

RSA Encryption

CS 5158/6058 Data Security and Privacy

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Exponentiation in Group

- Group G , group order m ,
 - A function $f_e: G \rightarrow G$: exponentiation on element g with integer e
 - If G is additive, $f_e(g) = e^*g$,
 - If G is multiplicative, $f_e(g) = g^e$
- If $\gcd(e, m) = 1$, then f_e is a **permutation** (bijection)

Exponentiation in Group

- A function $f_e: G \rightarrow G: f_e(g) = e * g$ (additive group)
 - If $\gcd(e, m) = 1$, then f_e is a **permutation** (bijection)
- E.g., $Z_5 = \{0, 1, 2, 3, 4\}$ is an additive group,
 - if $e = 2$, $\gcd(e, m) = \gcd(2, 5) = 1$
 - $f_e(0) = 2 * 0 = 0$; $f_e(1) = 2 * 1 = 2$; $f_e(2) = 2 * 2 = 4$;
 - $f_e(3) = 2 * 3 = 6 = 1 \pmod{5}$;
 - $f_e(4) = 2 * 4 = 8 = 3 \pmod{5}$
 - $f_e: \{0, 2, 4, 1, 3\}$ a permutation of Z_5

Exponentiation in Group

- A function $f_e: G \rightarrow G: f_e(g) = e * g$ (additive group)
 - If $\gcd(e, m) = 1$, then f_e is a **permutation** (bijection)
- E.g., $Z_6 = \{0, 1, 2, 3, 4, 5\}$ is an additive group,
 - if $e = 2$, $\gcd(e, m) = \gcd(2, 6) = 2 \neq 1$
 - $f_e(0) = 2 * 0 = 0$; $f_e(1) = 2 * 1 = 2$; $f_e(2) = 2 * 2 = 4$;
 - $f_e(3) = 2 * 3 = 6 = 0 \bmod Z_6$;
 - $f_e(4) = 2 * 4 = 8 = 2 \bmod Z_6$
 - $f_e(5) = 2 * 5 = 10 = 4 \bmod Z_6$
 - $f_e: \{0, 2, 4\}$ is not a permutation of Z_6

Exponentiation in Group

- Practice: $Z_7 = \{0, 1, 2, 3, 4, 5, 6\}$ is an additive group,
 - Group order $m = ??$ identity ??
 - If $e = 2$, is f_e a permutation of Z_7 ?
 - Compute $f_e(g)$, for all the elements in Z_7
 - $\gcd(e, m) = \gcd(2, 7) = 1$, f_e is a permutation
 - $f_e(0) = 2*0 = 0$; $f_e(1) = 2*1 = 2$; $f_e(2) = 2*2 = 4$;
 - $f_e(3) = 2*3 = 6$; $f_e(4) = 2*4 = 8 = 1 \bmod Z_7$
 - $f_e(5) = 2*5 = 10 = 3 \bmod Z_7$
 - $f_e(6) = 2*6 = 12 = 5 \bmod Z_7$
 - $f_e: \{0, 2, 4, 6, 1, 3, 5\}$

Exponentiation in Group

- A function $f_e: G \rightarrow G: f_e(g)$, group order is m
 - if $ed = 1 \pmod m$, then f_d is inverse of f_e
 - $f_d(f_e(g)) = g$

- If G is additive

$$f_d(f_e(g)) = f_d(eg) = deg = (ed \pmod m)g = g$$

- If G is multiplicative

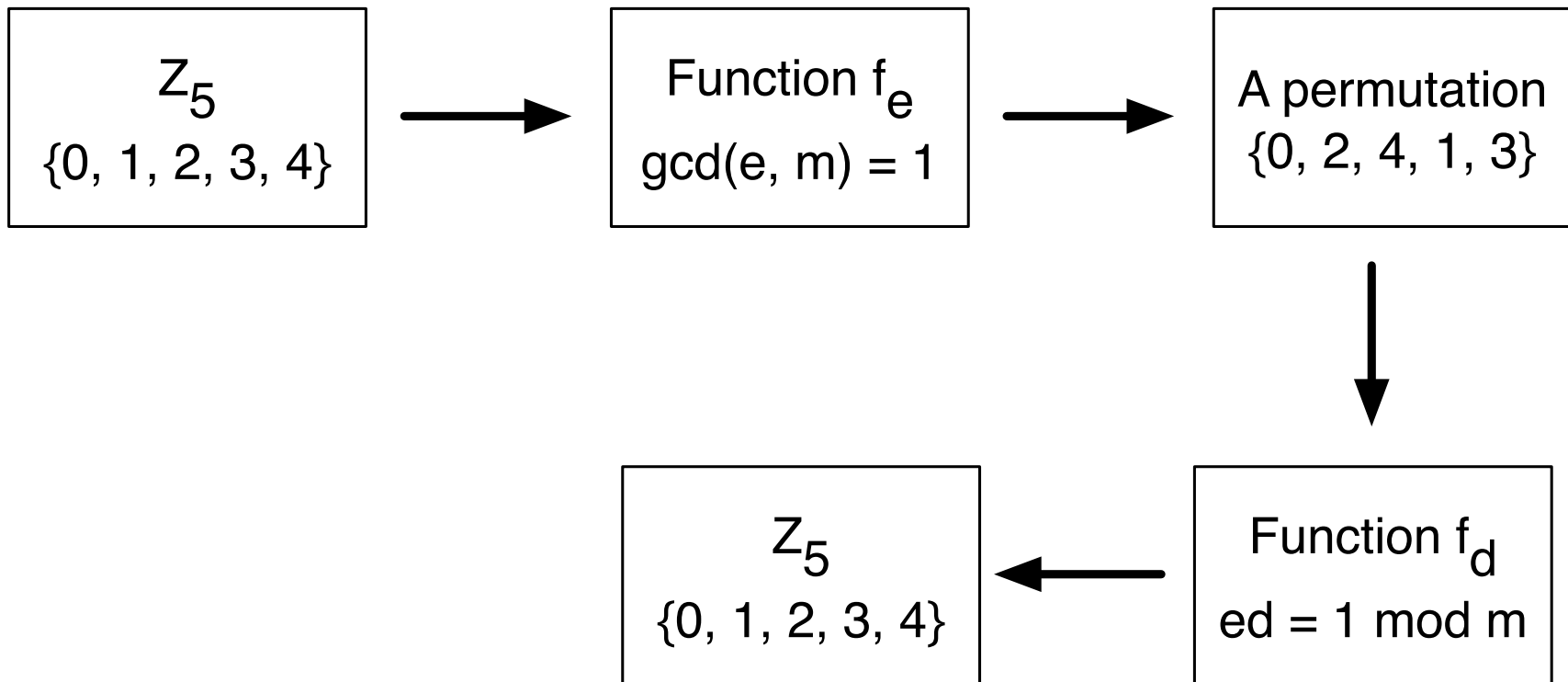
$$f_d(f_e(g)) = f_d(g^e) = (g^e)^d = g^{ed \pmod m} = g^1 = g$$

Exponentiation in Group

- Example: $Z_5 = \{0, 1, 2, 3, 4\}$ is an additive group,
 - if $e = 2$, $\gcd(e, m) = \gcd(2, 5) = 1$
 - $f_e : \{0, 2, 4, 1, 3\}$ a permutation of Z_5
- If $ed = 1 \pmod m$, then f_d is inverse of f_e
 - $e = 2$, $d = 3$, $ed = 3 \cdot 2 \pmod m = 1 \pmod 5$
 - $f_d(0) = 3 \cdot 0 = 0$; $f_d(2) = 3 \cdot 2 = 1$;
 - $f_d(4) = 3 \cdot 4 = 12 = 2 \pmod{Z_5}$; $f_d(1) = 3 \cdot 1 = 3$;
 - $f_d(3) = 3 \cdot 3 = 9 = 4 \pmod{Z_5}$
 - $f_d : \{0, 1, 2, 3, 4\}$ inverse of f_e

Exponentiation in Group

- A function $f_e: G \rightarrow G: f_e(g)$
 - If $ed = 1 \pmod m$, then f_d is inverse of f_e



Exponentiation in Group

- Practice: $Z_7 = \{0, 1, 2, 3, 4, 5, 6\}$ is an additive group,
 - $e = 2$, f_e is a permutation, $f_e: \{0, 2, 4, 6, 1, 3, 5\}$
 - what is $d = ??$ and compute f_d with all the inputs
 - $e * d = 2 * 4 = 1 \bmod 7$, $d = 4$
 - $f_d(0) = 4 * 0 = 0$; $f_d(2) = 4 * 2 = 1$;
 - $f_d(4) = 4 * 4 = 16 = 2 \bmod Z_7$;
 - $f_d(6) = 4 * 6 = 24 = 3 \bmod Z_7$; $f_d(1) = 4 * 1 = 4$
 - $f_e(3) = 4 * 3 = 12 = 5 \bmod Z_7$
 - $f_e(5) = 4 * 5 = 20 = 6 \bmod Z_7$
 - $f_d: \{0, 1, 2, 3, 4, 5, 6\}$, f_d is inverse of f_e

Group under Multiplication

- $Z_N = \{0, 1, 2, \dots, N-1\}$ is a group under addition but **not multiplication**
 - If identity is 1, then some do not have inverse
 - E.g., 2 does not have an inverse, since $1/2$ is not an element of Z_N , i.e. 2 is not invertible
- Find a group under multiplication mod N
 - Remove all the elements in Z_N that are **not invertible**
 - If b and N are **relatively prime**, i.e., $\gcd(b, N) = 1$, then b is invertible mod N

Group under Multiplication

- Find a group under multiplication mod N
 - Remove elements in \mathbb{Z}_N that are not invertible
 - If b and N are **relatively prime**, i.e., $\gcd(b, N) = 1$, then b is invertible mod N
- \mathbb{Z}_N^* has all the integers in $\{1, 2, \dots, N-1\}$ that are relatively prime to N
$$\mathbb{Z}_N^* = \{b \in \{1, \dots, N-1\} \mid \gcd(b, N) = 1\}$$
- \mathbb{Z}_N^* is a multiplicative group

Group under Multiplication

- Example: $N = 6$, what are the elements in Z_6^* ?
 - $Z_6 = \{0, 1, 2, 3, 4, 5\}$,
 - **remove 0** first
 - $\gcd(1, N) = 1$, relatively prime
 - $\gcd(2, N) = 2$, not relatively prime, **remove 2**
 - $\gcd(3, N) = 3$, not relatively prime, **remove 3**
 - $\gcd(4, N) = 2$, not relatively prime, **remove 4**
 - $\gcd(5, N) = 1$, relatively prime
 - $Z_6^* = \{1, 5\}$ is a multiplicative group

Group under Multiplication

- Practice: $N = 9$, what are the elements in Z_9^* ?
 - $Z_9 = \{0, 1, 2, 3, 4, 5, 6, 7, 8\}$,
 - **remove 0** first
 - $\gcd(1, N) = 1$, $\gcd(2, N) = 1$,
 - $\gcd(3, N) = 3 \neq 1$, **remove 3**
 - $\gcd(4, N) = 1$, $\gcd(5, N) = 1$,
 - $\gcd(6, N) = 3 \neq 1$, **remove 6**
 - $\gcd(7, N) = 1$, $\gcd(8, N) = 1$
 - $Z_9^* = \{1, 2, 4, 5, 7, 8\}$

Group under Multiplication

- Practice: $N = 11$, what are the elements in Z_{11}^* ?
 - $Z_{11} = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$
 - remove 0 first
 - $N=11$ is a prime, for any element $b > 0$ in Z_{11}
 - $\gcd(b, N) = 1$
 - no need to remove any b
- $Z_{11}^* = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$

Group under Multiplication

- Example: $N = 9$, what are the elements in Z_9^* ?
 - $Z_9 = \{0, 1, 2, 3, 4, 5, 6, 7, 8\}$; $Z_9^* = \{1, 2, 4, 5, 7, 8\}$
- $|Z_N^*|$: Order of group Z_N^*
 - $|Z_N^*| = \phi(N)$, $\phi(\cdot)$ is called Euler phi function
 - The number of the integers in $\{1, 2, \dots, N-1\}$ that are relatively prime to N
 - If N is a prime, then $\phi(N) = |Z_N^*| = N-1$
 - E.g., $N=7$ is a prime, $|Z_7^*|=6$, $Z_7^* = \{1, 2, 3, 4, 5, 6\}$

Group under Multiplication

- If $N = pq$, p, q are primes, $\phi(N) = |Z_N^*| = (p-1)(q-1)$
 - Why $|Z_N^*|$ is equal to $(p-1)(q-1)$??
- $|Z_N^*|$: the number of the integers in $\{1, 2, \dots, N-1\}$ that are relatively prime to N
- If a in $\{1, \dots, N-1\}$ is not relatively prime to N
 - $\gcd(a, N) \neq 1$
 - $\gcd(a, N) = p$ or $\gcd(a, N) = q$
 - $p \mid a$ or $q \mid a$

Group under Multiplication

- If $N = pq$, p, q are primes, then $|Z_N^*| = (p-1)(q-1)$
 - There are $(q-1)$ elements, s.t., $\gcd(a, N) = p$
 - Elements divided by p : $p, 2p, \dots, (q-1)p$
 - There are $(p-1)$ elements, s.t., $\gcd(a, N) = q$
 - Elements divided by q : $q, 2q, \dots, (p-1)q$
- No. of elements s.t., $\gcd(a, N) \neq 1$
 - $(q-1) + (p-1)$
- No. of elements s.t., $\gcd(a, N) = 1$
 - $N - 1 - (q-1) - (p-1) = (p-1)(q-1) = |Z_N^*|$

Group under Multiplication

- If $N = pq$, p, q are primes, then $|Z_N^*| = (p-1)(q-1)$
- Example: $N = 6 = 2 \cdot 3$, $Z_6 = \{0, 1, 2, 3, 4, 5\}$
 - $p = 2$, $q = 3$, p and q are primes
 - $(q-1) = 2$ elements, s.t., $\gcd(a, N) = p = 2$
 - $a = \{2, 4\}$
 - $(p-1) = 1$ elements, s.t., $\gcd(a, N) = q = 3$
 - $a = \{3\}$
 - Remove $\{0\}$, $\{2, 4\}$ and $\{3\}$
 - $Z_6^* = \{1, 5\}$ $|Z_6^*| = (2-1)(3-1) = 2$

- If $N = pq$, p, q are primes, then $|Z_N^*| = (p-1)(q-1)$
- Practice: $N = 21 = 3 \cdot 7$, $Z_{21} = \{0, 1, 2, \dots, 20, 21\}$
 - What are the elements in Z_{21}^* ?
 - How many elements in Z_{21}^* ? Or $\phi(N) = |Z_{21}^*| = ?$
 - $p = 3$, $q = 7$, p and q are primes
 - $(q-1) = 6$ elements, s.t., $\gcd(a, N) = p = 3$
 - $a = \{3, 6, 9, 12, 15, 18\}$
 - $(p-1) = 2$ elements, s.t., $\gcd(a, N) = q = 7$
 - $a = \{7, 14\}$
 - Remove $\{0\}$, $\{3, 6, 9, 12, 15, 18\}$ and $\{7, 14\}$
 - $Z_{21}^* = \{1, 2, 4, 5, 8, 11, 13, 16, 17, 19, 20\}$
 - $|Z_{21}^*| = (3-1)(7-1) = 2 \cdot 6 = 12$

Group under Multiplication

- If group order is m , then $g^m = 1$, for any g in G
 - Exponentiation on element g with integer m is equal to group identity 1
- \mathbb{Z}_N^* is a multiplicative group mod N
 - Exponentiation on element g with integer $|\mathbb{Z}_N^*|$ is equal to group identity 1

$$g^{|\mathbb{Z}_N^*|} = g^{\phi(N)} = 1 \pmod{N}$$

Group under Multiplication

- $Z_7 = \{0, 1, 2, 3, 4, 5, 6\}$, $N=7$, group order is $m=7$
- $Z_7^* = \{1, 2, 3, 4, 5, 6\}$, $N=7$, group order is $m=6$

$$g^{|Z_N^*|} = g^{\phi(N)} = 1 \pmod{N}$$

- $1^6 = 1 \pmod{N}$, $2^6 = 64 = 1 \pmod{N}$,
 - $3^6 = 729 = 1 \pmod{N}$, $4^6 = 4096 = 1 \pmod{N}$,
 - $5^6 = 15625 = 1 \pmod{N}$, $6^6 = 46656 = 1 \pmod{N}$
- Z_7 is additive, Z_7^* is multiplicative

Group under Multiplication

- $Z_6 = \{0, 1, 2, 3, 4, 5\}$, $N=2*3$, group order is $m=6$
- $Z_6^* = \{1, 5\}$, $N=2*3$, group order is $m=(2-1)(3-1)=2$

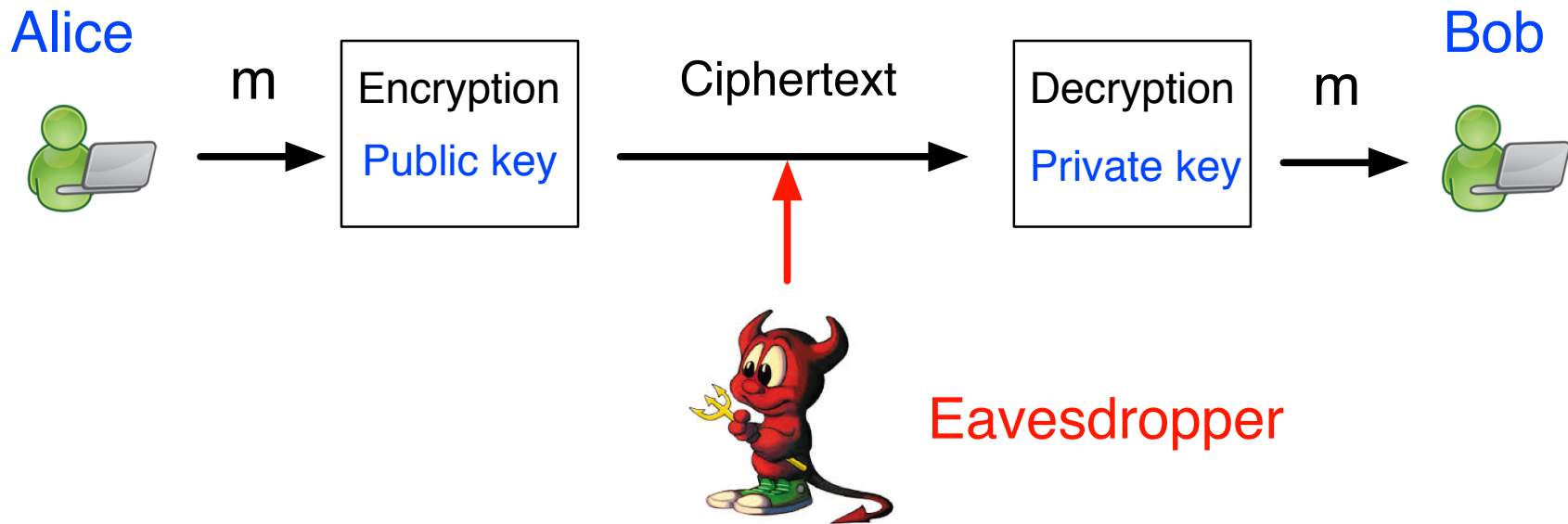
$$g^{|Z_N^*|} = g^{\phi(N)} = 1 \pmod{N}$$

- $1^2 = 1 \pmod{N}$, $5^2 = 25 = 1 \pmod{N}$,
- Z_6 is additive, Z_6^* is multiplicative

Group under Multiplication

- Group G , order m
 - A function $f_e: G \rightarrow G: f_e(g) = g^e$
 - If $\gcd(e, m) = 1$, then f_e is a permutation
 - $ed \equiv 1 \pmod m$, then f_d is inverse of f_e
- Group Z_N^* , order $m = |Z_N^*| = \phi(N)$
 - A function $f_e: Z_N^* \rightarrow Z_N^*: f_e(g) = g^e$
 - If $\gcd(e, |Z_N^*|) = 1$, then f_e is a permutation
 - $ed \equiv 1 \pmod{|Z_N^*|}$, then f_d is inverse of f_e

Public-Key Encryption



- Alice obtains Bob' public key from public channels
- Alice encrypts with public key
- Bob decrypts with private key

Public-Key Encryption

- PKE includes 3 algorithms
- KeyGen: given a security parameter 1^n , output a pair of keys (pk, sk) , where pk is public and sk is private.
- Enc: given a public key pk and a message m , output a ciphertext c
- Dec: given a private key sk and a ciphertext c , output a message m .

One-Way Function

- One-way function is used in public-key crypto
 - **Easy** to compute
 - **Hard** to invert: (with polynomial-time algorithm)
 - E.g., factoring, discrete-log problem
- **Factoring** (integer factorization)
 - Assume p, q are large primes
 - Easy: $p, q \longrightarrow N = p \cdot q$
 - Hard: $N \longrightarrow p, q$

RSA

- Rivest-Shamir-Adleman (RSA), 1978
 - Widely used today, SSL/TLS, email, HTTPS, etc.
 - ACM Turing award in 2002



Textbook RSA

- KeyGen: given a security parameter 1^n , generate two n -bit primes p, q , compute $N = pq$, choose e s.t. $\gcd(e, \phi(N)) = 1$, compute $d = e^{-1} \bmod \phi(N)$, output public key $pk = (N, e)$, private key $sk = d$
- Enc: given a message m and a public key $pk = (N, e)$, return $c = m^e \bmod N$
- Dec: given a ciphertext c and a private key $sk = d$, return $m = c^d \bmod N$

Correctness

- For group \mathbb{Z}_N^* $g^{|\mathbb{Z}_N^*|} = g^{\phi(N)} = 1 \pmod{N}$

- We know in RSA:

$$ed = 1 \pmod{\phi(N)} \quad \text{Enc}_{pk}(m) = c = m^e \pmod{N}$$

$$\begin{aligned} \text{Dec}_{sk}(c) = c^d &= (m^e)^d \\ &= m^{ed} \pmod{\phi(N)} \\ &= m \pmod{N} \end{aligned}$$

- m, c are elements of group \mathbb{Z}_N^*
- e, d are integers

Textbook RSA

- Example: $p \cdot q = 17 \cdot 23 = 391 = N$
 - $\phi(N) = |Z_N^*| = (p-1)(q-1) = 16 \cdot 22 = 352$
 - choose e , s.t. $\gcd(e, \phi(N)) = 1$
 - $e = 3$, $\gcd(3, 352) = 1$
 - $d = 235$, $e \cdot d = 3 \cdot 235 = 705 = 1 \pmod{352}$
 - $pk = (N, e) = (391, 3)$, $sk = d = 235$
- Given $m=158$, $c=m^e \pmod N = 158^3 = 295 \pmod{391}$
- Given $c=295$, $m=c^d \pmod N = 295^{235} = 158 \pmod{391}$

Textbook RSA

- Practice: $p \cdot q = 11 \cdot 7 = 77 = N$
 - $\phi(N) = |Z_N^*| = (p-1)(q-1) = 10 \cdot 6 = 60$
 - choose e , s.t. $\gcd(e, \phi(N)) = 1$
- Can we choose $e = 3$ in KeyGen?
 - **No**, $\gcd(e, \phi(N)) = \gcd(3, 60) = 3 \neq 1$
- Can we choose $e = 7$ in KeyGen?
 - **Yes**, $\gcd(e, \phi(N)) = \gcd(7, 60) = 1$

Textbook RSA

- Practice: $p \cdot q = 11 \cdot 7 = 77 = N$
 - $\phi(N) = |Z_N^*| = (p-1)(q-1) = 10 \cdot 6 = 60$
 - choose e , s.t. $\gcd(e, \phi(N)) = 1$
 - $e = 7$,
 - $d = 43$, $ed = 7 \cdot 43 = 1 \pmod{60}$
 - What is $pk = ?$ what is $sk = ?$
 - $pk = (N, e) = (77, 7)$, $sk = d = 43$
 - Given $m = 2$, what is $c = m^e$?
 - $c = m^e = 2^7 = 128 = 51 \pmod{77}$

Additional Reading

Chapter 8, *Introduction to Modern Cryptography*, Drs.
J. Katz and Y. Lindell, 2nd edition