

Statement (or Proposition): A sentence which is either always true or always false is called a statement.

Eg.	1. Sky is blue	T	Stat.
	2. Blood is red	T	Stat.
	3. It will rain tomorrow		Not stat.

Boolean Variable: It is a symbol (say  $p$ ), which takes the boolean values.

Eg.  $p$ : Sky is blue

Boolean Expression: The expression with Boolean variables and logical operations is called Boolean expression

Truth Table: It is table that gives the truth values corresponding to each possible combination of inputs for the logical operations.

Negation : Negation of a statement  $p$  gives  $T$  for the actual value  $F$  and  $F$  for actual value  $T$ .

e.g.  $p$ : sky is blue  $\quad T$   
 $\sim p$ : sky is not blue  $\quad F$

Truth Table

$p$	$\sim p$
T	F
F	T

Conjunction '1'

$$[p \wedge q = p]$$

$p \wedge q$  : T if both  $p$  and  $q$  are T, otherwise false.

Truth Table

$p$	$q$	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

Disjunction 'V'

$$p \vee q = p \text{ or } q$$

$p \vee q$  : T if at least one of  $p$  or  $q$  is true

$p$	$q$	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

## Conditional Operator (Implication operator)

$p \Rightarrow q$  : if  $p$  then  $q$

Hypothesis

Implication

Truth Table

$p$	$q$	$p \Rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

$p \Rightarrow q$  simply means that if hypothesis  $p$  is true and it is implying a true statement then  $p \Rightarrow q$  is true.

Again if hypothesis  $p$  is true and it is implying a false statement then  $p \Rightarrow q$  is false.

In case when the hypothesis is false then it is of no meaning whatever the implication is we cannot say  $p \Rightarrow q$  to be false, so in each case we take  $p \Rightarrow q$  is true.

Converse of  $p \Rightarrow q$ : is defined as  $q \Rightarrow p$

e.g.  $p$ : sky is blue

$q$ : parrot is red

→ Then  $p \Rightarrow q$ : If sky is blue then parrot is red.

its converse is  $q \Rightarrow p$ : If parrot is red then sky is blue.

Inverse of  $p \Rightarrow q$ : is defined as  $\sim p \Rightarrow \sim q$

e.g.  $p \Rightarrow q$ : If sky is blue then parrot is red

its inverse is  $\sim p \Rightarrow \sim q$ : If sky is not blue then parrot is not red.

Contrapositive of  $p \Rightarrow q$ : is defined as  $\sim q \Rightarrow \sim p$

e.g.  $p \Rightarrow q$ : If sky is blue then parrot is red

its contrapositive is  $\sim q \Rightarrow \sim p$ : If parrot is not red then sky is not blue.

# A Conditional expression is equivalent to its contrapositive.

$$\text{i.e. } p \Rightarrow q = \sim q \Rightarrow \sim p$$

→ This can be seen from their truth tables.

# But  $p \Rightarrow q$  and its inverse or converse are not equivalent.

## Biconditional operator

$p \Leftrightarrow q$  is defined as  $\neg p$  if and only if  $q$ .

$p \Leftrightarrow q$  means the both side implication

$$\text{ie } p \Leftrightarrow q = (\neg p \Rightarrow q) \wedge (q \Rightarrow \neg p)$$

Truth Table:

$p$	$q$	$p \Leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

eg:

$p$ : Blood is red

$q$ : sky is blue

$p \Leftrightarrow q$ : Blood is red if and only if  
sky is blue

#  $p \Leftrightarrow q$  means that either both are true  
at the same time or both are false  
at the same time.

Tautology: An expression which is true for all possible inputs is called a tautology.

Ex:

$p$ : sky is blue

$\sim p$ : sky is not blue

Then  $\sim(p \wedge \sim p)$  is a tautology

$p$	$\sim p$	$p \wedge \sim p$	$\sim(p \wedge \sim p)$
T	F	F	T
F	T	F	T
F	T	F	T
F	T	F	T

Contradiction: An expression which is false for all possible inputs is called a contradiction.

Ex:

$p$ : sky is blue

$\sim p$ : sky is not blue

Then  $p \wedge \sim p$  is a contradiction

$p$	$\sim p$	$p \wedge \sim p$
T	F	F
F	T	F
F	T	F
F	T	F

# De Morgan's Laws

$$1. \sim(p \vee q) = \sim p \wedge \sim q$$

$$2. \sim(p \wedge q) = \sim p \vee \sim q$$

1.

$p$	$\sim p$	$q$	$\sim q$	$p \vee q$	$\sim(p \vee q)$	$\sim p \wedge \sim q$
T	F	T	F	T	F	F
T	F	F	T	T	F	F
F	T	T	F	T	F	F
F	T	F	T	F	T	T

2.

$p$	$\sim p$	$q$	$\sim q$	$p \wedge q$	$\sim(p \wedge q)$	$\sim p \vee \sim q$
T	F	T	F	T	F	F
T	F	F	T	F	T	T
F	T	T	F	F	T	T
F	T	F	T	F	T	T

Dual Statement (Expression): Two expressions are said to be dual of each other if either of them can be obtained from the other by replacing  $\wedge$  by  $\vee$ ,  $\vee$  by  $\wedge$ ,  $\neg T$  by  $F$ , and  $\neg F$  by  $T$ .

- Example:
1.  $(\neg p \wedge q) \vee r$   
its dual is  $(\neg p \vee q) \wedge r$
  2.  $\neg(p \vee q) \wedge (p \vee \neg(p \wedge r))$   
its dual is  $\neg(p \wedge q) \vee (p \wedge \neg(p \vee r))$

Principle of Duality: It states that the dual of any two equal expressions will also be equal.

Example: We have two equal expressions  
 $\neg(p \vee q) = \neg p \wedge \neg q$

Then by the principle of duality, their duals will also be equal i.e.

$$\neg(p \wedge q) = \neg p \vee \neg q$$

Predicate : A predicate is an expression of one or more Boolean variable defined on some specific domain.

This becomes a statement for some specific values of variables.

Example : Domain :  $D = \text{set of all fruits}$

Predicate :  $P(x) : x \text{ is mango}$ ,

So,  $P(x)$  is now a statement which is true.

Example :  $P(x) : x \text{ is prime} , x \in \mathbb{N}$

$\therefore P(2)$  is True,  $P(4)$  is false

# Simply the truth value of  $P(x)$  depends on the value of  $x$ .

Example :  $P(x) : x \text{ is odd} , x \in \mathbb{N}$

$\therefore P(1)$  is T,  $P(2)$  is F.

Universal Quantifier: It is denoted by  $\forall$  and read as 'for all' or 'for every' declares that a statement must be true for all variables within some domain.

Let  $P(x)$  be a predicate and  $D$  the domain then the universal statement ' $\forall x : P(x)$ ' means that ' $P(x)$  is true for every  $x$  in  $D$ '.

Example:  $\mathbb{N}$  = Set of all natural numbers

$$P(x) : x+1 > 1$$

' $\forall x : P(x)$ ' means that " $x+1 > 1$  is true for every  $x$  in  $\mathbb{N}$ ".

Example 2.:  $\mathbb{N}$  = Set of all natural numbers

$$P(x) : 3x \text{ is odd.}$$

$\therefore \forall x : P(x)$  means that " for every  $x \in \mathbb{N}$   $3x$  is odd."

or "  $3x$  is odd is true for every  $x \in \mathbb{N}$ ."

Existential Quantifier: It is denoted by ' $\exists$ ' and read as 'there exists' or 'for at least one' declares that a statement must be true for at least one variable within the domain.

Let  $P(x)$  be a predicate and  $D$  be the domain of  $x$ . Then the existential statement ' $\exists x : P(x)$ ' means that " $P(x)$  is true for at least one  $x$  in  $D$ ".

Example 1.:  $N = \text{Set of natural numbers}$

$$P(x), x+1 > 1$$

' $\exists x : P(x)$ ' means that " $x+1 > 1$  is true for at least one  $x$  in  $N$ ".

Example 2.:  $N = \text{Set of natural numbers}$

$$P(x) = 3x+1 \text{ is odd}$$

' $\exists x : P(x)$ ' means that " $3x+1$  is odd" is true for at least one  $x$  in  $N$ ".

Example 3.: Domain = Set of all birds

$$P(x) : x \text{ is a crow}, Q(x) : x \text{ is black}$$

Statement : All crows are not black.

Boolean Expression :  $\exists x : P(x) \Rightarrow \sim Q(x)$

Nested Quantifier: If we use a quantifier that happens with the scope of another quantifier, it is called nested quantifier.

Example 1.

$N = \text{Set of all natural numbers}$ ,

$P(x) : x \text{ is prime}$

$\forall x \exists y : (y > x) \wedge P(y)$  means

"for every  $x$  there exists at least one  $y$  such that  $y$  is greater than  $x$  and  $y$  is prime."

Example 2.  $N = \text{Set of natural numbers}$

$P(x,y) : y > x$

$\forall x \exists y : P(x,y)$  means that "for each  $x \in N$ , there is at least one  $y \in N$  which is greater than  $x$ ".

Negating Quantified Expressions: To negate quantified statement, we first negate all the quantifiers from left to right (keeping in the same order), then we negate the statements.

Example: 1.  $\sim [\forall x : P(x)] = \exists x : \sim P(x)$

2.  $\sim [\exists x : P(x)] = \forall x : \sim P(x)$

3.  $\sim [\forall x, \exists y : P(x,y)] = \exists x, \forall y : \sim P(x,y)$

4.  $\sim [\exists x, \forall y : P(x,y)] = \forall x, \exists y : \sim P(x,y)$ .

# "2x is an even number for every natural number x"

$\forall x : P(x)$  where  $P(x) : 2x$  is an even no. and  $x \in \mathbb{N}$ .

$\therefore \sim [\forall x : P(x)] = \exists x : \sim P(x)$

"there exists at least one  $x$  in  $\mathbb{N}$  such that  $2x$  is not an even number."