

ECE437/CS481

# M07D: VIRTUAL MEMORY

CHAPTER 9.5-9.4

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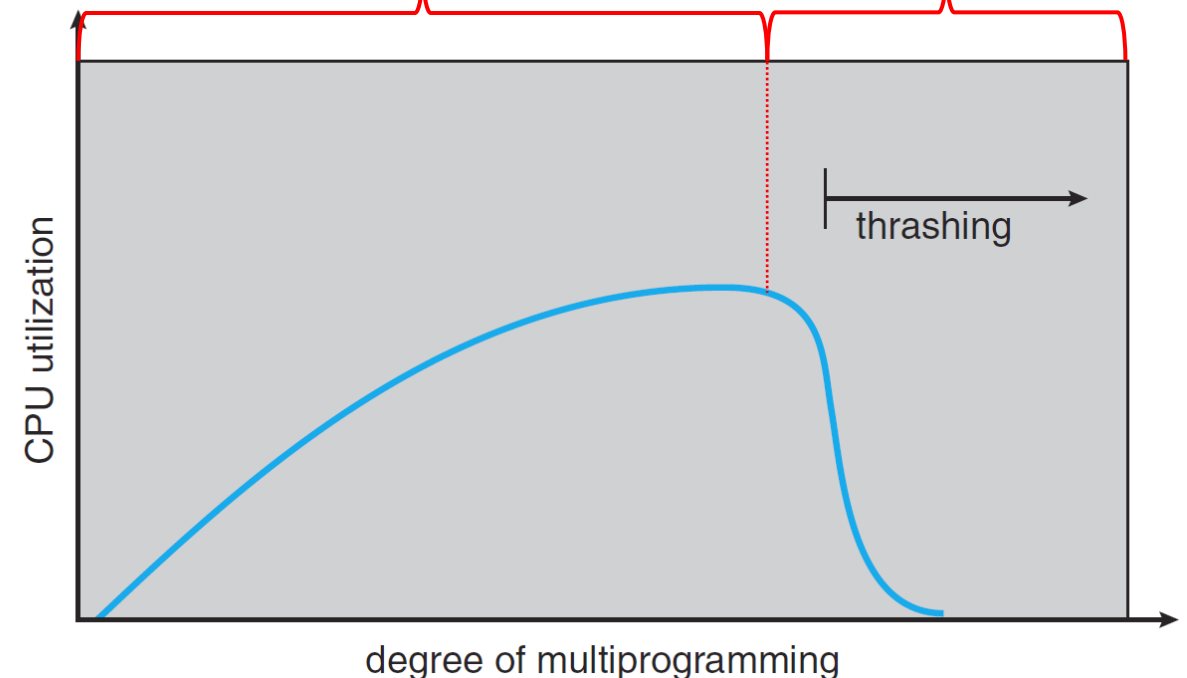
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A decorative blue wavy line that spans the width of the slide, starting with a small upward curve on the left, dipping into a V-shape in the center, and then curving back up on the right before continuing as a straight line to the edge.

# Thrashing

- ❑ As the degree of multiprogramming increases, more processes will be brought into the main memory
  - The average CPU utilization **may** increase.
  - The size of memory space allocated to each process **may** reduce.
    - If reduced, more pages of a process will be stored in the hard drive.
    - The page fault rate increases.
    - The OS spends more time on swapping in/out pages.
    - Reduce the CPU utilization.

❑ **Thrashing** → A process is busy in swapping pages in and out.



# Thrashing

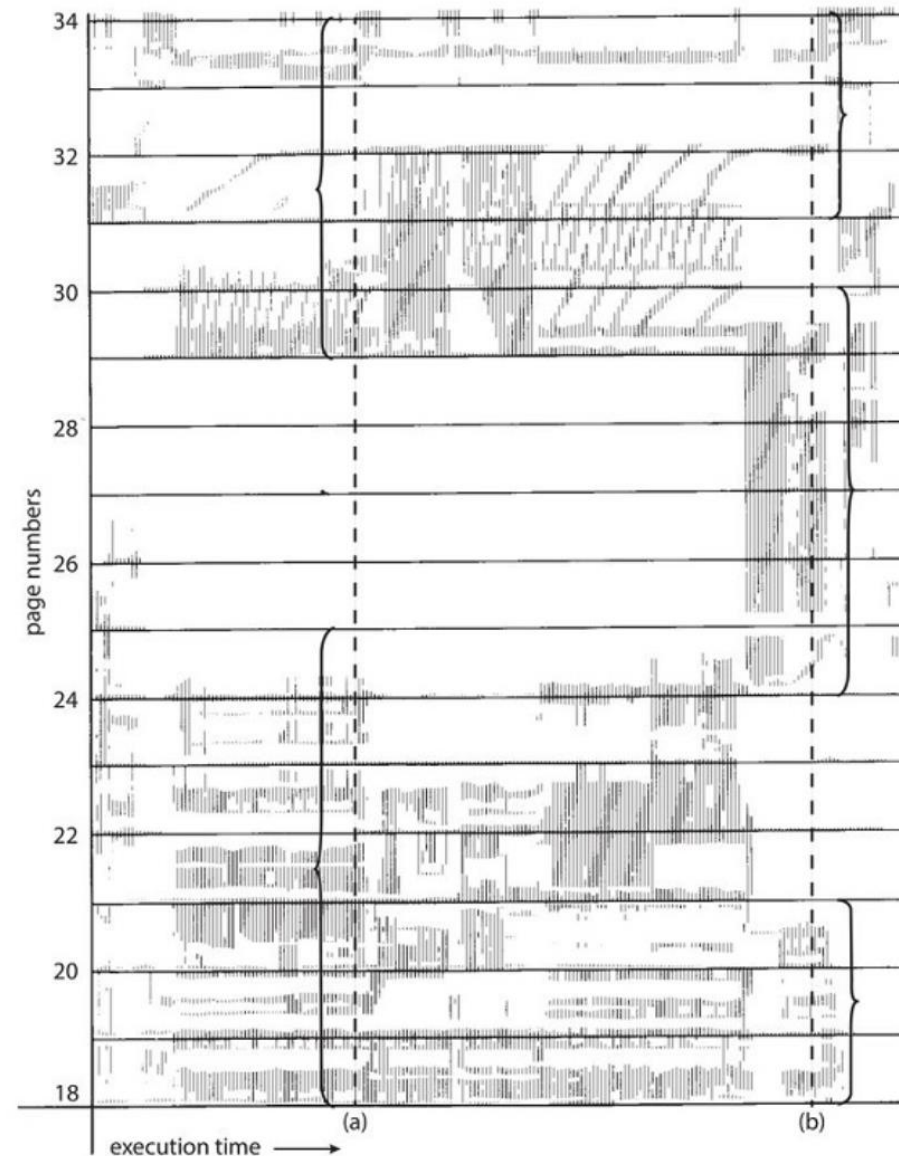
## ❑ How to mitigate process thrashing?

- Basic idea: 1) prefetch the memory pages of a process that will be referenced in the near future; 2) if the memory space is not enough, suspend the process.
  - ✓ Figuring out the memory pages of a process to be referenced soon is difficult.
- Locality feature of a process—a set of pages for a process are actively used together
- Two types of locality
  - ✓ **Temporal locality**
    - A process's pages, which are referenced recently, will likely be referenced in the near future.
    - E.g., looping, counting/reduction variables, etc.
  - ✓ **Spatial locality**
    - One page of a process is referenced, and its nearby pages will be referenced in the near future.
    - E.g., array, sequential code.

# Thrashing

## □ Locality in a memory-reference pattern

- In time slot (a), the spatial locality is [18,19,20,21,22,23,24,29,30,33]
- In time slot (b), the spatial locality is [18, 19, 20, 24, 25, 26, 27, 28, 29, 31, 32, 33].
- The locality may change over time.

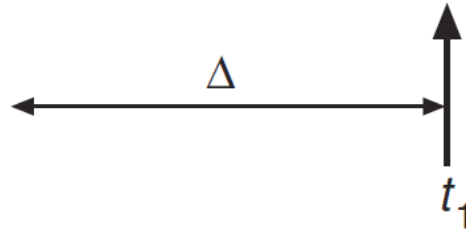


# Thrashing

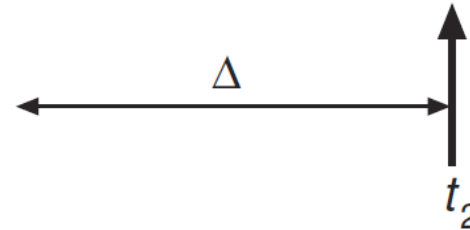
- ❑ Based on the locality feature, the **working set model** is designed to figure out the need of a process in the near future.
- One parameter, working set window  $\Delta$ , is used to monitor the recent referenced pages.

page reference table

. . . 2 6 1 5 7 7 7 7 5 1 6 2 3 4 1 2 3 4 4 4 3 4 3 4 4 4 1 3 2 3 4 4 4 3 4 4 4 . . .



$WS(t_1) = \{1, 2, 5, 6, 7\}$



$WS(t_2) = \{3, 4\}$

$\Delta=10$

- The accuracy of the working set depends on the value of  $\Delta$ .
  - ✓ If  $\Delta$  is too small, it will not cover the entire locality.
  - ✓ If  $\Delta$  is too large, it may overlap several localities.
  - ✓ If  $\Delta$  is infinite, the working set includes all the pages during the process execution.

# Thrashing

## □ Procedures of using working set to mitigate the process thrashing.

- The OS monitors the working set of each running process and tries to allocate the frames to each running process, i.e., # of frames allocated to process  $i = |WS_i|$ , where  $|WS_i|$  indicates the size of the working set for process  $i$ .
- Check if the total number of free frames (denoted as  $D$ ) is larger than the requirements, i.e.,  $D > \sum_i |WS_i|$  ?
  - ✓ If yes, another process can be initiated (by allocating frames to the process).
  - ✓ If no, the OS selects a process(es) to suspend.

# Thrashing

## ❑ Relationship between the working set of a process and its page fault rate.

- As mentioned before, the locality/working set of a process may vary over time.
- The page fault rate of a process transits between **peaks** and **valleys** over time. Here,
  - ✓ The peaks indicate the pages in working set changes, and thus the OS has to bring the pages from the hard disks into memory (demand paging)
  - ✓ The valleys indicate all the pages in the working set have swapped in the memory.

