# Fault Tolerant IVY Protocol (Integrated shared Virtual memory at Yale)

Go-based implementation of fault-tolerant Ivy Architecture with a backup central manager for consistent (meta)data handling during primary CM failures.

# How to run the code

There are 2 types of nodes, Central Manager and Client. Each node runs on separate terminals on separate processes and communicate with each other using RPC calls. For a fresh run of the protocol, ensure the /data folder is empty.

## Running the different types of Nodes

To start the primary CM:

- 1. Ensure you are in the root directory of the project.
- 2. Run go build && ./ivy
- 3. You will be prompted to choose the node type: "Enter Node type ('1': CM, '2': Client, 'restartCM', 'restartBackup')"
- 4. Type '1'. This will create a new cm. json file and add the primary CM object to the file.
- 5. The primary CM should now be running.

To start the Backup CM:

- 1. Run go build && ./ivy
- 2. You will be prompted to choose the node type: "Enter Node type ('1': CM, '2': Client, 'restartCM', 'restartBackup')"
- 3. Type '1'. This will see that a primary CM already exists in cm. json and add a Backup CM object to the file.
- 4. The Backup CM should now be running.

To start a Client:

- 1. Run go build && ./ivy
- 2. You will be prompted to choose the node type: "Enter Node type ('1': CM, '2': Client, 'restartCM', 'restartBackup')"
- 3. Type '1'. This will check client.json and add Client (currentHighestID + 1) to the file.
- 4. The Client should now be running

## How to kill any Node (PrimaryCM/BackupCM/Client)

To kill any node simply go to its terminal and press ctrl+c

# How to reboot a CM (PrimaryCM/BackupCM)

To reboot a PrimaryCM:

- 1. Run go build && ./ivy
- 2. You will be prompted to choose the node type: "Enter Node type ('1': CM, '2': Client, 'restartCM', 'restartBackup')"
- 3. Type restartCM. This checks cm. json for the IP of the original Primary CM and runs the CM on that IP.

To reboot a BackupCM:

- 1. Run go build && ./ivy
- 2. You will be prompted to choose the node type: "Enter Node type ('1': CM, '2': Client, 'restartCM', 'restartBackup')"
- 3. Type restartBackupCM. This checks cm. json for the IP of the Backup CM and runs the CM on that IP.

# How to send read/write requests

- Send a write request by typing writePage <pageNo> <content> . For example, you type writePage P1 Content1 .
- Send a read request by typinh readPage <pageNo> . For example, you type readPage P1

### Useful command

- CM: Type print to view the MetaData.
- Client: Type print to view the PageStore

# **Fault Tolerance**

# Syncing MetaData

When you have 2 CMs running: one Primary CM and one Backup CM, the Backup CM has runs goroutine ( go pulseCheck()). With this, it polls the Primary CM every 1s to check if it is alive. In response, the Primary CM sends a Payload containing its MetaData. This way, every second the Backup CM has synced up with the Primary CM.

# Transfer of Primary title

When the Primary CM is killed, the pulseCheck will detect the death of the Primary CM. The Backup CM now takes over as Primary and sends a CHANGE\_CM message to all Clients in client.json to notify them about the change in CM. All read and write requests are now routed to the Backup CM.

# **Rebooting Primary CM**

When you reboot the Primary CM, it sends a IM\_BACK message to the Backup CM to let them know who's the real boss. It again sends a CHANGE\_CM message to all the Clients to inform them about the change in CM. Read/Write requests are back to being routed to the Primary CM.

## Rebooting Backup CM

In the event of Backup CM death, you can reboot it. This will simply restart the goroutine to send PulseChecks to the PrimaryCM and resume syncing the MetaData every 1s.

# **Experiments**

For each test, 12 terminals were used for:

- 2 CMs (1 Primary + 1 Backup)
- 10 Clients (Client 1 10)
- One of the client initially seeds CM with Pages P1-P10
- For each scenario, all 10 Clients ran randomizeRWRequests() which sends a total of 10 requests randomized by a coin flip to either read or write. You can run the simulation on each Client by typing x.
- The time taken for all 10 requests to complete in each client was recorded (using time.Now().UnixMilli()).
- For each scenario the Average Time for all 10 clients to complete 10 requests is shown.

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## Test results

#### Scenario 1: No faults

• Type x on all Client terminals. No further interference.

Average Time/ms: 10065

### Scenario 2(a): Single Fault in Primary CM

- Type x on all Client terminals.
- Kill the primary CM by ctrl-c .

Average Time/ms: 10078

### Scenario 2(b): Fail and reboot Primary CM

- Type x on all Client terminals.
- Kill the primary CM by ctrl-c.
- Reboot the primary CM with go build && ./ivy and restartCM

Average Time/ms: 10082\*\*

### Scenario 3: Fail and reboot Primary CM multiple times

- Type x on all Client terminals.
- Kill the primary CM by ctrl-c .
- Reboot the primary CM with go build && ./ivy and restartCM
- Repeat the killing and rebooting multiple times

Average Time/ms: 10080

### Scenario 4: Fail and reboot both Primary CM and Backup CM multiple times

- Type x on all Client terminals.
- Kill and Reboot Primary CM
  - Kill the primary CM by ctrl-c.
  - Reboot the primary CM with go build && ./ivy and restartCM
- Kill and Reboot Backup CM
  - Kill the Backup CM by ctrl-c.
  - Reboot the Backup CM with go build && ./ivy and restartBackupCM

Average Time/ms: 10082

### Inferences

Due to the nature of the implementation particularly with the non blocking syncing process, the time taken for each simulation is approximately constant.

# Is Fault Tolerant Version of IVY sequentially consistent?

A sequentially consistent implementation should maintain **some total ordering of read and write requests amongst all clients**. Since all requests are routed to a single Central Manager, the incoming requests to the CM are automatically ordered. In situations where proceeding to the next flow in the logic needs to be blocked while waiting for a response from all Clients, the code has checks in place to deal with it.

For example, when a CM receive a WRITE\_REQUEST. It sends an INVALIDATE\_COPY to all the Clients in the CopySet for the particular page. In my implementation, a WRITE\_FORWARD is only sent when all INVALIDATE\_CONFIRMATION are received. The code below demonstrates how the function return s when a Client does not acknowledge a INVALIDATE COPY message.

```
for _, clientPointer := range pageInfo.CopySet {
  invalidateCopy := Message{
    Type: INVALIDATE_COPY,
    Payload: Payload{
      InvalidateCopy: InvalidateCopy{
        WriteRequesterID: writeRequesterID,
        PageNumber:
                          targetPageNo,
      },
    },
  }
  reply := cm.CallRPC(invalidateCopy, CLIENT, clientPointer.ID, clientPointer.IP)
  if !reply.Ack {
    logerror.Printf("Msg [%s] from CM not acknowledged by Client %d\n", invalidateCopy.Type, clientPointer.ID)
    logerror.Println("Cannot forward Write Request")
  }
}
// All InvalidateCopy responses have been received.
// Send WriteForward to Page Owner
writeForward := Message{
  Type: WRITE_FORWARD,
  Payload: Payload{
    WriteForward: WriteForward{
      WriteRequesterID: writeRequesterID,
      WriteRequesterIP: writeRequesterIP,
      PageNumber:
                        targetPageNo,
      Content:
                         content,
    },
  },
updatedPageInfo := cm.MetaData[targetPageNo]
ownerID := updatedPageInfo.Owner.ID
ownerIP := updatedPageInfo.Owner.IP
reply := cm.CallRPC(writeForward, CLIENT, ownerID, ownerIP)
```

### Sequential consistency with Primary CM failure

Here we have 4 cases:

- 1. Primary CM fails after last request is completed + Data sync with Backup CM is complete.
- 2. Primary CM fails before last request is completed + Data sync with Backup CM is complete.
- 3. Primary CM fails after last request is completed + Data sync with Backup CM is incomplete.
- 4. Primary CM fails before last request is completed + Data sync with Backup CM is incomplete.

#### Case 1

This is dealt with seamlessly as all requests are registered in MetaData and BackupCM can take over with most updated MetaData.

#### Case 2

In this case, the last request to the Primary CM is now lost because the the Backup CM only has MetaData complete up till the request before the last request to the Primary CM. This is not dealt with in my implementation but can easily be done. If a client does not receive a PAGE\_SEND after sending a request, it can resend the request after a timeout.

### Case 3

This is not dealt with at all. The client receives PAGE\_SEND and assumes the request is complete. But the MetaData of the Backup CM does not reflect the latest request, hence the request is lost forever.

### Case 4

This is similar to Case 2, where an incomplete request can trigger a re-request by the Client after a timeout.

### Conclusion

Overall, sequential consistency is still maintained as there is no case where a request which was ordered by the CM is re-ordered in the Backup CM. That means that *some* total order of the request is always maintained. Furthermore, the likelihood of lost requests can be significantly reduced with more frequent syncing of the data with the Backup CM.