# Lexical Analysis Programming Languages

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## Language Processing Compilers



Example – Natural Languages

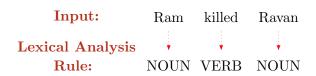
Input: Ram killed Ravan

Example – Natural Languages

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Rule: NOUN VERB NOUN

Example – Natural Languages



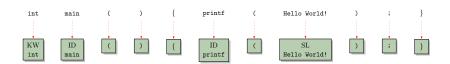
Example – Programming Languages

```
int main ( ) \left\{ \begin{array}{cccc} & printf \end{array} \right. ( Hello World! ) ; \left. \right\}
```

#### **Terminology**

- Token classes/Tokens. e.g. KW, ID, (, } etc.
- Lexemes. e.g. "int", "main", "(", "}" etc.

Example – Programming Languages

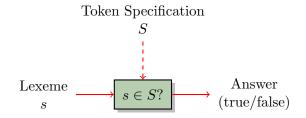


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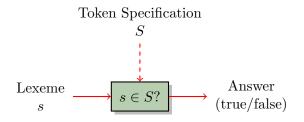
#### Specification of Token Classes

Example



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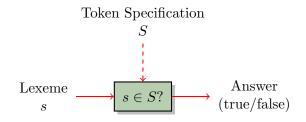
Example



Specification (S) of token classes

#### Specification of Token Classes

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Specification (S) of token classes –  $regular\ expressions$ 

Example – identifier

- **■** English:
  - Must comprise of only alphabetic character  $(\alpha)$

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Rules of Construction

- $\bullet$  (empty symbol)
- 2 Let r and s be two regular expressions:

Expression	Meaning
r s	Choice
rs	Concatenation
$r\star$	Zero or more
r+	One or more
r?	Zero or one

Rules of Construction

- $\bullet$  (empty symbol)
- 2 Let r and s be two regular expressions:

Expression	Meaning	Example	Instance
r s	Choice	$\alpha N$	a, b,, 1, 2,
rs	Concatenation	rs	a1, b1,, a2,
$r\star$	Zero or more	$\alpha\star$	$\epsilon$ , a, ab, aaa,
r+	One or more	N+	1, 11, 12,
r?	Zero or one	N?	$\epsilon$ , 1, 2,

Example – Identifier

#### **■** English:

- 1 Must start with an alphabetic character
- 2 Subsequently, may have either an alphabetic character  $(\alpha)$ , a numeric character (N), separated by zero or more underscores  $('_{-}'_{-})$ .

Example – Identifier

#### English:

- 1 Must start with an alphabetic character
- 2 Subsequently, may have either an alphabetic character  $(\alpha)$ , a numeric character (N), separated by zero or more underscores  $('_-)$ .
- Regular Expression:  $\alpha\{[(\alpha|N)(\dot{\ }_{-}\dot{\ }\star)]\star(\alpha|N)\}$ ?

## Regular Expressions Activity

Regular expression for floating point numbers in C programming language

#### Implementation of Lexical Analysis

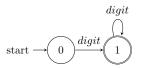
Finite State Automata

### **Example** NUMBER:

#### Implementation of Lexical Analysis

Finite State Automata

### Example NUMBER:



return NUMBER

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Overview

- Finite state automata
- Implementation of finite state automata
- Implementation of lexical analysers: manual and automated generation