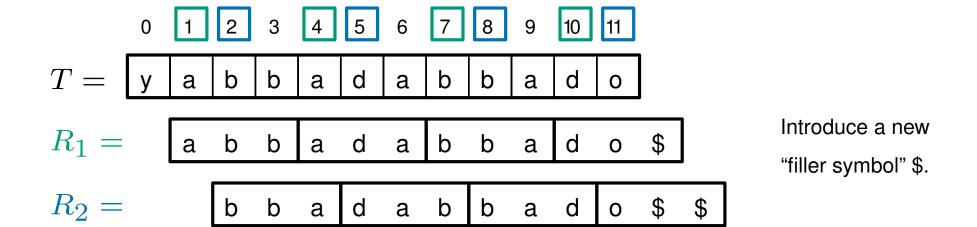
Concatenate R_1 and R_2 to obtain R:



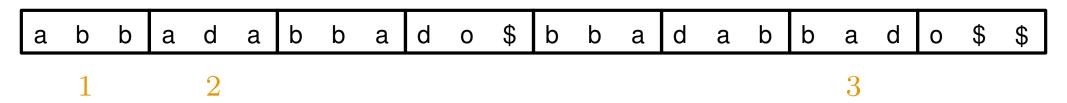
Number the **blocks** in lexicographical order

 B_2 contains indices with $i \mod 3 = 2$





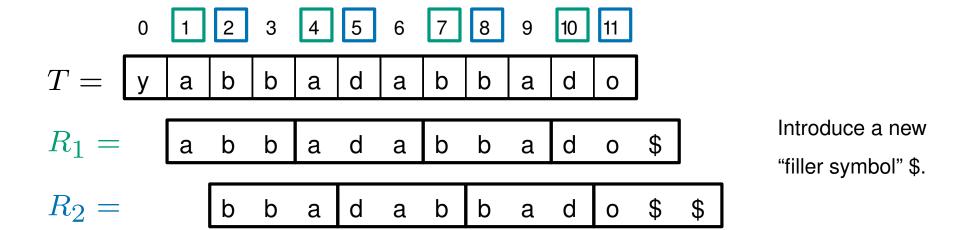
Concatenate R_1 and R_2 to obtain R:



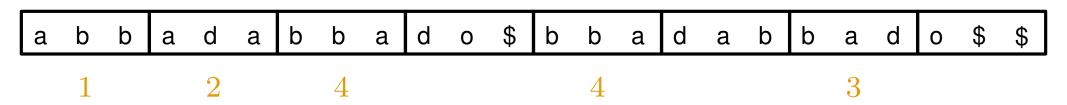
Number the **blocks** in lexicographical order

 B_2 contains indices with $i \mod 3 = 2$





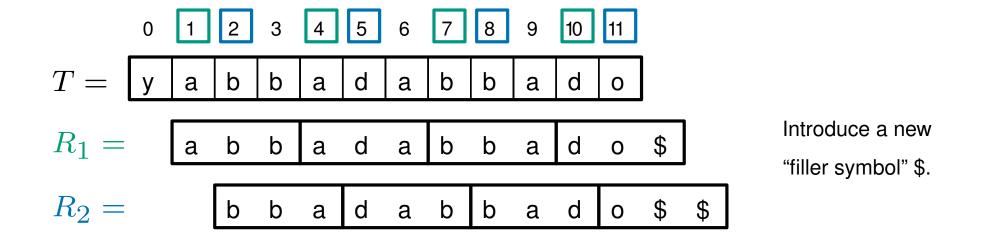
Concatenate R_1 and R_2 to obtain R:



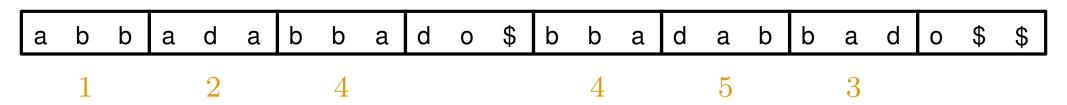
Number the **blocks** in lexicographical order

 B_2 contains indices with $i \mod 3 = 2$





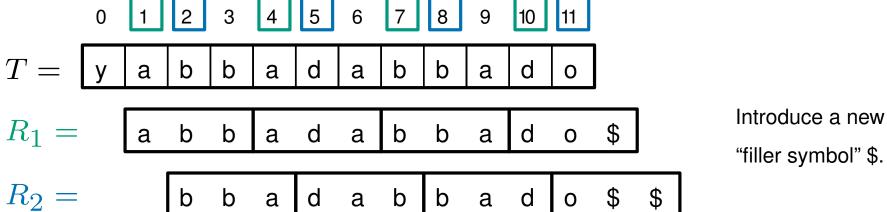
Concatenate R_1 and R_2 to obtain R:



Number the **blocks** in lexicographical order

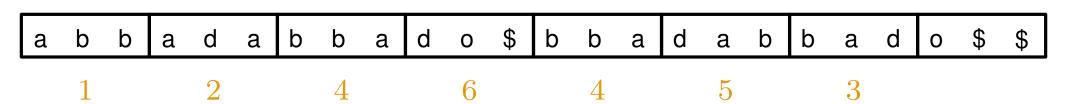
 B_2 contains indices with $i \mod 3 = 2$





Introduce a new

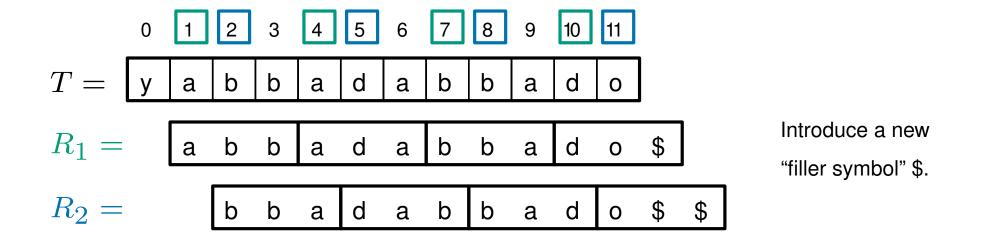
Concatenate R_1 and R_2 to obtain R:



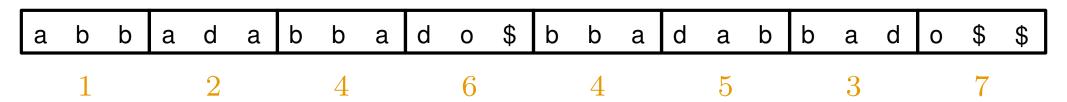
Number the blocks in lexicographical order

 B_2 contains indices with $i \mod 3 = 2$



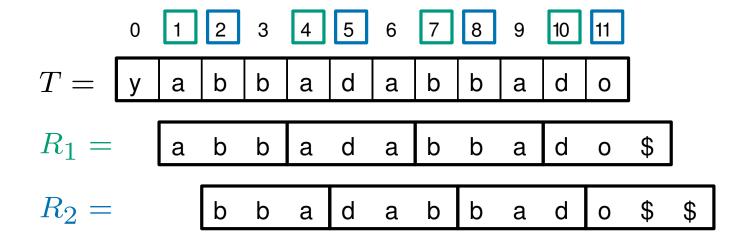


Concatenate R_1 and R_2 to obtain R:



Number the **blocks** in lexicographical order

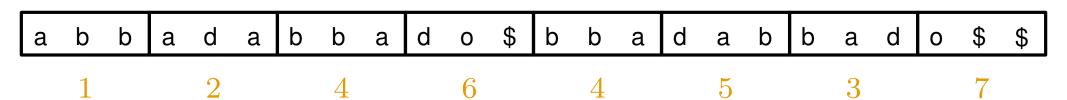
$$i \mod 3 = 2$$



Introduce a new

"filler symbol" \$.

Concatenate R_1 and R_2 to obtain R:



Number the **blocks** in lexicographical order

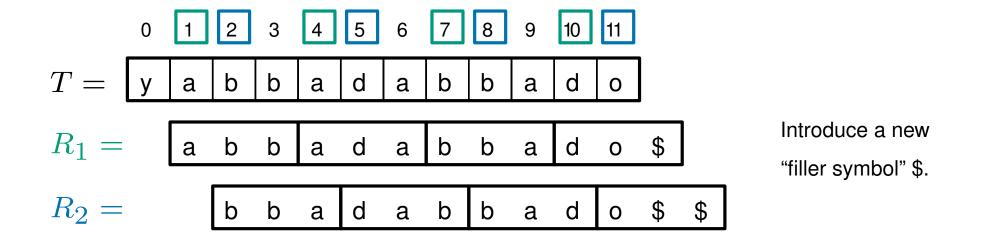
(\$ is the smallest symbol)

This can be done by sorting the blocks in O(n) time using radix sort we assume that the bit representation of each symbol uses $O(\log n)$ bits. (which is a common and realistic assumption)

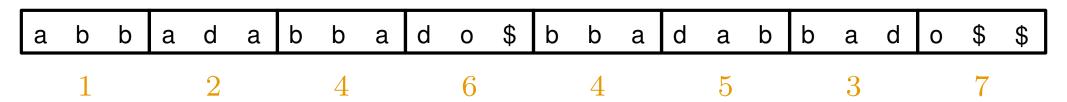


 B_2 contains indices with $i \mod 3 = 2$





Concatenate R_1 and R_2 to obtain R:

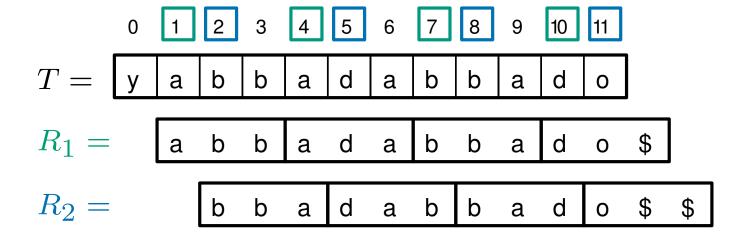


Number the **blocks** in lexicographical order

 $i \mod 3 = 2$



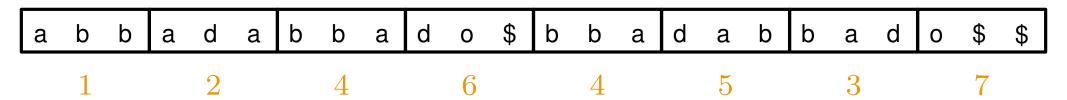
$$i \mod 3 = 1$$



Introduce a new

"filler symbol" \$.

Concatenate R_1 and R_2 to obtain R:



Number the blocks in lexicographical order

let
$$R' = \begin{bmatrix} 1 & 2 & 4 & 6 & 4 & 5 & 3 & 7 \end{bmatrix}$$

 $i \mod 3 = 2$

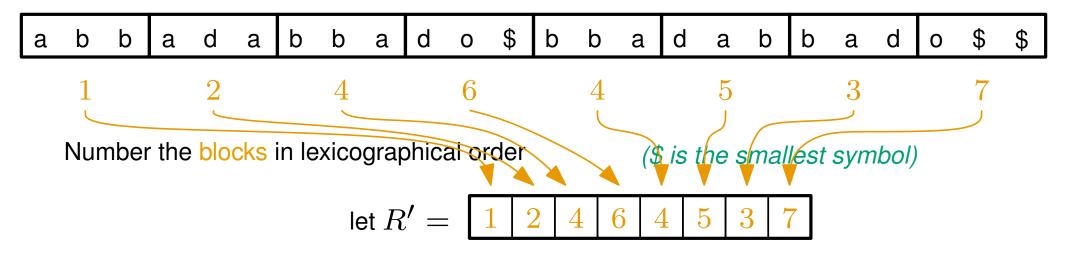


$$i \mod 3 = 1$$

Introduce a new

"filler symbol" \$.

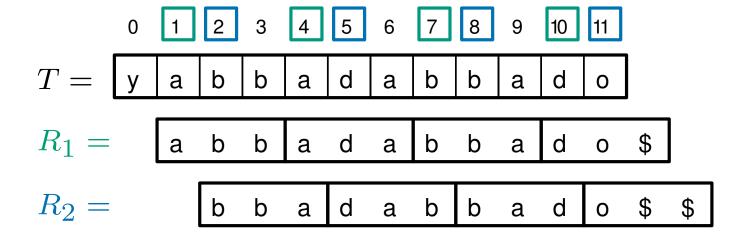
Concatenate R_1 and R_2 to obtain R:



 $i \mod 3 = 2$



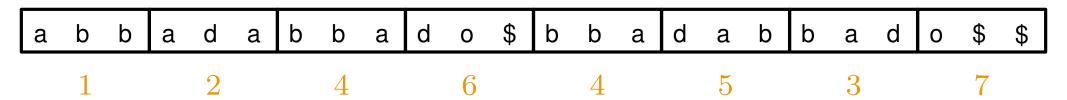
$$i \mod 3 = 1$$



Introduce a new

"filler symbol" \$.

Concatenate R_1 and R_2 to obtain R:



Number the blocks in lexicographical order

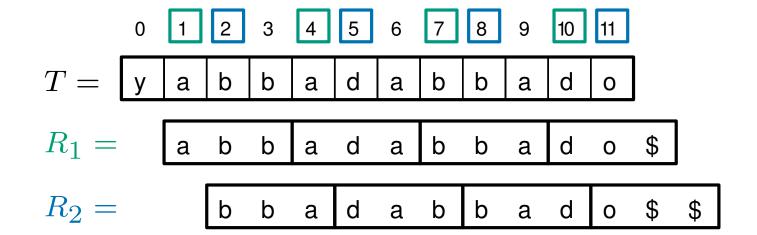
let
$$R' = \begin{bmatrix} 1 & 2 & 4 & 6 & 4 & 5 & 3 & 7 \end{bmatrix}$$

 B_2 contains indices with

 $i \mod 3 = 2$



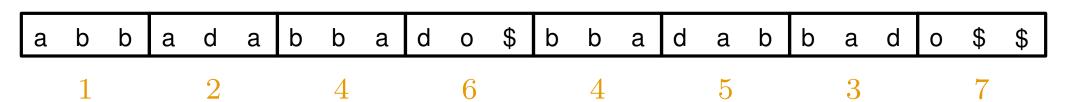
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Introduce a new

"filler symbol" \$.

Concatenate R_1 and R_2 to obtain R:



Number the **blocks** in lexicographical order

(\$ is the smallest symbol)

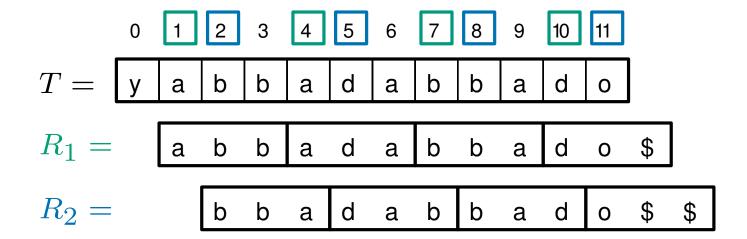
compute the suffix array of R':

 B_2 contains indices with

 $i \mod 3 = 2$



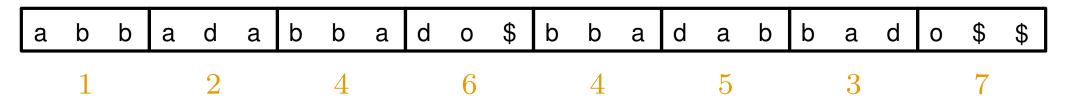
$$i \mod 3 = 1$$



Introduce a new

"filler symbol" \$.

Concatenate R_1 and R_2 to obtain R:



Number the blocks in lexicographical order

(\$ is the smallest symbol)

compute the suffix array of R':



 $R_2 =$

The DC3 method

 B_2 contains indices with

 $i \mod 3 = 2$



d

a

b

b

a

d

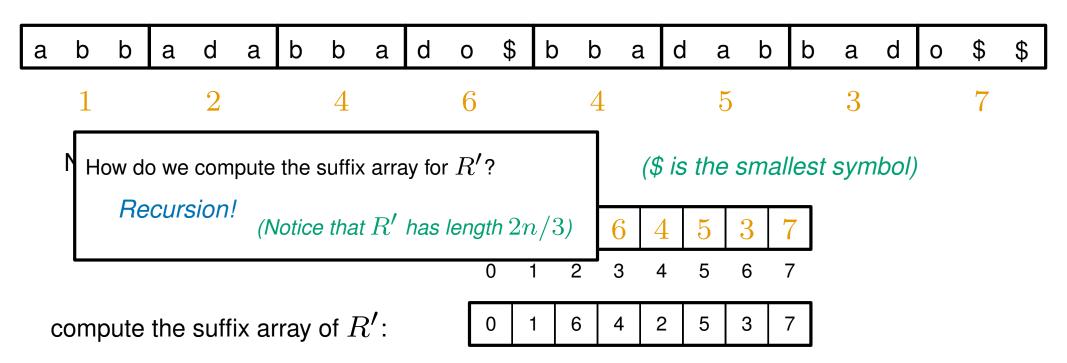
Introduce a new

"filler symbol" \$.

Concatenate R_1 and R_2 to obtain R:

b

a

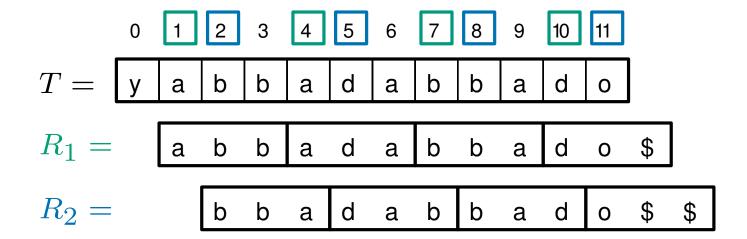


 B_2 contains indices with

 $i \mod 3 = 2$



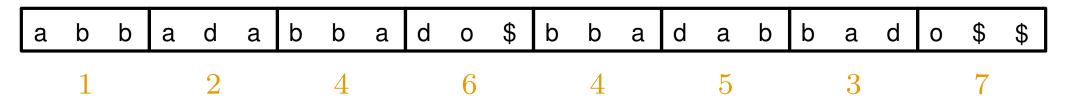
$$i \mod 3 = 1$$



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"filler symbol" \$.

Concatenate R_1 and R_2 to obtain R:



Number the blocks in lexicographical order

(\$ is the smallest symbol)

compute the suffix array of R':



 $R_2 =$

The DC3 method

 B_2 contains indices with

 $i \mod 3 = 2$



d

a

b

b

a

d

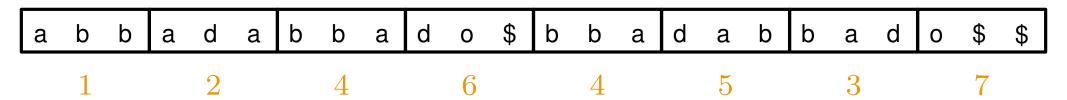
a

Introduce a new

"filler symbol" \$.

Concatenate R_1 and R_2 to obtain R:

b



Number the blocks in lexicographical order

(\$ is the smallest symbol)

let
$$R' = \begin{bmatrix} 1 & 2 & 4 & 6 & 4 & 5 & 3 & 7 \\ & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \end{bmatrix}$$

what use

compute the suffix array of R':

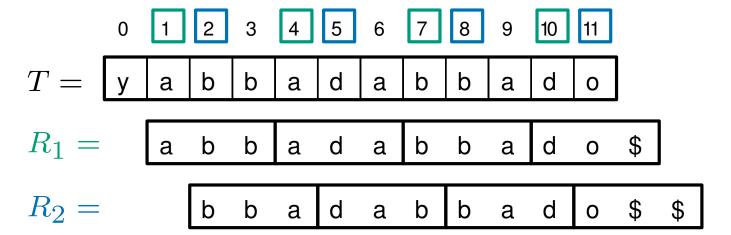
0 1 6 4 2 5 3 7

 B_2 contains indices with

 $i \mod 3 = 2$



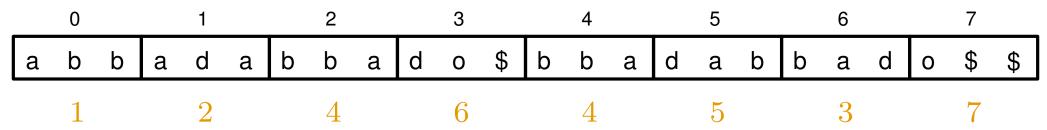
$$i \mod 3 = 1$$



Introduce a new

"filler symbol" \$.

Concatenate R_1 and R_2 to obtain R:



Number the blocks in lexicographical order

(\$ is the smallest symbol)

what use

compute the suffix array of R':

0 1 6 4 2 5 3 7

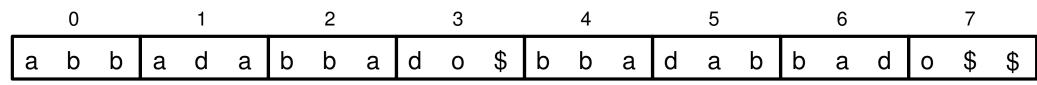
 B_2 contains indices with

 $i \mod 3 = 2$

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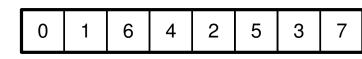
$$i \mod 3 = 1$$

Concatenate R_1 and R_2 to obtain R:



what use

compute the suffix array of R':



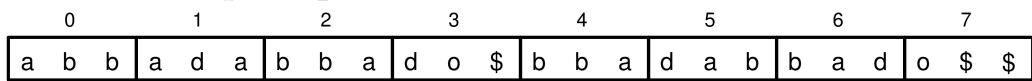
 B_2 contains indices with

 $i \mod 3 = 2$

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 $i \mod 3 = 1$

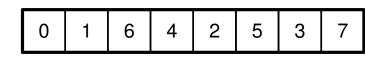
Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$

what use

compute the suffix array of R':



 $i \mod 3 = 1$

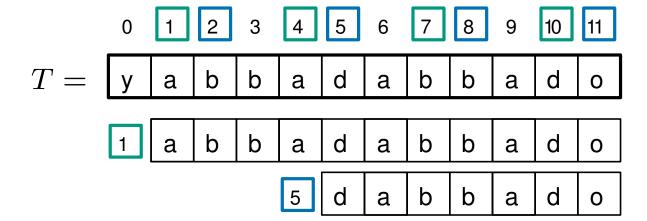
The DC3 method

 B_2 contains indices with

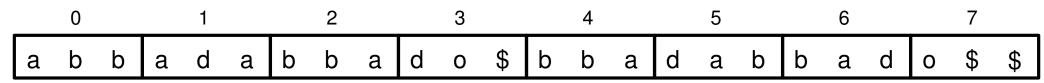
 $i \mod 3 = 2$

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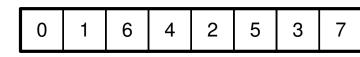
Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$

what use

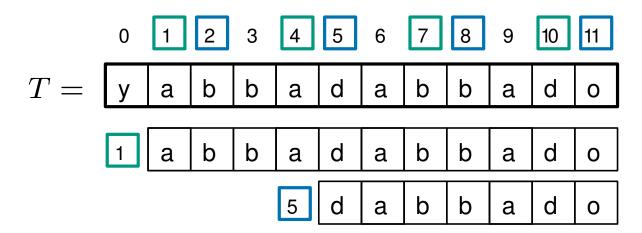
compute the suffix array of R':



 B_2 contains indices with

 $i \mod 3 = 2$



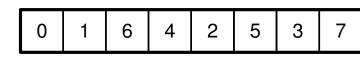


Concatenate R_1 and R_2 to obtain R:

Take any two suffixes in $B_1 \cup B_2$ and find them in R

what use

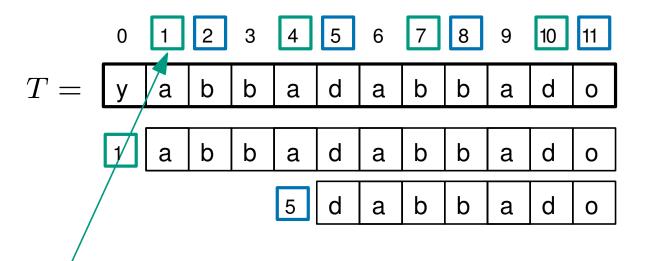
compute the suffix array of R':



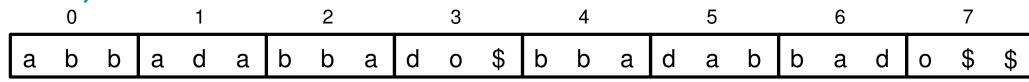
 B_2 contains indices with

 $i \mod 3 = 2$





Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$ and find them in R

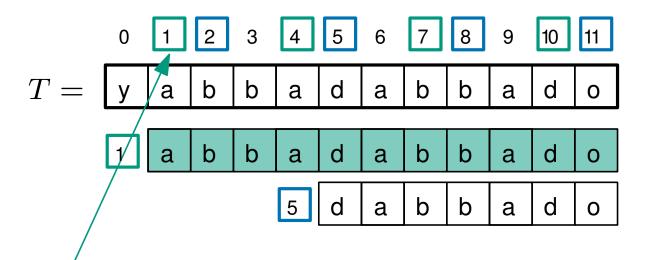
what use

compute the suffix array of R':

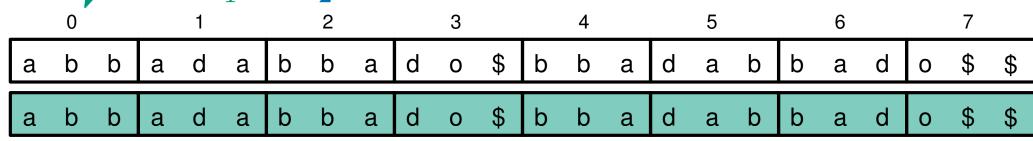
 B_2 contains indices with







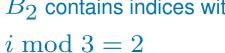
Concatenate R_1 and R_2 to obtain R:



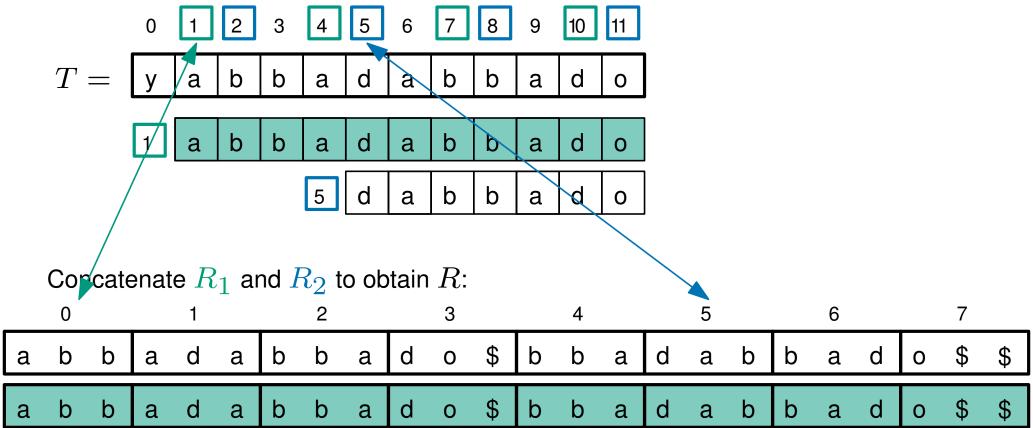
Take any two suffixes in $B_1 \cup B_2$ and find them in R

what use

compute the suffix array of R':



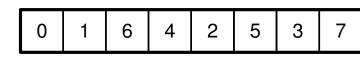
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Take any two suffixes in $B_1 \cup B_2$ and find them in R

what use

compute the suffix array of R':



Take any two suffixes in $B_1 \cup B_2$ and find them in R

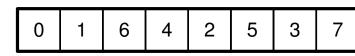
what use

a

b

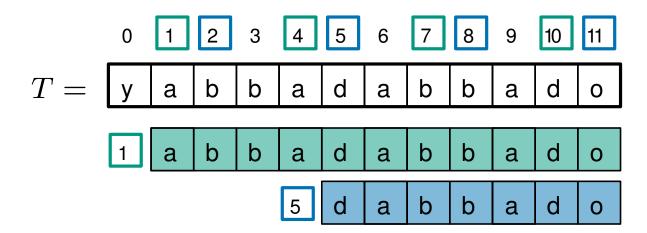
a

compute the suffix array of R':

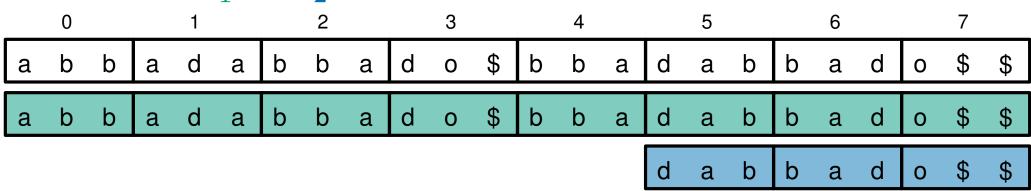


 $i \mod 3 = 2$





Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

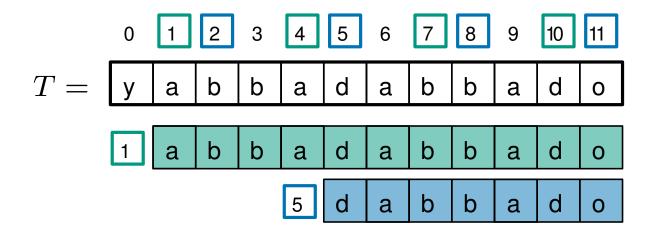
what use

compute the suffix array of R':

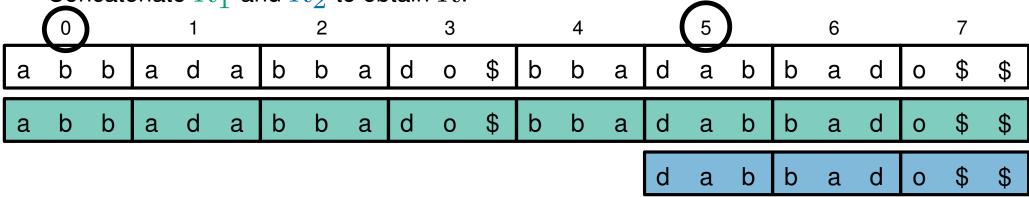
0 1 6 4 2 5 3	7
---------------	---







Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$ and find them in R

their order is given by the suffix array of R':

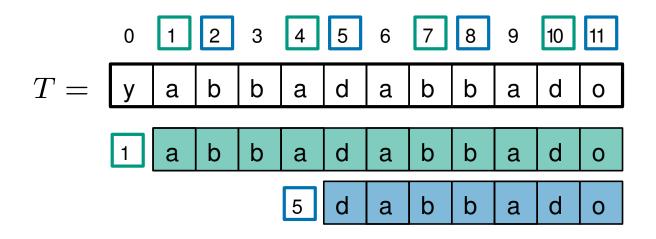
what use

compute the suffix array of R':

|--|

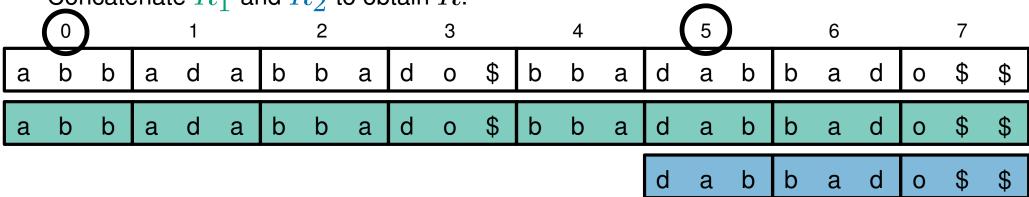
 B_2 contains indices with $i \mod 3 = 2$





Suffix 1 is smaller than suffix 5

Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

what use

compute the suffix array of R':

0 1 6 4 2 5 3 7

 B_2 contains indices with

 $i \mod 3 = 2$

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$$i \mod 3 = 1$$

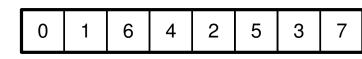
Concatenate R_1 and R_2 to obtain R:

	0			1			2			3		4			5			6			7	
а	b	b	а	d	а	b	b	а	d	0	\$ b	b	а	d	а	b	b	а	d	0	\$	\$

Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

what use

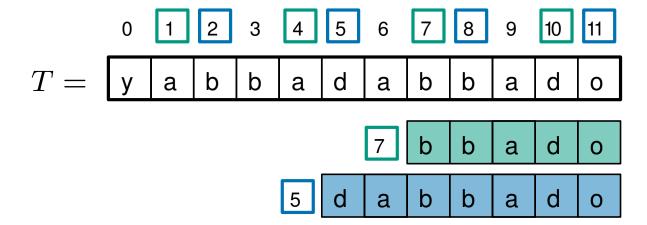
compute the suffix array of R':



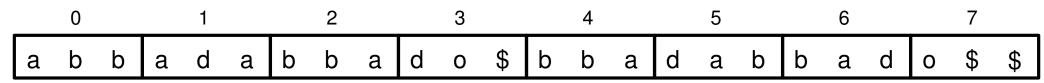
 B_2 contains indices with

 $i \mod 3 = 2$





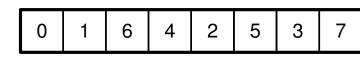
Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

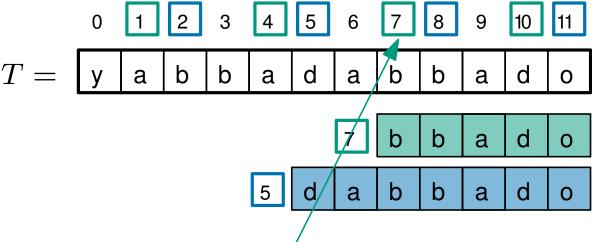
what use

compute the suffix array of R':

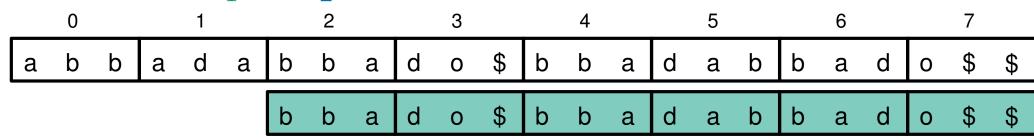


 B_2 contains indices with

 $i \mod 3 = 2$



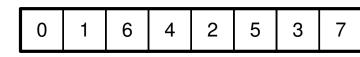
Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

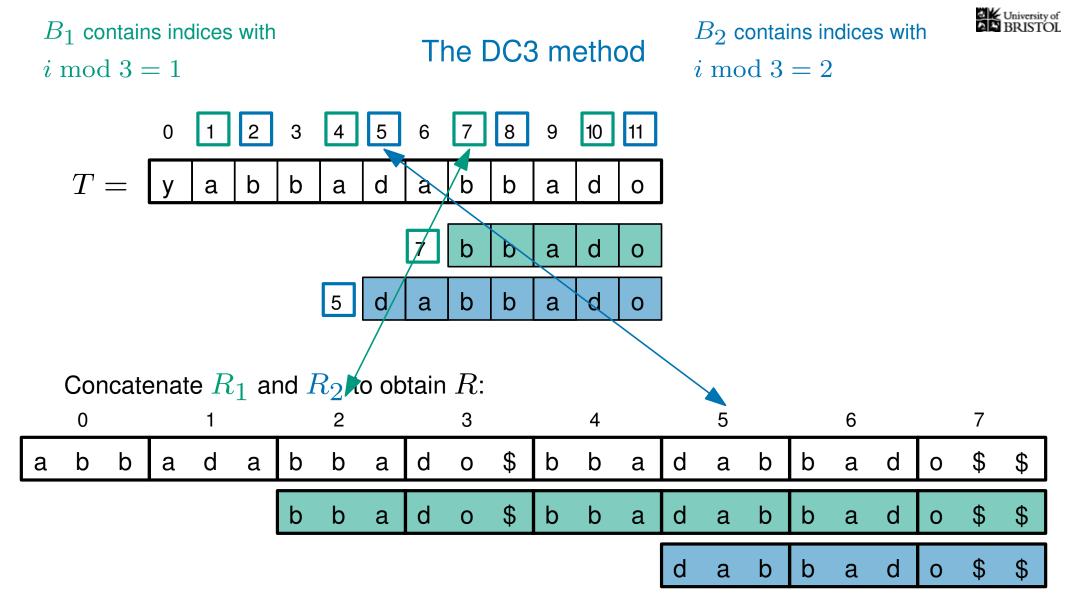
what use

compute the suffix array of R':



is this?!

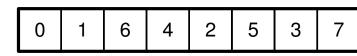
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Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

what use

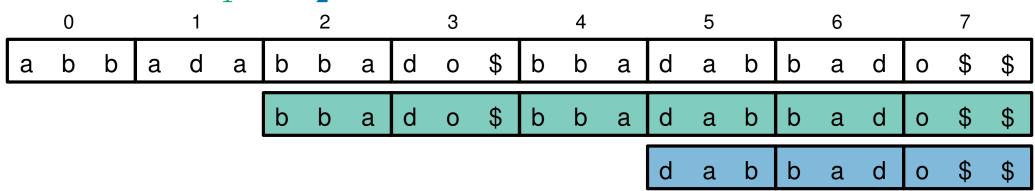
compute the suffix array of R':



 $i \mod 3 = 2$



Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

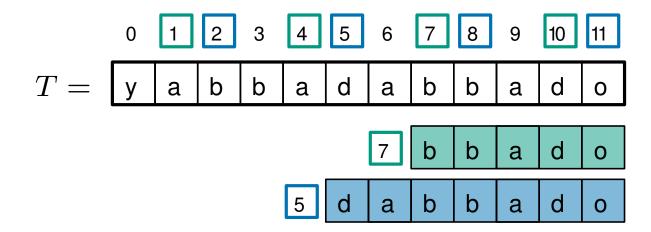
what use

compute the suffix array of R^{\prime} :

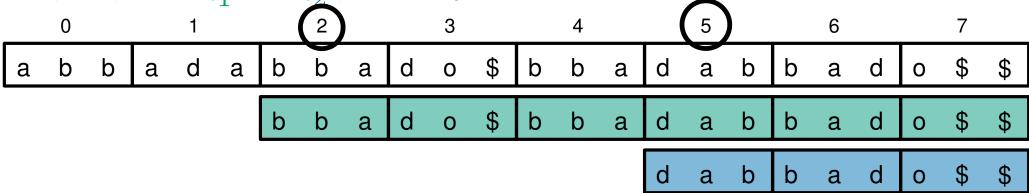
0 1 6 4 2 5 3	7
---------------	---

 $i \mod 3 = 2$





Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$ and find them in R

their order is given by the suffix array of R':

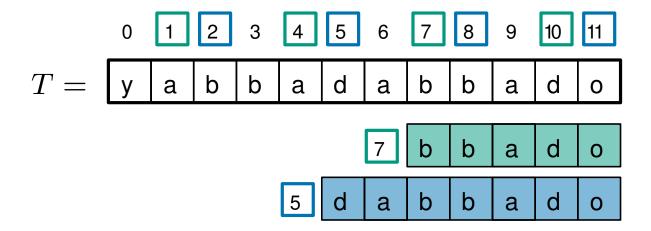
what use

compute the suffix array of R^{\prime} :

0 1 6 4 2 5 3 7

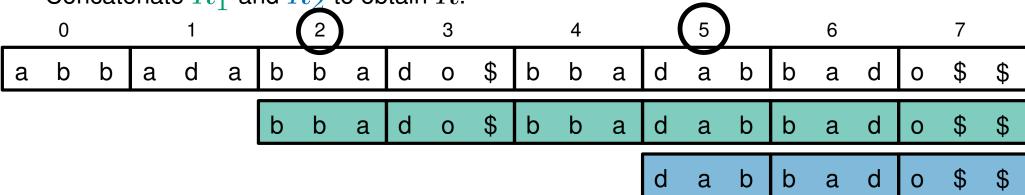
 B_2 contains indices with $i \mod 3 = 2$





Suffix 7 is smaller than suffix 5

Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

what use

compute the suffix array of R':

0 1 6 4 2 5 3 7

 B_2 contains indices with

 $i \mod 3 = 2$

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$$i \mod 3 = 1$$

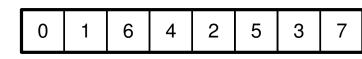
Concatenate R_1 and R_2 to obtain R:

	0			1			2			3		4			5			6			7	
а	b	b	а	d	а	b	b	а	d	0	\$ b	b	а	d	а	b	b	а	d	0	\$	\$

Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

what use

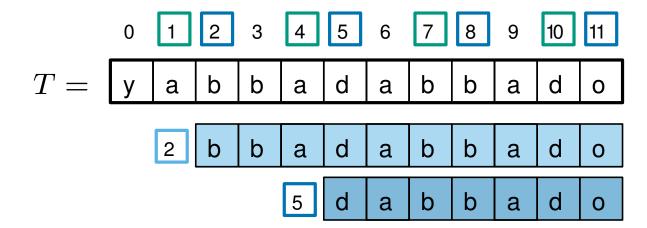
compute the suffix array of R':



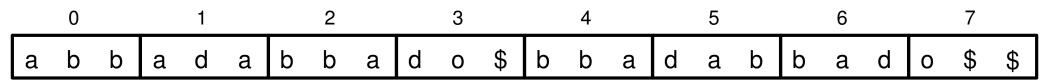
 B_2 contains indices with

 $i \mod 3 = 2$





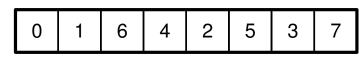
Concatenate R_1 and R_2 to obtain R:

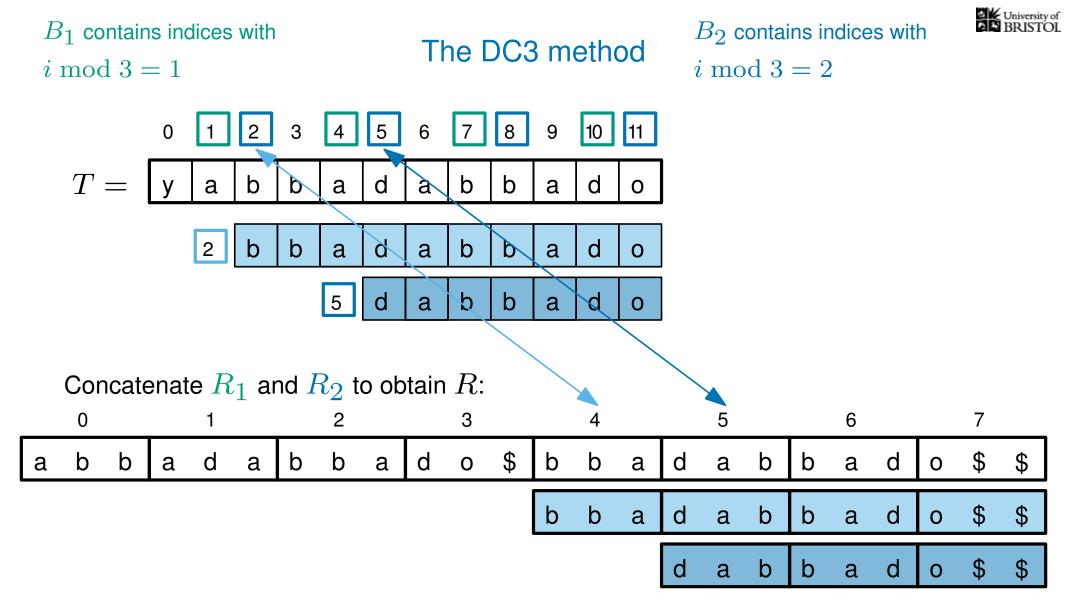


Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

what use

compute the suffix array of R':

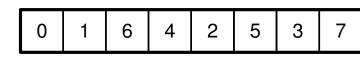




Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

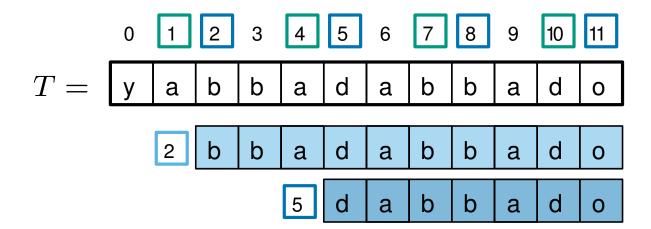
what use

compute the suffix array of R':

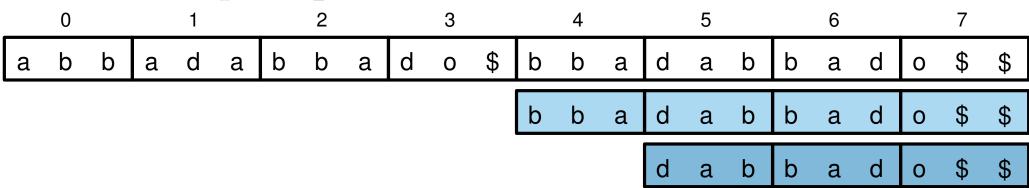


 $i \mod 3 = 2$





Concatenate R_1 and R_2 to obtain R:



Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

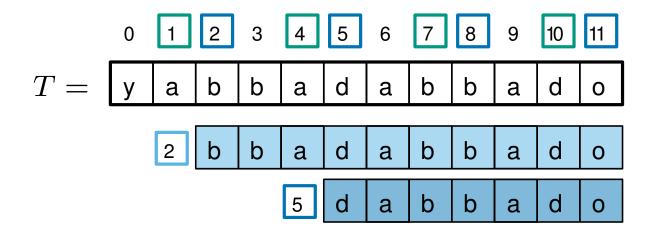
what use

compute the suffix array of R':

0 1 6 4	2	5	3	7
---------	---	---	---	---

 $i \mod 3 = 2$





Concatenate R_1 and R_2 to obtain R: 6 b b a d b b d a 0 b b b d b b d a a a 0 b a a

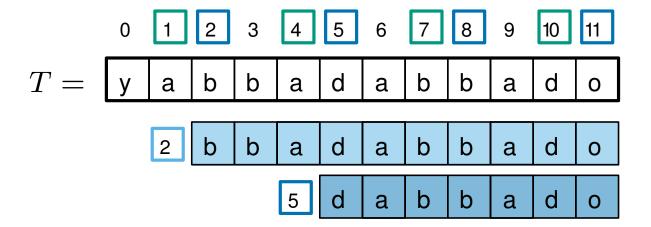
Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

what use

compute the suffix array of R':

0 1 6 4 2 5 3 7





Suffix 2 is smaller than suffix 5

Concatenate R_1 and R_2 to obtain R: b b a d b b d a 0 b b b d b b d a a a 0 b a a

Take any two suffixes in $B_1 \cup B_2$ and find them in R their order is given by the suffix array of R':

what use

compute the suffix array of R':

0 1 6 4 2 5 3 7

is this?!

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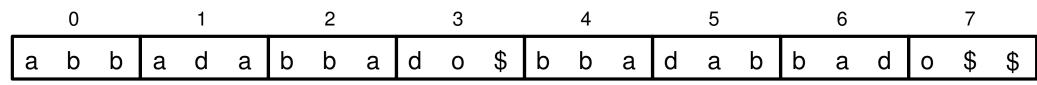
 B_2 contains indices with

 $i \mod 3 = 2$

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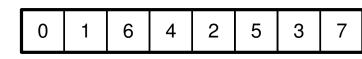
$$i \mod 3 = 1$$

Concatenate R_1 and R_2 to obtain R:



what use

compute the suffix array of R':

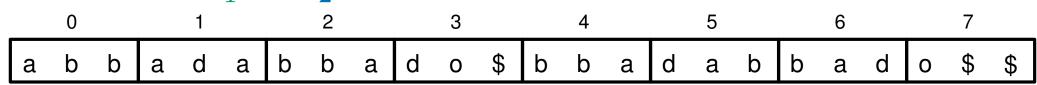


 B_2 contains indices with

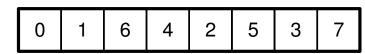
 $i \mod 3 = 2$



Concatenate R_1 and R_2 to obtain R:



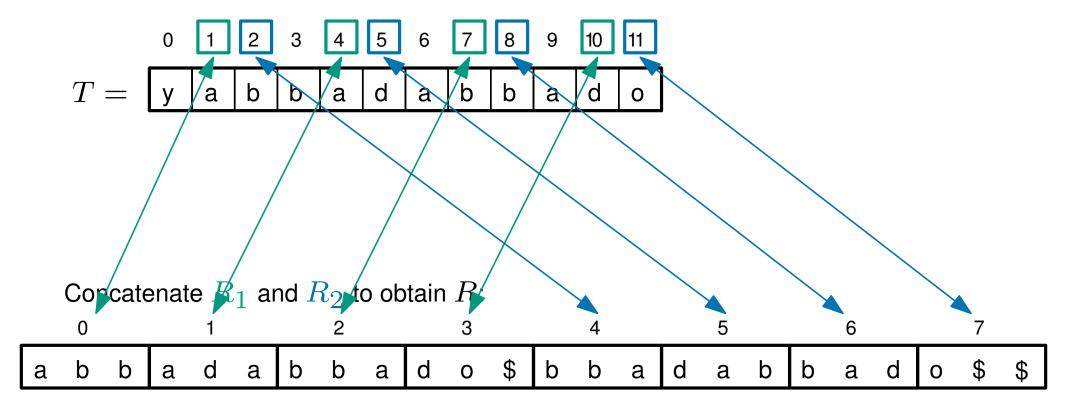
The suffix array of R':



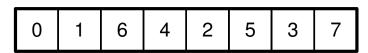








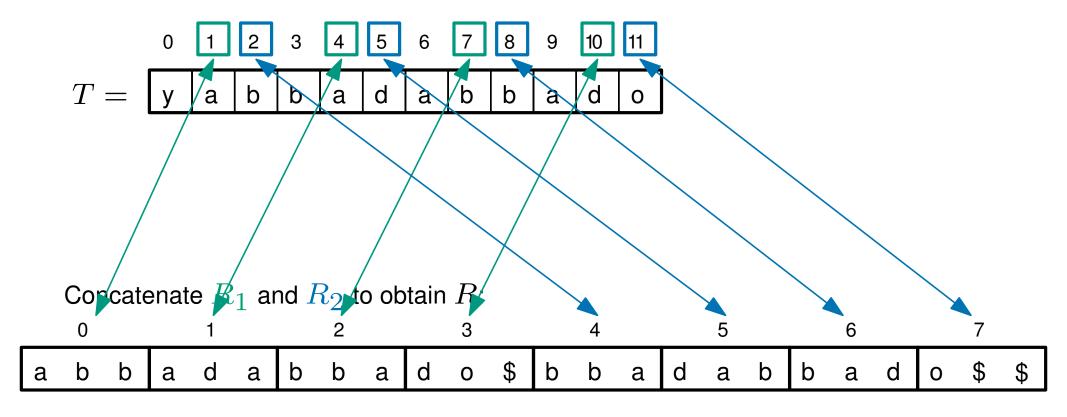
The suffix array of R':



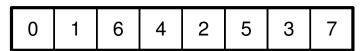
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 $i \mod 3 = 1$

 $i \mod 3 = 2$



The suffix array of R':

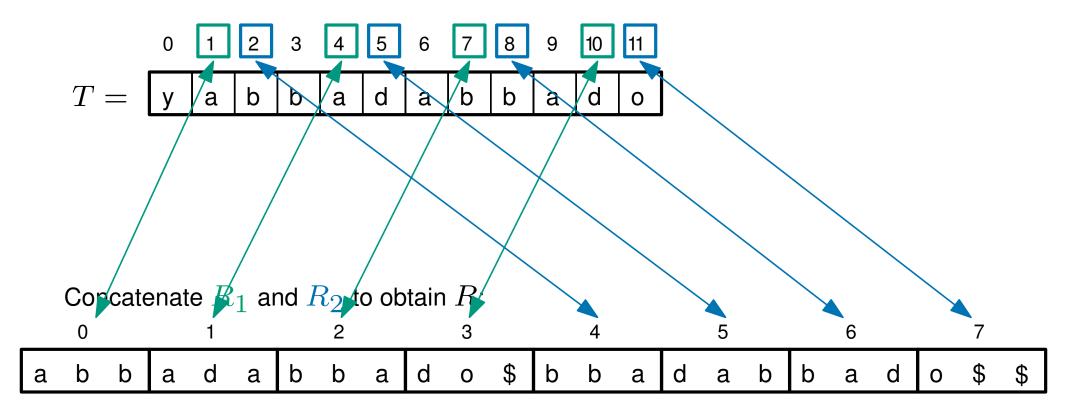


after relabelling,

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 $i \mod 3 = 1$

 $i \bmod 3 = 2$



The suffix array of R':

0 1 6 4 2 5 3 7

after relabelling,

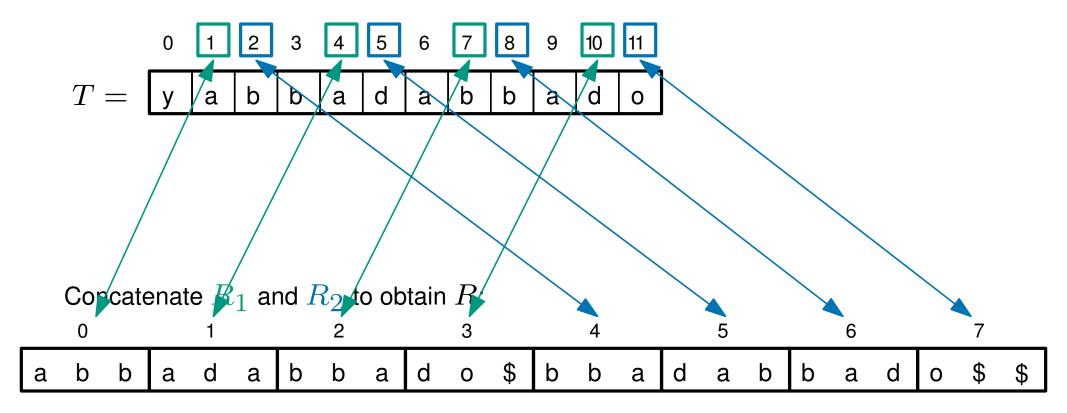
1 4 8 2 7 5 10 11

 B_2 contains indices with





 $i \mod 3 = 2$



The suffix array of R':

0
1
6
4
2
5
3
7

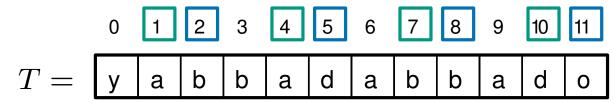
after relabelling,

1
4
8
2
7
5
10
11

we have the suffix array of just the suffixes from $B_1 \cup B_2$

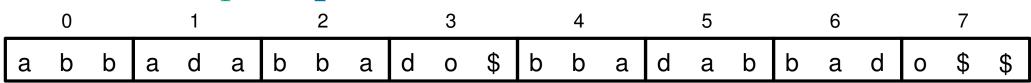
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 $i \mod 3 = 1$



 $i \mod 3 = 2$

Concatenate R_1 and R_2 to obtain R:



The suffix array of R':

0 1 6 4 2 5 3 7

after relabelling,

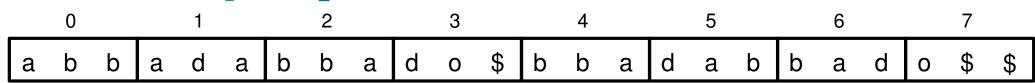
1 4 8 2 7 5 10 11

we have the suffix array of just the suffixes from $B_1 \cup B_2$



$$i \mod 3 = 1$$

Concatenate R_1 and R_2 to obtain R:



we have the suffix array of just the suffixes from $B_1 \cup B_2$

 B_2 contains indices with $i \mod 3 = 2$



 B_2 contains indices with $i \mod 3 = 2$



Suffix array for just $B_1 \cup B_2$

1 4 8 2 7 5 10 11

 B_2 contains indices with $i \mod 3 = 2$

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Suffix array for just $B_1 \cup B_2$



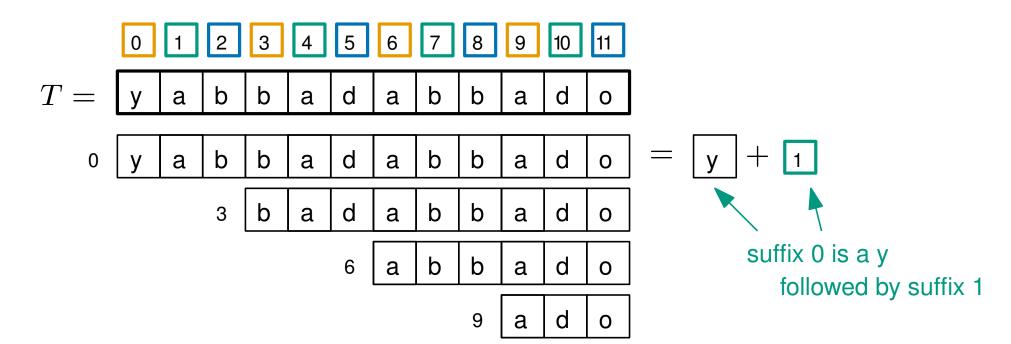
Suffix array for just $B_1 \cup B_2$

0	1	2	3	4	5	6	7
1	4	8	2	7	5	10	11



Suffix array for just $B_1 \cup B_2$





Suffix array for just $B_1 \cup B_2$



Suffix array for just $B_1 \cup B_2$



$$i \mod 3 = 1$$

Suffix array for just $B_1 \cup B_2$



Each suffix $i \in B_0$ is represented by (T[i], r) where r is the rank of suffix (i + 1) (the ranks are given by the array below)

rank: 0 1 2 3 4 5 6 7 Suffix array for just $B_1 \cup B_2$ 1 4 8 2 7 5 10 11



6

3 b a d a b b a d o = b + 4
$$(b,1)$$

Each suffix $i \in B_0$ is represented by (T[i], r) where r is the rank of suffix (i + 1) (the ranks are given by the array below)

rank: 0 1 2 3 4 5 6 7 Suffix array for just
$$B_1 \cup B_2$$
 1 4 8 2 7 5 10 11



Each suffix $i \in B_0$ is represented by (T[i], r) where r is the rank of suffix (i+1) (the ranks are given by the array below)

We then sort in O(n) time using radix sort



(b, 1)

(y, 0)

3
$$\begin{bmatrix} b & a & d & a & b & b & a & d & o \end{bmatrix} = \begin{bmatrix} b & + & 4 \end{bmatrix}$$

0
$$y$$
 a b b a d a b b a d o = y + 1

Each suffix $i \in B_0$ is represented by (T[i], r) where r is the rank of suffix (i + 1) (the ranks are given by the array below)

We then sort in O(n) time using radix sort

rank: 0 1 2 3 4 5 6 7

1 4 8 2 7 5 10 11

Suffix array for just $B_1 \cup B_2$

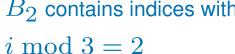


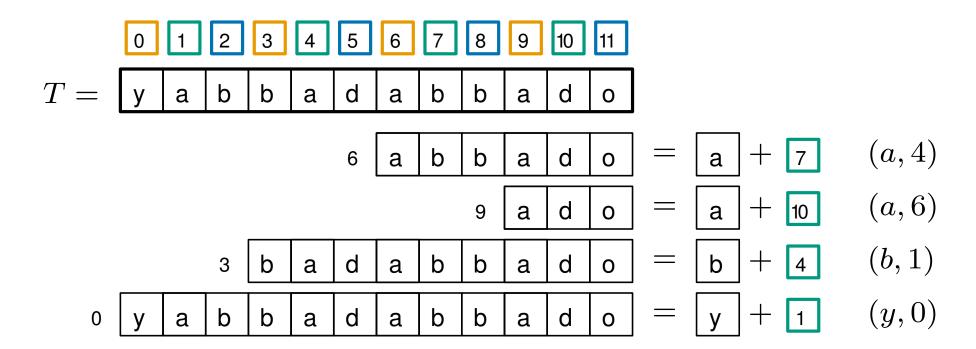
b b d b b d a a a 0 (a, 4)b 6 b d a (a, 6)10 9 a d 0 a (b, 1)3 b b d d a a a (y, 0)b b d b b d 0 0 a a a a

Each suffix $i \in B_0$ is represented by (T[i], r) where r is the rank of suffix (i + 1) (the ranks are given by the array below)

We then sort in O(n) time using radix sort

 B_2 contains indices with





Each suffix $i \in B_0$ is represented by (T[i], r) where r is the rank of suffix (i + 1)(the ranks are given by the array below)

We then sort in O(n) time using radix sort

rank: 0 5 7 6

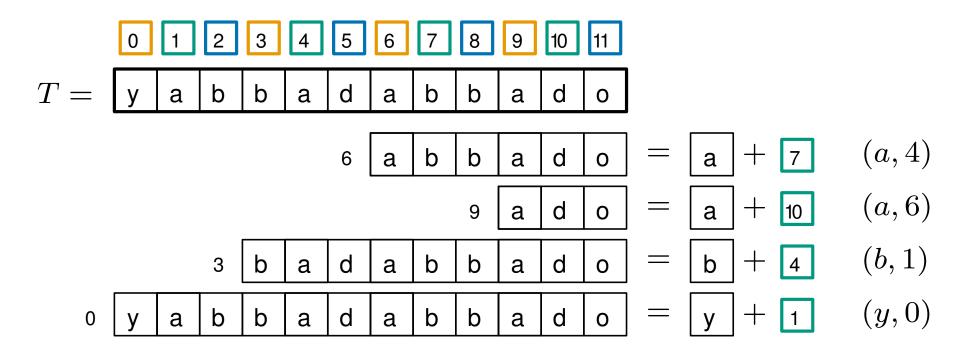
Suffix array for just $B_1 \cup B_2$



 B_2 contains indices with

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 $i \mod 3 = 2$



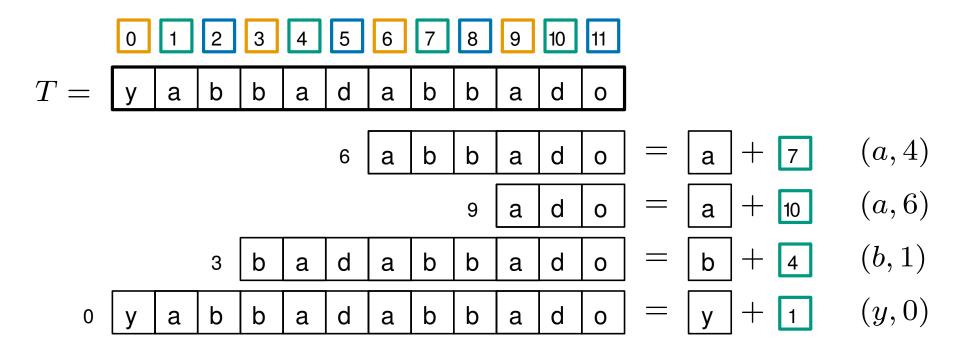
Each suffix $i \in B_0$ is represented by (T[i], r) where r is the rank of suffix (i + 1) (the ranks are given by the array below)

We then sort in O(n) time using radix sort

 B_2 contains indices with

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$$i \mod 3 = 2$$



Each suffix $i \in B_0$ is represented by (T[i], r) where r is the rank of suffix (i + 1) (the ranks are given by the array below)

We then sort in O(n) time using radix sort

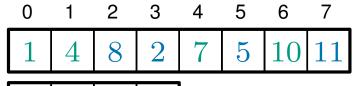
Suffix array for just B_0

6 9 3 0

 B_2 contains indices with $i \mod 3 = 2$



Suffix array for just $B_1 \cup B_2$ Suffix array for just B_0



 B_2 contains indices with $i \mod 3 = 2$



Merge them like in mergesort...

Suffix array for just $B_1 \cup B_2$ Suffix array for just B_0

```
    0
    1
    2
    3
    4
    5
    6
    7

    1
    4
    8
    2
    7
    5
    10
    11
```

 B_2 contains indices with $i \mod 3 = 2$



Merge them like in mergesort...

which is smaller, suffix 1 or 6?

Suffix array for just $B_1 \cup B_2$ Suffix array for just B_0

 $i \mod 3 = 2$

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$$i \mod 3 = 1$$

Merge them like in mergesort...

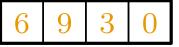
which is smaller, suffix 1 or 6?

$$|a| = |a| + |7|$$

$$_{1} = |_{a} + [_{2}]$$

Suffix array for just $B_1 \cup B_2$

Suffix array for just B_0



 $i \mod 3 = 2$



$$i \mod 3 = 1$$

Merge them like in mergesort...

which is smaller, suffix 1 or 6 ?

$$a = a + 7 \qquad (a,4)$$

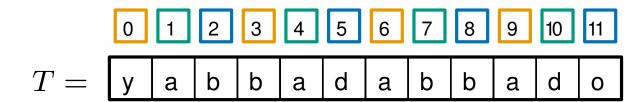
$$a = a + 2 \qquad (a,3)$$

Suffix array for just $B_1 \cup B_2$

Suffix array for just B_0

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$$i \mod 3 = 2$$



Merge them like in mergesort...

which is smaller, suffix 1 or 6 ?

$$[6] = [a] + [7] \quad (a,4)$$

$$1 = a + 2$$
 $(a,3)$

It takes O(1) time to decide that 1 is smaller

Suffix array for just $B_1 \cup B_2$

Suffix array for just B_0

 $i \mod 3 = 2$



$$i \mod 3 = 1$$

Merge them like in mergesort...

which is smaller, suffix 1

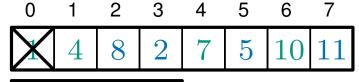
$$a = a + 7 \qquad (a,4)$$

$$a = a + 2 \qquad (a,3)$$

It takes O(1) time to decide is smaller that 1

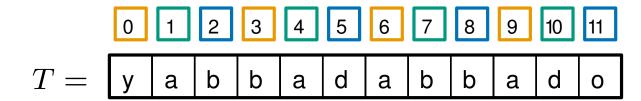
Suffix array for just $B_1 \cup B_2$

Suffix array for just B_0



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$$i \mod 3 = 2$$

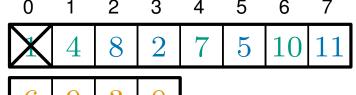


Merge them like in mergesort...

1

which is smaller, suffix 4 or 6?

Suffix array for just $B_1 \cup B_2$ Suffix array for just B_0



$$i \mod 3 = 2$$



Merge them like in mergesort...

1

which is smaller, suffix 4 or 6?

$$a = a + 7$$

$$a = a + 5$$

Suffix array for just $B_1 \cup B_2$

Suffix array for just B_0



 $i \mod 3 = 2$



 $i \mod 3 = 1$

Merge them like in mergesort...

1

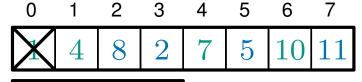
which is smaller, suffix 4 or 6?

$$[6] = [a] + [7] \quad (a,4)$$

$$a = a + 5 \qquad (a,5)$$

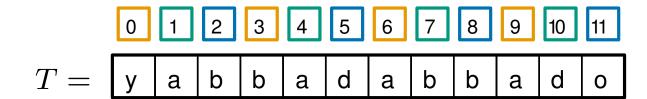
Suffix array for just $B_1 \cup B_2$

Suffix array for just B_0



 B_2 contains indices with $i \mod 3 = 2$





Merge them like in mergesort...

1

which is smaller, suffix 4 or 6?

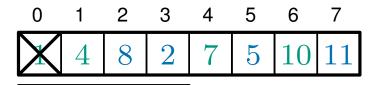
$$a = a + 7 \qquad (a,4)$$

$$a = a + 5 \qquad (a,5)$$

Again, it takes O(1) time to decide that 6 is smaller

Suffix array for just $B_1 \cup B_2$

Suffix array for just B_0



6 9 3 0

 B_2 contains indices with $i \mod 3 = 2$



Merge them like in mergesort...

Suffix array for just $B_1 \cup B_2$ Suffix array for just B_0



 B_2 contains indices with

 $i \mod 3 = 2$



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b d b b a

Merge them like in mergesort...

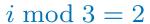
which is smaller, suffix 4 or 9 ?

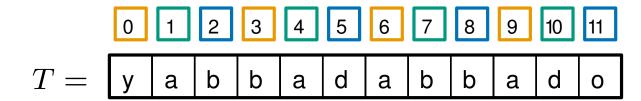
Suffix array for just $B_1 \cup B_2$ Suffix array for just B_0



 B_2 contains indices with

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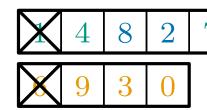


Merge them like in mergesort...

which is smaller, suffix 4 or 9 ?

4 is smaller)

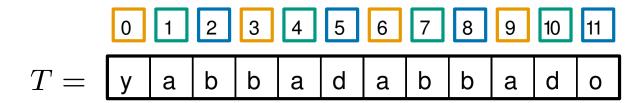
Suffix array for just $B_1 \cup B_2$ Suffix array for just B_0



 B_2 contains indices with

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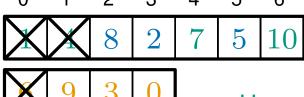


Merge them like in mergesort...

which is smaller, suffix 4 or 9 ?

4 is smaller)

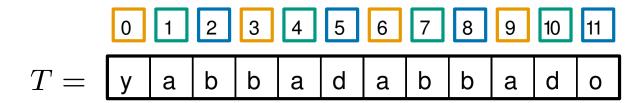
Suffix array for just $B_1 \cup B_2$ Suffix array for just B_0



 B_2 contains indices with

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$$i \mod 3 = 2$$



Merge them like in mergesort...

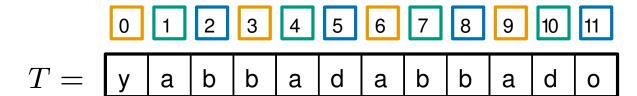
which is smaller, suffix 8 or 9 ?

Suffix array for just $B_1 \cup B_2$ Suffix array for just B_0



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$$i \mod 3 = 2$$



Merge them like in mergesort...

which is smaller, suffix 8 or 9 ?

$$9 = |a| + 10$$

$$a = b + 9$$

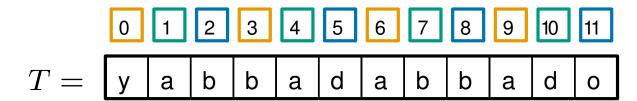
Suffix array for just $B_1 \cup B_2$

Suffix array for just B_0



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$$i \mod 3 = 2$$



Merge them like in mergesort...

which is smaller, suffix 8 or 9 ?

$$9 = a + 10$$

$$|a| = |b| + |9|$$

Uh oh! how do we compare 9 to 10 ?

Suffix array for just $B_1 \cup B_2$

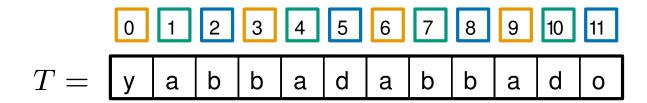
Suffix array for just B_0





 B_2 contains indices with

 $i \mod 3 = 2$



Merge them like in mergesort...

which is smaller, suffix 8 or 9 ?

$$9 = \boxed{a} + \boxed{d} + \boxed{11}$$

$$a = b + a + 10$$

It *still* takes O(1) time to decide that 9 is smaller

Uh oh! how do we compare 9 to 10 ?

Suffix array for just $B_1 \cup B_2$

Suffix array for just B_0



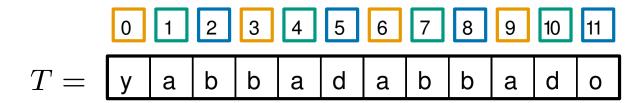




 B_2 contains indices with

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$$i \mod 3 = 2$$

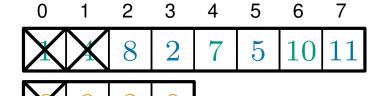


Merge them like in mergesort...

Overall this merging phase takes O(n) time

(because processing each suffix takes O(1) time)

Suffix array for just $B_1 \cup B_2$ Suffix array for just B_0





Theorem

The DC3 algorithm constructs a suffix array in O(n) time.



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Proof

Suppose T(n) is the running time. We have

$$T(n) = T(2n/3) + O(n)$$



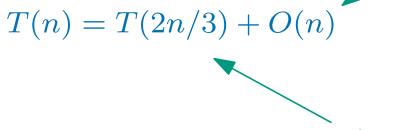
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radix sorting and merging



recursion to construct a suffix array of size 2n/3



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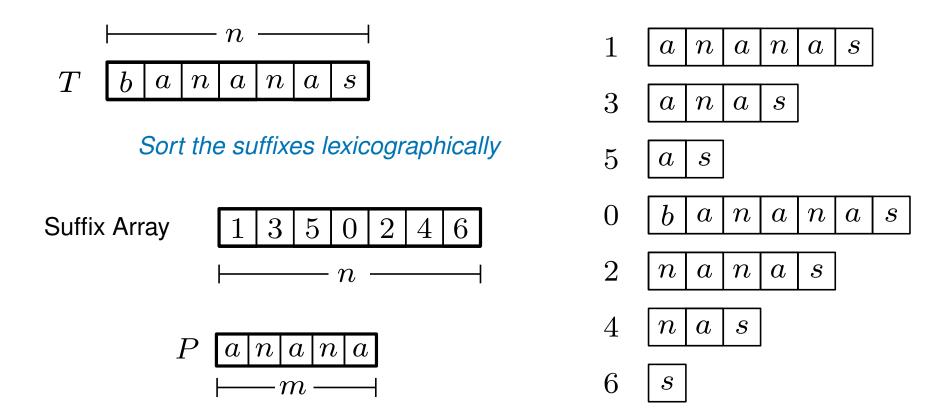
Solving this recurrence gives $T(n) \in O(n)$.

recursion to construct a suffix array of size 2n/3

radix sorting and merging



The suffix array



Finding an occurrence of a pattern (length m) takes $O(m \log n)$ time

Finding all occurrences takes $O(m \log n + occ)$ time

where occ is the number of occurences

This can be further improved to $O(m + \log n + \operatorname{occ})$ time (using LCP queries which we will see in a future lecture)

We can construct the suffix array in O(n) time