

## 5. Optimization problems

- optimization problems
- solving optimization problems
- problem transformations
- control example

## Optimization problem

$$\begin{array}{ll}\text{minimize} & f(x) \\ \text{subject to} & g_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_j(x) = 0, \quad j = 1, \dots, p\end{array}\tag{5.1}$$

- $x = (x_1, \dots, x_n)$  is the *optimization* or *decision variable*
- $f : \mathbb{R}^n \rightarrow \mathbb{R}$  is the *objective* or *cost* function
- $g_i : \mathbb{R}^n \rightarrow \mathbb{R}$  are the *inequality constraint* functions
- $h_j : \mathbb{R}^n \rightarrow \mathbb{R}$  are the *equality constraint* functions
- can be compactly written as

$$\begin{array}{ll}\text{minimize} & f(x) \\ \text{subject to} & g(x) \leq 0 \\ & h(x) = 0\end{array}\tag{5.2}$$

where  $g(x) = (g_1(x), \dots, g_m(x))$  and  $h(x) = (h_1(x), \dots, h_p(x))$

## Feasible and optimal points

**Feasible point:**  $\hat{x}$  is *feasible* if  $\hat{x} \in \text{dom } f$  and it satisfies the constraints

**Solution:** a point  $x^\star$  is an *optimal point* or a *solution* if it is feasible and

$$f(x^\star) \leq f(x) \quad \text{for any feasible } x$$

### Optimal value

greatest  $\rho$  such that  $\rho \leq f(x)$  for all feasible  $x$ , denoted by  $p^\star$

- if there exists an optimal point  $x^\star$ , then  $p^\star = f(x^\star)$ 
  - we say the optimal value is *attained* or *achieved* and the problem is *solvable*
- a minimization problem is *unbounded below* if  $p^\star = -\infty$
- if a minimization problem is infeasible, then we let  $p^\star = +\infty$

## Examples

- the unconstrained problem

$$\text{minimize} \quad (x_1 - 1)^2 + (x_2 - 1)^2 = \|x - \mathbf{1}\|^2$$

has optimal value  $p^\star = 0$ , which is attained at the optimal point  $x^\star = (1, 1) = \mathbf{1}$

- the problem

$$\begin{array}{ll}\text{minimize} & x_1 + x_2 \\ \text{subject to} & -x_1 \leq 10 \\ & x_2 \geq 0\end{array}$$

has solution  $x^\star = (-10, 0)$  and  $p^\star = -10$

- the problem

$$\text{minimize} \quad x_1^2 - x_2^2$$

is unbounded below since  $f(x) \rightarrow -\infty = p^\star$  as  $|x_2| \rightarrow \infty$

- consider the problem

$$\text{minimize } f(x) = e^{-x}$$

the optimal value is  $p^\star = 0$ , but it is not attained since it only holds as  $x \rightarrow \infty$

- for the problem

$$\text{minimize } f(x) = 1/x, \quad \text{dom } f = \{x \mid x > 0\}$$

we have  $p^\star = 0$  but is not attained by any feasible  $x$

- the problem

$$\begin{array}{ll} \text{minimize} & x_1^2 + x_2^2 \\ \text{subject to} & x_1 + x_2 \leq 1 \\ & x_1 + x_2 \geq 2 \end{array}$$

is infeasible; hence,  $p^\star = \infty$

## Domain and implicit constraints

the *domain* of an optimization problem is

$$\mathcal{D} = \text{dom } f \cap \bigcap_{i=1}^m \text{dom } g_i \cap \bigcap_{j=1}^p \text{dom } h_j$$

- the standard form problem (5.1) has an *implicit constraint*  $x \in \mathcal{D}$
- the *explicit constraints* are  $g(x) \leq 0$  and  $h(x) = 0$
- a problem is *unconstrained* if it has no explicit constraints ( $m = p = 0$ )
- for example, the unconstrained problem

$$\text{minimize} \quad -\log x_1 + \log(x_2 - x_1)$$

has implicit constraints  $x_1 > 0$ ,  $x_2 - x_1 > 0$ , which defines the domain  $\mathcal{D}$

## Maximization problems

$$\begin{array}{ll}\text{maximize} & f(x) \\ \text{subject to} & g(x) \leq 0 \\ & h(x) = 0\end{array}$$

- $f$  is often called *utility function* instead of cost
- note that  $\max f(x) = -\min -f(x)$
- thus, maximization problems can be written as minimization problems

$$\begin{array}{llll}\text{maximize} & f(x) & & \text{minimize} & -f(x) \\ \text{subject to} & g(x) \leq 0 & \Leftrightarrow & \text{subject to} & g(x) \leq 0 \\ & h(x) = 0 & & & h(x) = 0\end{array}$$

both problems have the same solutions

## Standard form

we refer to problem (5.1) as an optimization problem in *standard form*

- equality  $r_j(x) = \tilde{r}_j(x)$  is same as  $h_j(x) = 0$  with  $h_j(x) = r_j(x) - \tilde{r}_j(x)$
- inequality  $\tilde{g}_i(x) \geq 0$  is same as  $g_i(x) \leq 0$  with  $g_i(x) = -\tilde{g}_i(x)$
- maximization can be represented as minimization by changing the objective sign

**Example:** the problem

$$\begin{array}{ll}\text{maximize} & -x_1^2 + x_2^2 \\ \text{subject to} & -x_1 + x_2 \geq 10 \\ & x_2 = 2 - x_1\end{array}$$

can be expressed in standard form:

$$\begin{array}{ll}\text{minimize} & x_1^2 - x_2^2 \\ \text{subject to} & x_1 - x_2 + 10 \leq 0 \\ & x_1 + x_2 - 2 = 0\end{array}$$



# Outline

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## Set-constrained problems

$$\underset{x \in \mathcal{X}}{\text{minimize}} \quad f(x)$$

- find  $x$  that minimizes  $f(x)$  among all points in the *constraint set*  $\mathcal{X} \subseteq \mathbb{R}^n$
- for problem (5.1), the constraint set is described by *functional constraints*

$$\mathcal{X} = \{x \mid g_i(x) \leq 0, h_j(x) = 0, i = 1, \dots, p, j = 1, \dots, m\}$$

- this is not always the case, for example consider the integer set

$$\mathcal{X} = \{1, 2, 3\} \subset \mathbb{R}$$

## Existence of a solution

- existence of a solution is not always guaranteed
- can be guaranteed to exist under some conditions

**Solution existence:** a continuous  $f$  over  $\mathcal{X}$  has an optimal point over  $\mathcal{X}$  if either

- $\mathcal{X}$  is nonempty and compact (closed and bounded)
- $f$  is *coercive*:

$$\lim_{\|x\| \rightarrow \infty} f(x) = \infty$$

and  $\mathcal{X} \subseteq \mathbb{R}^n$  is a nonempty closed set

## Simple problems solution

- general optimization problems require sophisticated methods to solve them that utilize derivatives, linear equations, nonlinear operators,...etc
- that said, there are some simple optimization problems that can be solved by inspection or using some basic inequalities such as Cauchy-Schwarz

### Example

$$\begin{array}{ll}\text{minimize} & \|x - \mathbf{1}\| \\ \text{subject to} & -\mathbf{1} \leq x \leq 0\end{array}$$

- we seek to find a feasible  $x$  that is closest in distance to  $\mathbf{1}$
- we have

$$x^{\star} = 0 \quad \text{and} \quad p^{\star} = \|\mathbf{1}\| = \sqrt{n}$$

## Example

$$\begin{array}{ll}\text{minimize} & x_1 + x_2 \\ \text{subject to} & x_1^2 + x_2^2 \leq 1\end{array}$$

- using Cauchy-Schwarz, we can lower bound the objective by

$$x_1 + x_2 = \mathbf{1}^T x \geq -\|\mathbf{1}\| \|x\| \geq -\sqrt{2}$$

for all  $x_1^2 + x_2^2 \leq 1$

- the minimum value is attained at  $x = (-1/\sqrt{2}, -1/\sqrt{2})$ , which is feasible
- hence, the optimal point is  $x = (-1/\sqrt{2}, -1/\sqrt{2})$

# Optimization methods

- after formulating the problem, a suitable algorithm is applied to solve it
- an optimization *algorithm* is a set of calculations and rules that are followed to find a solution or an approximate solution to an optimization problem

**Iterative algorithms:** start from an initial guess  $x^{(0)}$  and computes

$$x^{(k+1)} = F(x^{(k)}), \quad k = 0, 1, \dots$$

- $F$  depends on  $f(x)$ ,  $g_i(x)$ ,  $h_j(x)$  to generate a new estimate  $x^{(k+1)}$
- moving from  $x^{(k)}$  to  $x^{(k+1)}$  is called an *iteration* of the algorithm
- stops when a good estimate of a solution is reached or  $k = k^{\max}$  for some  $k^{\max}$

# General iterative algorithm

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**given** a starting point  $x^{(0)}$ , tolerance  $\epsilon$ , stopping criteria, and  $k^{\max}$

**for**  $k \geq 1$

1. determine a search direction  $v^{(k)}$
2. **quit** if stopping criterion is met (depends on  $\epsilon$ ), and output  $x^{(k)}$
3. determine scalar  $\alpha_k$
4. update:

$$x^{(k+1)} = x^{(k)} + \alpha_k v^{(k)}$$

**until**  $k = k^{\max}$

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- $\alpha_k$  is called the *stepsize* or *learning rate*
- $v^{(k)}$  depends on  $f, g_i, h_j$  and their derivatives if differentiable

## Local minimum point

a point  $x^\circ \in \mathcal{X}$  is a *local minimizer* or *local minimum point* (locally optimal) of

$$\underset{x \in \mathcal{X}}{\text{minimize}} \quad f(x)$$

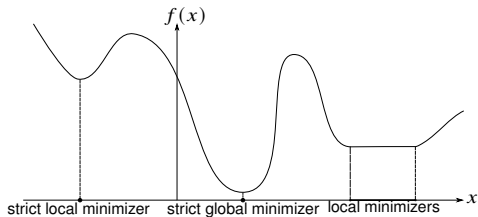
if there exists a scalar  $r > 0$  such that:

$$f(x^\circ) \leq f(x) \quad \text{for all } x \in \mathcal{X} \quad \text{and} \quad \|x - x^\circ\| \leq r$$

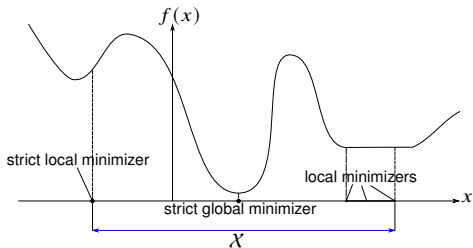
- if  $f(x^\circ) < f(x)$ , then the point  $x^\circ$  is called a *strict local minimizer*
- $x^\star \in \mathcal{X}$  is a *global minimizer* (*minimum point*) if  $f(x^\star) \leq f(x)$  for all  $x \in \mathcal{X}$
- ‘globally optimal’ is used for ‘optimal’ to distinguish from ‘locally optimal’
- a point is a maximum point of  $f$  if it is a minimum point of  $-f$



# Global and local optima

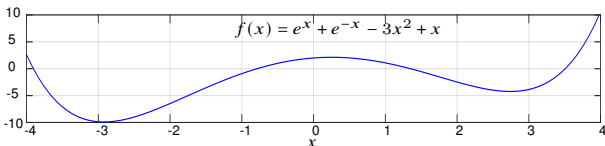
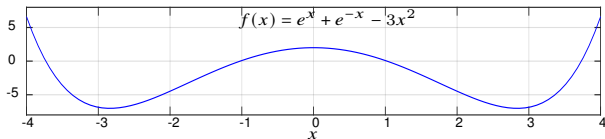
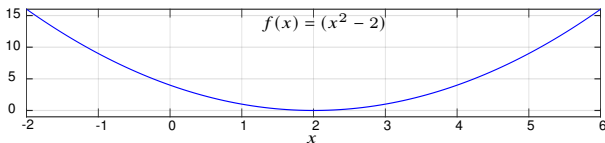


unconstrained case  $\mathcal{X} = \mathbb{R}^n$



constrained case  $x \in \mathcal{X}$

## Example



- $(x - 2)^2$ ;  $p^* = \min f(x) = 0$ ; global minimizer  $x^* = 2$
- $e^x + e^{-x} - 3x^2$ ;  $p^* = -7.02$ ; two global minima:  $x^* = \pm 2.84$
- $e^x + e^{-x} - 3x^2 + x$ ;  $p^* = -9.9$ ; global min.  $x^* = -2.92$ ; local min. at  $x = 2.74$

# Nonlinear optimization methods

## Local optimization methods

- find a locally optimal solution with no global optimality guarantees
- fast, can handle large-scale problems, and are widely applicable
- can be used to improve the performance of an engineering design obtained by manual, or other, design methods

## Global optimization methods

- true global solution is found with optimality guarantees
- difficult to find in general; even small problems, with a few tens of variables, can take a very long time (*e.g.*, hours or days) to solve
- usually seek the global optimum by finding local solutions to a sequence of approximate subproblems

# Efficiently solvable problem classes

## (linear) Least squares

$$\text{minimize} \quad \sum_{i=1}^m \left( \sum_{j=1}^n a_{ij} x_j - b_i \right)^2$$

where the coefficients  $a_{ij}$ ,  $b_i$  are given constants

- reliable and efficient algorithms and software
- least-squares problems are easy to recognize
- has many applications such as data-fitting and linear estimation

## Linear program (optimization)

$$\begin{array}{ll}\text{minimize} & \sum_{j=1}^n c_j x_j \\ \text{subject to} & \sum_{j=1}^n a_{ij} x_j \leq b_i, \quad i = 1, \dots, m \\ & \sum_{j=1}^n g_{ij} x_j = h_i, \quad i = 1, \dots, p\end{array}$$

the coefficients  $c_j, a_{ij}, g_{ij}, h_i, b_i$  are given constants

- there exist robust and efficient algorithms and software for solving LPs
- LPs isn't as immediately recognizable as that of least-squares problems
- common techniques are available to transform various problems into LPs

## Convex optimization

$$\begin{array}{ll}\text{minimize} & f(x) = g_0(x) \\ \text{subject to} & g_i(x) \leq 0, \quad i = 1, \dots, m, \\ & \sum_{j=1}^n a_{ij}x_j = b_i, \quad i = 1, \dots, p\end{array}$$

the coefficients  $a_{ij}$ ,  $b_i$  are given constants

- the objective and constraints functions are *convex*:

$$g_i(\theta x + (1 - \theta)y) \leq \theta g_i(x) + (1 - \theta)g_i(y), \quad 0 \leq \theta \leq 1$$

- if objective or one constraint is nonconvex, problem is called *nonconvex*

# Convex optimization

- include least-squares problems and linear programs as special cases
- has many of applications
- reliable and efficient algorithms
- difficult to recognize
- many tricks can be used to transform nonconvex problems into convex form
- basis for several heuristics for solving nonconvex problems

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## Equivalent optimization problems

two optimization problems are *equivalent* if from a solution of one, we can find a solution of the other, and vice versa

- for example, maximization problems are equivalent to minimization problems
- many optimization problems can be transformed into equivalent ones
- can be very useful if the equivalent problem is easier to solve

## Scaling and slack variables

**Scaling:** problem (5.1) is equivalent to

$$\begin{array}{ll}\text{minimize} & \alpha f(x) \\ \text{subject to} & \beta_i g_i(x) \leq 0, \quad i = 1, \dots, m \\ & \gamma_j h_j(x) = 0, \quad j = 1, \dots, p\end{array}$$

- $\alpha > 0$ ,  $\beta_i > 0$  and  $\gamma_j \neq 0$
- ‘scaling’ does not alter solutions

**Slack variables:** problem (5.1) is equivalent to

$$\begin{array}{ll}\text{minimize} & f(x) \\ \text{subject to} & s_i \geq 0, \quad i = 1, \dots, m \\ & g_i(x) + s_i = 0, \quad i = 1, \dots, m \\ & h_j(x) = 0, \quad j = 1, \dots, p\end{array}$$

- the variables are  $x \in \mathbb{R}^n$  and  $s \in \mathbb{R}^m$
- we replaced  $g_i(x) \leq 0$  with  $g_i(x) + s_i = 0$  for some  $s_i \geq 0$
- variable  $s_i$  is called *slack variable* associated with inequality constraint  $g_i(x) \leq 0$

## Monotone transformations

$$\begin{array}{ll}\text{minimize} & \phi(f(x)) \\ \text{subject to} & g_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_j(x) = 0, \quad j = 1, \dots, p\end{array}$$

- $\phi : \mathbb{R} \rightarrow \mathbb{R}$  is a continuous and monotone increasing function, *i.e.*,

$$\phi(a) > \phi(b) \quad \text{for all } a > b \text{ over the optimization domain}$$

this implies that  $\phi$  is one-to-one and its inverse  $\phi^{-1}$  is well defined

- this problem is equivalent to (5.1)

**Constraint transformation:** find  $\psi_i : \mathbb{R} \rightarrow \mathbb{R}$  and  $\varphi_j : \mathbb{R} \rightarrow \mathbb{R}$  so that

- $\psi_i(g_i(x)) \leq 0$  if and only if  $g_i(x) \leq 0$
- $\varphi_j(h_j(x)) = 0$  if and only if  $h_j(x) = 0$

## Example

$$\begin{array}{ll}\text{minimize} & \|x\| \\ \text{subject to} & g(x) \leq 0\end{array}$$

- norm is non-differentiable and we prefer differentiable objectives
- norm is nonnegative and  $\phi(\cdot) = (\cdot)^2$  is monotone increasing over nonneg. no.
- hence, we can transform the problem into

$$\begin{array}{ll}\text{minimize} & \|x\|^2 \\ \text{subject to} & g(x) \leq 0\end{array}$$

- the new objective function is differentiable, which simplifies the problem

## Change of variables

$$\begin{array}{ll}\text{minimize} & f(F(y)) \\ \text{subject to} & g_i(F(y)) \leq 0, \quad i = 1, \dots, m \\ & h_j(F(y)) = 0, \quad j = 1, \dots, p\end{array}$$

- $F : \mathbb{R}^n \rightarrow \mathbb{R}^n$  is 1-to-1 with image covering problem domain ( $\mathcal{D} \subseteq F(\text{dom } F)$ )
- this implies for each  $x \in \mathcal{D}$ , there's a unique  $y \in \text{dom } F$  such that

$$x = F(y) \iff y = F^{-1}(x)$$

- if  $x$  solves (5.1), then  $y = F^{-1}(x)$  solves the above problem
- if  $y$  solves the above problem, then  $x = F(y)$  solves (5.1)

## Example

$$\begin{array}{ll}\text{minimize} & x_1 x_2 x_3^2 \\ \text{subject to} & x_1 x_2 \leq 2 \\ & x_1, x_2, x_3 > 0\end{array}$$

- $\log(\cdot)$  is strictly increasing (for non-negative argument),
- hence, we can use monotone transformations on objective and constraints

$$\log(x_1 x_2 x_3^2) \quad \text{and} \quad \log(x_1 x_2) \leq \log(2)$$

- also use the change of variable  $y_i = \log x_i$  to transform the problem into

$$\begin{array}{ll}\text{minimize} & y_1 + y_2 + 2y_3 \\ \text{subject to} & y_1 + y_2 \leq \log 2\end{array}$$

this is a linear program, which is easier to solve

## Example

$$\begin{array}{ll}\text{minimize} & x_1 x_2 - x_3^2 \\ \text{subject to} & x_1 + x_2 + x_3 \leq 20 \\ & x_2 \geq 10\end{array}$$

- let  $y_1 = (x_1 + x_2)/2$  and  $y_2 = (x_1 - x_2)/2$  so that

$$x_1 = y_1 + y_2, \quad x_2 = y_1 - y_2$$

- thus, we can transform the problem into

$$\begin{array}{ll}\text{minimize} & y_1^2 - y_2^2 - x_3^2 \\ \text{subject to} & 2y_1 + x_3 \leq 20 \\ & y_1 - y_2 \geq 10\end{array}$$

- the objective is now separable in the new variables
- for separable problems, there exist efficient specialized optimization methods

## Eliminating equality constraints

suppose  $\phi : \mathbb{R}^k \rightarrow \mathbb{R}^n$  and  $h(x) = 0$  iff there is some  $z$  such that  $x = \phi(z)$

then problem (5.1) is equivalent to

$$\begin{array}{ll}\text{minimize} & f(\phi(z)) \\ \text{subject to} & g_i(\phi(z)) \leq 0, \quad i = 1, \dots, m\end{array}$$

- for optimal  $z$ ,  $x = \phi(z)$  is optimal for the original problem
- for optimal  $x$ , any  $z$  such that  $x = \phi(z)$  is optimal for the transformed problem

### Example

$$\begin{array}{ll}\text{minimize} & x_1 + x_2 + x_3^2 \\ \text{subject to} & x_1 - x_2x_3 = 1\end{array}$$

we use  $x_1 = 1 + x_2x_3$  to remove the equality constraint and get

$$\text{minimize} \quad 1 + x_2x_3 + x_2 + x_3^2$$

in this case, we have  $\phi(z_1, z_2) = (1 + z_1z_2, z_1, z_2)$



## Example

$$\begin{array}{ll}\text{minimize} & x_1 + 4x_2 + x_3 \\ \text{subject to} & 2x_1 - 2x_2 + x_3 = 4 \\ & x_1 - x_3 = 1 \\ & x_2 \geq 0, x_3 \geq 0\end{array}$$

- using the equality constraints, we have  $x_1 = 1 + x_3$  and

$$2x_1 - 2x_2 + x_3 = 2(1 + x_3) - 2x_2 + x_3 = 4 \implies x_2 = \frac{3}{2}x_3 - 1$$

- hence, the problem can be simplified to

$$\begin{array}{ll}\text{minimize} & 8x_3 + 3 \\ \text{subject to} & x_3 \geq 2/3\end{array}$$

with solution  $x_3 = 2/3$  from which we can find  $x^\star = (5/3, 0, 2/3)$

## Eliminating linear constraints

let  $A \in \mathbb{R}^{p \times n}$ ,  $b \in \mathbb{R}^p$ , and consider the constraints

$$h(x) = Ax - b = 0$$

- solution of  $Ax = b$  can be parametrized by (see page 2.41)

$$x = \hat{x} + Fz \quad \text{for any arbitrary } z \in \mathbb{R}^{(n-p)}$$

where columns of  $F$  form a basis for the nullspace of  $A$  ( $\text{range}(F) = \text{null}(A)$ )

- we can use change of variable  $x = \hat{x} + Fz$ , to transform (5.1) into

$$\begin{array}{ll} \text{minimize} & f(\hat{x} + Fz) \\ \text{subject to} & g_i(\hat{x} + Fz) \leq 0, \quad i = 1, \dots, m \end{array}$$

with variable  $z$

## Example

$$\begin{array}{ll}\text{minimize} & f(x_1, \dots, x_n) \\ \text{subject to} & x_1 + \dots + x_n = b\end{array}$$

- we can eliminate any  $x_i$ , we choose  $x_n$ :

$$x_n = b - x_1 - \dots - x_{n-1}$$

- the above corresponds to the choice

$$\hat{x} = be_n, \quad F = \begin{bmatrix} I \\ -\mathbf{1}^T \end{bmatrix} \in \mathbb{R}^{n \times (n-1)}$$

- the transformed problem is

$$\text{minimize} \quad f(x_1, x_2, \dots, b - x_1 - \dots - x_{n-1})$$

## Adding equality constraints

- sometimes it is useful to introduce equality constraints
- for example, consider

$$\text{minimize } f(Ax + b)$$

where  $A \in \mathbb{R}^{m \times n}$  and  $b \in \mathbb{R}^m$

- the above problem is equivalent to

$$\begin{array}{ll}\text{minimize} & f(z) \\ \text{subject to} & z = Ax + b\end{array}$$

with variables  $x \in \mathbb{R}^n$  and  $z \in \mathbb{R}^m$

## Optimizing over some variables

it holds that (with abuse of inf notation):

$$\min_{x,y} f(x,y) = \min_x \tilde{f}(x)$$

where  $\tilde{f}(x) = \min_y f(x,y)$

### Example

$$\begin{array}{ll} \text{minimize} & x_1^T Q_{11} x_1 + 2x_1^T Q_{12} x_2 + x_2^T Q_{22} x_2 \\ \text{subject to} & g_i(x_1) \leq 0, \quad i = 1, \dots, m, \end{array}$$

where  $Q_{11}$  and  $Q_{22}$  are symmetric; we can analytically minimize over  $x_2$ :

$$\min_{x_2} (x_1^T Q_{11} x_1 + 2x_1^T Q_{12} x_2 + x_2^T Q_{22} x_2) = x_1^T (Q_{11} - Q_{12} Q_{22}^{-1} Q_{12}^T) x_1$$

thus, the original problem is equivalent to

$$\begin{array}{ll} \text{minimize} & x_1^T (Q_{11} - Q_{12} Q_{22}^{-1} Q_{12}^T) x_1 \\ \text{subject to} & g_i(x_1) \leq 0, \quad i = 1, \dots, m \end{array}$$

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# Dynamical system

a nonlinear dynamical system has the form

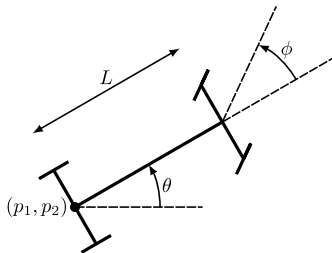
$$x_{k+1} = h(x_k, u_k), \quad k = 0, 1, \dots, K$$

- $x_k \in \mathbb{R}^n$  is the *state vector* at instant  $k$
- $u_k \in \mathbb{R}^m$  is the *input* or *control* at instant  $k$
- $h : \mathbb{R}^{n+m} \rightarrow \mathbb{R}^n$  describes evolution of the system (system dynamics)
- examples: vehicle dynamics, robots, chemical plants evolution...

## Optimal control

- initial state  $x_1 = x_{\text{initial}}$  is known
- choose the inputs  $u_1, \dots, u_K$  to achieve some goal for the states/inputs

## Car control example



$$\frac{dp_1}{dt}(t) = s(t) \cos \theta(t)$$

$$\frac{dp_2}{dt}(t) = s(t) \sin \theta(t)$$

$$\frac{d\theta}{dt}(t) = (s(t)/L) \tan \phi(t)$$

- $L$  wheelbase (length)
- $p(t)$  position
- $\theta(t)$  orientation (angle)
- $\phi(t)$  steering angle
- $s(t)$  speed
- we control speed  $s$  and steering angle  $\phi$



## Discretized car dynamics

$$p_1(t + \tau) \approx p_1(t) + \tau s(t) \cos \theta(t)$$

$$p_2(t + \tau) \approx p_2(t) + \tau s(t) \sin \theta(t)$$

$$\theta(t + \tau) \approx \theta(t) + \tau(s(t)/L) \tan \phi(t)$$

- $\tau$  is a small time interval
- let state vector  $x_k = (p_1(k\tau), p_2(k\tau), \theta(k\tau))$
- input vector  $u_k = (s(k\tau), \phi(k\tau))$
- discretized model is

$$x_{k+1} = h(x_k, u_k)$$

with

$$h(x_k, u_k) = x_k + \tau (u_k)_1 \begin{bmatrix} \cos(x_k)_3 \\ \sin(x_k)_3 \\ (\tan(u_k)_2)/L \end{bmatrix}$$

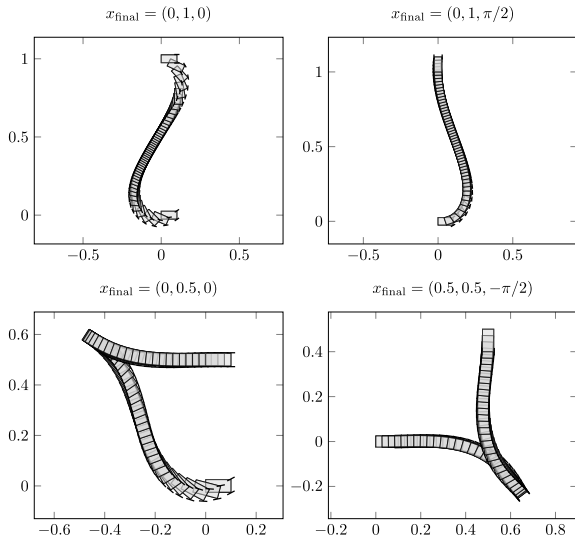
## Car control problem

- move car from given initial to desired final position and orientation
- using a small and slowly varying input sequence

### Problem formulation

$$\begin{array}{ll}\text{minimize} & \sum_{k=0}^K \|u_k\|^2 + \rho \sum_{k=0}^{K-1} \|u_{k+1} - u_k\|^2 \\ \text{subject to} & x_1 = h(0, u_0) \\ & x_{k+1} = h(x_k, u_k), \quad k = 1, \dots, K-1 \\ & x_{\text{final}} = h(x_K, u_K)\end{array}$$

- variables  $u_0, \dots, u_N$ , and  $x_1, \dots, x_N$
- the initial state is assumed to be zero
- the objective ensures the input is small with little variation
- $\rho > 0$  is an input variation trade-off parameter



solution trajectories with different final states; the outline of the car shows the position  $(p_1(k\tau); p_2(k\tau))$ , orientation  $\theta(k\tau)$ , and the steering angle  $\phi(k\tau)$  at time  $k\tau$

## References and further readings

- S. Boyd and L. Vandenberghe. *Convex Optimization*. Cambridge University Press, 2004. (ch 4.1)
- S. Boyd and L. Vandenberghe. *Introduction to Applied Linear Algebra: Vectors, Matrices, and Least Squares*. Cambridge University Press, 2018. (ch 19.4, car control example)