



# **CS-2001**

## **DATA STRUCTURE**

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**National University of Computer  
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# BINARY TREE

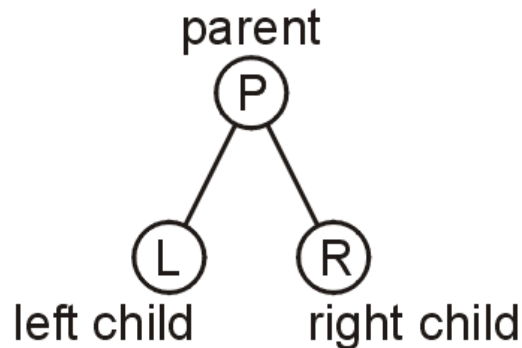


# Binary Trees

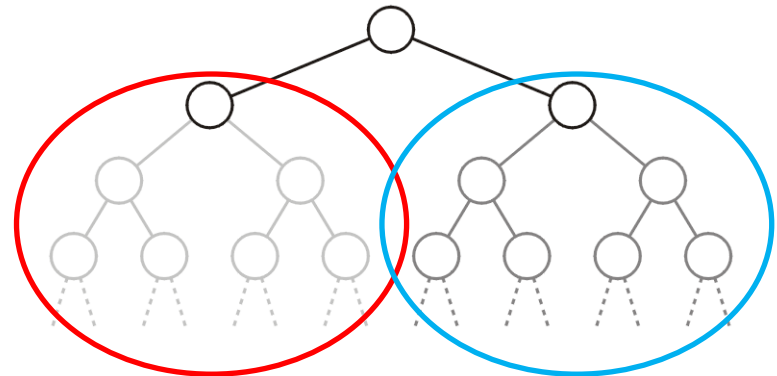
3

## Binary Tree

- In a binary tree, **each node has at most two children**
  - ▣ Allows to label the children as left and right



- Likewise, the two sub-trees are referred as
  - ▣ **Left sub-tree**
  - ▣ **Right sub-tree**



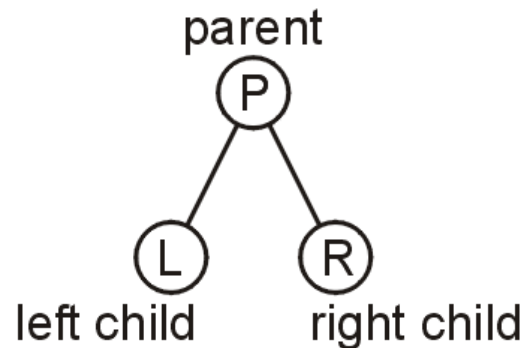
# Binary Trees

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## Binary Tree

The mathematical definition of a binary tree is

“A binary tree is a finite set of elements that is either empty or is partitioned into three disjoint subsets. The first subset contains a single element called the root of the tree. The other two subsets are themselves binary trees called the left and right sub-trees”. Each element of a binary tree is called a node of the tree. Following figure shows a binary tree.

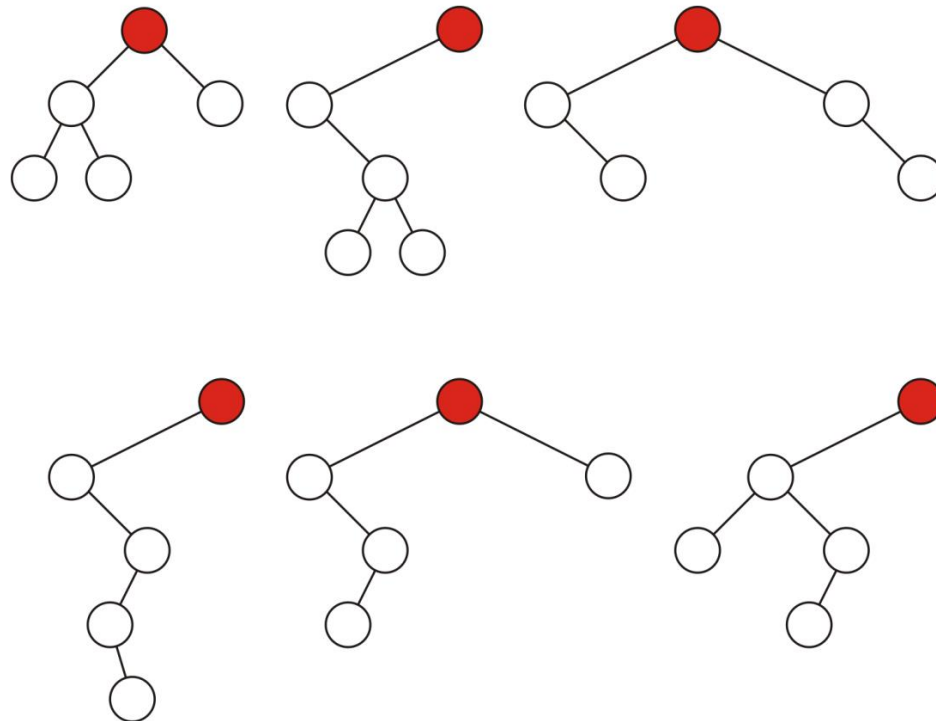


# Binary Trees

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## Binary Tree

Some variations on binary trees with five nodes

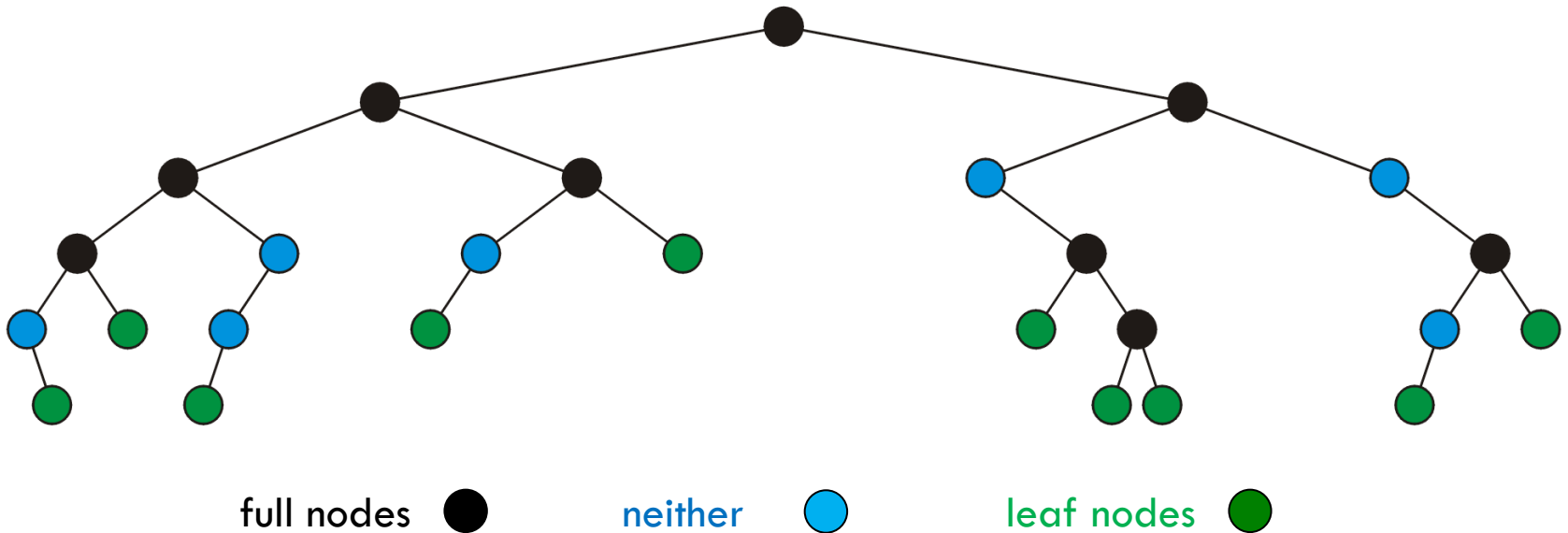


# Binary Trees

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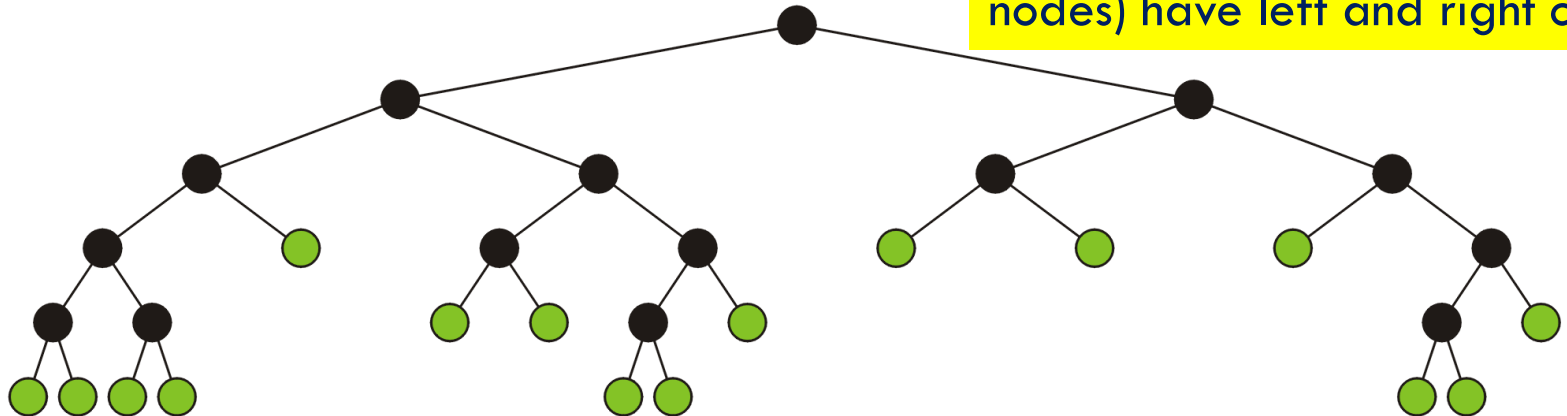
## Full Node

A **full node** is a node where both the left and right sub-trees are non-empty trees



## 7

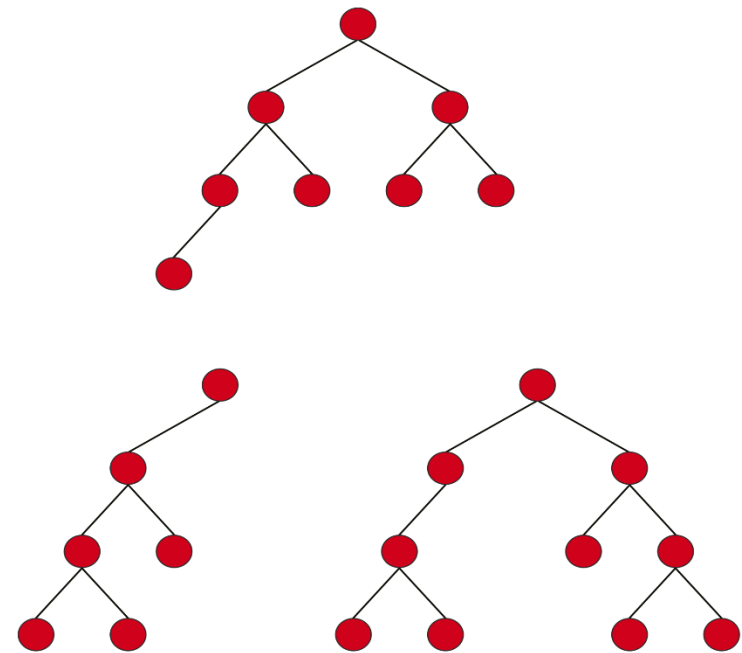
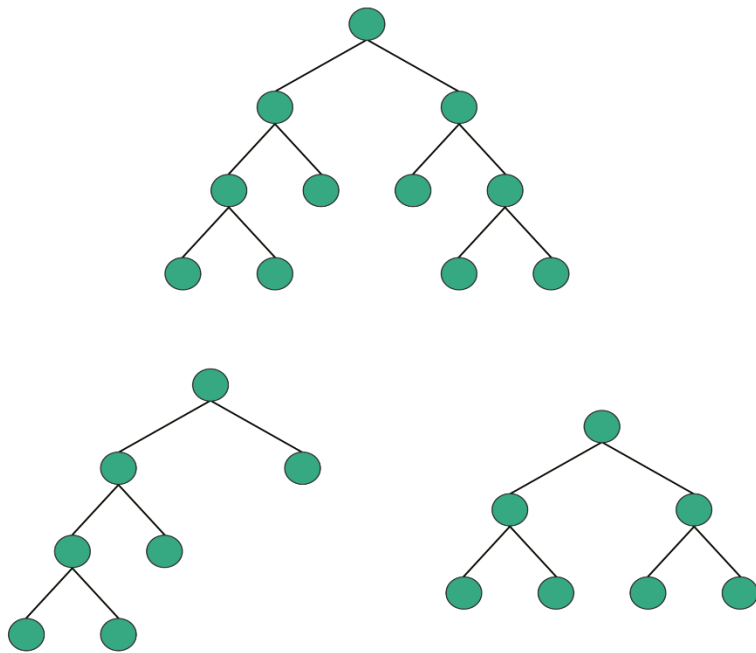
All the nodes (other than leaf nodes) have left and right child



# Full Binary Trees

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## Full Binary Tree



## Valid and Invalid Structure of Full Binary Tree

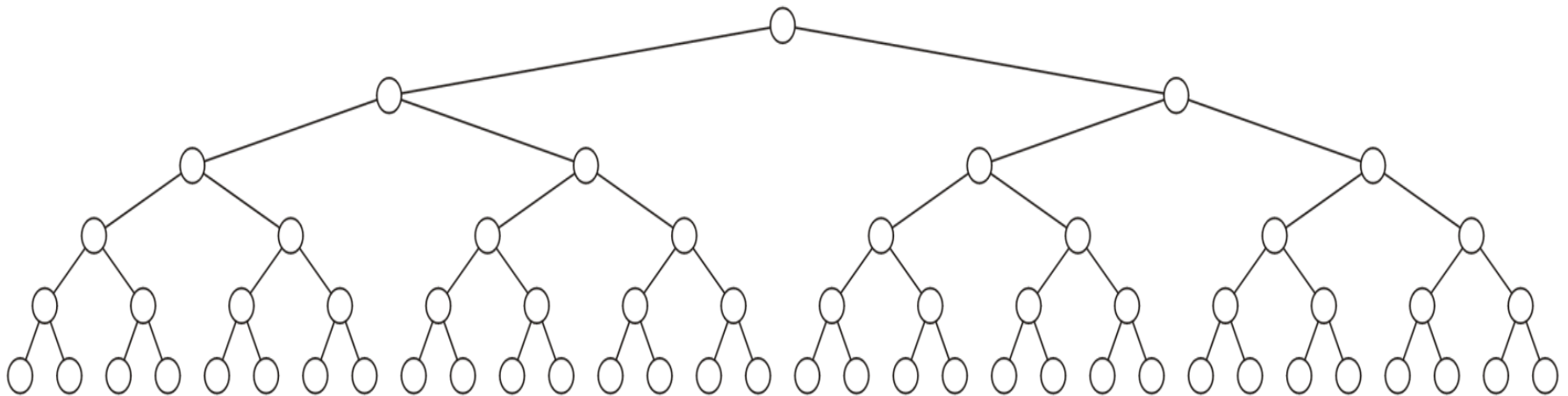


# Complete Binary Trees

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## Complete (Perfect) Binary Tree

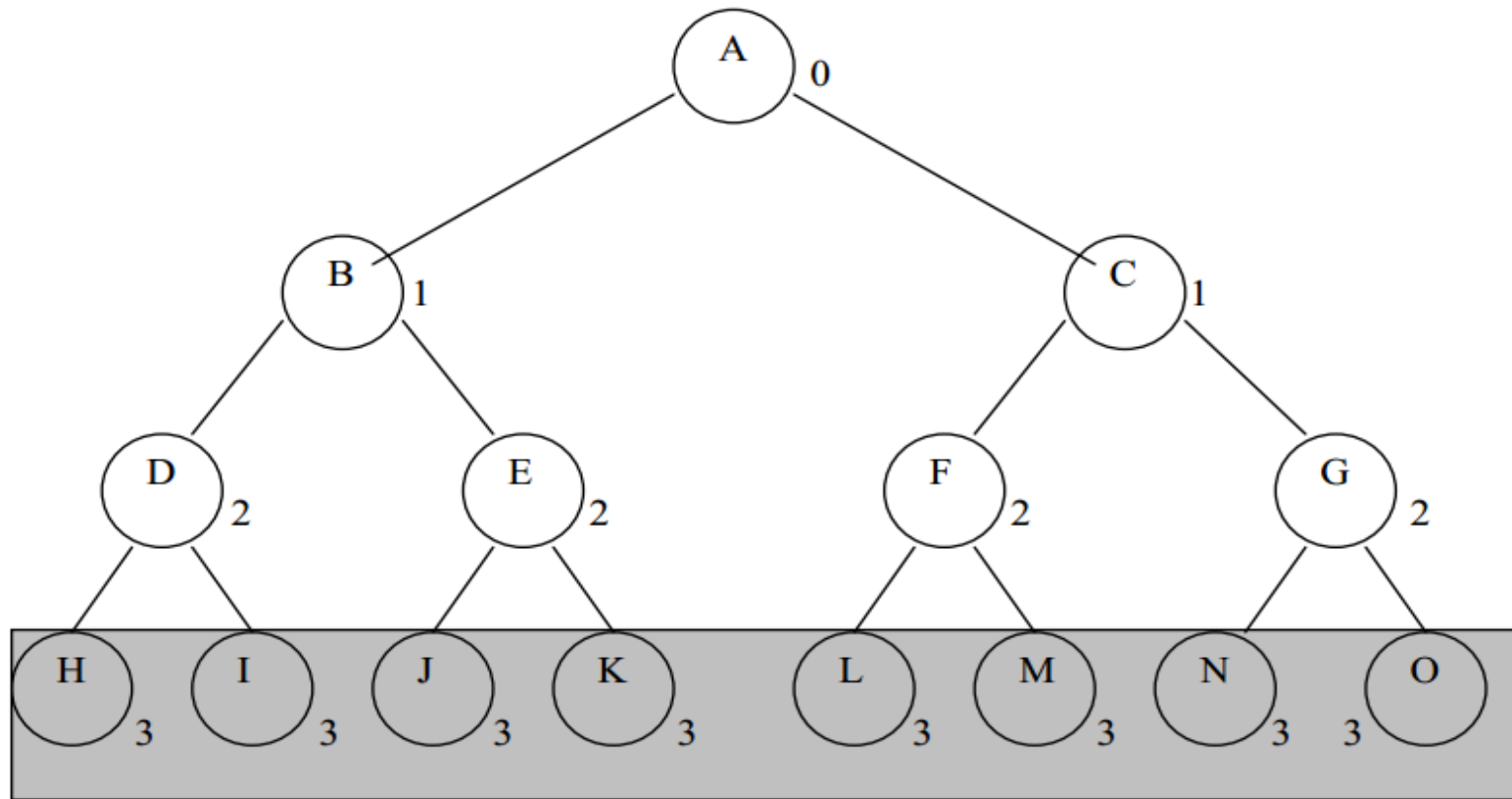
- A complete binary tree of height  $h$  is a binary tree, where,
  - All leaf nodes have the same depth  $h$
  - All other nodes are full nodes



# Complete Binary Trees

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## Complete (Perfect) Binary Tree ... Examples



# Complete Binary Trees

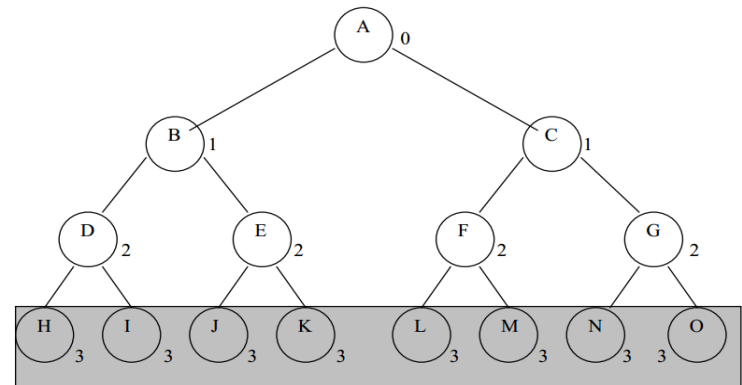
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## Complete (Perfect) Binary Tree ... Examples

The definition of the complete binary tree is

“A complete binary tree of depth  $d$  is the strictly binary tree all of whose leaves are at level  $d$ ”.

Now look at the tree, the leaf nodes of the tree are at level 3 and are H, I, J, K, L, M, N and O. There is no such a leaf node that is at some level other than the depth level  $d$  i.e. 3. All the leaf nodes of this tree are at level 3 (which is the depth of the tree i.e.  $d$ ). So this is a complete binary tree. In the figure, all the nodes at level 3 are highlighted.



# Complete Binary Trees

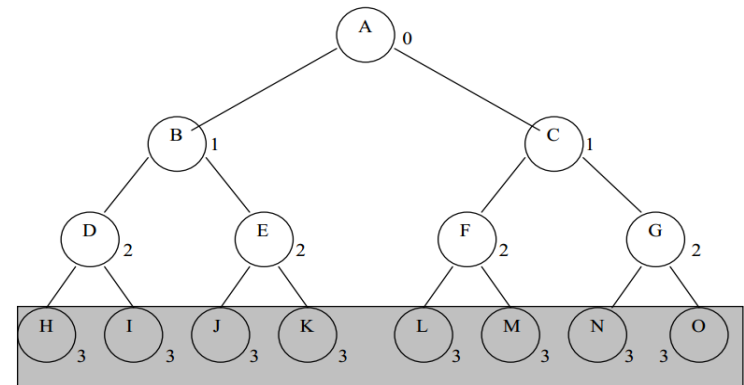
12

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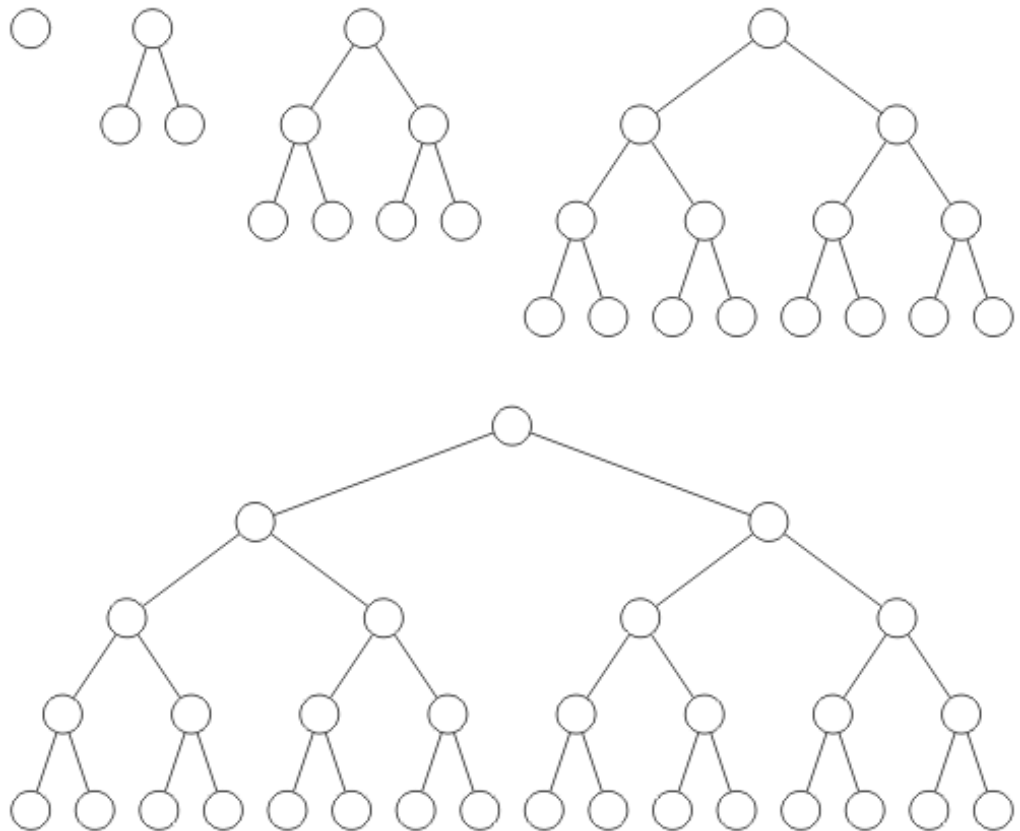
We can say every Strictly Binary Tree may or may not be a complete binary tree but every complete or perfect binary tree will always Strictly as well

# Complete Binary Trees

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## Complete (Perfect) Binary Tree ... Examples

Complete binary trees  
of height  $h = 0, 1, 2, 3$   
and 4

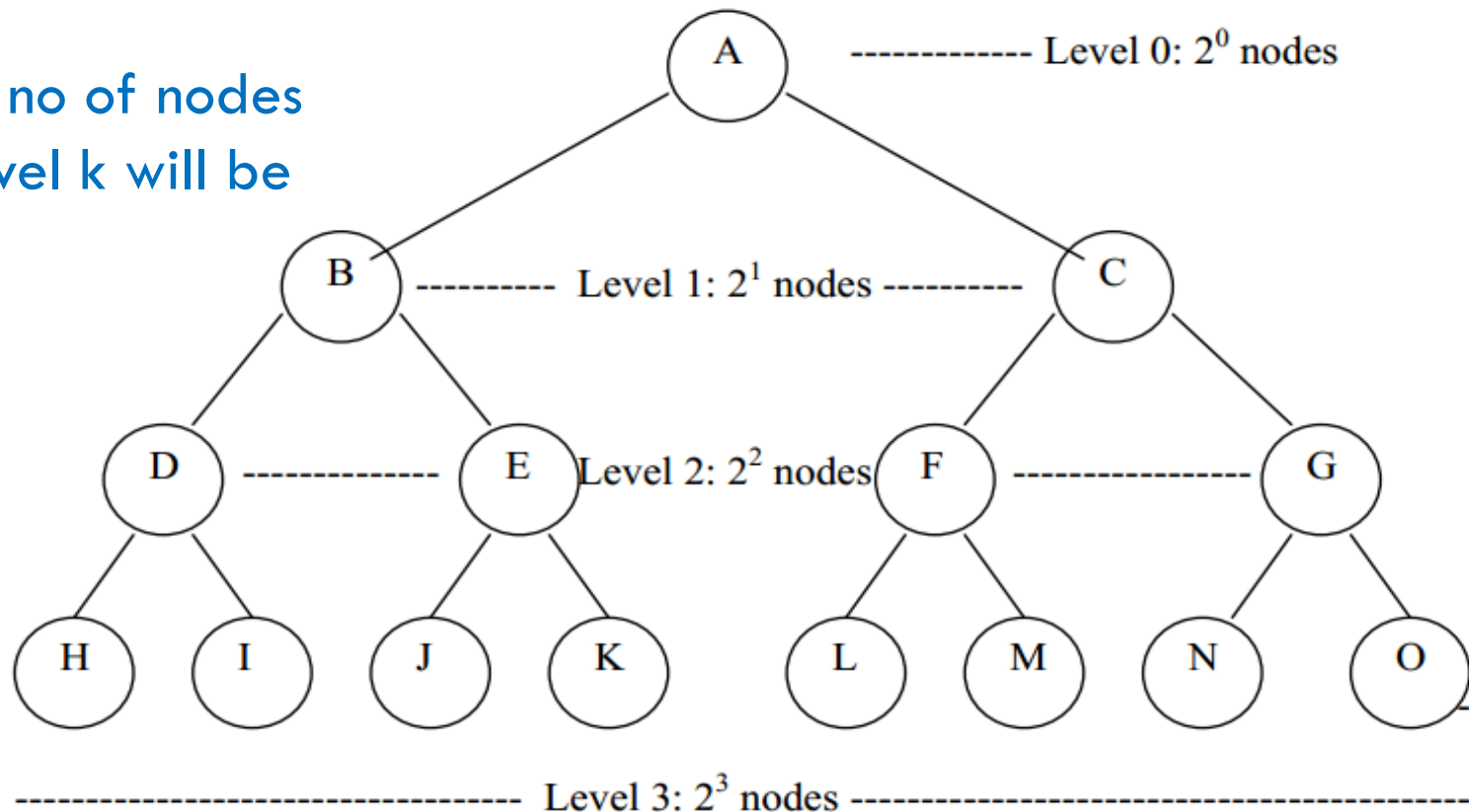


# Complete Binary Trees

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## Complete (Perfect) Binary Tree ... Examples

Max no of nodes  
at level k will be  
 $2^k$

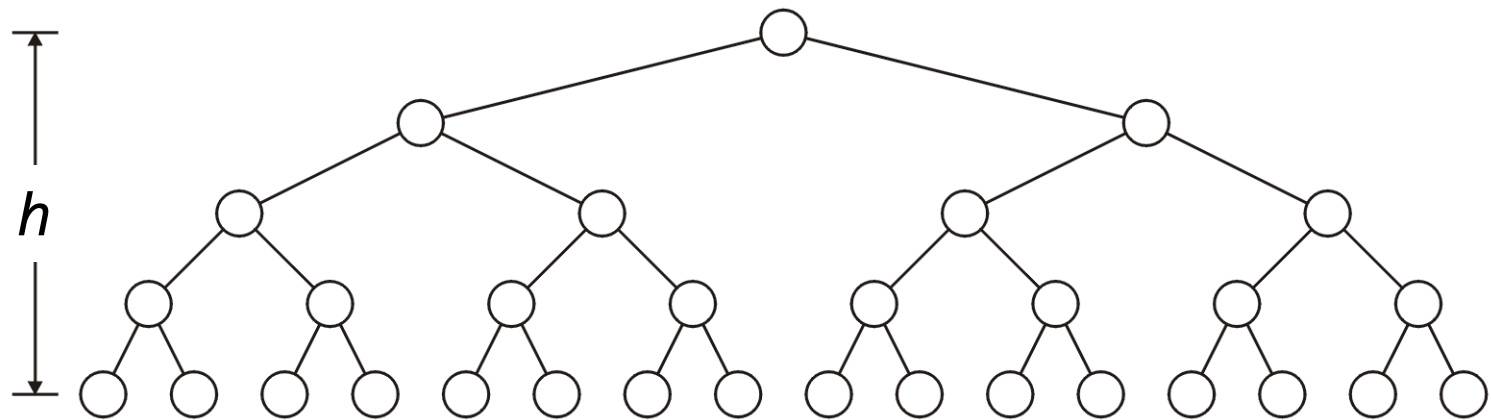


# Complete Binary Trees

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## Complete (Perfect) Binary Tree ... Properties

- A complete binary tree with height  $h$  has  $2^h$  leaf nodes



# Complete Binary Trees

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## Complete (Perfect) Binary Tree ... Properties

Depth d	# nodes at depth d	# of child nodes
0	1 = $2^0$	2 (each node has 2 children)
1	2 = $2^1$	4 (each node has 2 children)
2	4 = $2^2$	8 (each node has 2 children)
...		

- The **number** of nodes **doubles** every time the **depth** increases by 1 !

$$\# \text{ nodes at depth } d = 2^d$$

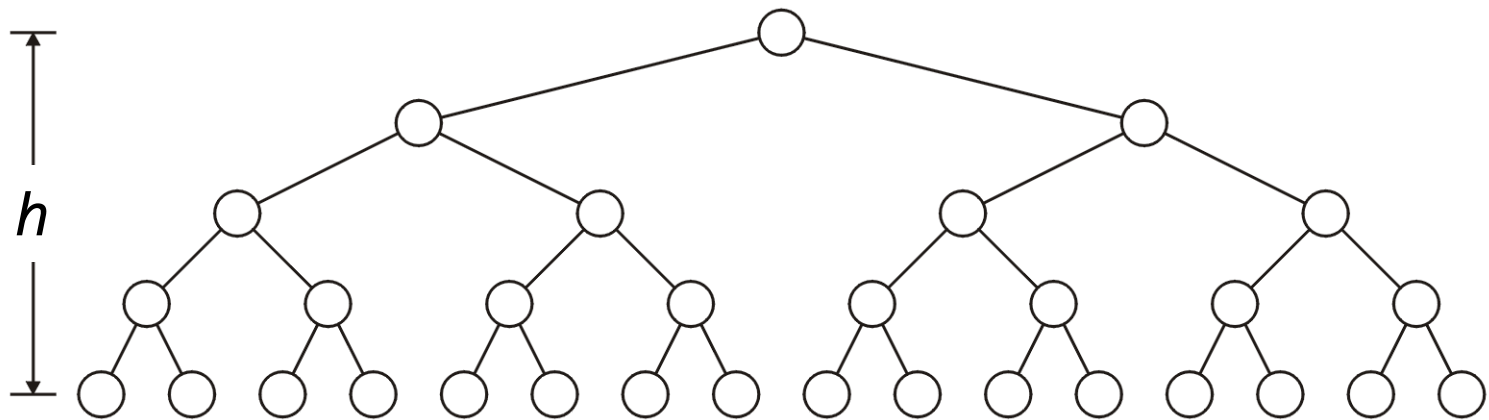


# Complete Binary Trees

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## Complete (Perfect) Binary Tree ... Properties

- A complete binary tree with height  $h$  has  $2^h$  leaf nodes
- A complete binary tree of height  $h$  has  $2^{h+1} - 1$  nodes

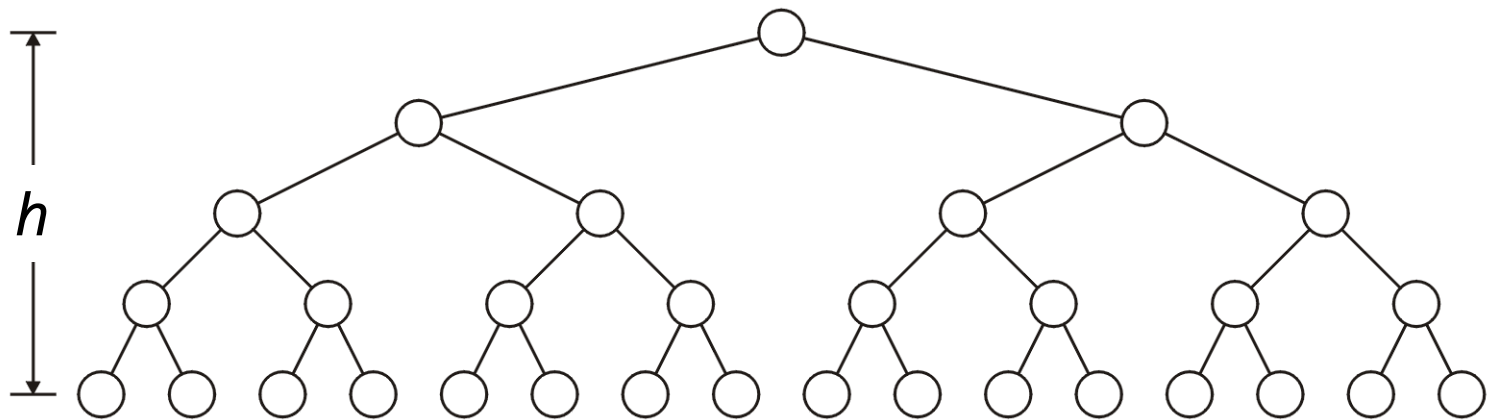


# Complete Binary Trees

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## Complete (Perfect) Binary Tree ... Properties

- A complete binary tree with height  $h$  has  $2^h$  leaf nodes
- A complete binary tree of height  $h$  has  $2^{h+1} - 1$  nodes
  - Number of leaf nodes:  $L = 2^h$
  - Number of internal nodes:  $2^h - 1$
  - Total number of nodes:  $2^{h+1} - 1$



# Complete Binary Trees

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## Complete (Perfect) Binary Tree ... Properties

- A complete binary tree with height  $h$  has  $2^h$  leaf nodes
- A complete binary tree of height  $h$  has  $2^{h+1} - 1$  nodes
  - Number of leaf nodes:  $L = 2^h$
  - Number of internal nodes:  $2^h - 1$
  - Total number of nodes:  $2^{h+1} - 1$
- A complete binary tree with  $n$  nodes has height  $\log_2(n + 1) - 1$

$$n = 2^{h+1} - 1$$

$$2^{h+1} = n + 1$$

$$h + 1 = \log_2(n + 1)$$

$$\Rightarrow h = \log_2(n + 1) - 1$$

# Almost Complete Binary Trees

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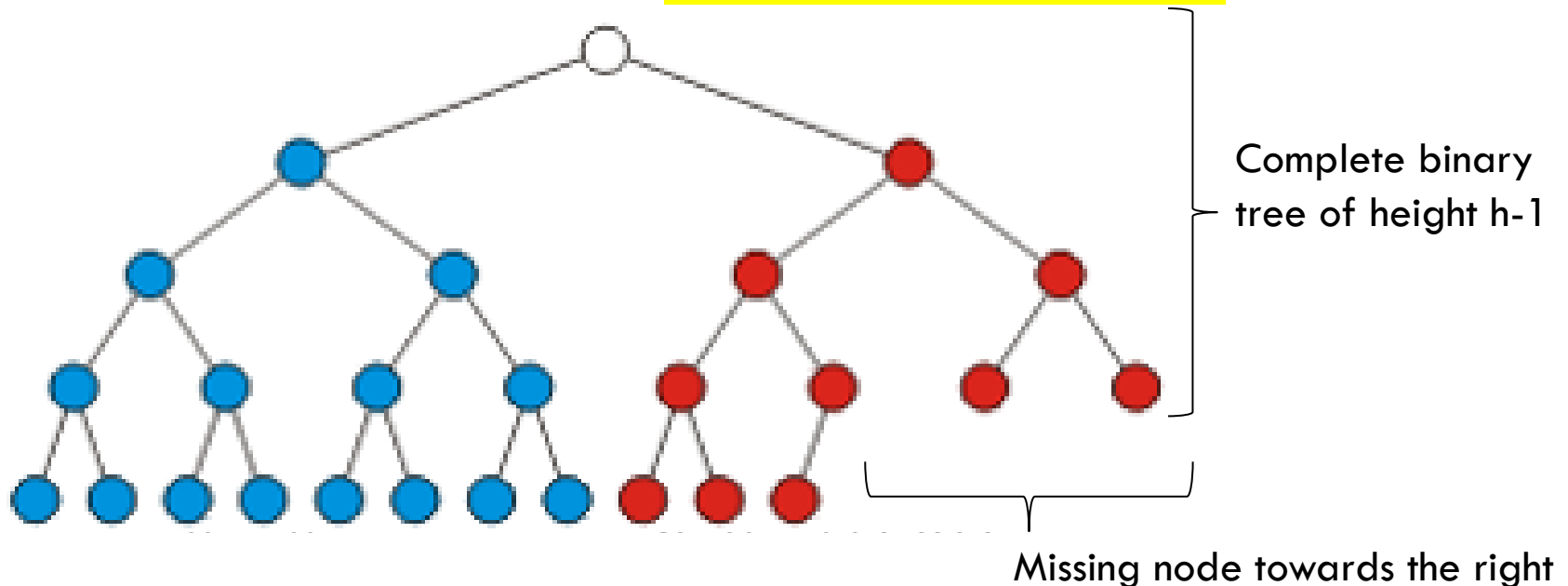
## Almost Complete Binary Tree

Almost complete binary tree of height  $h$  is a binary tree in which

1. There are  $2^d$  nodes at depth  $d$  for  $d = 1, 2, \dots, h-1$ .

Each leaf in the tree is either at level  $h$  or at level  $h-1$ .

2. The nodes at depth  $h$  are as far left as possible



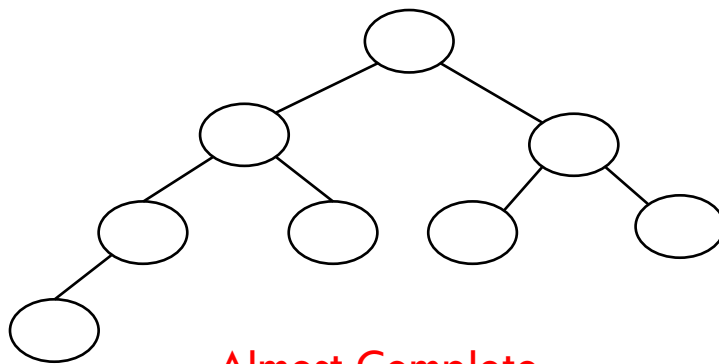
# Almost Complete Binary Trees

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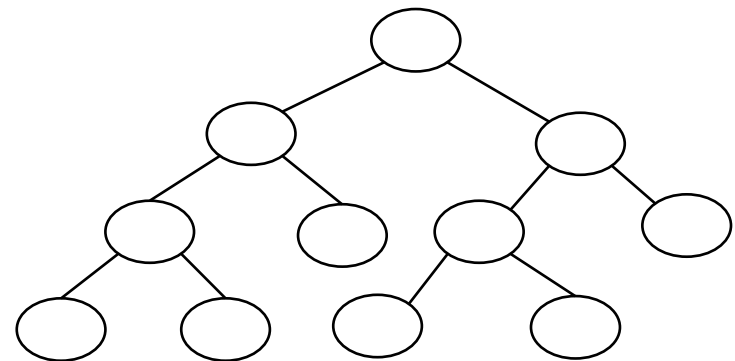
## Almost Complete Binary Tree

*The nodes at depth  $h$  are as far left as possible*

- If a node  $p$  at depth  $h-1$  has a left child
  - Every node at depth  $h-1$  to the left of  $p$  has 2 children
- If a node at depth  $h-1$  has a right child
  - It also has a left child



Almost Complete  
binary tree



Not Almost Complete binary tree  
(condition 2 violated)

# Almost Complete Binary Trees

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## Almost Complete Binary Tree ... Properties

Total number of nodes  $n$  is between

A complete binary tree of height  $h-1$ , i.e.,  $2^h$  nodes

A complete binary tree of height  $h$ , i.e.,  $2^{h+1} - 1$  nodes

Height  $h$  is the largest integer less than or equal to  $\log_2(n)$

# Balanced Binary Trees

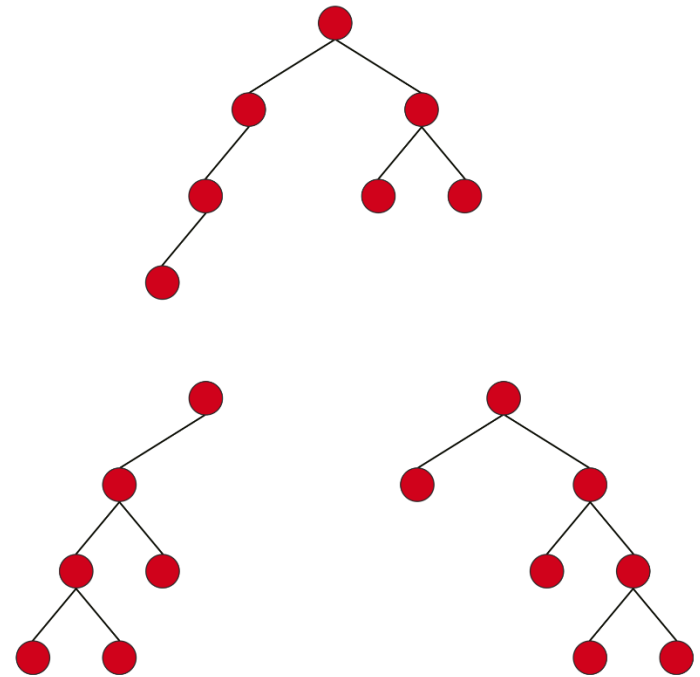
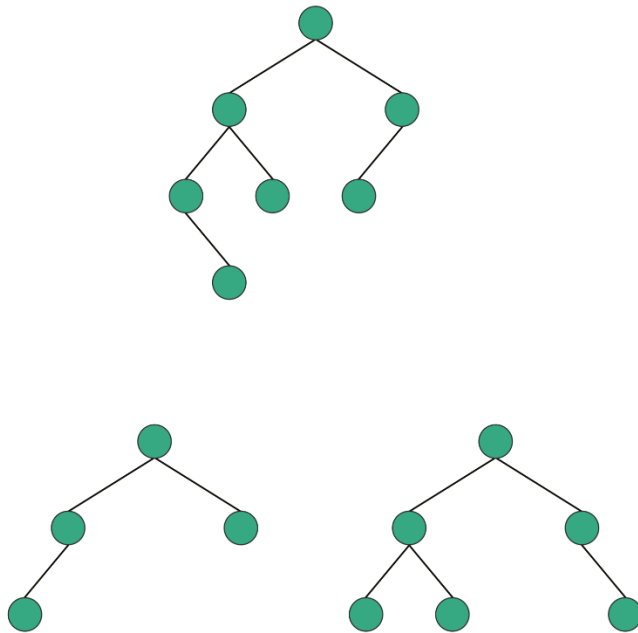
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## Balanced Binary tree

- is a Binary tree in which height of the left and the right sub-trees of every node may differ by at most 1

# Balanced Binary Trees

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**Valid and Invalid Structure of Balanced Binary Tree**



# BINARY TREE IMPLEMENTATION

# Binary Tree ... Implementation

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- Binary Tree ... Implementation
  - ❖ Array based implementation
  - ❖ Linked List based implementation

# Binary Tree ... Implementation

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- Objects: any type of objects can be stored in a tree data structure
- Accessor methods
  - ▣ `root()` – return the root of the tree
  - ▣ `parent(p)` – return the parent of a node
  - ▣ `children(p)` – returns the children of a node

# Binary Tree ... Implementation

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## □ query methods

- `size()` – returns the number of nodes in the tree
- `isEmpty()` - returns true if the tree is empty
- `elements()` – returns all elements
- `isRoot(p)`, `isInternal(p)`, `isExternal(p)`

## □ other methods

- Tree traversal, Node addition/deletion, create/destroy

# Binary Tree ... Implementation

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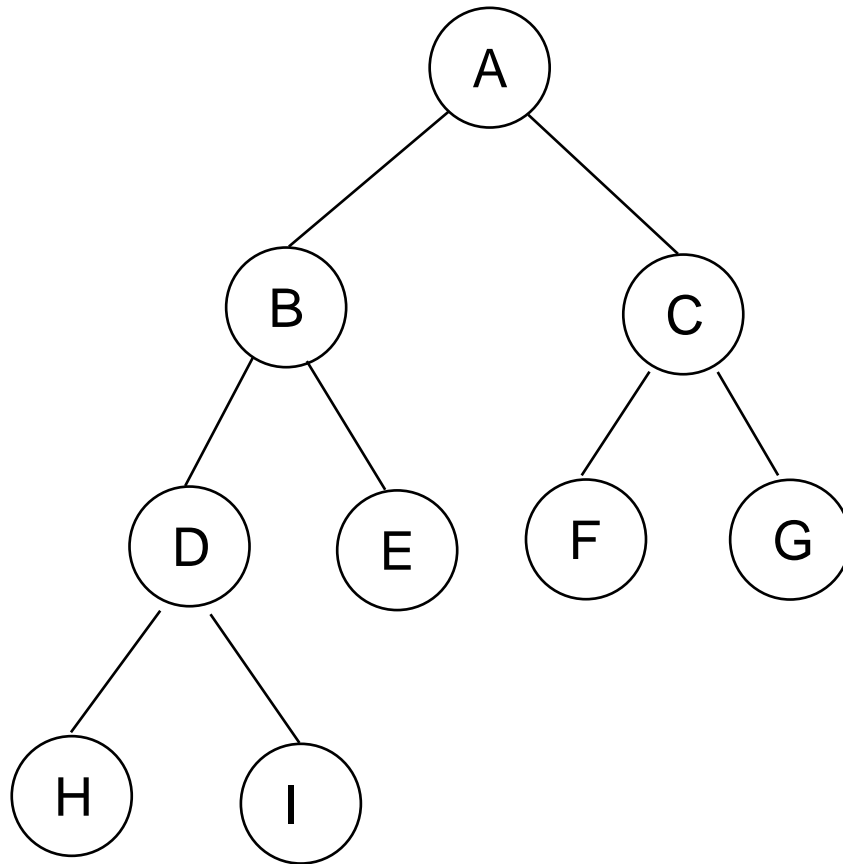
## Array-based Implementation:

- Value in root node stored first, followed by left child, then right child
- Each successive level in the tree stored left to right; unused nodes in tree represented by a bit pattern to indicate nothing stored there
- Children of any given node  $n$  is stored in cells  $2n$  and  $2n + 1$  (If array index starts at 1)
- *Storage allocated as for full tree, even if many nodes empty*
- For a tree of depth  $h$  we need array of  $2^{h+1}-1$  cells

# Binary Tree ... Implementation

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## Array-based Implementation:

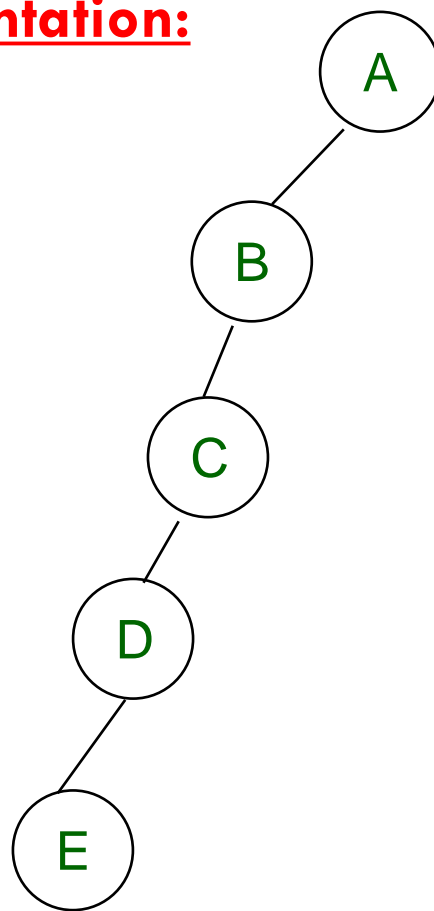
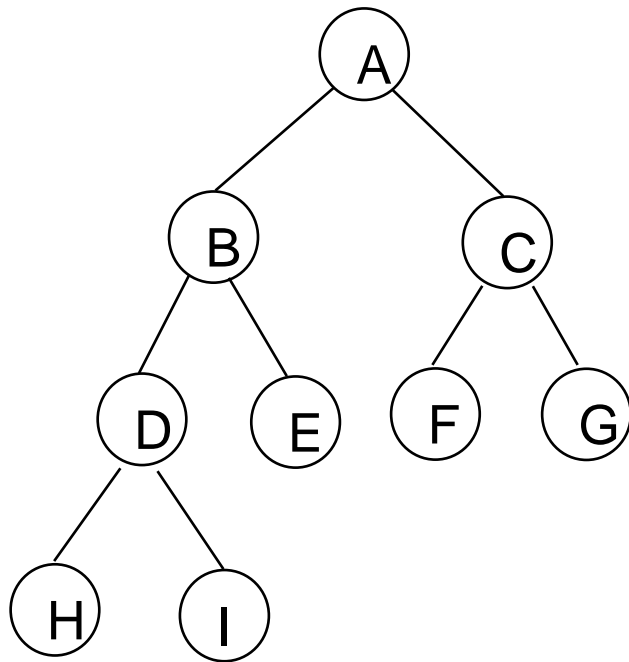


[1]	A
[2]	B
[3]	C
[4]	D
[5]	E
[6]	F
[7]	G
[8]	H
[9]	I

# Binary Tree ... Implementation

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## Array-based Implementation:



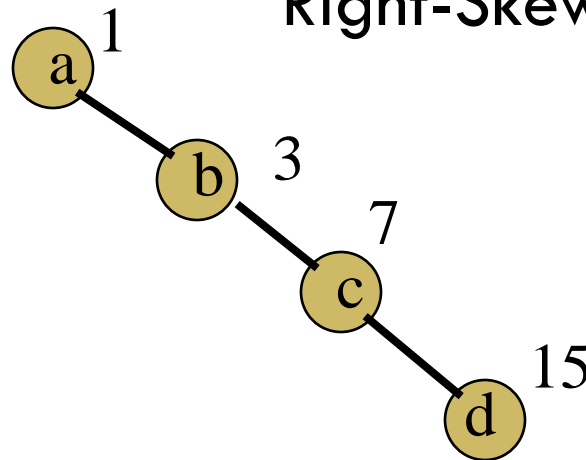
[1]	A
[2]	B
[3]	--
[4]	C
[5]	--
[6]	--
[7]	--
[8]	D
[9]	--
.	.
[16]	E

# Binary Tree ... Implementation

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## Array-based Implementation:

Right-Skewed Binary Tree



tree[]

a	-	b	-	-	-	c	-	-	-	-	-	-	-	d
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---



# Binary Tree ... Implementation

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## Linked List-based Implementation:

- Implementation of Binary tree can be done using dynamic creation of node having data and two pointers.

```
Struct Node {  
    int data;  
    Node* left;  
    Node * right;  
}
```

# Binary Tree ... Implementation

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- A binary tree is composed of **zero or more nodes**
- Each node contains:
  - ▣ A **value** (some sort of data item)
  - ▣ A **reference or pointer** to a left child (may be **null**),  
and
  - ▣ A **reference or pointer** to a right child (may be **null**)
- A binary tree may be *empty* (contain no nodes)

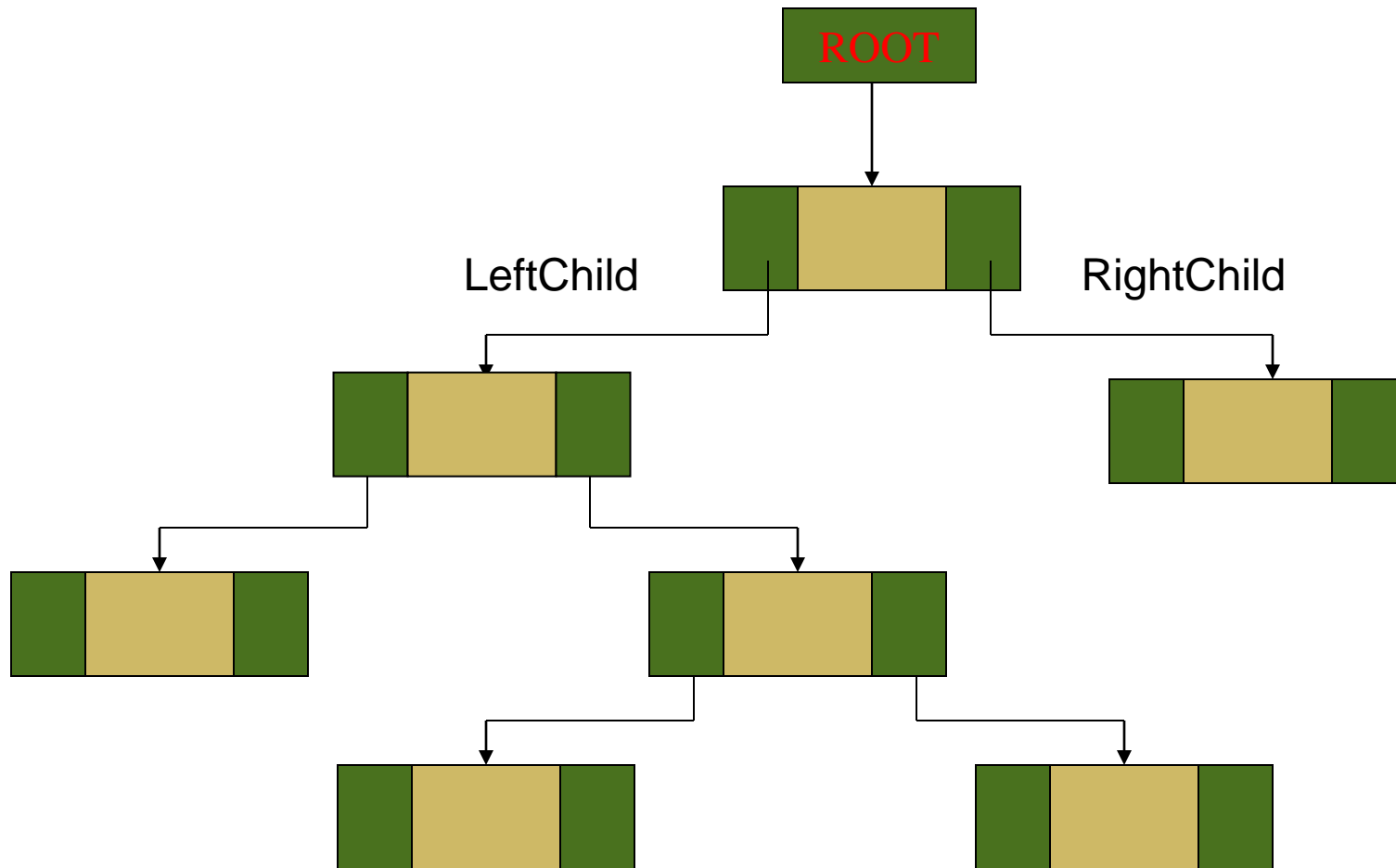
# Binary Tree ... Implementation

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- If not empty, a binary tree has a **root node**
  - ▣ Every node in the binary tree is reachable from the root node by a **unique** path
- A node with neither a left child nor a right child is called a **leaf**
  - ▣ In some binary trees, only the leaves contain a value

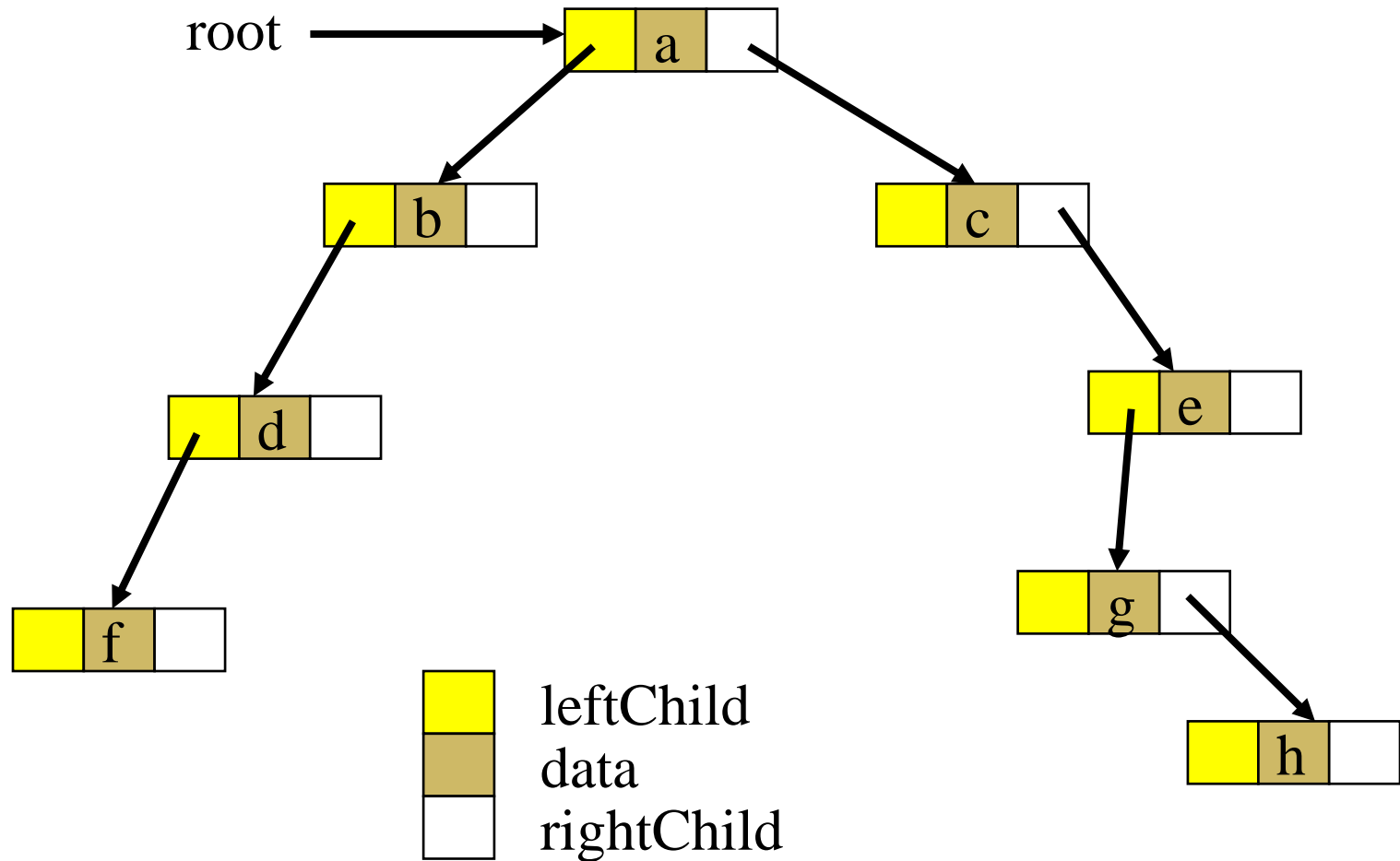
# Binary Tree ... Implementation

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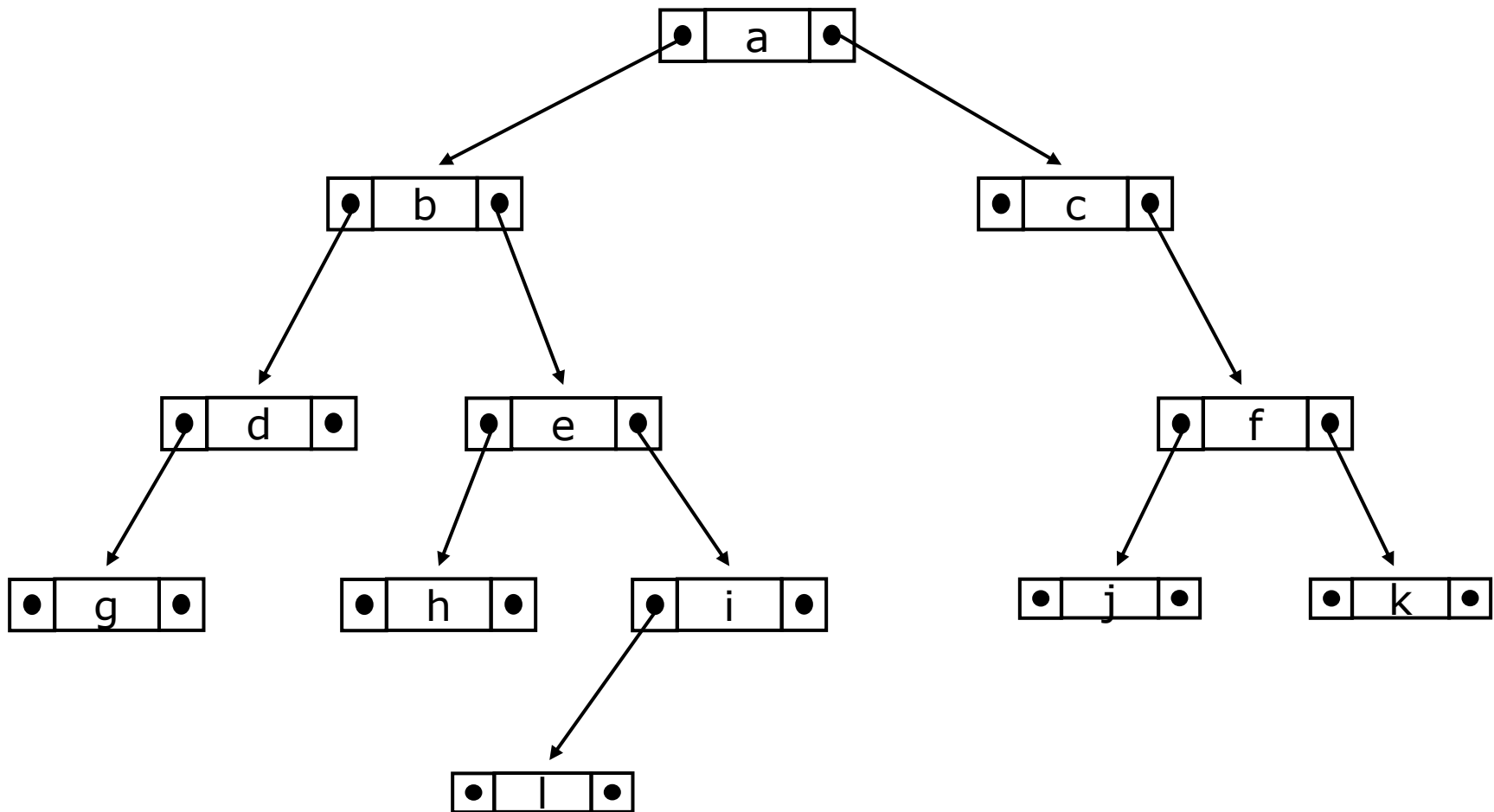
# Binary Tree ... Implementation

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# Binary Tree ... Implementation

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# Binary Tree ... Implementation

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## Implementation: with Linked List

- **Each node** in the tree consists of:
  - ▣ The **data**, or value contained in the element
  - ▣ A **left child pointer** (pointer to first child)
  - ▣ A **right child pointer** (pointer to second child)
- A **root pointer** points to the root node
  - ▣ Follow pointers to find every other element in the tree
- *Add and remove nodes by manipulating pointers*
- **Leaf nodes** have child pointers set to null

# Binary Tree ... Implementation

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- **Left(node):** Gives index/pointer of left child
- **Right(node):** Gives index/pointer of right child
- **Parent(node):** Returns index/pointer of parent
- **Brother(node):** Returns index/pointer of brother
- **Root:** Gives index/pointer of root node
- **Info(Node):** Data/Info stored at node
- **IsLeft(node):** Is node left child? Yes/No
- **IsRight(node):** Is node right child? Yes/No



# Reading Materials

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- Nell Dale Chapter#8
- Schaum's Outlines Chapter#7
- D. S. Malik Chapter#11

