

Homework 2

1.) Heap-Delete(A, i):

if ($A[i] < A[A.\text{heap-size}]$)

Heap-Increase-Key($A, i, A[A.\text{heap-size}]$)

$A.\text{heap-size} \leftarrow A.\text{heap-size} - 1$

else

$A[i] = A[A.\text{heap-size}]$

$A.\text{heap-size} \leftarrow A.\text{heap-size} - 1$

Max-Heapify(A, i)

2.) QuickSort

part 1: When all elements have the same value,
 r is returned.

Modifying partition,

~~Partition(A, p, r)~~

~~$x = A[r]$
 $i = p - 1$
 $\text{count} = 0$
for $j = p$ to $j = r$~~

Partition(A, p, r)

$\{ x = A[r]$

$i = p - 1$

$\text{count} = 0$

for $j = p$ to $j = r$

$\{ \text{if } A[j] == x$

$\text{count}++$

$\{ \text{if } A[j] < x$

$i++$

swap($A[i], A[j]$)

$\}$

swap($A[i+1], A[r]$)

return $(i+1 - \text{count}) / 2$

$\}$

Exercise 2:

part 2:

The run time of Quicksort when all elements of A have the same value is $\Theta(n^2)$. The Partition would return index

~~part 3:~~ $q=r$, changing a problem with size n to size $n-1$.

Therefore ~~for~~ $T(n) = \cancel{T(n-1)} + n$

Using the iteration method, $T(n) = \Theta(n^2)$

~~part 3:~~

the

part 3:

The run time of Randomized-Quicksort would be $O(n)$ when all elements of array A have the same value.

For example if partitioned into 2 equal subproblems,
 $T(n) = 2T(n/2) = \Theta(n)$

part 4:

Partition(A, P, r) {

int x = A[r]

swap(A[r], A[p])

int i = p - 1

int k = p

for (int j = p + 1 to r - 1)

{ if (A[j] < x)

i++;

k = i + 2;

swap(A[i], A[j])

swap(A[k], A[j])

if (A[j] = x)

k++;

swap(A[k], A[j])

} swap(A[i+1], A[r])

return (k+1) (i+1)

Modified so that elements will be moved after k

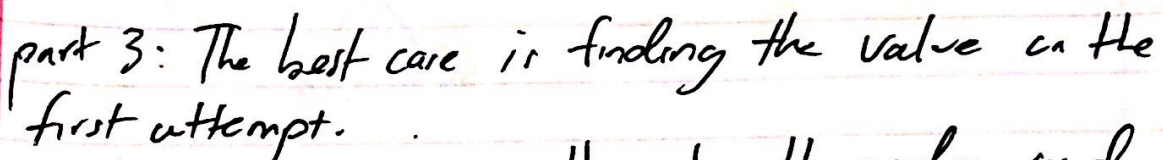
5.) * loop invariant

1) x is equal to $A[i]$

2) i is equal to $p - 1$

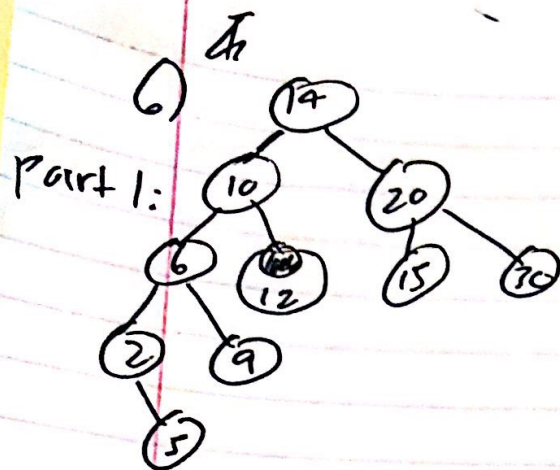
3.) k is equal to p

3) part 1: $a_1 \ a_2$



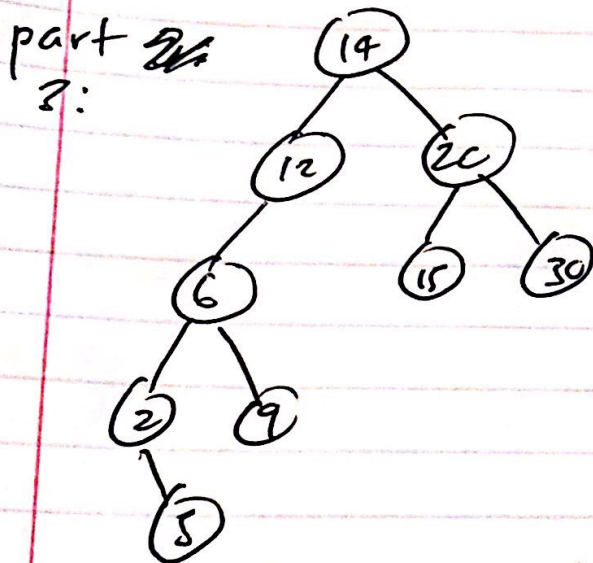
part 4:

Worst case depth = 8



part 2: ~~2, 5, 6, 9, 10, 12, 14, 15,~~
~~20, 30~~

2, 5, 6, 9, 10, 12, 14, 15, 20, 30



deleting 10

deleting 12

