

## Socket Programming

Unit 13 - Hands-On Networking - 2018

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### Recap

- IP addresses identify **hosts**.
- Ports identify **processes** on a host.
- TCP and UDP are **transport layer** protocols.
- Services to the user are provided by the application layer.
- Client-Server and Peer-to-Peer architectures can be used.



### What is a Socket?

• Intuition: A in the one could with.

Wall: Operating System

Hole: API

 Do Stuff: Communicate Information (Send/Recv)

Socket: OS Object



Source

Sockets are **the** network programming interface.

- Originally developed for UNIX, now everywhere.
- Well-defined network API for common network operations.
- Developing application layer protocols requires socket programming.

**A** Lazy Programmers: As TCP sockets are a bit complicated, devs tend to use HTTP for writing endpoints. It is easier, provides messages and there are frameworks.



### Socket API

#### **Abstractions**

- API provides primitives (bind, send).
- Sockets abstract nasty details away:
  - Buffers incoming data.
  - Reliably transmit data.
  - Establish / teardown connections.
- Concepts shared amongst implementations.

#### **Language Support**

- Socket syntax differs between languages.
- We are using Python Low-Level Socket API.
- Python syntax similar to glibc.
  - Nice to use, as we can deal with objects in Python.
  - Don't have to mess around with file descriptors.



### Learning Objectives

• How to use UDP, TCP, and raw sockets.

#### Why low-level API and no frameworks?

Dealing directly with socket API lets you understand...

- ... the different states a socket can be in.
- ... how different parameters of sockets can be manipulated.
- ... what frameworks take over for you.
- ... what may be the reason when things don't work.

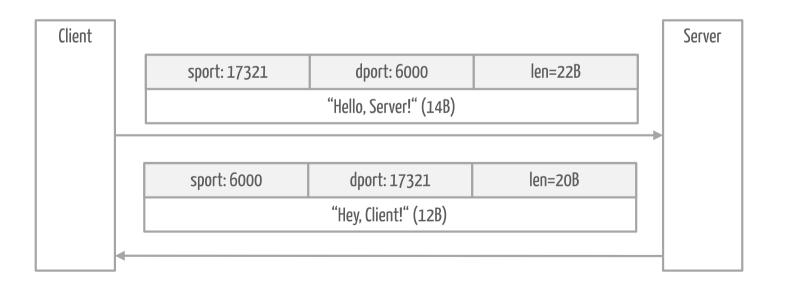


### **UDP Sockets**



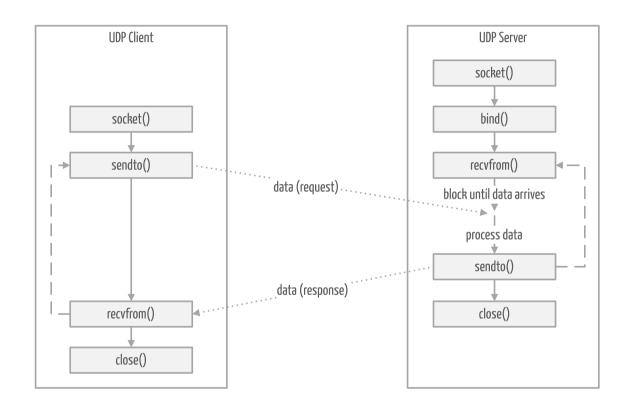
### Motivating Example

- **© Goal:** Simple UDP client/server programs.
  - Client sends data (request).
  - Server waits for request and responds (reply).





# UDP | State Graph





# UDP | Socket Creation (Step 1)

```
socket.socket(family, type, proto)
```

- Client and server create sockets.
- Returns: socket object
- family:
  - AF\_INET: IPv4 socket
  - AF INET6: IPv6 socket
  - AF\_PACKET: device-level socket (receive raw packets)
  - AF\_UNIX: UNIX socket
- type:
  - SOCK STREAM: TCP socket
  - SOCK\_DGRAM: UDP socket
  - SOCK\_RAW: Raw socket
- proto: usually zero (0)

#### **Protocols**

- The default of 0 lets the OS pick the appropriate **protocol**.
   See IP protocol field in U10.
- Examples:
  - IPPROTO\_TCP: 6
  - IPPROTO\_UDP: 17
- All protocols can be found in /etc/procotols on UNIX.



# UDP | Binding (Step 2)

#### s.bind(address\_tuple)

- Server determines service port. (mandatory step)
- Client can chose source port, otherwise random (optional step)
- address\_tuple: Address and port:
  - IP address: specific IP, or0.0.0.0 to bind to all IPs.
  - Port: specific port, or 0 for random port assignment.
- bind() can only be called once per socket.
- Binding to ports <= 1023 requires root privileges.



# UDP | Receiving (Step 3a)

s.recvfrom(bufsize[, flags])

- Returns: (buffer, address) tuple
  - buffer: byte string with received data
  - address: address of sender (tuple of IP and port)
- bufsize: Maximum number of bytes to receive.
- flags: options (see later)
- Blocking, unless otherwise specified.
- Returns entire single UDP segment.
- Specify maximum UDP payload size as bufsize (round up to next  $2^x$ )
- bufsize =  $2^{16}$  = 65536
- or: pass MSG\_TRUNC flag option

- Segment payload larger than bufsize is discarded.
- Example: recvfrom(1024) with payload of 2024 bytes will return first 1024B and discard remaining 1000B.



# UDP | Send (Step 3b)

```
s.sendto(buf, flags, addr)
```

- Returns: number of bytes sent (len(buf))
  - sendto atomic for UDP: entire buffer is sent (but no guarantee on reception)
- buf: byte string to be sent.
  - Has to fit into UDP payload of single IP packet.
  - Maximum size (IPv4): 65507
- flags: optional control flags.
- addr: address tuple of recipient.

Multicast: addr is specified when sending data and returned when receiving data, making UDP sockets usable for one-to-many transmissions.



# UDP | Maximum Segment Size (IPv4)

**UDP Length** field is 16 bit large, hence maximum theoretical length is:

$$UDP_{max} = 2^{16} - 1 = 65535.$$

**IP Total Length** field is 16 bit large, hence maximum theoretical length is:

$$IP_{max}=2^{16}-1=65535$$
 (too).

UDP datagram must fit into one IP packet. Probably fragmented, but maximum length IPmax holds.

Operating systems (Linux, etc.) know that you are putting UDP segments into IP packets.

IP ( $IP_{hdr} \geq 20$ ) and UDP ( $UDP_{hdr} = 8$ ) have a minimal / fixed header size.

Consequently, the practical maximum size of a UDP segment is:

$$IP_{max} - IP_{hdr} - UDP_{hdr} = 65535 - 20 - 8 = 65507$$



# UDP | Close (Step 4)

#### s.close()

- Flushes all data if pending.
- Closing UDP sockets:
  - Free system resources.
  - Undo binding.
  - No explicit notification of remote side.

```
import socket
s = socket.socket(
        socket.AF_INET,
                                    # family
        socket.SOCK DRAM,
        0)
                                    # proto
s.bind(("192.168.0.1", 4567))
                                    # (IP, port)
(data, address) = s.recvfrom(65536) # bufsize
data = data.decode("utf-8").toupper()
c = s.sendto(data.encode("utf-8"),
                                    # flags
            address)
                                    # addr
s.close()
```



### TCP Sockets



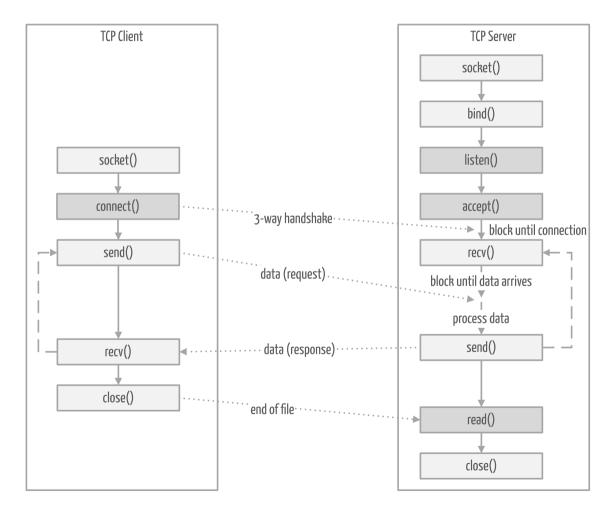
# TCP | Motivation

- TCP is
  - o connection-oriented.
  - o stream-oriented.
  - o reliable.

- TCP sockets provide
  - similar functions to UDP.
  - additional connection management.
  - additionally message segmentation.



# TCP | State Graph





# TCP Server | Preparation (#1)

- A) Create socket
  - Command: socket.socket(socket.AF\_INET, socket.SOCK\_STREAM, 0)
- B) Bind socket to port (similar to UDP)
  - Command: s.bind(address\_tuple)
  - address\_tuple: IP and port at which to wait for clients.
- C) Listen for connections made to socket
  - Starting at this command, kernel queues connections
  - Command: s.listen(backlog)
  - backlog: maximum number of queued connections



# TCP Server | Connection Backlog

#### Concept

- After listen, kernel automatically completes TCP handshakes.
- OS kernel has two queues:
  - Incomplete connections: SYN received.
  - Completed connections: Handshake completed.
- backlog: maximum length of total entries (sum of both queues).
- Additional incoming TCP connections are ignored.

#### **Design Decision**

- Backlog of 0 is allowed, but can trigger unspecified behaviour.
- Backlog of 5 is a common value, but only suitable for small servers.
- Modern TCP-based servers have much larger values.



# TCP Server | Accepting Connections (#2)

#### s.accept()

- Remove connection from completed listen queue.
  - First in, first out (FIFO) order.
- Returns: (conn, address) tuple.
  - conn: socket object of TCP connection.
  - address: address of sender (tuple of IP and port).
- accept is a blocking call:
  - Wait until client connects.
  - Blocks forever if nobody connects.



### Quiz

### **②** Assume a socket has a backlog of 5 connections. What is the state after the following sequence of events?

- SYNs from 3 sources.
- Handshake completed with additional 2 sources.
- 1 socket accept()ed.

A: Backlog is full.

C: Two slots free.

B: One slot free.

**D:** Three slots free.

#### Answer:

- A: Wrong, accept() removes one.
- C: Completed connections are not automagically removed.
- B: True, because 3 incompletes remain and one completed remains.
- D: Why? Off-by-two, maybe?



### TCP Client | Establish Connection

#### A) Create Socket

 Command: socket.socket(socket.AF\_INET, socket.SOCK\_STREAM, 0)

#### B) Connect to server

- Command:s.connect(address\_tuple)
- address\_tuple: IP and port of destination server
- Socket connected when connect returns.



# TCP Sockets | Recv / Send (#3a/b)

#### Receiving

- Command:
  - s.recv(bufsize[, flags])
- Returns: Buffer that was received (up to bufsize bytes).
  - Empty buffer indicates other side **terminated** the connection.
- bufsize: Maximum number of bytes to receive.
- flags: options (see later)
- Blocking, unless otherwise specified.

#### Sending

- Command:
  - s.send(buf[, flags])
- Returns: Number of bytes sent.
  - send is **not** atomic for TCP:
     partial data might be sent.
- buf: Byte string to be sent.
  - Can be of any size, will be fragmented accordingly.
  - If too large, send() will only send partial data.
- flags: options (see later)

For UDP, recvfrom() returned source address tuple. For TCP recv() does not. Why?



# TCP Send/Recv | Correct Usage

- Issues with recv and send:
  - May return less than bufsize bytes.
  - recv may return multiple messages/packets.
  - send may not send all data passed as argument.
- Corner Cases:
  - You may need multiple send calls to transmit all data.
  - Multiple send calls may be handeled by a single recv call.
  - A single send call may need multiple calls to recv.
- Cases may not be encountered under perfect conditions (e.g. testing), but it is quite common under high load.

#### Reliably Send all data

• sendall() in Python returns only if all data was sent (more convenient than C calls).

#### Reliably Recv all data

- Not trivial.
- Depends on message separation (e.g. \n, \0 or \n\r).



# TCP Send/Recv | Correct Usage II

• Idea: Reliably receive all data by using fixed message length.

Approach: Use MSG\_WAITALL.

- Optional flag for recv().
- recv() waits until bufsize bytes have been received.
- (... or a signal was caught.)
- (... or the connection was terminated)

**Problem:** Protocols with fixed length are usually bad.

**②** Why is fixed length bad?

**A:** Sending larger messages than predefined size is not easy.

**B:** Fixed size messages waste space, when the message is shorter than the fixed length.

**C:** Sender and recipient have to agree on length, but might have different opinions.

A,B,C: True!



### TCP Message Separation Options

Sender and receiver have to agree on message separation.

#### A) Fixed message length (e.g 1024 bytes).

- Easy to program, but lots of padding overhead.
- Large messages may not fit. **Bad design** in most cases.

#### B) Only one message per socket.

- Connect, send, recv, disconnect (e.g. basic HTTP used this).
- Connection establishment overhead.

#### C) Delimit messages by well-defined string/character (e.g. \n).

- Need escaping (or well-defined delimiter forbidden in message).
- Unclear how much needs to be read to reach delimiter.

#### D) Message length indicated at start of message.

- Acceptable overhead, very flexible.
- Message length field must have fixed size (therefore limits overall size).



# TCP Sockets | Closing (#4)

#### s.close()

- Flushes all data if pending.
- Closing TCP sockets:
  - Send FIN segment ("I am done sending data.")
  - Other side calling recv() will get empty buffer.
  - At later point, remote will also close and send FIN.



### **Raw Sockets**



### Raw | Motivation

- TCP, UDP sockets:
  - Do most of the work for you (put data into a socket, give an address and the rest is figured out by the OS).
  - Required for transport layer networking.
- Sometimes, it is necessary to work with protocol on lowest layers:
  - ICMP: directly in IP without UDP/TCP.
  - ARP: Ethernet frame, without IP.
  - o ...
- Security testing requires you to work on raw layer, so that packets can be crafted that would never be created in such a way by an OS.
- For more control and low-level networking: Use raw sockets!
  - Sender **can** generate arbitrary packets (including IP and Ethernet).
  - Sender **must** specify every detail of packets (including checksums).



# Raw | Sending Packets

- A) Create a socket with socket.SOCK\_RAW.
  - socket.socket(socket.AF\_PACKET, socket.SOCK\_RAW, 0)
  - Note: Requires root/admin privileges (as we are circumventing the kernel/OS).
- B) Bind to an interface (instead of IPv4 address).
  - s.bind(("eth0", 0))
- C) Generate and send packet.
  - s.send("A" \* 20) (malformed packet)
  - Include Ethernet and IP header if needed.



# Raw | Sending Packets (Convenient)

Generate packets using third-party library dpkt.

```
import dpkt
# create Ethernet, IP and UDP objects
eth = dpkt.ethernet.Ethernet(src="08:00:27:bf:17:b4", dst="ff:ff:ff:ff:ff:ff")
ip = dpkt.ip.IP(src="127.0.0.1", dst="127.6.6.6")
udp = dpkt.udp.UDP(sport=1234, dport=5678)
eth.data = ip  # set IP pkt as Ethernet payload
ip.data = udp  # set UDP segment to be IP pkt payload
ip.len = len(ip)  # set IP length
udp.ulen = len(udp) # set UDP length
repr(ip)
                       # implicitly computes IP checksum
# Output:
# "IP(len=28, src='127.0.0.1', dst='127.6.6.6', data=UDP(sport=1234, dport=5678))"
repr(eth)
                       # show Ethernet <u>packet</u>
# Output:
# "Ethernet(dst='ff:ff:ff:ff:ff:ff', src='08:00:27:bf:17:b4',
             data=IP(len=28, src='127.0.0.1', dst='127.6.6.6',
             data=UDP(sport=1234, dport=5678)))"
```



# Raw | Receiving Packets

- **© Goal:** Receive all packets on interface.
- Create socket with socket.SOCK\_RAW:

```
ETH_P_ALL = socket.ntohs(0x0003)  # all Ethernet protos
socket.socket(socket.AF_PACKET,  # domain
socket.SOCK_RAW,  # type
ETH_P_ALL)  # protocol (non-empty / non-default)
```

- Put NIC to promiscuous mode and even get packets not intended for you.
- Receive packet:
  - o s.recv(4096)
  - Return Ethernet frame as captured on interface.
  - Note: captures packets even if they are filtered by firewall.



# Raw | Receiving Packets (Convenient)

Parse packets using third-party library dpkt.

```
import dpkt
pkt = s.recv(4096)
                             # receive via raw socket
dpkt.ethernet.Ethernet(pkt)
                              # print dpkt parsing result
# Output
# Ethernet(dst="\x08\x00'\xbf\x17\xb4", src='RT\x00\x125\x02',
           data=IP(len=40, id=17571, p=6, sum=7709, src='\n\x00\x02\x02', dst='\n\x00\x02\x0f', opts=''.
           data=TCP(sport=54945, dport=22, seq=126587142, ack=4007547311, flags=16, sum=59372, opts='')))
eth = dpkt.ethernet.Ethernet(pkt)
                                    # get Ethernet frame object
ip = eth.data
                                    # get IP packet object
tcp = ip.data
                                    # get TCP packet object
tcp.dport # access member variables
# Output:
# 22
```



# Raw | Power

**"With Great Power Comes Great Responsibility" -**

#### **Q** Network Link Sniffing

- See every packet on the wire.
   And many of those in the air.
- Passive attacks such as local eavesdropping become easy.

#### Arbitrary Packet Injection

- You can inject completely arbitrary packets.
- "Spoof" any arbitrary MAC or IP address(es).
- Low resource requirements (in contrast to full TCP/IP stack).



### Wrap-Up



#### **Take-Home Messages**

- Sockets are the API for network programming and an important concept.
- **UDP socket** offer atomic send/receive operations.
- TCP sockets require more attention (connection handling, partial receive/send).
- Raw sockets give your complete control, but also put burden on you.

#### **E** Further Reading

- Python 3.6 Low-level networking interface
- Python 3.6 Struct module
- Stevens, Fenner, Rudoff: UNIX Network Programming, Volume 1
  - Book examples in C, we use Python.