

# CS375 HW5

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## 1. Cost of Gaussian Elimination

Assume that an ancient computer solves a 1000-variable, upper-triangular, linear system by back substitution in 0.5 seconds. Estimate the time needed to solve a general (full) system by Gaussian Elimination (forward elimination + back-substitution). Use the counts from the lectures:  $\frac{2n^3}{3} + O(n^2)$  for elimination and  $n^2$  for back substitution.

Given a  $m \times n$  matrix, the number of operations on row  $m$  of an upper triangular matrix is 1, as we divide  $\frac{x_{m,n-1}}{x_{m,n}}$ .

We then move up one row. Here, we will do

1. plug in (multiply) prev value
2. Subtract value from pt. 1 to RHS
3. Divide to determine second term

This illustrates that solving an upper-right triangular matrix takes  $O(n^2)$  operations. Therefore, the total operations from back-substitution and elimination will be given by

$$\frac{2}{3}(1000)^3 + (1000)^2$$

total operations. We know that our computer can solve a 1000 upper-triangular matrix in 0.5 seconds. This tells us the computer can solve at a rate of  $2(1000^2) = 2,000,000$  FLOPS, or 2 Mega-FLOPS.

Simply dividing our total operations with the speed per second will give us the total time in seconds;

$$\begin{aligned} &= \frac{\frac{2}{3}(1000)^3 + (1000)^2}{2,000,000} \\ &= 333.83 \text{ seconds} \\ &\approx 5.6 \text{ minutes} \end{aligned}$$

## 2. Gaussian Elimination with and without partial pivoting

- (a) Write a Matlab script that uses naive Gaussian Elimination to solve the linear system

$$\begin{bmatrix} a & 1 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1+a \\ 2 \end{bmatrix}$$

for  $a = 10^{-2k}$ ,  $k = 1, 2, \dots, 10$ . The exact solution is  $x = \begin{bmatrix} 1 & 1 \end{bmatrix}^T$ , regardless of the value  $a$ . Place the solutions your script produces in a table. How does the accuracy of your numerical solution behave as  $k$  gets bigger. Explain.

```
k_values = 1:10;
exact_solution = [1; 1];

solutions = zeros(2, length(k_values));
errors = zeros(1, length(k_values));

for k = k_values
    a = 10^(-2*k);

    A = [a, 1;
         1, 1];
    b = [1 + a;
         2];

    n = length(b);
    x = zeros(n,1);

    % Forward Elimination
    for i = 1:n-1
        factor = A(i+1,i) / A(i,i);
        A(i+1,:) = A(i+1,:) - factor * A(i,:);
        b(i+1) = b(i+1) - factor * b(i);
    end

    % Back Substitution
    x(n) = b(n) / A(n,n);
    for i = n-1:-1:1
        x(i) = (b(i) - A(i,i+1:n) * x(i+1:n)) / A(i,i);
    end

    solutions(:, k) = x;
    errors(k) = norm(exact_solution - x);
end
```

$k$	$a$	$x_1$	$x_2$	Error
1	$1.000000 \times 10^{-2}$	1.000000	1.000000	$8.881784 \times 10^{-16}$
2	$1.000000 \times 10^{-4}$	1.000000	1.000000	$1.101341 \times 10^{-13}$
3	$1.000000 \times 10^{-6}$	1.000000	1.000000	$2.875566 \times 10^{-11}$
4	$1.000000 \times 10^{-8}$	1.000000	1.000000	$6.077471 \times 10^{-9}$
5	$1.000000 \times 10^{-10}$	1.000000	1.000000	$8.274037 \times 10^{-8}$
6	$1.000000 \times 10^{-12}$	0.999867	1.000000	$1.331440 \times 10^{-4}$
7	$1.000000 \times 10^{-14}$	0.999201	1.000000	$7.992778 \times 10^{-4}$
8	$1.000000 \times 10^{-16}$	2.220446	1.000000	1.220446
9	$1.000000 \times 10^{-18}$	0.000000	1.000000	1.000000
10	$1.000000 \times 10^{-20}$	0.000000	1.000000	1.000000

As  $k$  increases, the param  $a$  decreases, approaching closer and closer to 0. As this happens, we the naive approach we implemented allows for very small values of  $a$ , which quickly decrease the accuracy of the result.

- (b) Next, change your Matlab code to carry out partial-pivoting. What is the solution now? Give a table of solution values. Explain your result and the accuracy relative to part (a).

$k$	$a$	$x_1$	$x_2$	Error
1	$1.000000 \times 10^{-2}$	1.000000	1.000000	$0.000000 \times 10^0$
2	$1.000000 \times 10^{-4}$	1.000000	1.000000	$0.000000 \times 10^0$
3	$1.000000 \times 10^{-6}$	1.000000	1.000000	$0.000000 \times 10^0$
4	$1.000000 \times 10^{-8}$	1.000000	1.000000	$0.000000 \times 10^0$
5	$1.000000 \times 10^{-10}$	1.000000	1.000000	$0.000000 \times 10^0$
6	$1.000000 \times 10^{-12}$	1.000000	1.000000	$3.140185 \times 10^{-16}$
7	$1.000000 \times 10^{-14}$	1.000000	1.000000	$0.000000 \times 10^0$
8	$1.000000 \times 10^{-16}$	1.000000	1.000000	$1.110223 \times 10^{-16}$
9	$1.000000 \times 10^{-18}$	1.000000	1.000000	$0.000000 \times 10^0$
10	$1.000000 \times 10^{-20}$	1.000000	1.000000	$0.000000 \times 10^0$

This table shows how partial pivoting can significantly decrease our error across a similar problem set.

3. Gaussian Elimination with scaled partial pivoting - Using scaled partial pivoting without actually moving data in the matrix, show the steps required to solve the following system of equations. Calculate the scale vector (called  $s$  in the lecture). Show how the pivot row is selected at each step, and carry out the computations. At each step, include the index vector (called  $\ell$  in lecture).

For the augmented matrix,ple

$$\begin{bmatrix} 2 & -1 & 3 & 4 & 15 \\ 4 & 2 & 0 & 7 & 11 \\ 2 & 1 & 1 & 3 & 7 \\ 6 & 5 & 4 & 17 & 31 \end{bmatrix}, s = \begin{bmatrix} 4 \\ 7 \\ 3 \\ 17 \end{bmatrix}, \ell = \begin{bmatrix} 1 \\ 2 \\ 3 \\ 4 \end{bmatrix}$$

$$\begin{bmatrix} 2 & 1 & 1 & 3 & 7 \\ 0 & -2 & 2 & 1 & 8 \\ 0 & 0 & -2 & 1 & -3 \\ 0 & 0 & 3 & 9 & 18 \end{bmatrix}, s = \begin{bmatrix} 3 \\ 4 \\ 7 \\ 17 \end{bmatrix}, \ell = \begin{bmatrix} 3 \\ 1 \\ 2 \\ 4 \end{bmatrix}$$

$$\begin{bmatrix} 2 & 1 & 1 & 3 & 7 \\ 0 & -2 & 2 & 1 & 8 \\ 0 & 0 & -2 & 1 & -3 \\ 0 & 0 & 0 & 7.5 & 22.5 \end{bmatrix}, s = \begin{bmatrix} 3 \\ 4 \\ 7 \\ 17 \end{bmatrix}, \ell = \begin{bmatrix} 3 \\ 1 \\ 2 \\ 4 \end{bmatrix}$$

...finish?-

Answer looks like it will be a little off too