Department I - C Plus Plus

Modern and Lucid C++ Advanced for Professional Programmers

Week 8 - Compile-Time Computation

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Recap Week 7



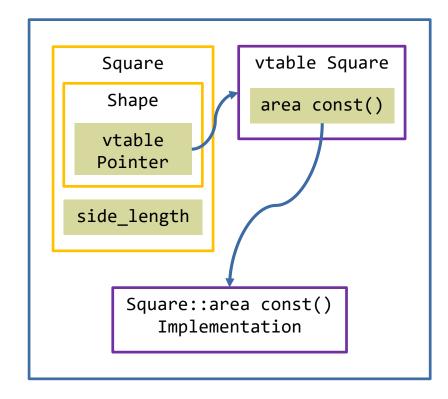




A polymorphic call of a virtual function requires lookup of the target function

```
struct Shape {
  virtual unsigned area() const = 0;
 virtual ~Shape();
};
struct Square : Shape {
  Square(unsigned side length)
    : side length{side length} {}
  unsigned area() const {
    return side length * side length;
  unsigned const side length;
};
```

```
decltype(auto) amountOfSeeds(Shape const & shape) {
  auto area = shape.area();
  return area * seedsPerSquareMeter;
};
```



• Article on this topic: http://eli.thegreenplace.net/2013/12/05/the-cost-of-dynamic-virtual-calls-vs-static-crtp-dispatch-in-c

```
struct SpaceDriveTag {};
struct HyperspaceDriveTag : SpaceDriveTag {};
struct InfniteProbabilityDriveTag : SpaceDriveTag {};
```

```
struct MultiPurposeCrewVehicle;
struct GalaxyClassShip;
struct HeartOfGoldPrototype;
```

- Approach 1: Derive space ship from the associated tag typ
 - This is not applicable for all types (e.g. for primitive types, as we will see later)
 - This is not extensible (i.e. you cannot specify new kinds of tag kinds as a user of the API)
- Approach 2: SpaceshipTraits template

```
template<typename>
struct SpaceshipTraits {
  using Drive = SpaceDriveTag;
};
```

```
template<>
struct SpaceshipTraits<GalaxyClassShip> {
  using Drive = HyperspaceDriveTag;
};
```

```
struct IntIterator { /* Member Types Omitted */
  explicit IntIterator(int const start = 0) :
    value { start } {}
  bool operator==(IntIterator const & r) const {
    return value == r.value;
  bool operator!=(IntIterator const & r) const {
    return !(*this == r);
  value_type operator*() const {
    return value;
  IntIterator & operator++() {
    ++value;
    return *this;
  IntIterator operator++(int) {
    auto old = *this;
    ++(*this);
    return old;
private:
 value_type value;
```

Explicit constructor

Implement != through operator ==

Implement postfix through prefix operators

Reuse pre-defined type

- You can write C++ code that is evaluated by the compiler
- You know the restrictions of constexpr functions
- You can write your own literal types
- You know how SFINAE works and can apply it to your code

Constexpr







- ROMable data (guaranteed!)
 - Smaller binary
- Better run-time performance
- Plain C++
- Not yet everything possible provided
 - Standard library not made up for everything desirable
- Compiletimes can increase horrendously
 - Example that is supposed to take 24h to be compiled http://cpptruths.blogspot.ch/2005/11/c-templates-are-turing-complete.html

- (static) const variables in namespace scope of built-in types initialized with constant expression are usually put into ROMable memory, if at all.
- Allowed in constant expression context
- No complicated computations (except with macros)
- No guarantee to be done at compile-time in all cases

Non-type template arguments

```
std::array<Element, 5> arr{};
```

Array bounds

```
double matrix[ROWS][COLS]{};
```

Case expressions

```
switch (value) {
case 42:
   //...
}
```

Enumerator initializers

```
enum Light {
    Off = 0, On = 1
};
```

static_assert

```
static_assert(order == 66);
```

constexpr variables

```
constexpr unsigned pi = 3;
```

• constexpr if statements

```
if constexpr (size > 0) {
}
```

noexcept

```
Blob(Blob &&) noexcept(true);
```

• . . .

```
static_assert(sizeof(int) == 4, "unexpected size of int");
static_assert(isGreaterThanZero(Capacity));
```

static_assert checked at compile-time

- Compilation fails if it evaluates to false
- The compiler needs to be able to evaluate the expression

Syntax:

- static_assert(condition, message);
- static_assert(condition);

constexpr unsigned pi = 3;

- Evaluated at compile-time (mandatory)
- Initialized by a constant expression
 - Literal value
 - Expression computable by the compiler
 - constexpr function calls
- Require literal type
- Can be used in constant expression contexts

- Possible contexts
 - Local scope
 - Namespace scope
 - static data members
- constexpr variables are const

```
constexpr auto factorial(unsigned n) {
   ...
}
```

- constexpr functions can...
 - ... have local variables of "literal" type. The variables must be initialized before used.

```
int local;
LiteralType local{};
int local;
f(local);
std::string local{};
```

- ... use loops, recursion, arrays, references
- ... even contain branches that rely on run-time features, if branch is not executed during compile-time computations, e.g., throw
- ... but can only call constexpr functions

- constexpr evaluation cannot...
 - ... allocate dynamic memory (new, delete)

```
constexpr int * allocate() {
  return new int{};
}
```

■ ... use exception handling (throw, try/catch)

```
constexpr void throwError() {
  throw std::logic_error{""};
}
```

■ ... be virtual member functions (until C++20)

```
struct Base {
  constexpr virtual void modify();
};
```

Constexpr functions are usable in constexpr and non-constexpr contexts

```
constexpr auto factorial(unsigned n) {
  auto result = 1u;
  for (auto i = 2u; i <= n; i++) {
    result *= i;
  return result;
constexpr auto factorial0f5 = factorial(5);
int main() {
  static_assert(factorialOf5 == 120);
  std::cout << factorial(5);</pre>
```

- The compiler will prevent Undefined Behavior
 - Leads to compilation error

```
constexpr int divide(int n, int d) {
   return n / d;
}
constexpr auto surprise = divide(0, 0);
int main() {
   std::cout << surprise;
}</pre>
```

```
..\CTUB.cpp:7:33: in 'constexpr' expansion of 'divide(0, 0)'
..\CTUB.cpp:4:12: error: '(0 / 0)' is not a constant expression
  return n / d;
      ~~^~~
```

If constexpr evaluation does not reach invalid statement the code is valid

```
constexpr void throwIfZero(int value) {
  if (value == 0) {
    throw std::logic_error{""};
  }
}

constexpr int divide(int n, int d) {
  throwIfZero(d);
  return n / d;
}

constexpr auto five = divide(120, 24);
constexpr auto failure = divide(120, 0);
```

```
?
```

```
int whatIsTheAnswer() {
  return 42;
}
static_assert(whatIsTheAnswer() == 42);
```

```
constexpr int global = 42;
constexpr int const * allocate(bool useGlobal) {
  if (useGlobal) {
    return &global;
  } else {
    return new int{};
  }
}
constexpr int const * ptr = allocate(true);
```

```
?
```

```
int whatIsTheAnswer() {
  return 42;
}
static_assert(whatIsTheAnswer() == 42);
```

Incorrect

The expression in static_assert requires a compile-time expression. whatIsTheAnswer() is not a constexpr function.

```
constexpr int global = 42;
constexpr int const * allocate(bool useGlobal) {
   if (useGlobal) {
      return &global;
   } else {
      return new int{};
   }
}
constexpr int const * ptr = allocate(true);
```

Correct

As long as the path with the new expression is not taken, it is valid to have the code in a constexpr function. allocate(false); could not be used to initialize a compile-time constant.

- Built-in scalar types, like int, double, pointers, enumerations, etc.
- Structs with some restrictions¹
 - Trivial destructor (non-user-defined)
 - With a constexpr constructor and no virtual members
- Lambdas
- References
- Arrays of literal types
- void
- Literal Types can be used in constexpr functions, but only constexpr member functions can be called on values of literal type

- Trivial Destructor
- Constexpr Constructor
 - At least one
- Constexpr Member Functions
 - const & non-const
 - Only constexpr useable in constexpr context
 - All are non-virtual
- Can be a template
- It can still contain non-constexpr constructors and member functions

```
template <typename T>
class Vector {
 constexpr static size t dimensions = 3;
  std::array<T, dimensions> values{};
public:
  constexpr Vector(T x, T y, T z)
    : values{x, y, z}{}
  constexpr T length() const {
    auto squares = x() * x() +
                   y() * y() +
                   z() * z();
    return std::sqrt(squares);
  constexpr T & x() {
    return values[0];
 constexpr T const & x() const {
 return values[0];
```

- Can be used in constexpr and non-constexpr contexts
- Non-const member functions can be used to modify the object
- constexpr variables are const

```
constexpr Vector<double> create() {
  Vector<double> v{1.0, 1.0, 1.0};
  v.x() = 2.0;
  return v;
constexpr auto v = create();
static assert(doubleEqual(v.length(), 2.4495));
int main() {
  //v.x() = 1.0;
  auto v2 = create;
  v2.x() = 2.0;
```

 Note on Vector: Has hardcoded three dimensions (x/y/z).

```
template <size t n>
struct fact {
  static size_t const value{(n > 1)? n * fact<n-1>::value : 1};
};
template <>
struct fact<0> { // recursion base case: template specialization
  static size_t const value = 1;
};
void testFactorialCompiletime() {
  constexpr auto result = fact<5>::value;
 ASSERT_EQUAL(result, 2 * 3 * 4 * 5);
```

"Integer" only (almost) through non-type template parameters

Capture types (the types returned by lambda expressions) are literal types as well

- They can be used as types of constexpr variables
- They can be used in constexpr functions
- Restrictions to constexpr functions and variables apply as well

Examples for demonstration purposes

```
constexpr double pi = 3.14159;

constexpr auto area = [](double r) {
  return pi * r * r;
};

constexpr auto circleArea = area(2.0);
```

```
constexpr auto cubeVolume(double x) {
  auto area = [x] {return pi * x * x;};
  return area() * x;
}
constexpr auto cV = cubeVolume(5.0);
```

Syntax

```
template<[Parameters]>
Type name [= initialization];
```

- Can be specialized
- Usually constexpr

Purpose

- Compile-time predicates and properties of types
- Usually applied in template meta programming
- Before C++14 it was necessary to create a class template with a static member variable
 - Now less code is required for the same effect

```
template<typename T>
constexpr T pi = T(3.1415926535897932385);

template<typename T>
constexpr bool is_integer = false;

template<>
constexpr bool is_integer<int> = true;
```

```
template <size_t N>
constexpr size_t factorial = factorial<N - 1> * N;

template <> //Base case
constexpr size_t factorial<0> = 1;
```

Variable templates...

- ... can be constexpr
- ... can be defined recursively -> specialization to define the base case

Useage

Template-ID (Name and template arguments)

```
static_assert(factorial<0> == 1);
static_assert(factorial<5> == 120);
```

Base case required for printAll because of the recursion

```
void printAll() {
}

template<typename First, typename...Types>
void printAll(First const & first, Types const &...rest) {
   std::cout << first;
   if (sizeof...(Types)) {
      std::cout << ", ";
   }
   printAll(rest...);
}</pre>
```

Compile-time conditional inclusion statement

```
void printAll() {
}

template<typename First, typename...Types>
void printAll(First const & first, Types const &...rest) {
   std::cout << first;
   if constexpr (sizeof...(Types)) {
     std::cout << ", ";
     printAll(rest...);
   }
}</pre>
```

Requires compile-time expression

• Instance for printAll("Hello"s);

```
void printAll(std::string const & first) {
  std::cout << first;
  if constexpr (0) {
    std::cout << ", ";
    printAll(); //rest... expansion
  }
}</pre>
void printAll(std::string const & first) {
  std::cout << first;
  std::cout << first;
}
```

We don't need a base case anymore!

```
?
```

```
template <std::size_t N>
constexpr size_t McCarthy91 = [] {
   if constexpr (N <= 100) {
     return McCarthy91<McCarthy91<(N + 11)>>;
   }
   return N - 10;
   }();
```

```
struct Point {
  constexpr Point(double x, double y)
    : x{x}, y{y}{}
  Point() : Point{0.0, 0.0}{}
  private:
    double x;
    double y;
};
constexpr Point origin{0.0, 0.0};
```

```
?
```

```
template <std::size_t N>
constexpr size_t McCarthy91 = [] {
   if constexpr (N <= 100) {
      return McCarthy91<McCarthy91<(N + 11)>>;
   }
   return N - 10;
   }();
```

Correct

Variable templates can be initialized by lambdas, which are literal types, as long as there is a base case. With constexpr if we reach that base case.

```
struct Point {
  constexpr Point(double x, double y)
    : x{x}, y{y}{}
  Point() : Point{0.0, 0.0}{}
private:
  double x;
  double y;
};

constexpr Point origin{0.0, 0.0};
```

Correct

Point is a literal type. Not all constructors need to be constexpr. However, in compile-time constant expressions only the constexpr constructors can be used.

- Only preprocessor macros and templates (class & function)
- Templates allow Turing-complete computation
 - But it looks ugly and can be hard to understand
 - Compiler error-messages can be lengthy
 - Compile-time long and recursion depth limited (depends on the compiler and its settings)
- Almost only single values can be computed
 - Arrays/structs possible, but ugly
- Observation: static const member variables of class templates are variables that depend on template parameters

Substitution Failure Is Not An Error (SFINAE)







- What do you expect is the return value of the program on the right?
 - None, the program does not compile!

```
Name increment is looked up
increment(unsinged)
                     increment Template
  Template arguments are deduced
increment(unsigned)
                    increment<int>(int)
      Best overload is selected
            increment(int)
```

```
unsigned increment(unsigned i) {
  return i++;
}

template<typename T>
T increment(T value) {
  return value.increment();
}

int main() {
  return increment(42);
}
```

```
error: request for member 'increment' in 'value',
which is of non-class type 'int'
```

- During overload resolution the template parameters in a template declaration are substituted with the deduced types
 - This may result in template instances that cannot be compiled
 - Or otherwise suboptimal selection
- If the substitution of template parameter fails that overload candidate is discarded
- Substitution failure might happen in
 - Function return type
 - Function parameter
 - Template parameter declaration
 - And expressions in the above
- Errors in the instance body are still errors

```
Substitution

template<typename T>

T increment(T value) {
  return value.increment();
}

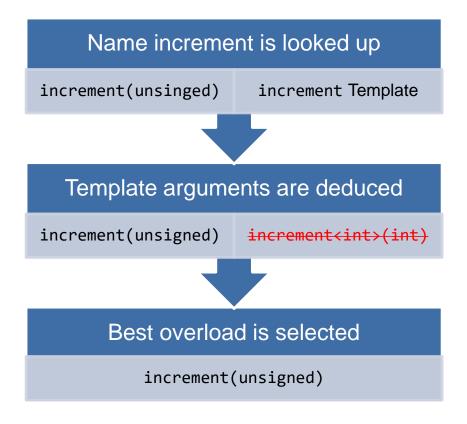
int increment(int value) {
  return value.increment();
}

//T = int
```

```
template<typename T>
auto increment(T value) -> decltype(value.increment()) {
  return value.increment();
}
```

- We can break the return type
- If we tell the compiler to use the type of value.increment() as return type for increment<int>
 - That type cannot be determined during substitution

 Since there is a problem during substitution that overload is discarded



Now the result is 42

```
unsigned increment(unsigned i) {
  return i++;
template<typename T>
auto increment(T value) ->
    decltype(value.increment()) {
  return value.increment();
int main() {
  return increment(42);
```

- This approach, using decltype(...) as trailing return type, is infeasible in general
 - Function might have return type void
 - It is not elegant for complex bodies

Let's examine our example again

We want the increment template to be selected only for class type arguments

There exists a template std::is_class<T>

- contains static constexpr bool value;
- value is true if T is a class, false otherwise
- Variable template: std::is_class_v<T>
 - Type bool (direct access of ::value)

Can we apply this to the increment template directly?

- No, either value (true or false) is still valid
- We need something to create an error in the type of the function

```
template<typename T>
T increment(T value) {
  return value.increment();
}
int main() {
  return increment(42);
}
```

```
#include <type_traits>
struct S {};
int main() {
   std::is_class<S>::value; //true
   std::is_class<int>::value; //false
}
```

```
template<bool expr, typename T = void>
struct enable_if;
```

- The std::enable_if_t template takes an expression and a type
 - If the expression evaluates to true std::enable if t represents the given type
 - Otherwise it does NOT represent a type

• Inside std::enable_if

```
template<typename T>
struct enable_if<true, T> {
  using type = T;
};
```

```
template<typename T, >

T increment( T value) {
  return value.increment();
}
```

Spots to apply enable_if (SFINAE)

Possibilities

```
template<typename T>
std::enable if t<std::is class<T>::value, T> increment(T value) {
  return value.increment();
template<typename T>
                                                                     enable_if as parameter
T increment(std::enable_if_t<std::is_class<T>::value, T> value) {
                                                                      impairs type deduction
  return value.increment();
template<typename T, typename = std::enable_if_t<std::is_class<T>::value, void>>
T increment(T value) {
  return value.increment();
                                                                    would be void per default
```

Example: Box-Container with

- Default constructor
- Copy constructor
- Move constructor
- Size constructor

Box<MemoryOperationCounter> b{1};

```
template <typename T>
struct Box {
    Box() = default;
    Box(Box const & box)
        : items{box.items}{}
    Box(Box && box)
        : items{std::move(box.items)} {}
    explicit Box(size_t size)
        : items(size) {}
    //...
private:
    std::vector<T> items{};
};
```

 What if we replace the copy/move constructors with a forwarding constructor?

Box<MemoryOperationCounter> b{1};

```
template <typename T>
struct Box {
  Box() = default;
  template <typename BoxType>
explicit Box(BoxType && other)
    : items(std::forward<BoxType>(other).items) {}
  explicit Box(size t size)
    : items(size) {}
  //...
private:
  std::vector<T> items{};
};
```

- We don't want the forwarding constructor to match anything else than Boxes
 - Type traits can be used to narrow down the valid calls.

```
template <typename T>
struct Box {
   Box() = default;
   template <typename BoxType, typename = std::enable_if_t<std::is_same_v<Box, BoxType>>>
   explicit Box(BoxType && other)
     : items(std::forward<BoxType>(other).items) {}
   explicit Box(size_t size)
     : items(size) {}
   //...
private:
   std::vector<T> items{};
};
```

- The standard library provides many predefined checks for type traits¹
- A trait contains a boolean value
- Usually they are available in two versions
- Example:

```
std::is_same<T, U>
```

std::is_same_v<T, U>

```
template <typename T, typename U>
struct is same : false type {
 // inherits;
 // static constexpr bool value = false;
};
template <typename T>
struct is_same<T, T> : true_type {
 // inherits;
 // static constexpr bool value = true;
};
template <typename T, typename U>
constexpr bool is_same_v = is_same<T, U>::value;
```

¹ <u>https://en.cppreference.com/w/cpp/header/type_traits</u>

```
?
```

```
template <typename T>
std::enable_if_t<
    std::negation_v<
        std::is_reference<T>
    >
    consume(T && value) { /*...*/ }

//This function can only be called with rvalues
```

```
template <typename T>
int convert(T value) {
  using namespace std;
  enable_if_t<
    is_constructible_v<int, T>, int
  > converted{value};
  return converted;
}
//This function can be eliminated by SFINAE
```

```
template <typename T>
std::enable_if_t<
    std::negation_v<
        std::is_reference<T>
    >
    consume(T && value) { /*...*/ }

//This function can only be called with rvalues
```

Correct

std::is_reference is std::true_type if T is a reference otherwise std::false_type. If the forwarding reference is initialized with an Ivalue T is a reference. std::negation_v negates the true/false_type. The overload is only enabled for rvalues.

```
template <typename T>
int convert(T value) {
  using namespace std;
  enable_if_t<
    is_constructible_v<int, T>, int
  > converted{value};
  return converted;
}
//This function can be eliminated by SFINAE
```

Incorrect

This is not SFINAE! It will compile if called with something that can be used to construct an int, but the overload will be taken anyway.

- Many computations for which the arguments are known upfront can be computed at compile time
- Until C++20 dynamic memory allocation is not possible
- All literal types can be used in constexpr contexts
- SFINAE is used to eliminate overload candidates

Self Study



Compile time code





Have a look at the following code on Godbolt and the binary that is generated

```
constexpr auto factorial(unsigned n) {
  auto result = 1u;
 for (auto i = 2u; i <= n; i++) {
    result *= i;
  return result;
constexpr auto factorial0f5 = factorial(5);
int main() {
  static_assert(factorialOf5 == 120);
  return factorialOf5;
```

```
1 main:
2 push rbp
3 mov rbp, rsp
4 mov eax, 120
5 pop rbp
6 ret
```

https://www.godbolt.org/z/suMDYp

No calls to the factorial function are present anymore