CS3243 : Introduction to Artificial Intelligence

Tutorial 6

NUS School of Computing

July 21, 2022

Admin

► Finals next week

Review

- ► Knowledge Base
- ► Propositional Logic
- ► Inferences using what we know
 - Entailment
 - Truth Table Enumeration
 - ► Inference Algorithm

- Verify the logical equivalences
- ▶ Idea : To use the rules of logic to proceed

Tutorial Question 1(a)

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- Let's solve the LHS
- $(\neg p \land q) \lor (\neg p \land \neg q)$
- $ightharpoonup \neg p \land (q \lor \neg q)$
- $ightharpoonup \neg p \wedge \mathbf{true}$
- **▶** ¬p
- ▶ Which is logically equivalent to RHS

Tutorial Question 1(b)

$$\blacktriangleright \ (p \land \neg (\neg p \lor q)) \lor (p \land q) \equiv p$$

Tutorial Question 1(b)

- $(p \land \neg (\neg p \lor q)) \lor (p \land q) \equiv p$
- Let's solve the LHS
- $\blacktriangleright (p \land (p \land \neg q)) \lor (p \land q)$
- $((p \land p) \land \neg q) \lor (p \land q)$
- $\blacktriangleright (p \land \neg q) \lor (p \land q)$
- $ightharpoonup p \wedge (\neg q \vee q)$
- $ightharpoonup p \wedge \mathbf{true}$
- ▶ Which is logically equivalent to RHS

- Three friends : Alice, Ben, and Cindy
- Constraints given
- Cindy comes to the party only if Alice does not come
- Alice comes to the party if either Ben or Cindy (or both) comes
- Cindy comes to the party if Ben does not come
- Using propositional logic, determine who attends the party

- ► Translate the constraints to propositional logic statements (like statements in the Knowledge Base)
- Let's say we define these binary variables
- ► C: Cindy comes to the party, A: Alice comes to the party, B: Ben comes to the party, and use them to formulate the constraints

- ► Translate the constraints to propositional logic statements (like statements in the Knowledge Base)
- Let's say we define these binary variables
- ► C: Cindy comes to the party, A: Alice comes to the party, B: Ben comes to the party, and use them to formulate the constraints
- ightharpoonup C o
 eg A
- $ightharpoonup (C \lor B) \to A$ translates to $C \to A$, $B \to A$
- ightharpoonup $\neg B \rightarrow C$

- $ightharpoonup (C o \neg A) \wedge (C o A)$
- $(\neg C \lor \neg A) \land (\neg C \lor A)$
- $ightharpoonup \neg C \lor (\neg A \land A)$
- ightharpoonup $\neg C \lor \mathbf{false}$
- ightharpoonup
- Cindy ain't coming to the party! Now onto the rest...

Thank you!

If you have any questions, please don't hesitate. Feel free to ask! We are here to learn together!