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**CS 170, Fall 2022**

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## **1. Algorithms with numbers**

## 2. Divide-and-conquer Algorithms

### 2.a. Multiplication

#### Definition 2.1 (Integer Multiplication)

A divide-and-conquer algorithm for integer multiplication is defined as follows:

```

1 function mul(x(0b[1...k]), y(0b[1...h]))
2 %Input: Positive integers x, y in binary
3 %Output: x times y
4
5 n = max(size of x, size of y)
6 if n == 1: return x × y
7
8 xL, xR = x(0b[1...⌊n/2⌋]), x(0b[⌊n/2⌋...n])
9 yL, yR = y(0b[1...⌊n/2⌋]), y(0b[⌊n/2⌋...n])
10
11 P1 = multiply(xL, yL)
12 P2 = multiply(xR, yR)
13 P3 = multiply(xL + xR, yL + yR)
14 return P1 × 2n + (P3 - P1 - P2) × 2n/2 + P2

```

Where  $0b[1...k]$  denotes the binary string representing a number

### 2.b. Recurrence Relations

#### Theorem 2.2 (Master Algorithm)

If  $T(n) = aT(n/b) + cn^k$  and  $T(1) = c$  for some constants  $a, b, c$  and  $k$ , then

$$T(n) \in \begin{cases} O(n^k) & \text{if } a < b^k \\ O(n^k \log n) & \text{if } a = b^k \\ O(n^{\log_b a}) & \text{if } a > b^k \end{cases}$$

### 2.c. Mergesort

**Definition 2.3** (Mergesort)

The Mergesort algorithm is defined as follows:

```

1 function mergesort(a[1...n])
2   %Input: An array of numbers a[1...n]
3   %Output: Sorted array a
4
5   if n>1:
6     return merge(mergesort(a[1...[n/2]]), mergesort(a[[n/2]+1...n]))
7   else:
8     return a
9
10 function merge(x[1...k], y[1...h])
11   %Input: Two arrays of numbers (x[1...k], y[1...h])
12   %Output: An array of numbers in x and y in ascending order
13
14   if k=0: return y
15   if l=0: return x
16   if x[1] <= y[1]:
17     return x[1] ◦ merge(x[2...k], y[1...h])
18   else:
19     return y[1] ◦ merge(x[1...k], y[2...h])

```

Where  $\circ$  denotes concatenation.

The *merge* function above does a constant amount of work (concatenating two arrays) per recursive call, for a total running time of  $O(k + l)$ . Thus the calls to *merge* in *mergesort* are linear, we conclude that the overall time taken by *mergesort* is

$$T(n) = 2T(n/2) + O(n)$$

Recall the Master Algorithm, here  $a = b^k$ , and therefore

$$T(n) \in O(n \log n)$$

**2.d. Medians****2.e. Matrix Multiplication****Definition 2.4****2.f. Fast Fourier Transform**