

CSE221 ALGORITHMS

Topic: Graph (Directed,
Undirected, Unweighted),
Adjacency Matrix, Adjacency
List, Tree

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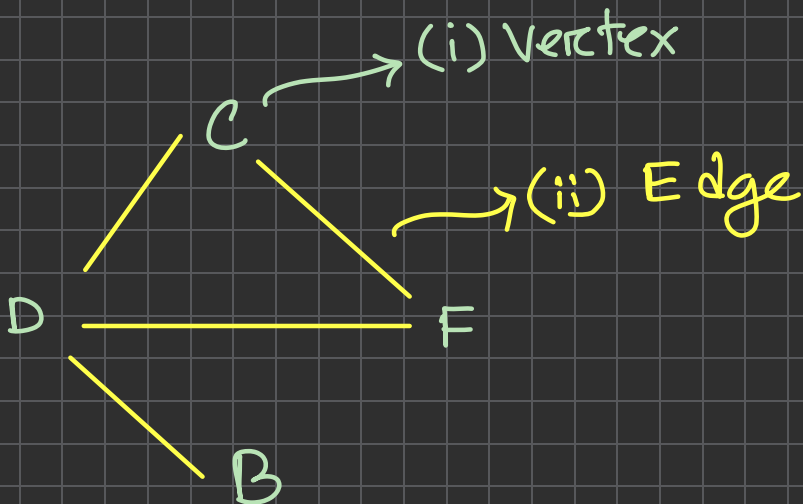
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Topic: Graph

Graph

Basic two parts of a graph data structure:



Imagine 'vertex' as cities and
'Edge' as roads connecting
them or linking them.

Types of Graph

Primarily two types:

(i) Undirected

(ii) Directed

Graph can also be categorized as

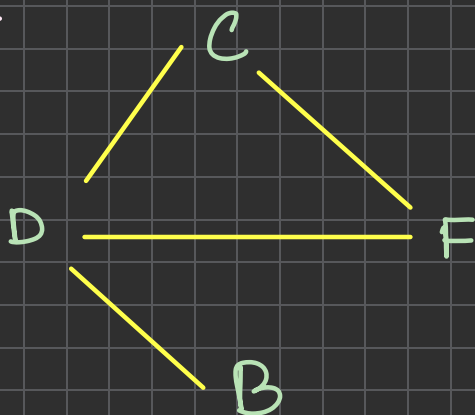
(a) weighted

(b) unweighted

Undirected Graph

No particular direction is given from one vertex (city) to another vertex. Both direction between possible.

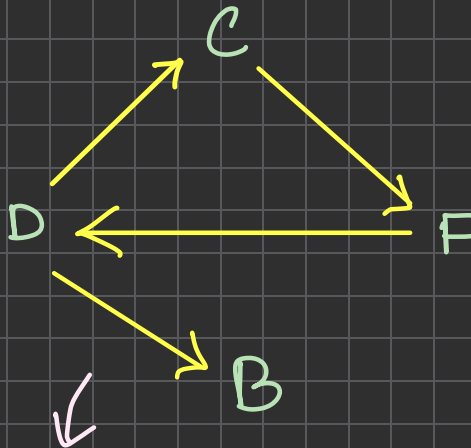
उत्तर:



Directed Graph

→ Directions from one vertex to adjacent vertices are given. And it allows only one way traverse (एकतरफ़ा).

उदाहरण:



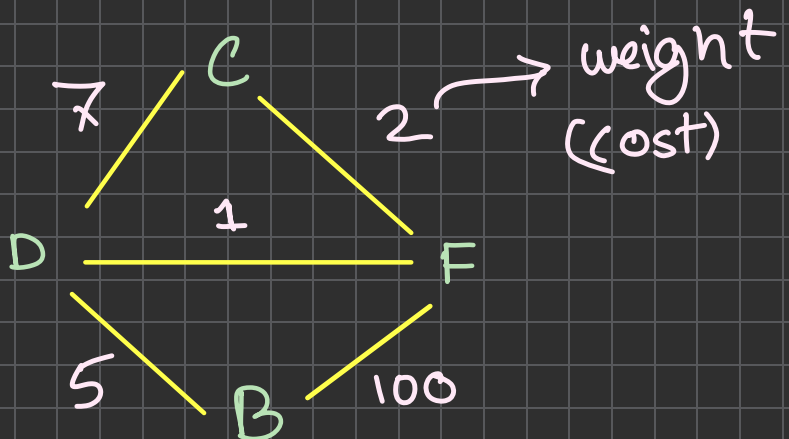
it can also be called unweighted graph

NB: → Both way direction ટફડી
જાત ના.

→ Direction ટફડી ના ટાકમે means
both way બાતબાત

Unweighted Graph

→ Weight ବଳା ନାହିଁ, e.g. ଏକ city (vertex) ଅନ୍ୟ city (vertex) ଯିବା ପାଇଁ cost (weight) କିଛି ନାହିଁ।



→ either મહત્તમ weight આ
થાકતો or લઘુત્તમ weight આ
થાકતો ના.

Important properties of Graph:

Properties of Undirected Graph:

(i) Degree: number of edges connected to the vertex.

उदा:

Vertex

Degree

B

1

C

2

D

3

F

2

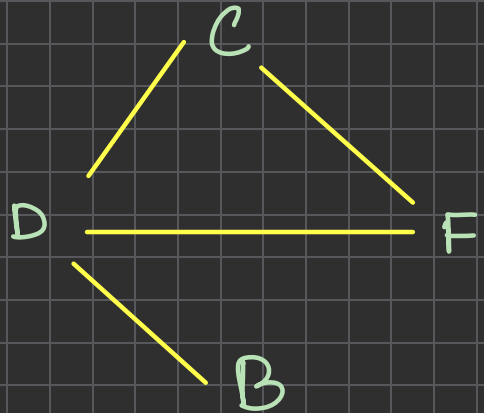
sum of
degrees = 8



notice something?

Degree sum of all vertices,

= $2 \times$ number of edges



reason: each edge connects
exactly two vertices. that's
why increase two times

Properties of

Directed Graph:

(i) Indegree: number of
incoming edges

(ii) Outdegree: number of
outgoing edges

| <u>vertex</u> | <u>Indegree</u> | <u>Outdegree</u> |
|---------------|-----------------|------------------|
|---------------|-----------------|------------------|

B

1

1

C

2

0

D

1

2

F

1

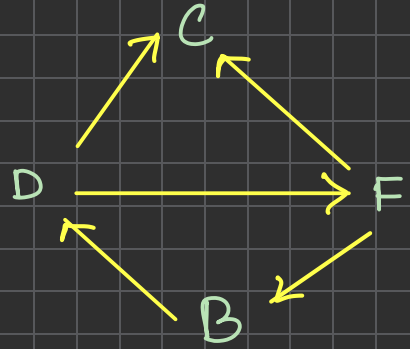
2

sum = 5

sum = 5



notice something?



$$\sum (\text{in degree}) = \sum (\text{out degree})$$

sum of indegree = sum of outdegree

Representation of Graph

Two methods to represent graph:

- Adjacency Matrix
- Adjacency List

Adjacency Matrix

↳ Table / 2d array

Vertex $n \in \mathbb{Z}^n$,

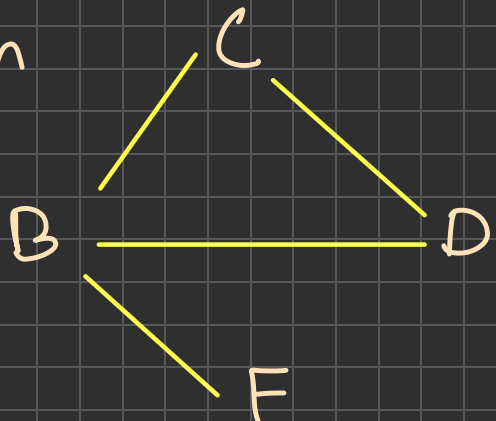
matrix size = $n \times n$

→ start with a $n \times n$ array where

column title and

row title is each

vertex one by one



→ then, if there is an edge between a vertex 'x' (row) and vertex 'y' (column) then insert '1' in the mutual box of vertex 'x' and 'y'

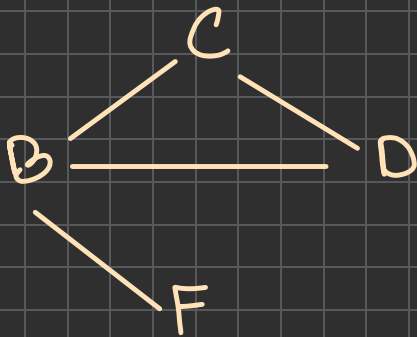
And if there is no edge between the two vertices, then insert '0' (zero) in the box.

0 → no edge

1 → edge exists

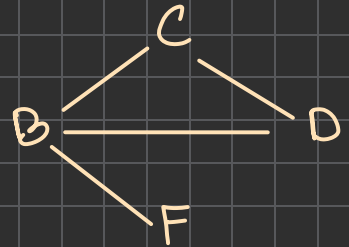
And you have a matrix
representation of Graph

For example




This graph can be represented
as follows,

| | B | C | D | F |
|---|---|---|---|---|
| B | 0 | 1 | 1 | 1 |
| C | 1 | 0 | 1 | 0 |
| D | 1 | 1 | 0 | 0 |
| F | 1 | 0 | 0 | 0 |



Reason for $B \times B$ being zero is we can't traverse from B to B in the given example.

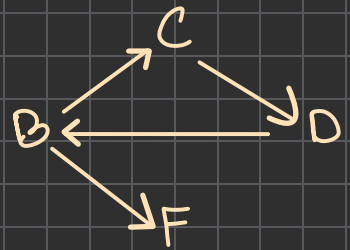
And if it was like  then $B \times B$ would be 1.

Directed graphs can also be represented in matrix. But the direction will be:

row \rightarrow column
always

For example,

| | B | C | D | F |
|---|---|---|---|---|
| B | 0 | 1 | 0 | 1 |
| C | 0 | 0 | 1 | 0 |
| D | 1 | 0 | 0 | 0 |
| F | 0 | 0 | 0 | 0 |



Adjacency List

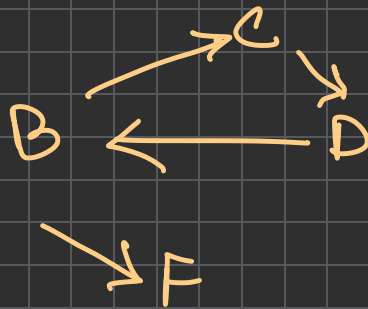
↘ list of lists

→ start with a list with all individual vertices

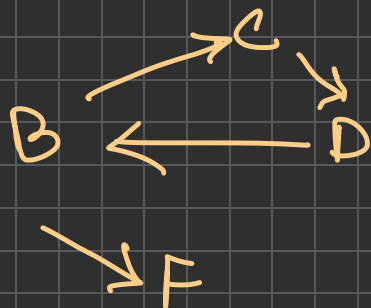
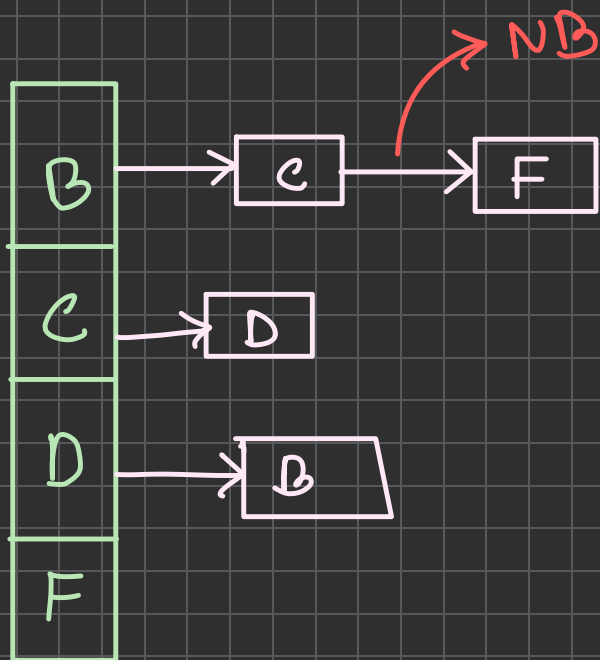
→ then the vertices in that list will be followed by all the vertices that have edges with that vertex from the main list and the

following vertices (with connecting edges) will be saved as list.

For example,



can be represented in Adjacency List as follows,



NB: There $c \rightarrow F$ doesn't mean we can go from c to F or that C and F have mutual edge. Because in the graph we see they don't. What that

means is we can go from

B to both C and F directly.

Meaning C and F both have
edges from B.

Tree

Tree is a graph which maintains 3 conditions.

Conditions:

(i) Connected \Rightarrow \Rightarrow :

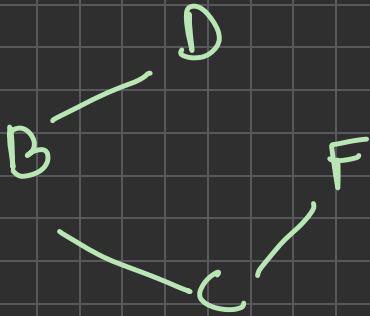
সেখানে

একটা vertex থেকে অন্য vertex-এ

somehow অন্য একটা vertex-এ

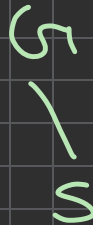
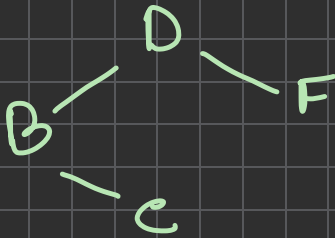
গিয়ে without lifting the pen.

उदाहरण:



→ it's a tree

but



} not a tree

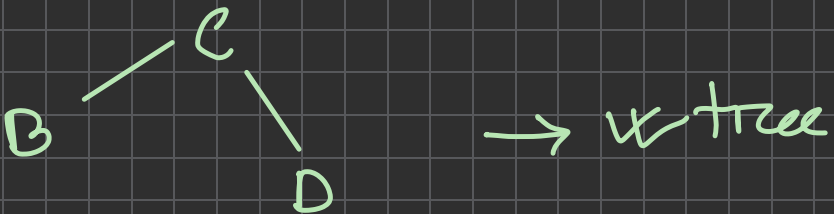
no connection

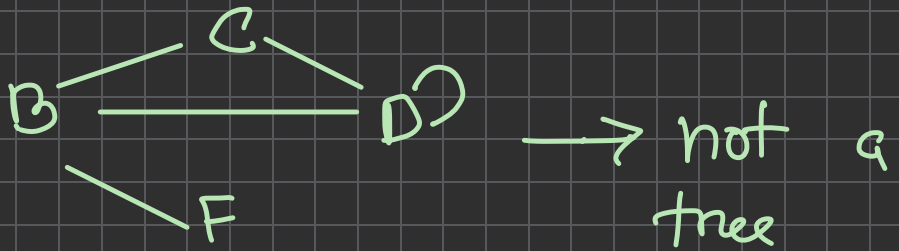
reason

we can't go from D to S without lifting the pen.

That's still a graph but not a tree.
~~~~~  
disconnected graph

condition (ii) No cycle whatsoever!





condition (iii):

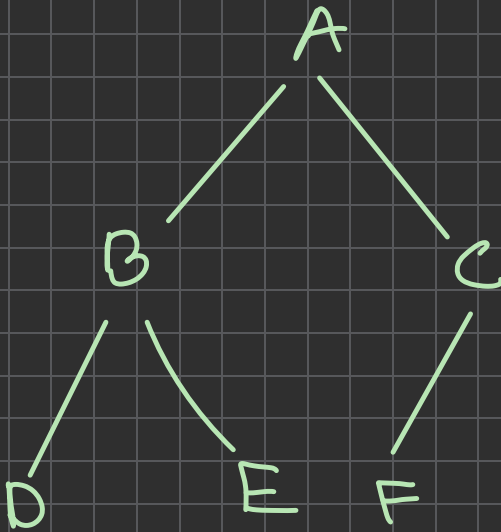
Graph  $\hookrightarrow$

exactly  $(n-1)$  edges

where

$n$  = number of vertices.

תוצא:



$$n=6$$

$$\text{edges} = 5$$