

# Communication Concepts, Information Theory, Math and Patterns

Unit 03 - Hands-On Networking - 2018

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# Recap

- Historical and modern **telecommunication systems**.
- Transmissions use certain **symbols** to transfer meaning.
- **Encoding** used to transmit written language.
- In **networking**, there are **challenges** galore.

# At the Train Station



Through computer scientist's glasses...

- *Party A* communicates
- with *Party B*
- with *Purpose P*
- in *Environment E*
- and under *Requirements R*

**Example:** Alice (*A*) asks Bob (*B*) for the time (*P*) at the train station (*E*).

## ❓ What are the Requirements (*R*)?

- *B*'s answer should be valid (R1).
- *E* should not be too noisy (R2).
- *A* and *B* should talk same language (R3).
- *A* has to say "Good day, may I ask you a question" first (R4).

Because otherwise *B* will think he is impolite and refrain from answering.

# Communicate

**Meaning (information)** is transmitted through **something else (symbols)**, which is built using e.g. **words (alphabet)**.

## Examples

- Information.  
Time of day. Sensor readings. Thought: "This is interesting."
- Symbols.  
11 o'clock. Tomorrow. In general: Sentences from a language.
- Alphabet.  
All words/numbers used to describe time information. In general: Language XYZ

# Time flies by...

## Processing Delay ( $D_s$ )

- Enquirer: Phrase the question.
- Responder: Understand question. Look at the watch. Phrase answer.

## Transmission Delay ( $D_t = \frac{\text{Data Length}}{\text{Data Rate}} = \frac{L}{R}$ )

- Enquirer: What time is it?
- Responder: Eleven o'clock.

## Propagation Delay ( $D_p = \frac{\text{Distance}}{\text{Speed of Propagation}} = \frac{d}{c}$ )

- Natural bound: Speed of sound / light etc.
- Distance between mouths and ears.

## Round-Trip Time ( $RTT$ )

- Time between asking the question and receiving the answer.
- $RTT = 2 \cdot D_p + 2 \cdot D_t + D_s$

# Loss

## ❓ What can cause misunderstandings?

- A cannot understand B because a train is approaching.
- B has a horrible accent.

## ❓ How can this be solved?

- B repeats what he said before.  
Later: *Automated Repeat reQuest (ARQ)*.
- B says something like "it's eleven o'clock on the minute".  
Later: *Forwarding Error Coding (FEC)*.

## 📦 Formal Sizes:

- Packet Loss Rate
- Bit Error Rate

# Data Rate

❓ What are the possible answers( $A$ )?

- e.g. precision using hours per day, and quarter-hours of day.

❓ How many bits would I need to encode a certain answer ( $s$ )?

- $s = \text{ld}(|A|)$

❓ How long does it take to transmit one symbol ( $t$ )?

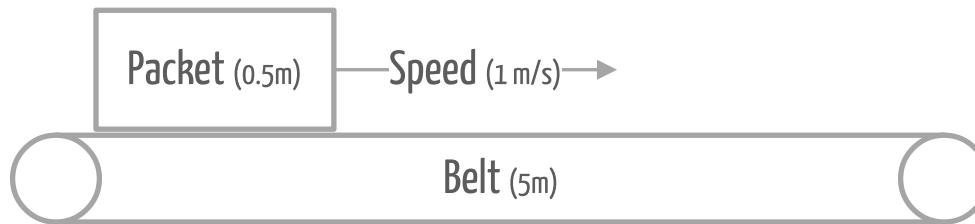
- See last slide.

📦 Formal Sizes:

- Data Rate:  $r = \frac{s}{t}$ ,  $[s] = \text{bit}$ ,  $[t] = \text{sec}$
- Throughput
- Goodput

# Links

💡 **Analogy:** Packet Conveyor Belt.



## ❓ Questions:

- How long does it take for a packet to traverse the belt?
- How large is the throughput (packets per second) if we send a single packet?
- How many packets can fit the belt / can be sent back to back?
- What happens if we send more?



# Data Network Links

## Base Sizes:

- Length  $\equiv$  **delay**  $D, [D] = sec.$
- Speed  $\equiv$  **throughput**  $B, [B] = bps.$   
Also called bandwidth or data rate.  
**Careful:** Bandwidth is most of the times given in Hertz and a frequency size.
- Packet size  $\equiv$  **packet length**  $L, [L] = byte.$   
Largest possible packet size: **Maximum Segment Size** ( $MSS, [MSS] = byte$ ).

## Composed Sizes:

- Number of *bytes* a link can contain at any time:  
*Bandwidth-Delay Product* ( $BDP = B \cdot D, [BDP] = byte$ ).
- Number of *packets* a link can contain at any time:  
*Window* ( $N = \frac{BDP}{MSS}, [N] = 1$ )

# Data Network Link - By Example

❶ Given:  $MSS = 1280\text{byte}$ ,  $D = 200\text{ms}$ ,  $B = 50\text{Mbps}$ .

❷ How big is the BDP? How many packets can be sent?

$$BDP = B \cdot D = 50\text{Mbps} \cdot 200\text{ms} = 10\text{Mbit} = 1280\text{kByte}$$

$$N = \frac{BDP}{MSS} = 1024$$

# Making Connections

Communication can be **connectionless** or **connection-based**.

- Consider our train station example again!
- Bob wanted Alice to establish a connection first.
  - Alice had to say "Hello" first, so that Bob isn't offended (and paying attention).
  - Usually people say "Bye" to each other when stopping conversation.  
Closing communication. Free resources.
  - Speaking with people involves polite start and end sentences.  
→ **Connection-Based** transmission
- There are regular announcements (*"Please mind the gap"*):
  - Everyone who is interested can tune in and listen.
  - There is neither an establishment nor a teardown routine.  
→ **Connectionless** transmission

# Symbol Rate

⚡ **Potential Problem:** Consider Alice speaking incredibly fast (think of chipmunks). Bob cannot understand her, because his hearing system is overwhelmed.

- For many physical transmission systems, the rate at which symbols are generated is important and has to be agreed on:
- Rate automatically selected / conveyed by the transmission systems or standardized among interoperating systems.
- Has to be specified when tuning in to a communication.  
e.g. baud rate for a serial connection between PC and embedded board
- Mathematically speaking:
  - Symbol Duration:  $T_s$
  - Symbol Rate:  $f_s = \frac{1}{T_s}$

# In-Order

- Assume Alice asks Bob for three different things using three separate letters (answers in brackets):
  - When did you receive this letter? (*02nd March*)
  - When did the HON course start? (*19th February*)
  - When is the HON exam? (*27th March*)
- Bob answers, but does not incorporate the original question asked.

## ⚡ Problems

- Post office might reorder answers.
- Alice might associate wrong answer with her questions.

## ❓ How to solve this?

💡 **Solution:** Include request. Identify messages (e.g. sequence numbers).

# Communication Building Blocks

# Technologies

## Media

 Air

 Water

 Copper

 Glass

 Plastic

## Approaches

 Optical

 Acoustic

 Electrical (wired and wireless)

## Communication Systems

- Combine *Media* and *Approach* to a general process able to transmit a symbol.
- Examples: Telegraph, Telephone, Copper Cable, Fibre Optics, Fire Signals, ...

More on this in detail in later lectures.

# Networks

- Communication Systems ...
  - are itself useful for transmitting.
  - require an established network for access and connectivity.
- Examples:
  - Fire Signal Network
  - PSTN
  - Internet

# Services

- Users want something to be done (e.g. telling someone something).
- Doing something is implemented by services.
- Examples:
  - Tell Troy has fallen.
  - Configure my router in Troy.
  - Book a hotel in Troy.
  - Check the weather in Troy.



# Protocols

**Definition (Wikipedia):** *In telecommunications, a communication protocol is a **system of rules** that allow **two or more entities** of a communications system to transmit information via any kind of variation of a physical quantity. These are the rules or standard that defines the **syntax, semantics and synchronization** of communication and possible **error recovery methods**. Protocols may be implemented by hardware, software, or a combination of both.*

## Protocol Specification:

- *Service* to be provided to a user of the protocol.
- *Assumptions* to be made about environment and participants.
- *Vocabulary* to be used by participants.
- *Encoding* to map elements of the vocabulary to transmittable symbols (flags etc.).
- *Behavioural Rules* to be followed by participants.

# Train Station Time Query Protocol

- **Service:** Get to know the current time by asking someone.
- **Assumptions:**
  - Participants can hear each other and talk the same language.
- **Vocabulary:**
  - Question words.
  - Time information bits (numbers (0-24), modifiers (half past), well-known times (noon, midnight), imperial specialities (a.m., p.m.)).
- **Encoding:**
  - Sounds (*W-H-A-T T-I-M-E I-S I-T, N-O-O-M*).
  - Gestures (*point at wrist*).
- **Behavioural Rules:** Ask politely!

 **Important:** Define an acronym, e.g. Train Station Time Query Protocol (TSTQP).

# Textual vs. Binary Protocols

## Textual


`GET /resource HTTP/1.1`

- Humans can read and understand the transmitted data **easily**.
- Usually done using 7-bit ASCII characters ([RFC20](#)).  
Published 1969, but just recently became a standard in 2015.
- Inefficient (information-theory wise).
- Straightforward to extend.  
Depending on protocol grammar.

## Binary


`0x450005dc` (Version: 4, Length: 20 bytes)

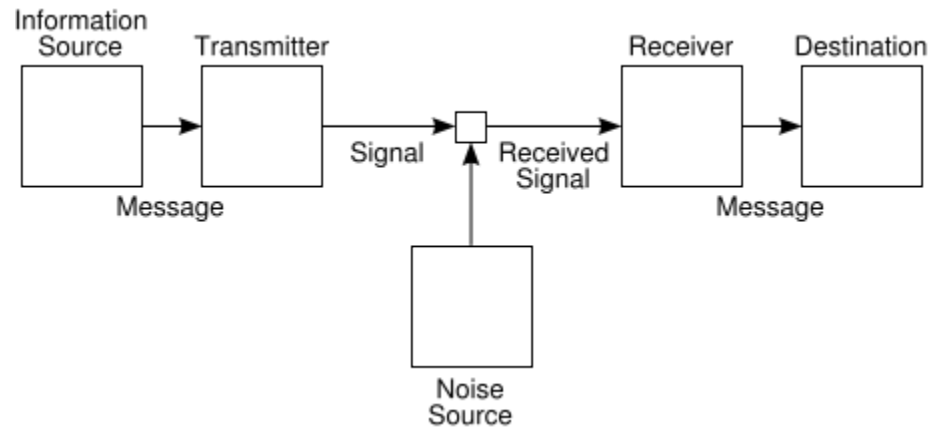
- Most humans **cannot** understand the transmitted data from the bits.
- Fields have differing length (number of bits) and type.
- Efficient.
- Potentially hard to extend.  
One way will be covered later.

 Both are needed! None is generally **better** than the other.

# A Quantum of Information Theory

# Communication Systems

 General description by Claude Elwood Shannon.



Source: [Wikipedia](#)

# Just a Little Bit

**Bit:** Basic unit of information in *computing* and *digital communications*. Portmanteu of *binary digit*.

- Only two values: 0, 1.
- Compare Light Bulb: Either on or off.
- In Information Theory:  
1 *Sh* (Shannon)  
Bit: Data ↔ Shannon: Information



[Source](#)

# Entropy

**Origin:** Thermodynamics. How irreversible is a system? *Ultimate chaos is inevitable!*

## Shannon (Information Theory)

Receiver attempts to infer which message has been sent. Receiver is uncertain about what is sent, but anticipates certain things.

### Intuition

- Likely Option  $\rightarrow$  Low Entropy.  
e.g. white pixel of a document scan
- Unlikely Option  $\rightarrow$  High Entropy.  
e.g. traffic light's yellow light
- More Options  $\rightarrow$  Higher Entropy.  
e.g. traffic light vs. 7-segm. display

### Quantification (in bits)

- Consider all options and call this  $\Omega$ .
- Count them as  $N = |\Omega|$ .
- Calculate  $\log_2(N)$ .  
The number of bits required to identify these options by a unique binary number.
- Entropy of Option  $x$ :  $-p_x \cdot \log_2(p_x)$

# Inefficiency of Textual Protocols

Consider protocol involving a *command* field, which can be one of the following:

Retrieve (GET), Create (ADD), Modify (MOD), Delete (DEL)

❓ How many bits are used for the textual and binary solution?

## Textual

- 3 characters for each command.
- One ASCII char needs 7 bits.  
Often even 1 byte, but let's be fair.
- **Result:**  $3 \cdot 7bit = 21bit$

## Binary

- 4 different commands (0,1,2,3).
- 4 values require two bits.  
(00, 01, 10, 11)
- **Result:**  $2bit$

❓ Why are textual protocols used at all?

Efficiency is not your **only** parameter! **Compression** (see later) can bring efficiency.



# Networking Math

# Number Systems

## Binary

- Symbols [0,1] (for EEs: High and Low)
- Base: 2
- Hardware uses this.
- Tedious to write.
- Protocol headers might have single bits of information.

## Hexadecimal

- Symbols [0,1,2,3,4,5,6,7,8,9,A,B,C,D,E,F]
- Base: 16
- Lazy humans that have to write binary numbers use this.
- Condensed presentation.

## Decimal

- Symbols [0,1,2,3,4,5,6,7,8,9]
- Humans use this.

## Example

- Binary: 101010
- Hexadecimal: 2A
- Decimal: 42

## Python

- Binary 0b???
- Hexadecimal 0x???
- Decimal ???

# Conversions

## Hex <-> Bin

- One hex value <-> 4-bit sequence.
- Convert blockwise.
- Examples:
  - `0xF8 = 0b11111000`  
(F = 1111, 8 = 1000)
  - `0b01101010 = 0x6A`  
(0110 = 6, 1010 = A)

## Hex / Bin -> Dec

- Multiply the symbols by the respective decimal basis and sum up.
- Example:
  - `0x1A` =  $16 + 10 = 26$
  - `0b1010` =  $8 + 2 = 10$

## Dec -> Hex (n=16) / Bin (n=2)

- Process:
  1. Divide the decimal number by  $n$   
Integer division.
  2. Write down the remainder.  
In hex/bin notation.
  3. Divide the result again by  $n$   
(integer division).
  4. Repeat 2. and 3. until result is 0.
  5. The value is the digit sequence  
of the remainders from the last  
to first.
- Example:
  - $42 / 16 = 2$  (Rest 10 = A)
  - $2 / 16 = 0$  (Rest 2)
  - $42 = 0x2A$

# Two's Complement

- Humans can decide that certain binary numbers are signed (have a sign +/-) or unsigned (always non-negative).
- Numbers are interpreted as two's complement as follows:
  - First bit indicates sign value (0 = positive, 1 = negative).

## Example

`0b11010110` ->  $1 * (-128) + 1 * 64 + 0 * 32 + 1 * 16 + 0 * 8 + 1 * 4 + 1 * 2 + 0 * 1 = -42$

# Quiz

❓ What is **0xCAFE** in binary?

A: 0b1110 1010 1110 1110

B: 0b1100 1010 1111 1110

C: 0b1100 1010 1100 1111

D: 0b110 1101 0101 10111

# A Means to an End

- Endianness = byte order of multi-byte numbers
- Little-endian: LSB first, MSB last
- Big-endian: MSB first, LSB last

## Example

- 4-byte value: 01 02 03 04 (hex, MSB to LSB)
- Big-Endian: 0x01 0x02 0x03 0x04
- Little-Endian: 0x04 0x03 0x02 0x01

 Big-Endianness is also the **Network-Byte Order**

# Quiz

❓ Value of `0x00010000`? Assume little-endian.

A: 1

B: 256

C: 4096

D: 65536

# Binary Operators

Bit A	Bit B	A & B	A   B	A ^ B
0	0	0	0	0
0	1	0	1	1
1	0	0	1	1
1	1	1	1	0

## Examples

And	Or	Xor
$\begin{array}{l} 0x1010 \\ \& 0x1100 \\ = 0x1000 \end{array}$	$\begin{array}{l} 0x1010 \\   0x1100 \\ = 0x1110 \end{array}$	$\begin{array}{l} 0x1010 \\ \wedge 0x1100 \\ = 0x0110 \end{array}$



# Shifts

## Logical Shift

- Does not preserve the sign bit.
- All bits moved in direction.
- Vacant (now empty) bit positions are filled with zeros.
- Python: `<<` and `>>`

## Arithmetic Shift / Signed Shift

- For right shifts: Instead of filling with 0s (logical), fill with sign bit.
- Not available in Python directly.

## Circular Shift

- Numbers wrap instead of filling.
- Not available in Python directly.

## Examples

- `0b0001 << 1 = 0b0010` (left logical shift)
- `0b1010 << 1 = 0b0101` (left circular shift)
- `0b0001 << 4 = 0b0000` (left logical shift, if size is 4 bit)
- `0b1001 >> 1 = 0b0100` (right logical shift)

# Quiz

❓ What is the result of `0b11001 >> 3`?

A: `0b00111`

B: `0b11111`

C: `0b00011`

D: `0b10101`

⚠️ Answer:

✗ A: That's a circular shift.

✗ B: That's a right arithmetic shift.

✓ C: Correct!

✗ D: That's nonsense.

# Masking and Setting Bits

Check for last 4 bits of a bit string b

```
b & 0b00001111
```

Set Bit at position x (from right) to y in Bit-String b

- If  $y=1$ :  $b \mid= (1 \ll x)$
- If  $y=0$ :  $b \&= \sim(1 \ll x)$

## Examples

- Set bit at position 2 to 1 in 0b1001:  
 $0b1001 \mid 0b0100 = 0b1101$
- Check if bit at position 4 is set:  
 $((0b01010100 \& 0b00010000) \gg 4) == 1$

# Metric Prefixes

- Usually orders of 1000:
  - Kilo ( $10^3$ ), Mega ( $10^6$ ), Giga ( $10^9$ )
  - Milli ( $10^{-3}$ ), Micro ( $10^{-6}$ ), Nano ( $10^{-9}$ )
- In computer science (especially with bytes):
  - Kibi =  $1024$
  - Mebi =  $1024^2$

**⚠ Careful:** The industry likes to use metric prefixes but calculates with binary prefixes (factor 1024).

- For more information:
  - [https://en.wikipedia.org/wiki/Metric\\_prefix](https://en.wikipedia.org/wiki/Metric_prefix)
  - <https://en.wikipedia.org/wiki/Kibibyte>

# Units

## Bit

- Information.
- 8 bits are one byte.

## Packet

- (More) Information.
- Sequence of bytes (seldomly bits).
- Variable length, even though parts have to be fixed.

## Second

- Time.
- Used for: Round-Trip Time, Forward-Trip Time, Jitter, ...

## Hertz

- Frequency.
- Used for: Packet Rate, Baud Rate, ...

# Size(s) matter... even for computer scientists

## Time Measures

- Latency/Delay: How long something takes in seconds.
- Forward-Trip Time (FTT): How long from source to sink?
- Round-Trip Time: FTT + way back. Sometimes  $2 \times \text{FTT}$ .
- Jitter: variance of delay.

## Data Rate

- $R = \frac{\text{TransmittedData}}{\text{Time}} \rightarrow [R] = \text{bits/s}$
- Throughput: Overall data rate.
- Goodput: Useful data rate (application data per time).
- Maximum Bandwidth: Maximally achievable data rate of a channel. Careful: Not Hertz!

## Loss

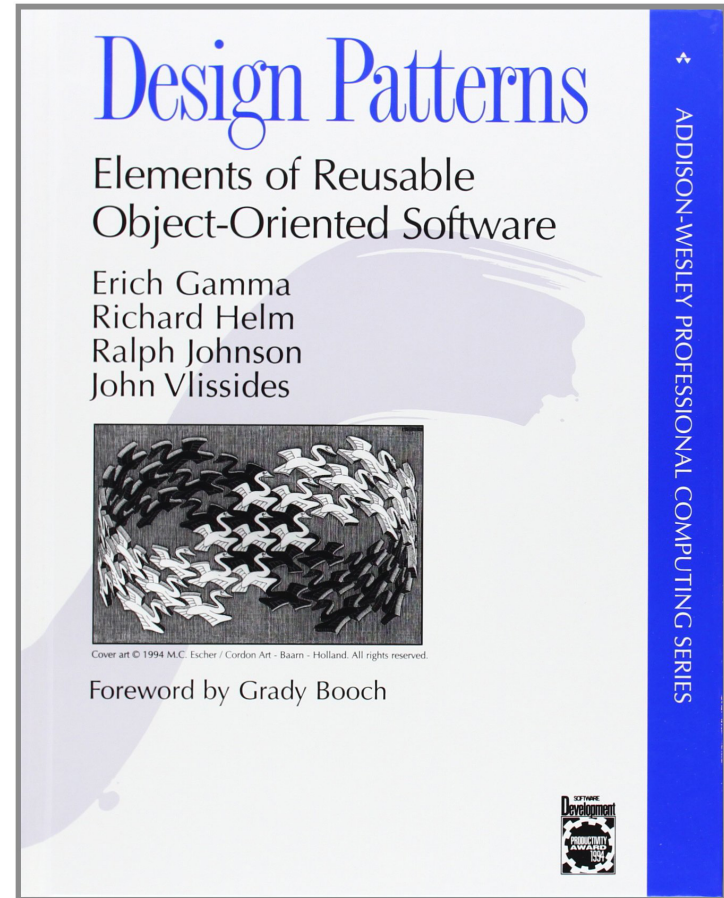
- Packet Loss Rate:  $plr = \frac{\text{LostPackets}}{\text{TransmittedPackets}} \rightarrow [plr] = 1$  (usually in %)
- Bit Error Rate:  $ber = \frac{\text{WrongBits}}{\text{TotalBits}} \rightarrow [ber] = 1$  (usually in %)



# Patterns

# Software Design Patterns

- **Motivation:** Computer Scientists tend to reinvent the wheel again and again.. and again.
- Mistakes and design weaknesses sneak in when reinventing.
- Other fields, e.g. architecture, have patterns for houses, bridges, etc.
- Proposed by the Gang-of-Four (GoF).






# Pattern

## Definition and Usage Advice

- In general: *Approach that can be used to solve common problems.*
- Should always be applied appropriately and only if required.  
Don't apply a pattern because you feel smart about knowing it.
- Nearly always has to be adapted to match the specific requirements.

## Profile


 **Name:** Suitable name to refer to it.


 **Context:** When to apply?

 **Implementation:** How to apply?

 **Benefits:** Why is it good?

 **Drawbacks:** What might be problems?

 **Variants:** How is the pattern usually modified?

 **Similar Patterns:** Which patterns can be used with this together?

# Request-Response Pattern / Poll Pattern

■ **Context:** Simple, often stateless way to ask questions and get answers.

⚙️ **Implementation:**

- Define a format for questions and answers.
- Enquirer sends requests, responder waits for it and answers accordingly.

👍 **Benefits:**

- Simple protocol.

📄 **Variants:**

- Long-Polling: Responses not sent immediately, but delayed until the responder has something to send (avoids poll-to-refresh).

🗨️ **Drawbacks:**

- If the responder wants to tell the requester something, he has to wait for the request.
- Realtime updates require constantly asking the responder:  
*"Has something changed." - "No." - "How about now?"*

📄 **Inverse Pattern:** Push Pattern

# Buffer Pattern

## ■ Context:

- Persist data.
- Enforce order.
- Smoothing data rate.

## 👍 Benefits:

- Simple.
- Universal.

## 📄 Variants: Queue, Stack, Pipe, ...

## ⚙️ Implementation:

- List, array, ordered-something implementation.
- Provide put() and get() methods.

## 🗨️ Drawbacks:

- Requires storage.
- Requires fill-level management.

# The Three-Letter Acronym Pattern

■ **Context:** Whenever you have to name a protocol, tool, project, ...

👍 **Benefits:** Nobody will know what it is when first encountering your TLA.

🗨️ **Drawbacks:** Nobody will know what it is when first encountering your TLA.

⚙️ **Implementation:** Pick 3 letters (or digits) from ASCII and invent suitable words for it.

📄 **Variants:**

- 4LAD (Four-Letter Acronym Definitions).

📄 **Similar Patterns:**

- "Let it end with a P"-pattern.
- Recursive acronyms (GNU - GNU not Unix).

📝 More information: [RFC5513](#).

# Wrap-Up

Questions?

## Take-Home Messages

- Party A wants to communicate with Party B with Purpose P in Environment E and with Requirements R.
- **Information** is transmitted through **symbols over an alphabet**.
- **Delay, loss** and **data rate** are important quantifiable parameters for communication.
- Networks, composed of multiple communication systems, are used to provide a service. Participants have to follow a **protocol**.
- No networking without **math and numbers**.
- **Patterns** help to structure and recognize approaches.

## Further Reading

- Kurose-Ross "Computer Networking" (Sec. 1.1, 1.3., 1.4)
- Gamma et al. "Design Patterns - Elements of Reusable Object-Oriented Software".