Stack Operations

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Book Chapter

- "Assembly Language for x86 Processors"
- Author "Kip R. Irvine"
- 6th Edition
- Chapter 5
 - Section 5.4

Stack

- A LIFO (Last In First Out) data structure
- New value is added to the top of stack
- Existing values are removed from the top of stack
- An essential part of calling from and returning to the procedures
- Real life example
 - A stack of plates

Runtime Stack (1/3)

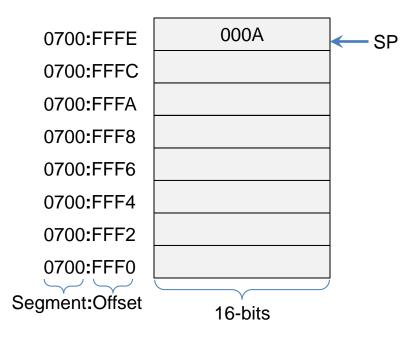
- A memory array managed by CPU using ESP/SP (Extended Stack Pointer) register and SS (Stack Segment)
- ESP/SP always points to the last value pushed on the top of stack and holds offset of that value in Stack Segment
- ESP/SP cannot be manipulated directly instead it can be modified indirectly by instructions such as PUSH, POP, CALL, RET

Runtime Stack (2/3)

- In protected mode i.e. 32-bit mode, size of each stack location is 32-bits
- In real-address mode i.e. 16-bit mode, size of each stack location is 16-bits
- emu8086 uses real-address mode
- Runtime Stack is different from Stack Abstract
 Data Type which is typically written in a HLL

Runtime Stack (3/3)

- SS contains the base address of Stack Segment
- SP contains the offset of value at the top of stack



Push Operation (1/3)

- A push operation in stack puts the value on the top location available in stack and decrements the stack pointer by size of stack element
- Size of each stack element is 32 bits in protected address mode
- Size of each stack element is 16 bits in realaddress mode

Push Operation (2/3)

- SP is decremented by 2 with each push operation
- These values are pushed on stack
 - **1000h**
 - 2000h
 - **3000h**

000011	Before			After	
0700:FFFE	000A	← SP	0700:FFFE	000A	
0700:FFFC			0700:FFFC	1000	
0700:FFFA			0700:FFFA	2000	
0700:FFF8			0700:FFF8	3000	← SP
0700:FFF6			0700:FFF6		
0700:FFF4			0700:FFF4		
0700:FFF2			0700:FFF2		
0700:FFF0			0700:FFF0		

Push Operation (3/3)

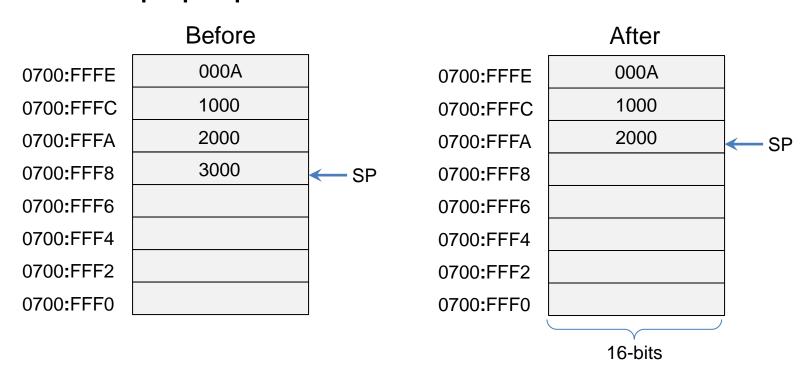
This value is pushed on stack

1000 0000h



Pop Operation

- Removes value from top of stack
- SP is incremented by stack element size with each pop operation



PUSH Instruction

- PUSH instruction is executed in two steps
 - First decrements SP by the size of stack element
 - Then copies the source operand on top of stack
- PUSH instruction formats
 - PUSH reg/mem16 → contents of 16-bit register or 16-bit memory location is pushed on stack
 - PUSH imm16 → 16-bit immediate value is pushed on stack
- Examples are
 - PUSH AX
 - PUSH 10h

POP Instruction

- POP instruction is executed in two steps
 - First the contents of stack element pointed to by SP are copied into destination operand
 - Then SP is incremented by the size of stack element
- Only one POP instruction formats
 - POP reg/mem16 → copies the value pointed to by SP into 16-bit register or 16-bit memory location
- Examples are
 - POP AX
 - POP var ;where var is a 16-bit memory location

PUSHF and POPF Instructions

- PUSHF is used to push EFLAGS register on the stack
- POPF pops the stack into EFLAGS register
- When using these instructions, make sure program's execution path does not skip over POPF instruction
- Syntax is
 - PUSHF
 - POPF

PUSHA and POPA Instructions

- PUSHA instruction pushes all 16-bit general purpose register on the stack in given order
 - AX, CX, DX, BX, SP, BP, SI, DI
- POPA instruction pops the same registers in the revers order
- Useful when modifying many general purpose registers inside a procedure

Stack Applications

- Registers can be saved temporarily when used for more than one purpose
- When CALL instruction executed, return address is saved on the stack
- Arguments are passed to a subroutine by pushing them on the stack
- Stack can be used as temporary storage for local variables inside a subroutine