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Seminar Approximation Algorithms

Approximating Nash Social Welfare under Submodular Valuations through (Un)Matchings

Based on a paper of the same title by J. Garg, P. Kulkarni, and R. Kulkarni

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Supervised by Dr Giovanna Varricchio

28th July 2023 · Algorithms and Complexity (Prof. Dr Martin Hoefer)

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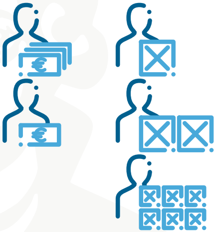

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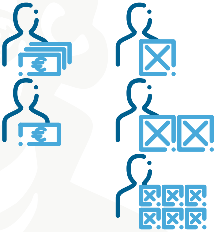

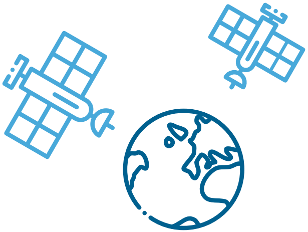

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- industrial procurement
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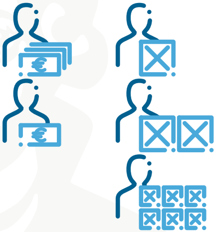
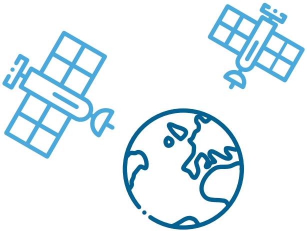
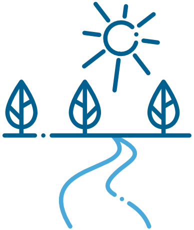
- water withdrawal







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But how to measure its efficiency and fairness?

Valuation Functions

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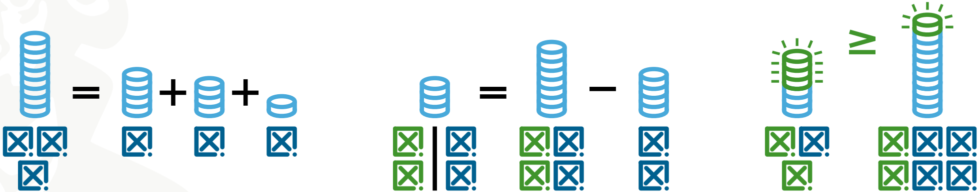
  - diminishing returns



Asymmetric Maximum Nash Social Welfare Problem

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Problem

$$\boldsymbol{x}^* \overset{!}{=} \underset{\boldsymbol{x} \in X_{\mathscr{A}}(\mathscr{G})}{\arg\max} \{\mathrm{NSW}(\boldsymbol{x})\}$$

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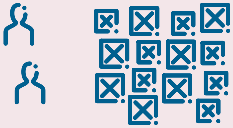

$n \lll m$

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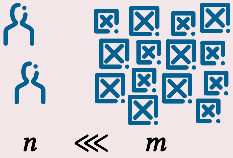

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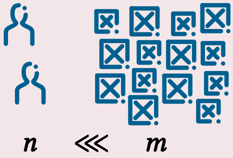

RepReMatch

2

Naïve Approach

Greedy algorithm:

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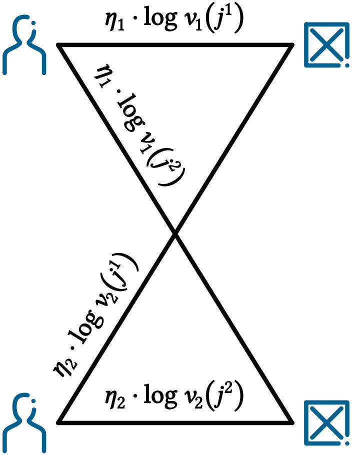

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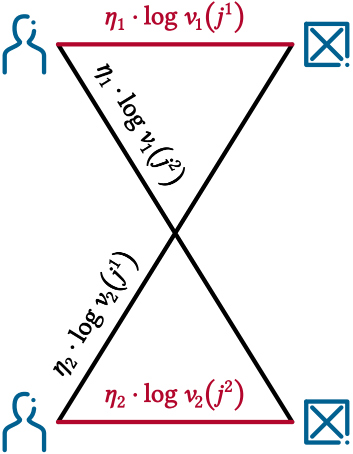

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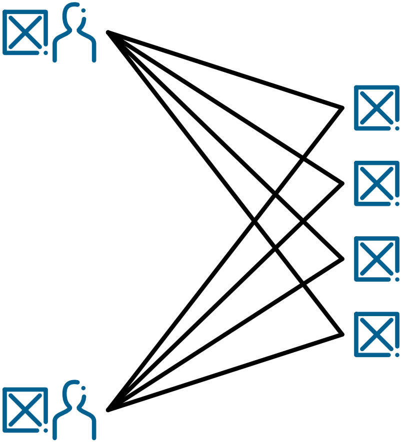
Greedy algorithm:

- repeatedly use maximum weight matchings



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Greedy algorithm:

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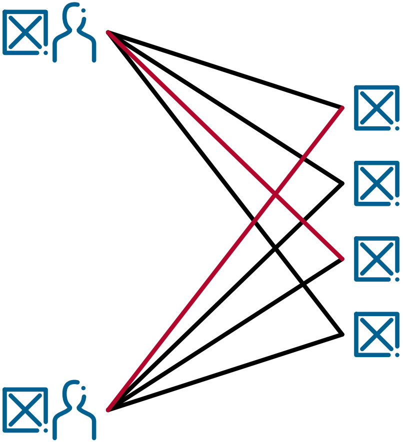

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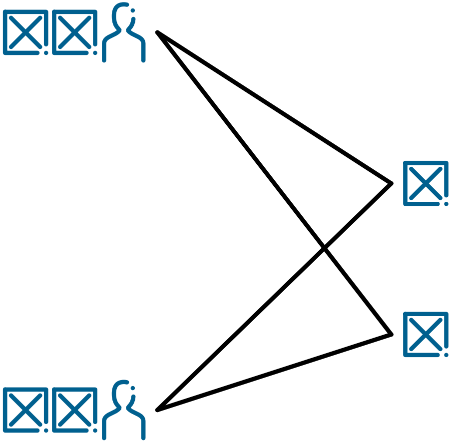

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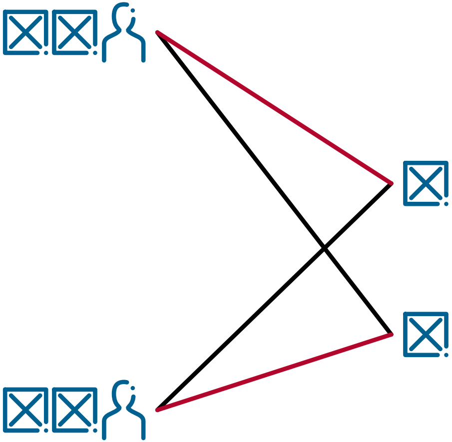

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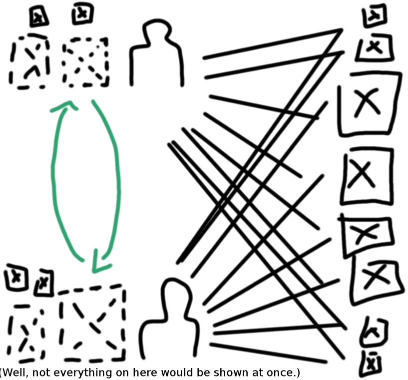
Key Ideas of the Algorithm

We need change the past in three phases:

Phase I Assign enough high-value items temporarily.

Phase II Assign the remaining items definitely.

Phase III Re-assign the items of phase I definitely.



(Well, not everything on here would be shown at once.)

Theorem

RepReMatch guarantees a $2n(\log_2 n + 3)$-approximation under submodular valuations.

The Algorithm

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1. repeat $\lceil \log_2 n \rceil + 1$ times

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5. create bundles $\boldsymbol{x}_i^{\mathrm{III}}$

Analysing Phases I & III (1/2)

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$\Rightarrow$ Are enough outstanding items reserved?

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Each agent can be matched with an outstanding item in phase III.

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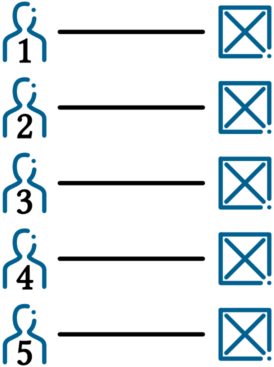

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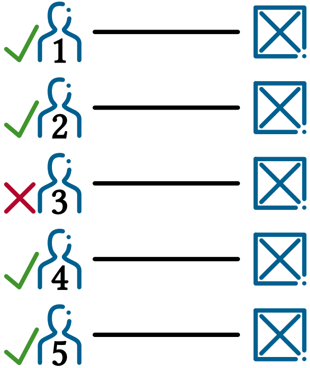

Analysing Phases I & III (2/2)

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Each agent can be matched with an outstanding item in phase III.

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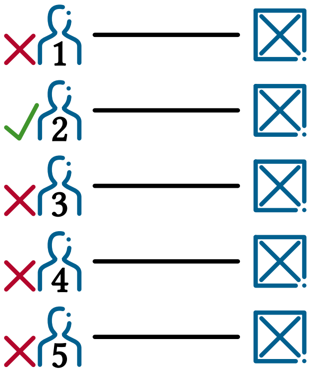

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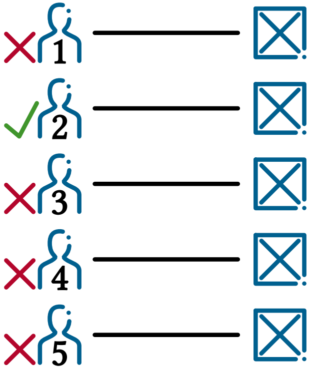

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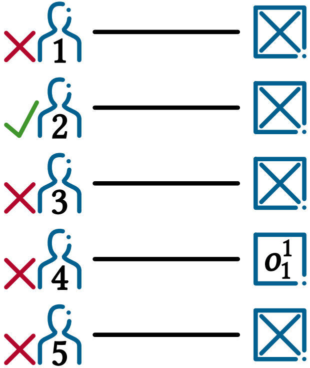

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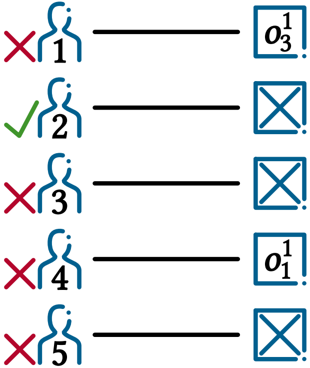

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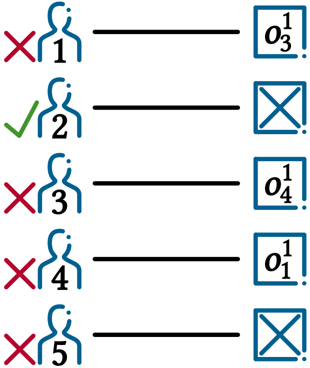

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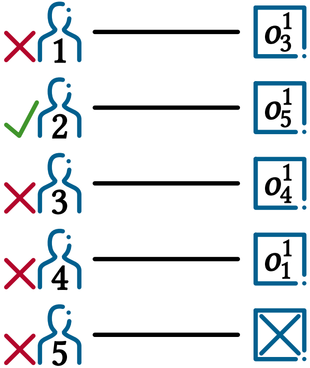

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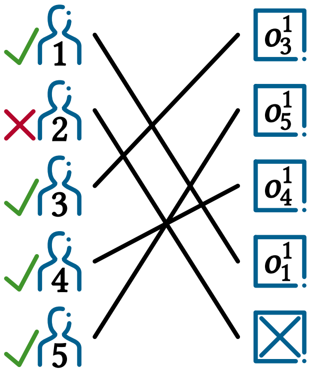

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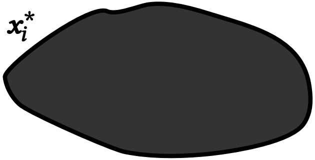

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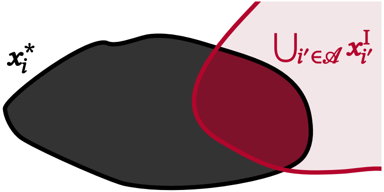

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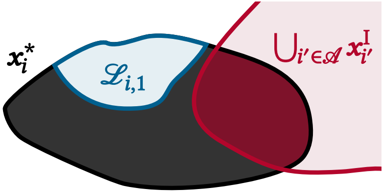

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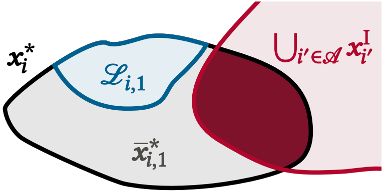

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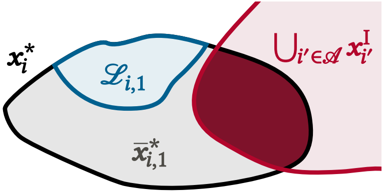

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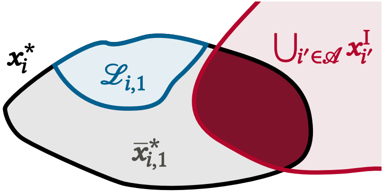

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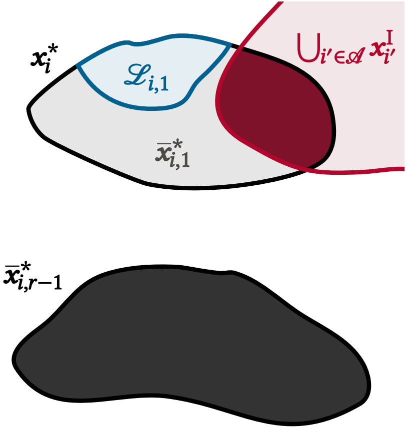

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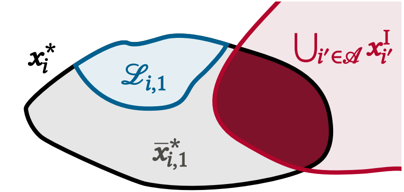

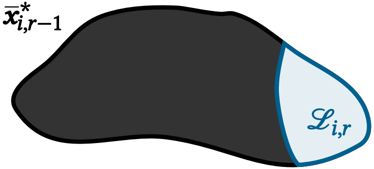

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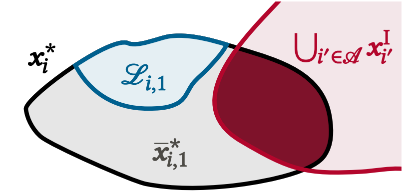

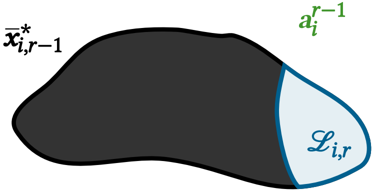

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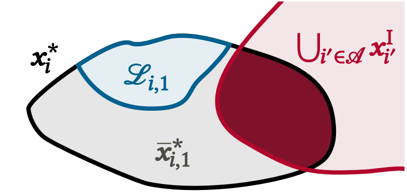

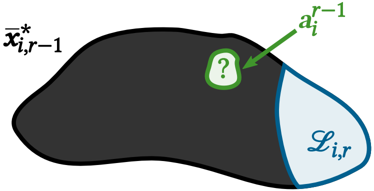

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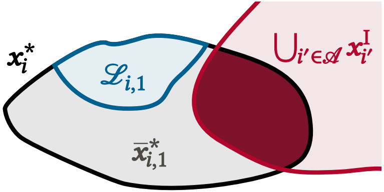

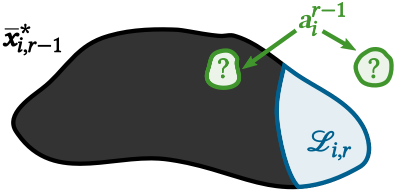

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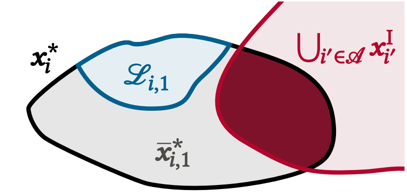

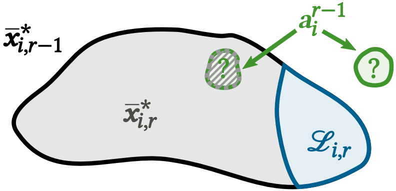

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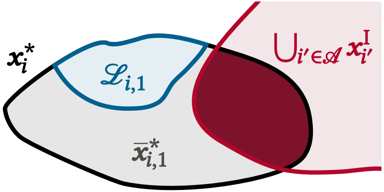

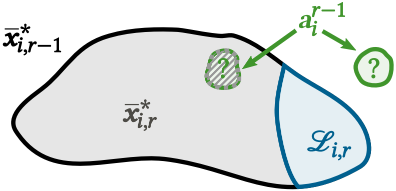

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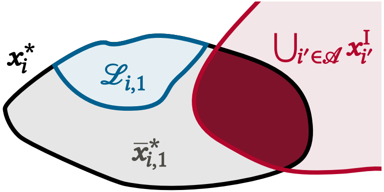

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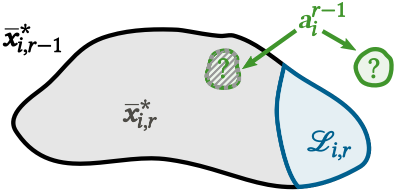

$\Rightarrow$ What is the valuation of the remaining items?

Analysing Phase II (2/2)

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$v_i\left(\overline{x}^*_{i,r} \mid a_i^1, \ldots, a_i^{r-1}\right)$

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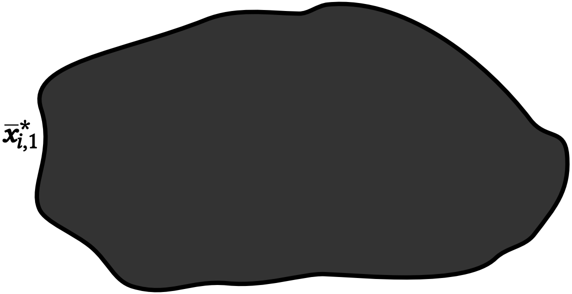

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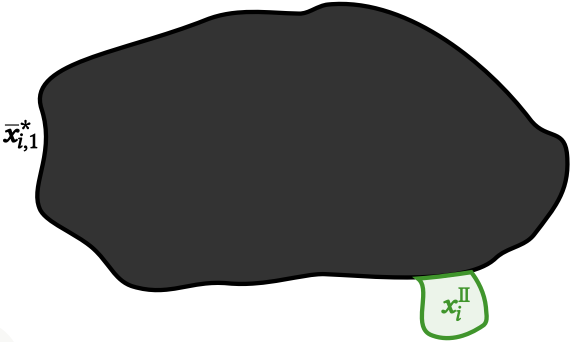

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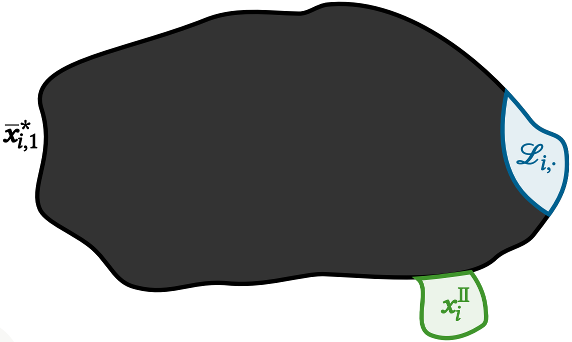

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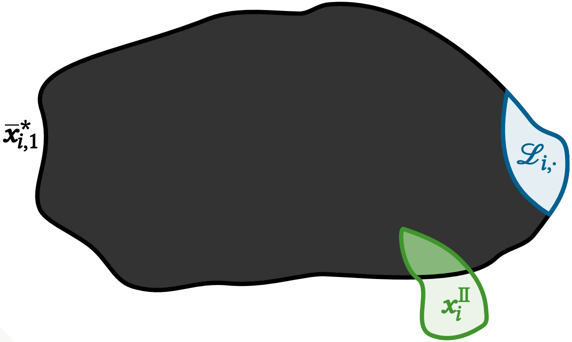

Analysing Phase II (2/2)

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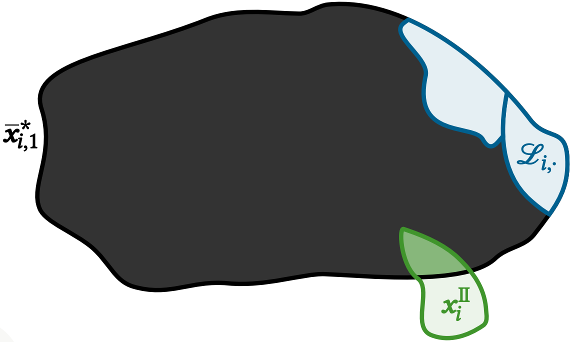

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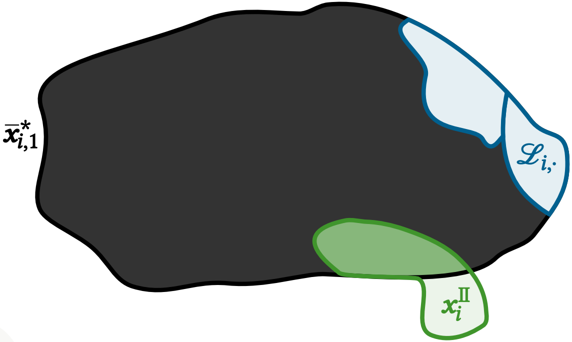

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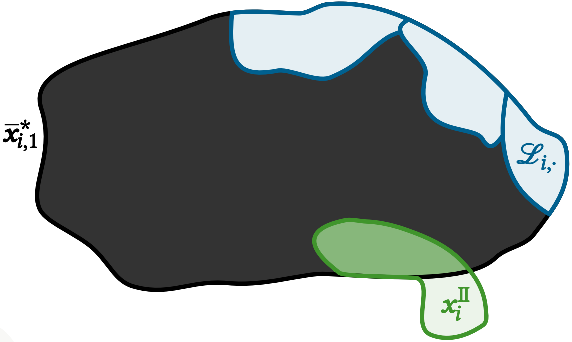

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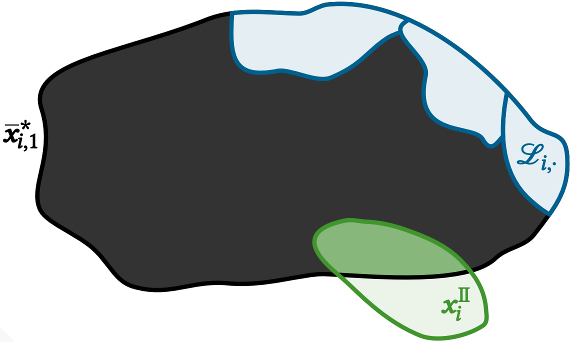

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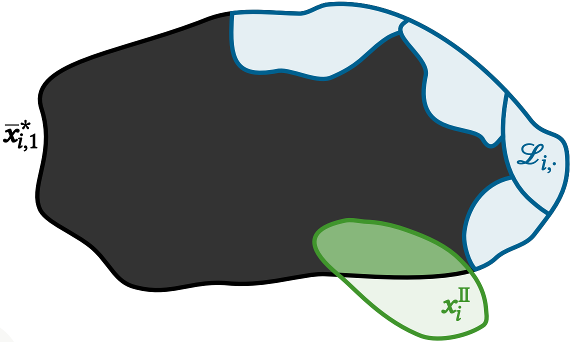

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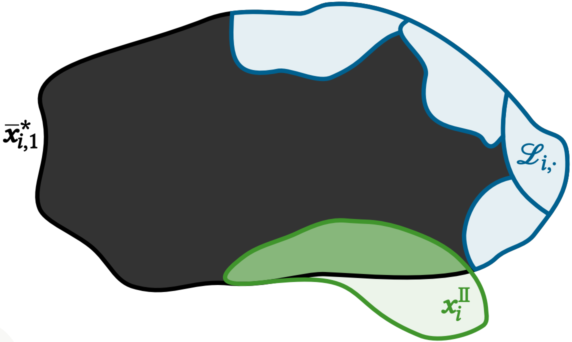

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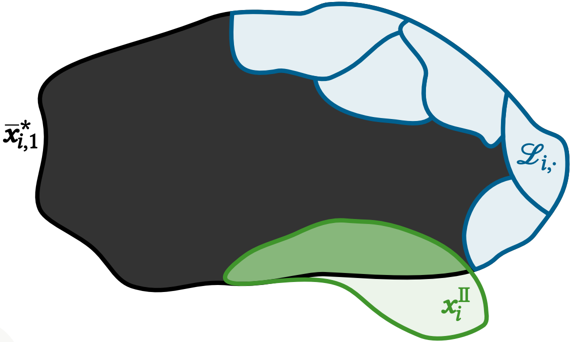

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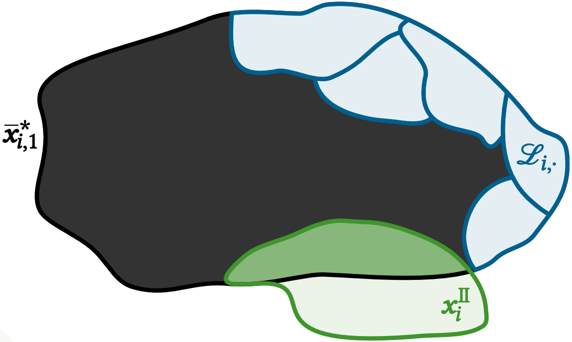

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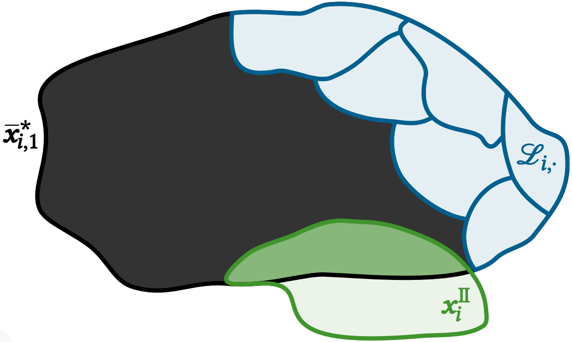

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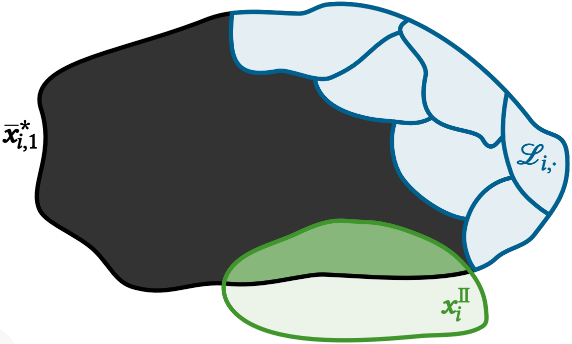

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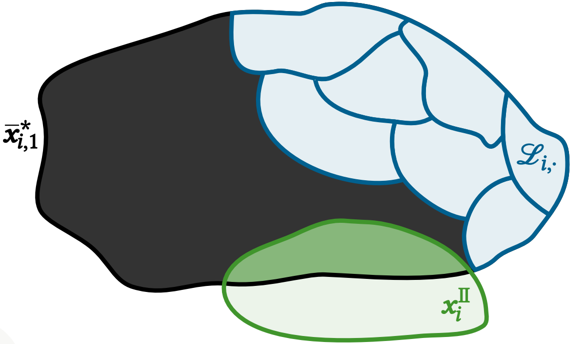

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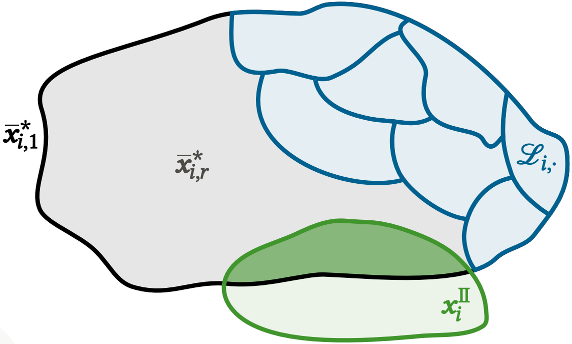

Analysing Phase II (2/2)

Lemma

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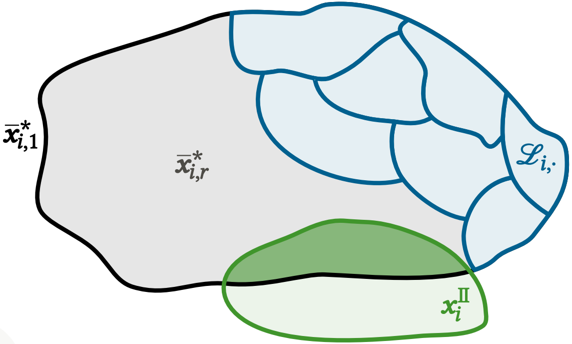

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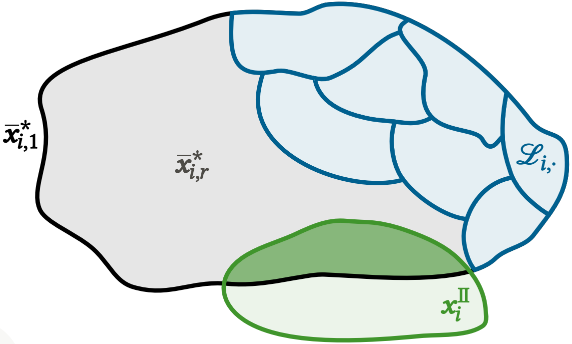

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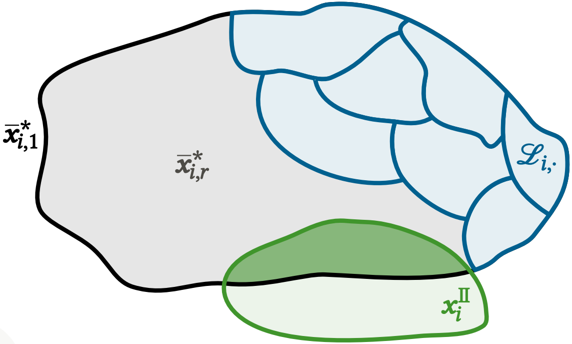

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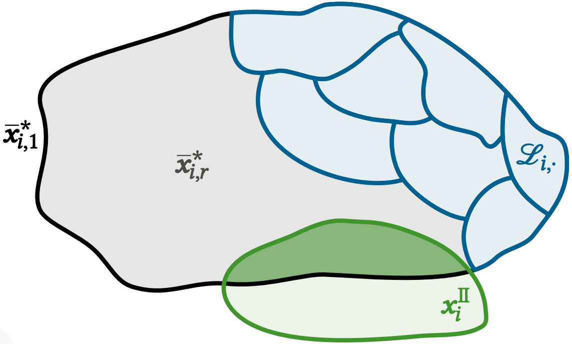

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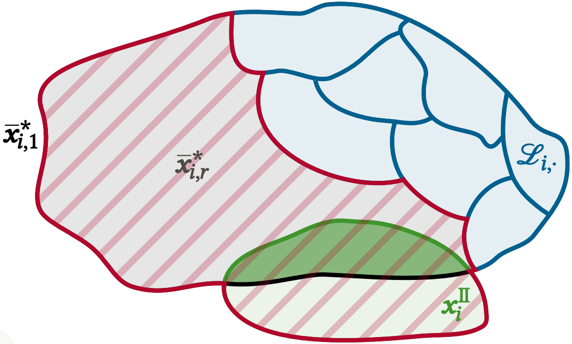

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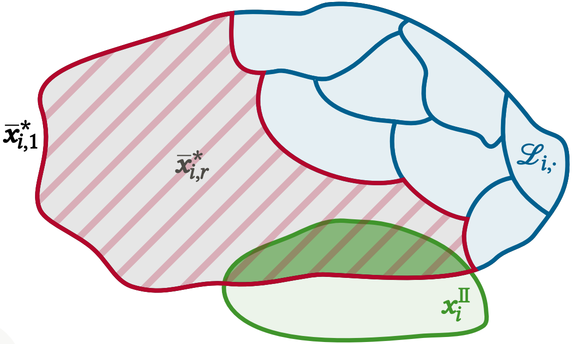

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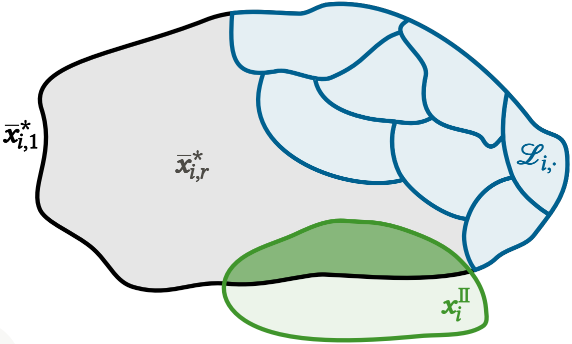

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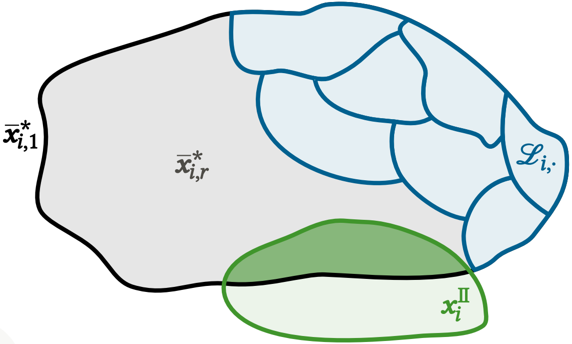

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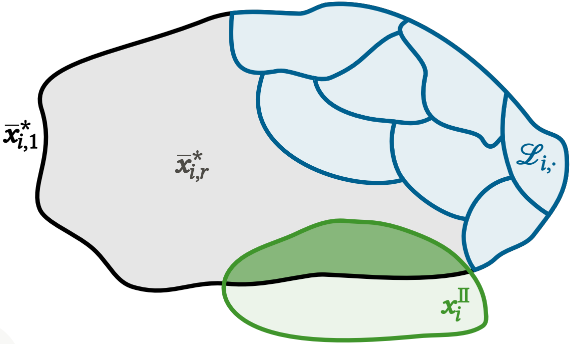

Analysing Phase II (2/2)

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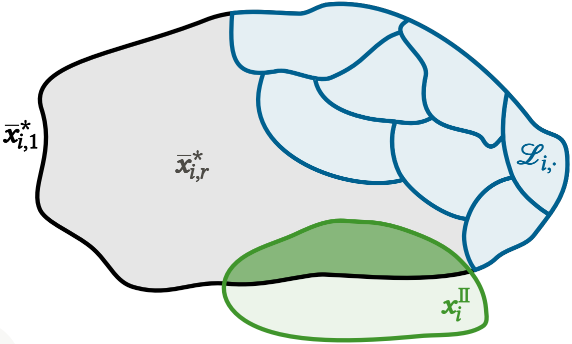

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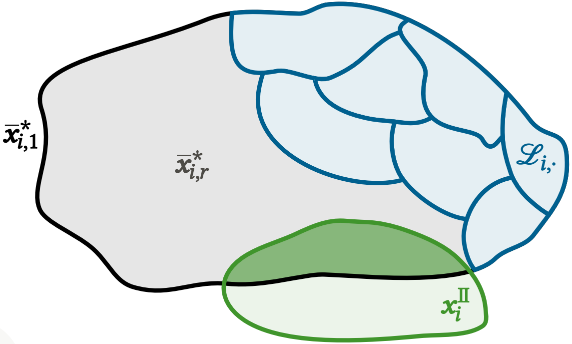

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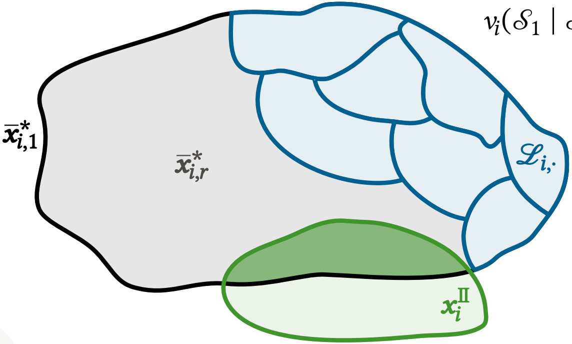

$v_i(\mathcal{S}_1 \mid \mathcal{S}_2) =$

Analysing Phase II (2/2)

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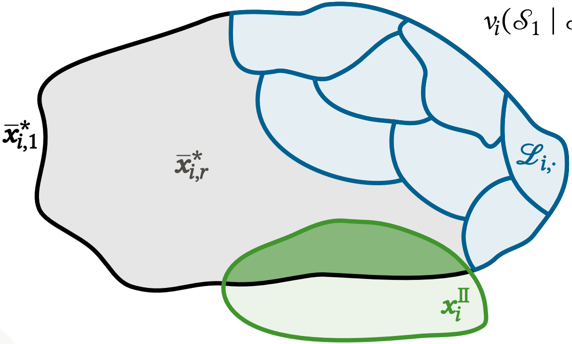
$$v_i(\mathcal{S}_1 \mid \mathcal{S}_2) = v_i\left(j^1 \mid \mathcal{S}_2\right)$$



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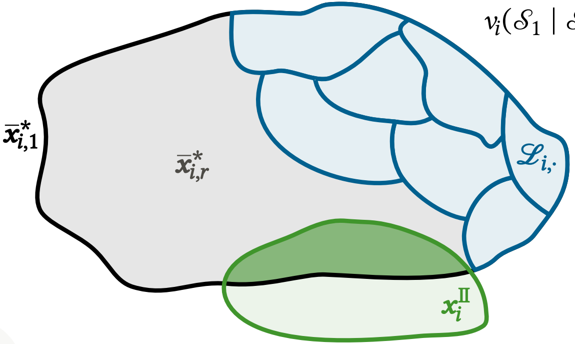

$$v_i(\mathcal{S}_1 \mid \mathcal{S}_2) = v_i\left(j^1 \mid \mathcal{S}_2\right) + v_i\left(j^2 \mid \mathcal{S}_2 \cup \left\{j^1\right\}\right)$$

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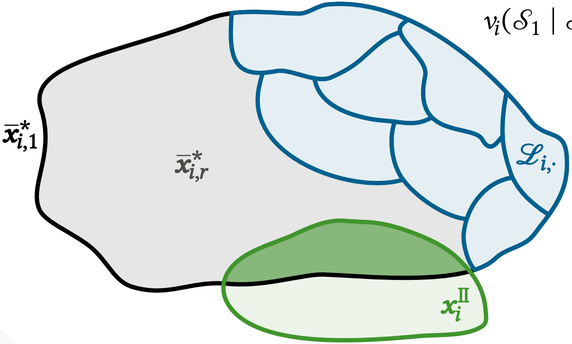

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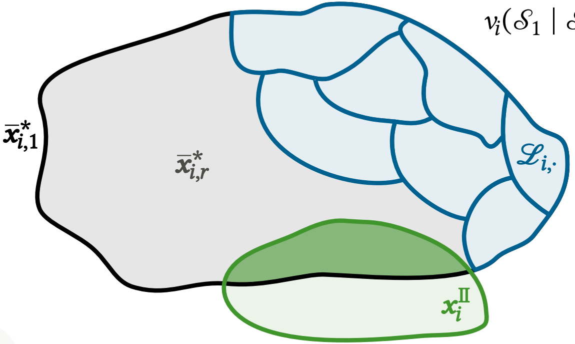

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Analysing Phase II (2/2)

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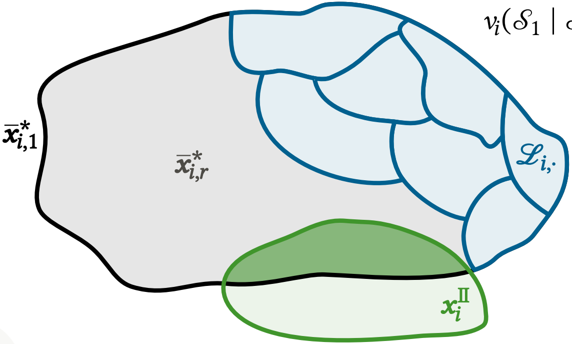

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Analysing Phase II (2/2)

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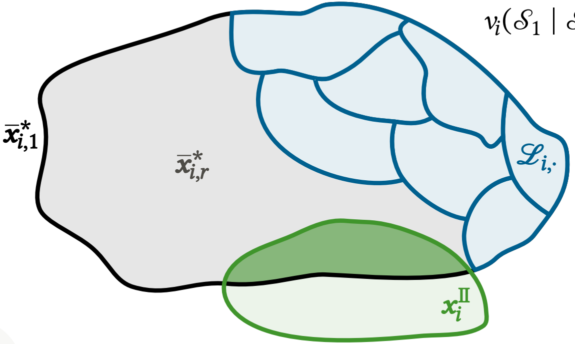
Lemma

$$v_i\left(\overline{\boldsymbol{x}}^*_{i,r} \mid a_i^1, \dots, a_i^{r-1}\right) \geq -v_i\left(a_i^1, \dots, a_i^{r-1}\right) + v_i\left(\overline{\boldsymbol{x}}^*_{i,1}\right) - \sum_{l=2}^{r} v_i\left(\mathscr{L}_{i,l} \mid a_i^1, \dots, a_i^{l-1}\right)$$



$$\begin{aligned} v_i(\mathcal{S}_1 \mid \mathcal{S}_2) \leq v_i\left(j^1 \mid \mathcal{S}_2\right) &+ \\ v_i\left(j^2 \mid \mathcal{S}_2\right) &+ \\ v_i\left(j^3 \mid \mathcal{S}_2\right) &+ \\ \vdots \end{aligned}$$

Analysing Phase II (2/2)

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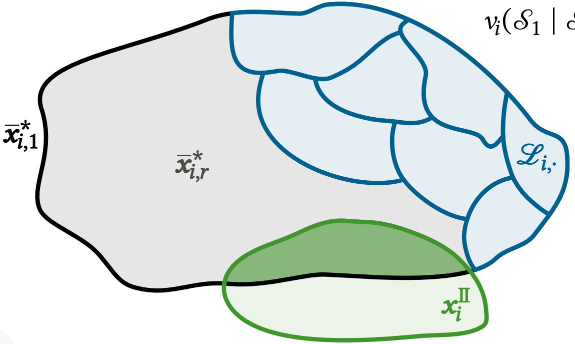
$$v_i(\mathcal{S}_1 \mid \mathcal{S}_2) \leq \sum_{j \in \mathcal{S}_1} v_i(j \mid \mathcal{S}_2)$$



Analysing Phase II (2/2)

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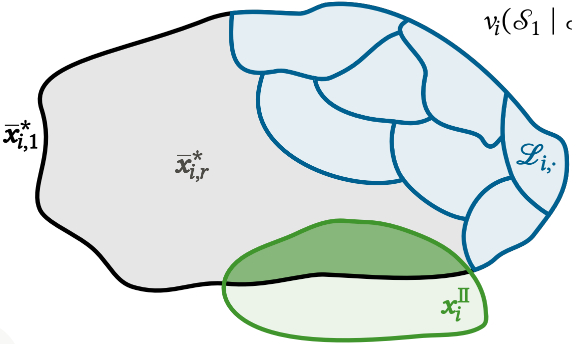

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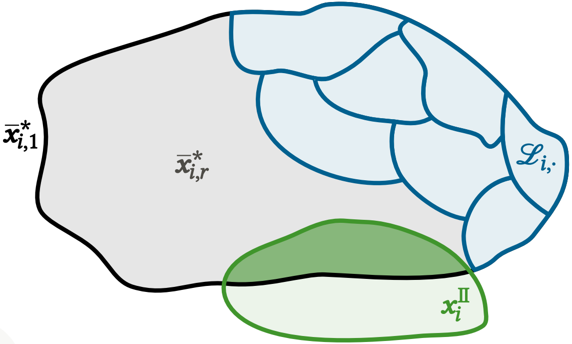

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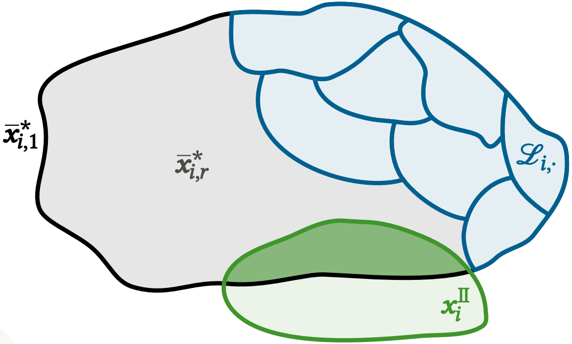

Analysing Phase II (2/2)

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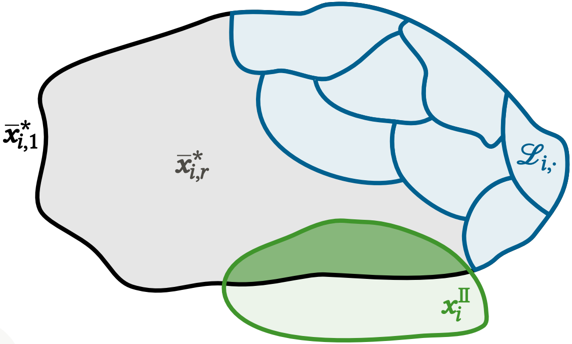

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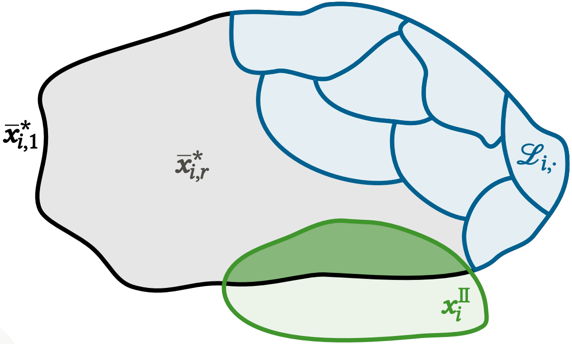

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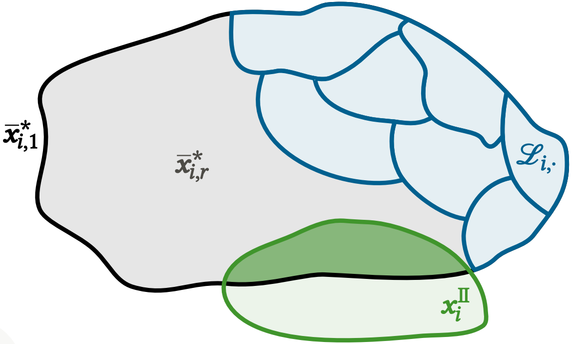

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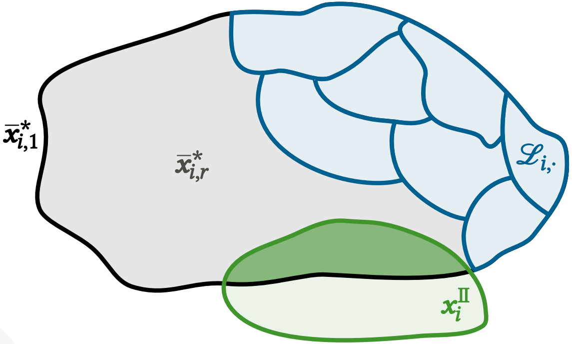

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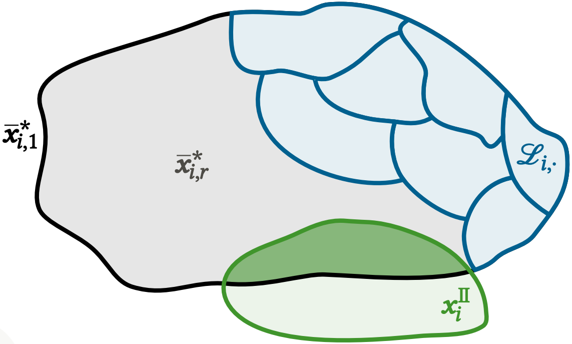

3

Conclusion

Summary & Outlook

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- allocation: partition of items amongst agents

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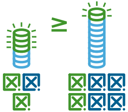

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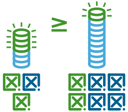

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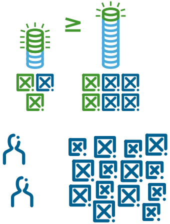

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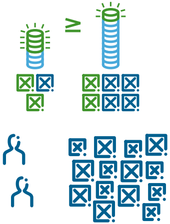

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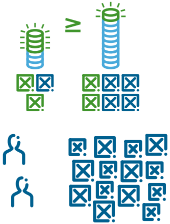

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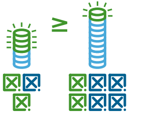

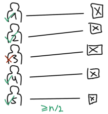

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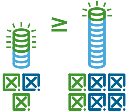

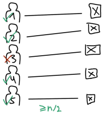

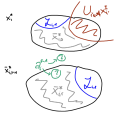

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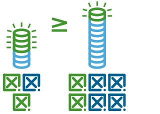

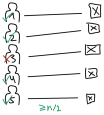

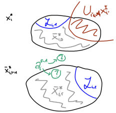

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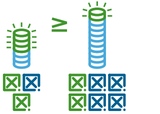

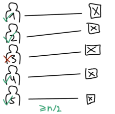

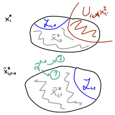

Any Room for Improvement?

Possibly! Lower bound of $\frac{e}{e-1} \approx 1.58$

End of Talk