70分:

```
plaintext = input("Please input plaintext: ")
P1 = input("Please input Prime1: ")
P2 = input("Please input Prime2: ")
#while(1):
#    P1 = 1
#    while not(miller_rabin(P1,512)):
#        P1 = random.getrandbits(1024)
#    P2 = 1
#    while not(miller_rabin(P2,512)):
#        P2 = random.getrandbits(1024)
#    if(P1*P2>=int(plaintext)and P1!=P2 and P1*P2>=1.340781e+154):
#        break
E,D,CRT_D = RSA_GO(int(P1),int(P2),int(plaintext))
```

```
Please input plaintext: 2018
Please input Prime1: 71
Please input Prime2: 83
Cipher is 2221
DeCipher is 2018
CRT DeCipher is 2018
```

P1&P2:

產生兩個大質數,驗證次數為 512 次,並且兩質數相乘 N 要大於等於 1024bit

```
while(1):
    P1 = 1
    while not(miller_rabin(P1,512)):
        P1 = random.getrandbits(1024)
    P2 = 1
    while not(miller_rabin(P2,512)):
        P2 = random.getrandbits(1024)
    if(P1*P2>=int(plaintext)and P1!=P2 and P1*P2>=1.340781e+154):
        break
```

```
def miller rabin(N, validate count):
    if (N<2):
        return False
    if(N==3 or N==2):
        return True
    if(N%2==0):
        return False
    u_{x} = 1, N-1
    r = x
    while x \% pow(2,u) != 0:
        r = x / pow(2,u)
        u+=1
    for i in range(validate count):
        candidate = random.randrange(2,N-1)
        b = pow(candidate,r,N)
        if(b==1 \text{ or } b==N-1):
            continue
        for j in range(u-1):
            b = pow(b, 2, N)
            if(b==N-1):
                 break
            elif(b==1):
                 return False
        else:
            return False
    return True
```

P3:

```
def square_mul(x,y,N):#x^y
   output = x
   for i in y[1:]:
        output = pow(output,2) % N
        if(i=='1'):
             output = output * x % N
        return output
```

Р4

```
def CRT(D,P1,P2,Cipher):
    Dp = D % (P1-1)
    Dq = D % (P2-1)
    (Xq,_,_) = ext_GCD(P2,P1)
    return (Xq*P2*(square_mul(Cipher,bin(Dp)[2:],P1))+(1-Xq*P2)*(square_mul(Cipher,bin(Dq)[2:],P2))) % (P1*P2)
```

整個 RSA 流程:

找到兩質數後,算 N 和 PHI,再找出和 PHI 互質的數 e,透過 Extended Euclidean algorithm,得出 e_inverse d,透過 Square & multiply 加速運算,分別

使用一般的和 Chinese Remainder Theorem 進行解密,得出結果

```
def RSA_GO(P1,P2,text):
    N = P1*P2
    PHI = (P1-1)*(P2-1)
    for i in range(2,PHI):
        if(gcd(i,PHI)==1):
            e = i
            break
    (d,_,_) = ext_GCD(e,PHI)

    cipher = square_mul(text,bin(e)[2:],N)
    Decipher = square_mul(cipher,bin(d % PHI)[2:],N)
    CRT_Decipher = CRT(d % PHI,P1,P2,cipher)
    return cipher,Decipher,CRT_Decipher
```

```
plaintext = input("Please input plaintext: ")
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while(1):
   P1 = 1
    while not(miller_rabin(P1,512)):
        P1 = random.getrandbits(1024)
    P2 = 1
    while not(miller_rabin(P2,512)):
        P2 = random.getrandbits(1024)
    if(P1*P2>=int(plaintext)and P1!=P2 and P1*P2>=1.340781e+154):
       break
E,D,CRT D = RSA GO(int(P1),int(P2),int(plaintext))
Please input plaintext: 2018
Cipher is 33466154331649568
DeCipher is 2018
```

CRT DeCipher is 2018