

# **Discrete Structures**

## **Lecture # 13**

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# TREE

A **tree** is a **connected graph** that does not contain any **nontrivial circuit**. (it is **circuit-free**).

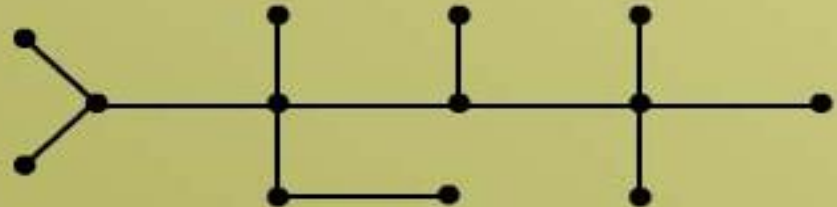
A **trivial circuit** is one that consists of a **single vertex**.

# EXAMPLE



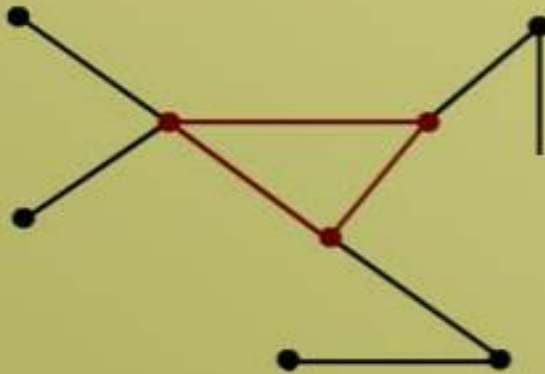
TREE

TREE

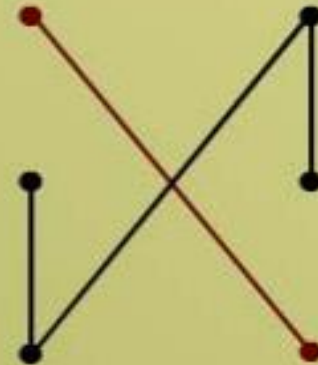


TREE

## EXAMPLES OF NON TREES



(a) Graph with a circuit.



(b) Disconnected graph.



(c) Graph with a circuit.

## SOME SPECIAL TREES

### 1. TRIVIAL TREE

A **graph** that consists of a **single vertex** is called a **trivial tree** or **degenerate tree**.

### 2. EMPTY TREE

A **tree** that does not have any **vertices or edges** is called an **empty tree**.

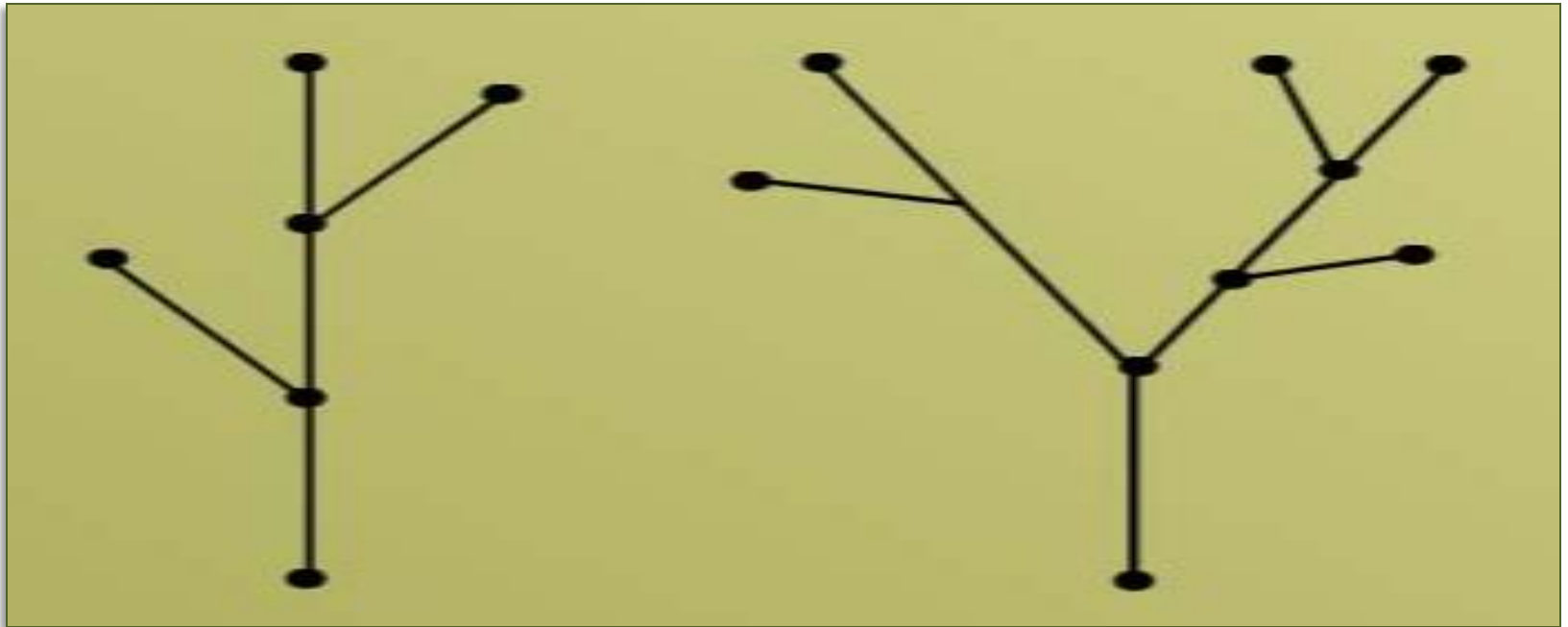
## SOME SPECIAL TREES

### FOREST:

A **graph** is called a **forest** if, and only if, it is **circuit-free**.

Hence, the **connected components** of a **forest** are **trees**.

# FOREST



Forest has only **two trees** in it.

# PROPERTIES OF TREES

1. A **tree** with  **$n$  vertices** has  **$n - 1$  edges**.
2. Any **connected graph** with  **$n$  vertices** and  **$n - 1$  edges** is a **tree**.
3. A **tree** has no **nontrivial circuit**; but if **one** new **edge** is added to it, then the resulting **graph** has exactly **one nontrivial circuit**.



## PROPERTIES OF TREES

4. A **tree** is connected, but if any **edge** is deleted from it, then the **resulting graph** is **not connected**.
5. Any **tree** that has more than one **vertex** has at least **two vertices** of **degree 1**.
6. A **graph** is a **tree** iff there is a **unique path** between any two of its **vertices**.

## EXERCISE

Explain why **graphs (Tree)** with the given specification do not **exist**.

1. Tree with **twelve vertices** and **fifteen edges**.
2. Trees with **five vertices** and total **degree 10**.

### SOLUTION:

1. Any tree with 12 vertices will have  $12 - 1 = 11$  edges.

Hence solution of first is not possible.

## SOLUTION

Tree with **five vertices** and total degree **10**.

Any **tree** with **five vertices** will have  $5 - 1 = 4$  **edges**.

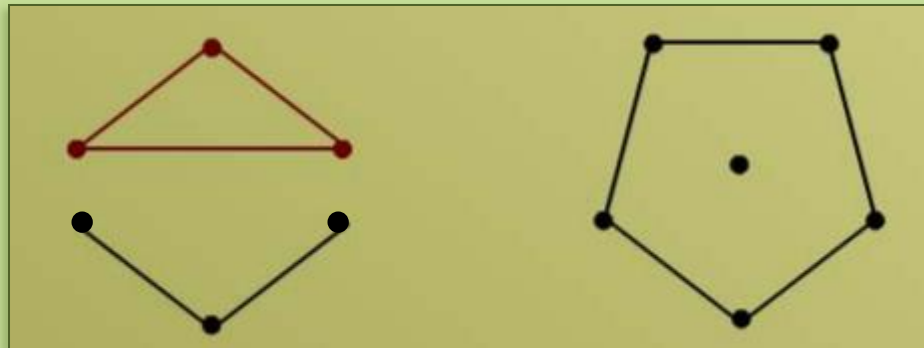
We are given **total degree** of graph is **10**. So it must have **edges**  $10/2 = 5$ .

The two conditions contradict each other.

## SOLUTION

Draw a graph with **six vertices**, **five edges** that is not a **tree**.

**SOLUTION:**

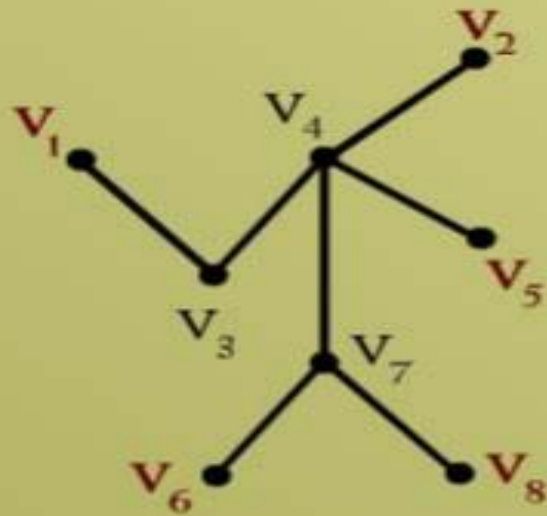


First is not **tree** because it is **not connected** and also has a **circuit** similarly for second.

## TERMINAL AND INTERNAL VERTEX

A vertex of degree '1' in a tree is called terminal vertex or a leaf (leaf has no children) and a vertex of degree greater than '1' in a tree is called an internal vertex or a branch vertex.

## EXAMPLE



$v_1, v_2, v_5, v_6$  and  $v_8$  are **terminal vertices**.

$v_3, v_4, v_7$  are **internal vertices**.

# ROOTED TREE

A **rooted tree** is a **tree** in which **one vertex** is **distinguished** from the others and is called the **root**.

The **level of a vertex** is the number of **edges** along the **unique path** between it and the **root**.

The **height of a rooted tree** is the maximum level to **any vertex** of the **tree**.



# ROOTED TREE

The **children** of any **internal vertex  $v$**  are all those **vertices** that are **adjacent to  $v$**  and are **one level farther** away from the **root** than  **$v$** . If  **$w$**  is a **child** of  **$v$** , then  **$v$**  is called the **parent** of  **$w$** .

Two **vertices** that are **both children** of the **same parent** are called **siblings**.

Given **vertices  $v$**  and  **$w$** , if  **$v$**  lies on the unique path between  **$w$**  and the **root**, then  **$v$**  is an **ancestor** of  **$w$**  and  **$w$**  is a **descendant** of  **$v$** .



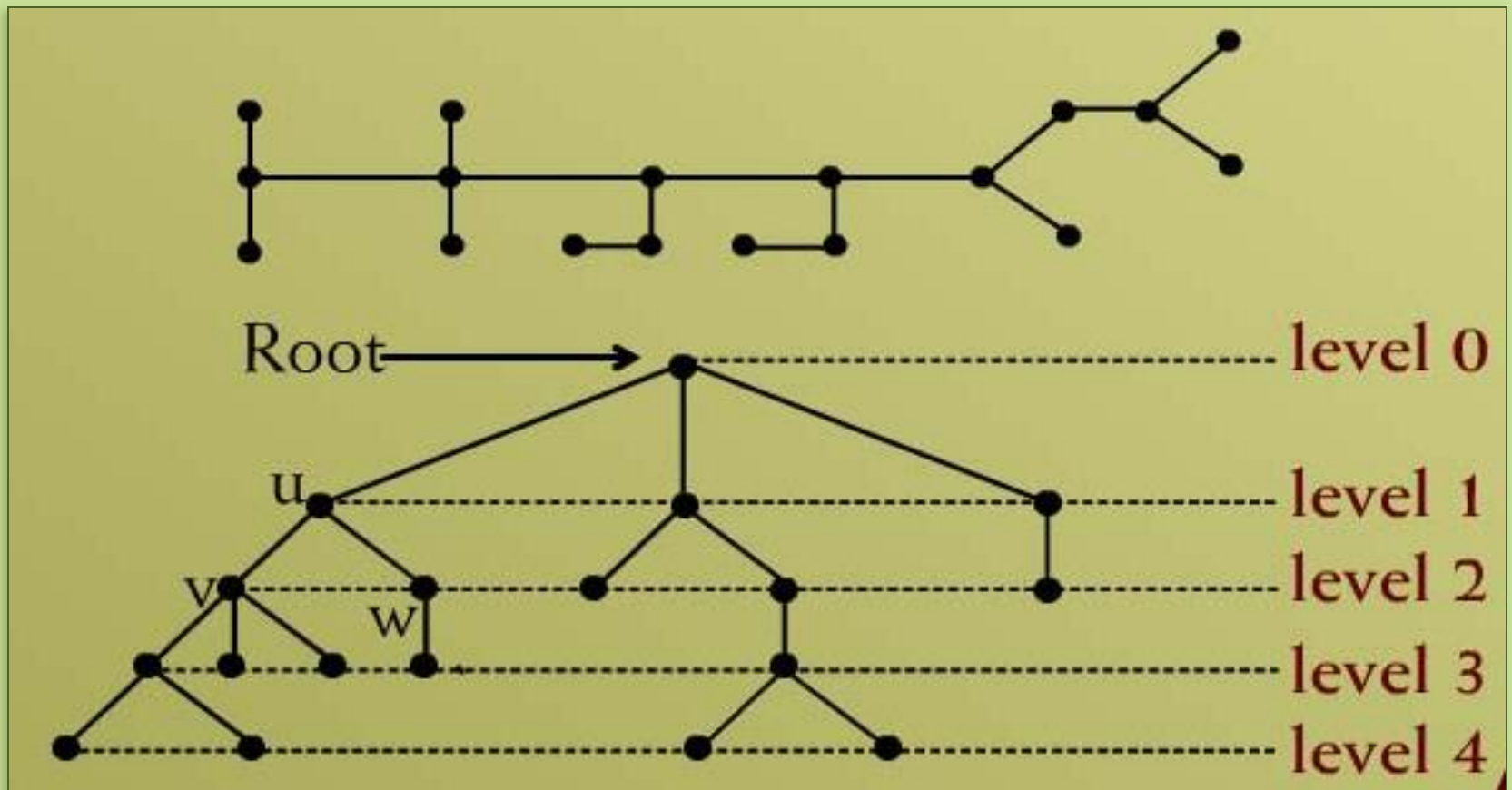
## ROOTED TREE

The root is an internal vertex unless it is the only vertex in the graph, and in that case it will be a leaf.

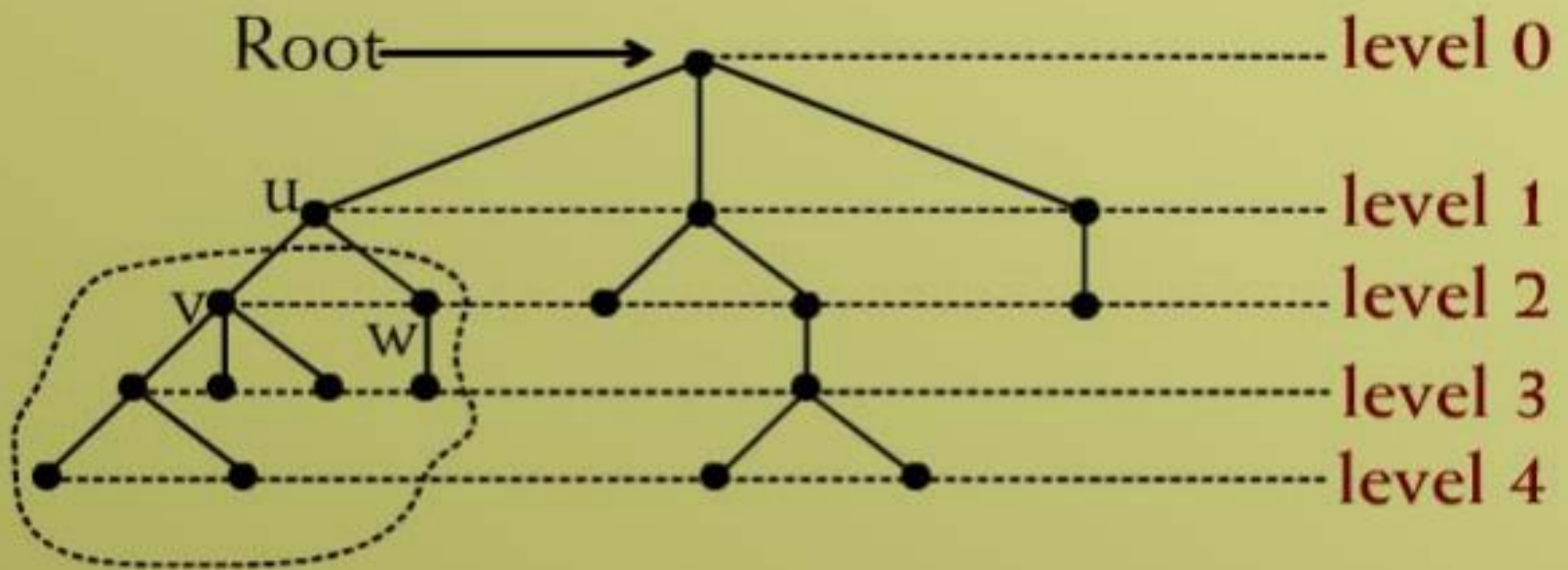
## Subtree

If  $a$  is a vertex in a tree, the subtree with  $a$  as its root is the subgraph of the tree consisting of  $a$  and its descendants and all edges incident to these descendants.

# EXAMPLE

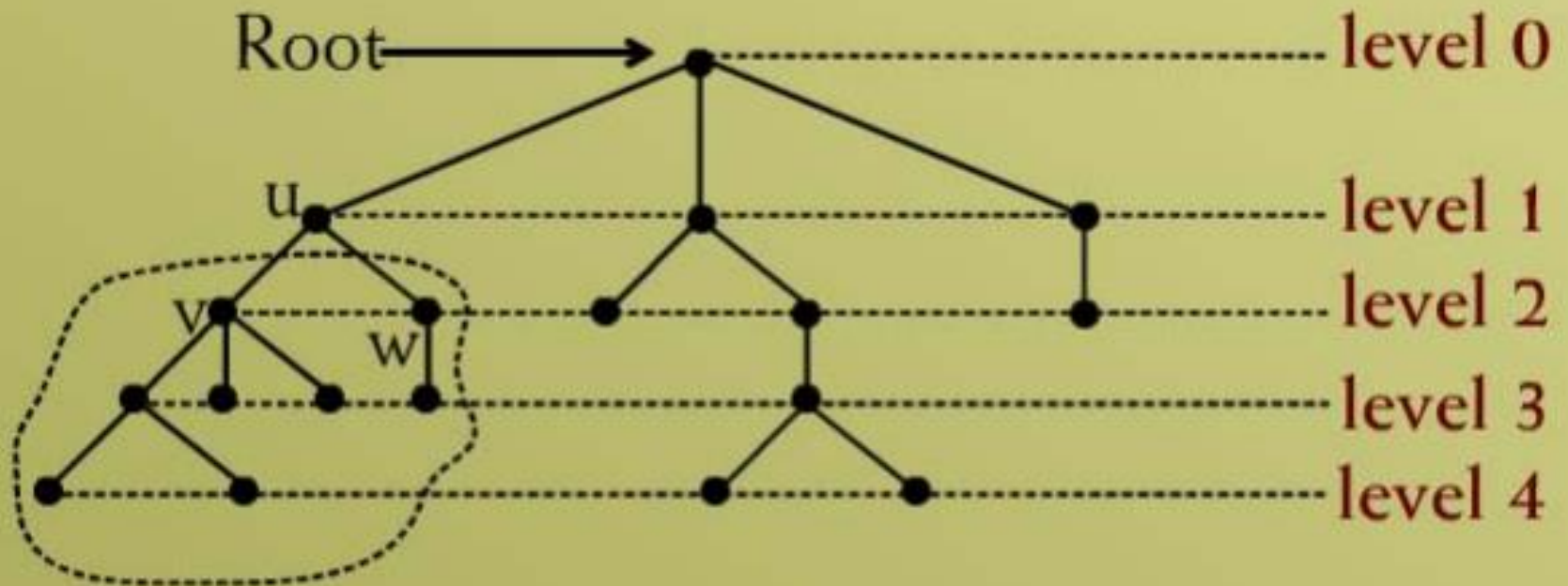


# EXAMPLE



Vertices in enclosed region are descendants of  $u$ , which is an ancestor of each.

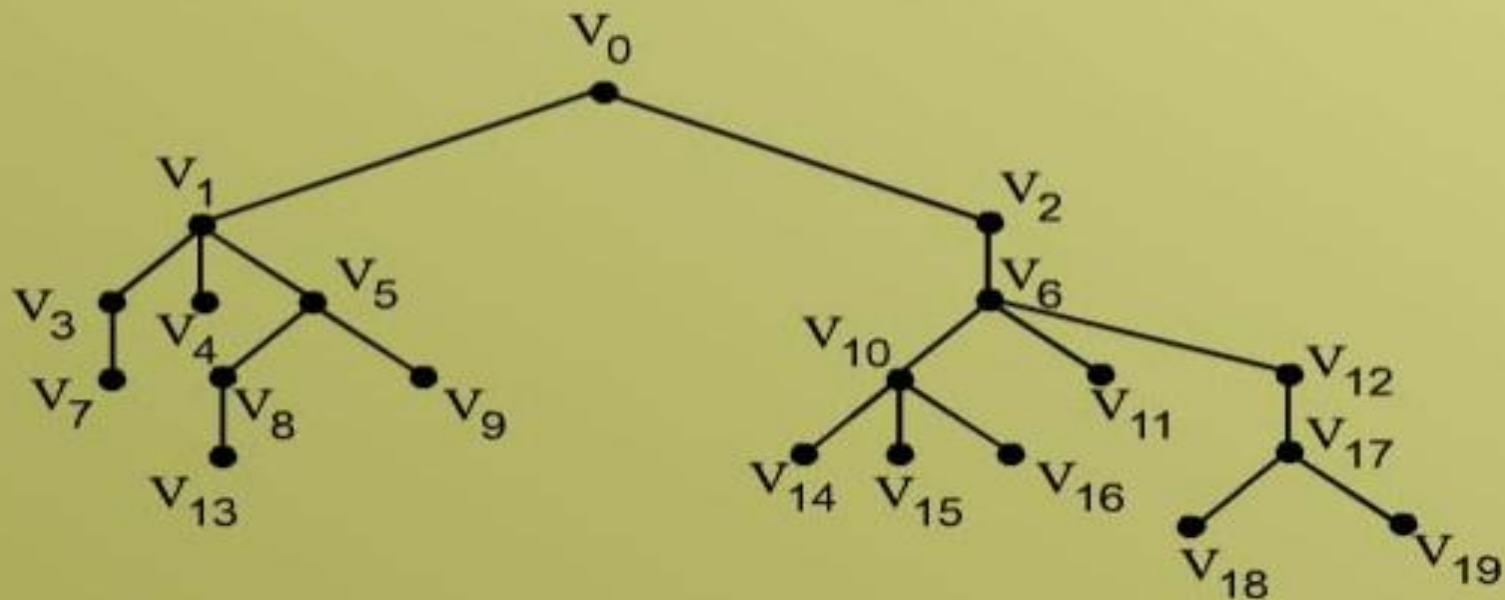
# EXAMPLE



$v$  is a **child** of  $u$ ,  $u$  is the **parent** of  $v$ ,  $v$  and  $w$  are **siblings**,  $\text{height} = 4$ .

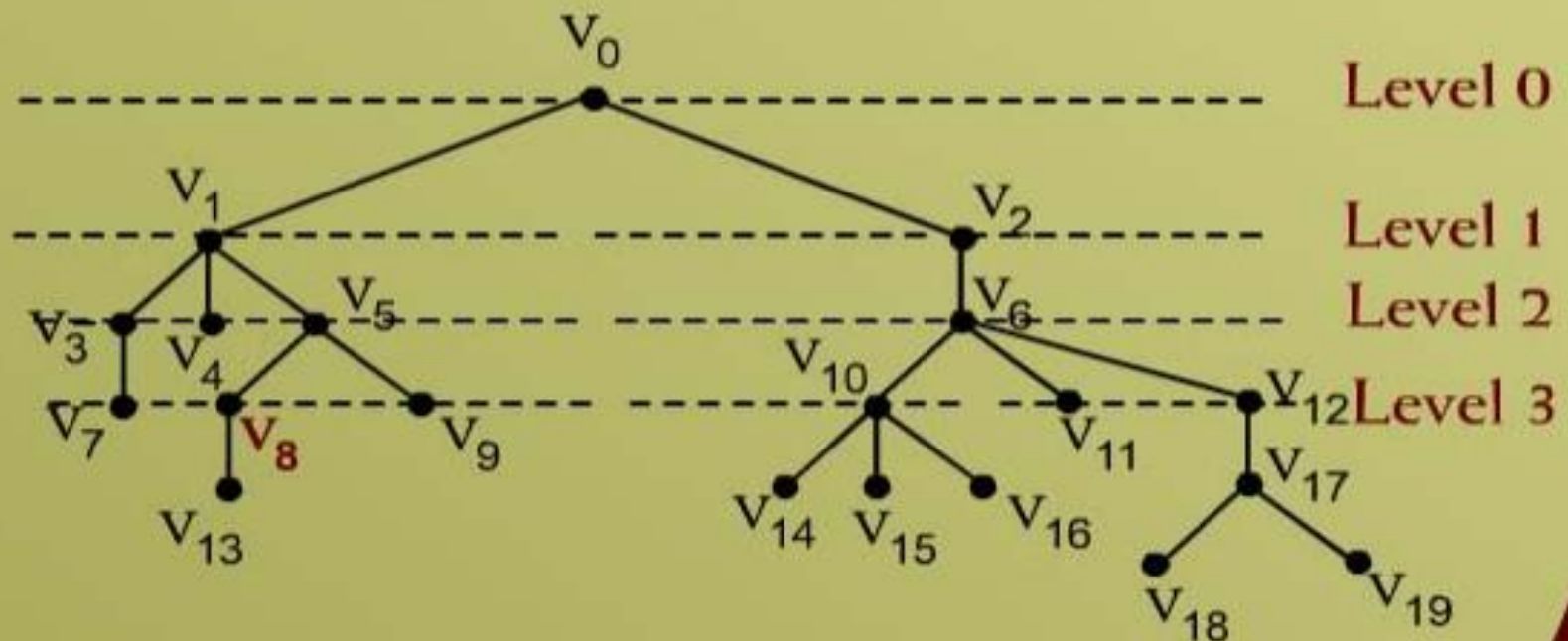
## EXERCISE

Consider the **rooted tree** shown below with root  $v_0$ .



## EXERCISE

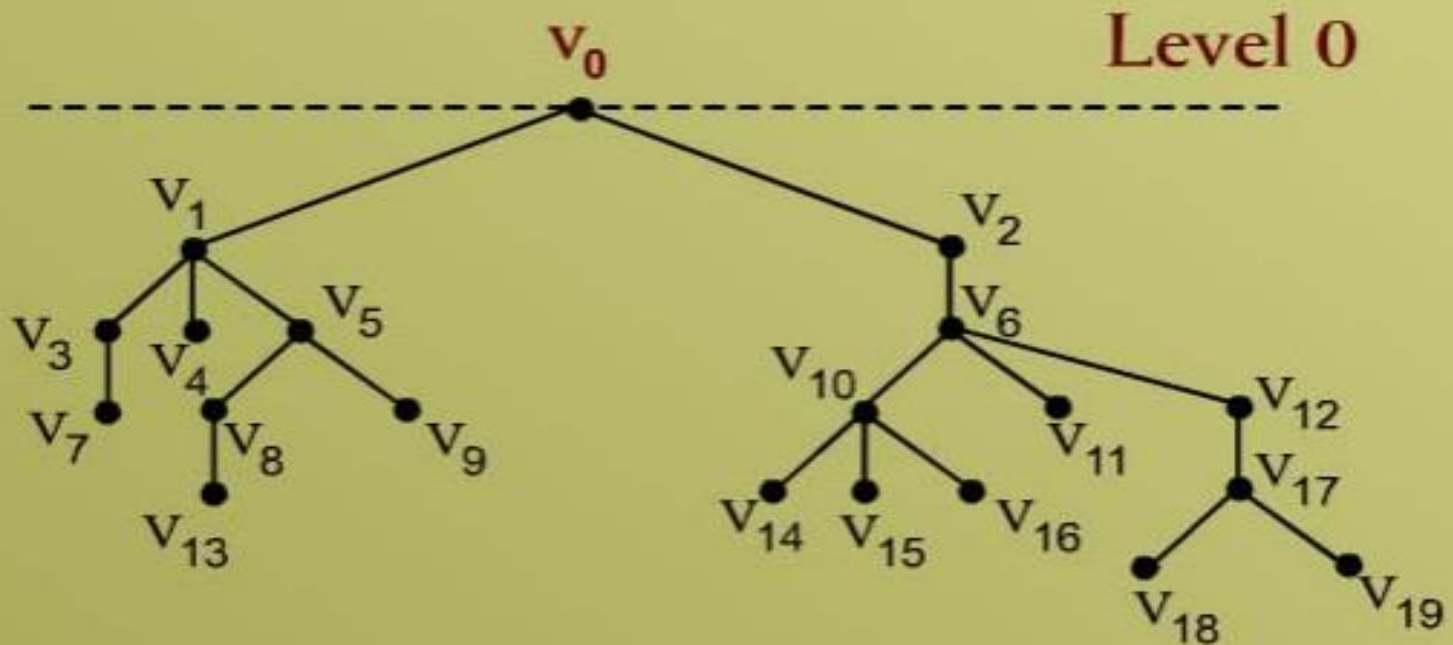
a. What is the **level** of  $v_8$ ?





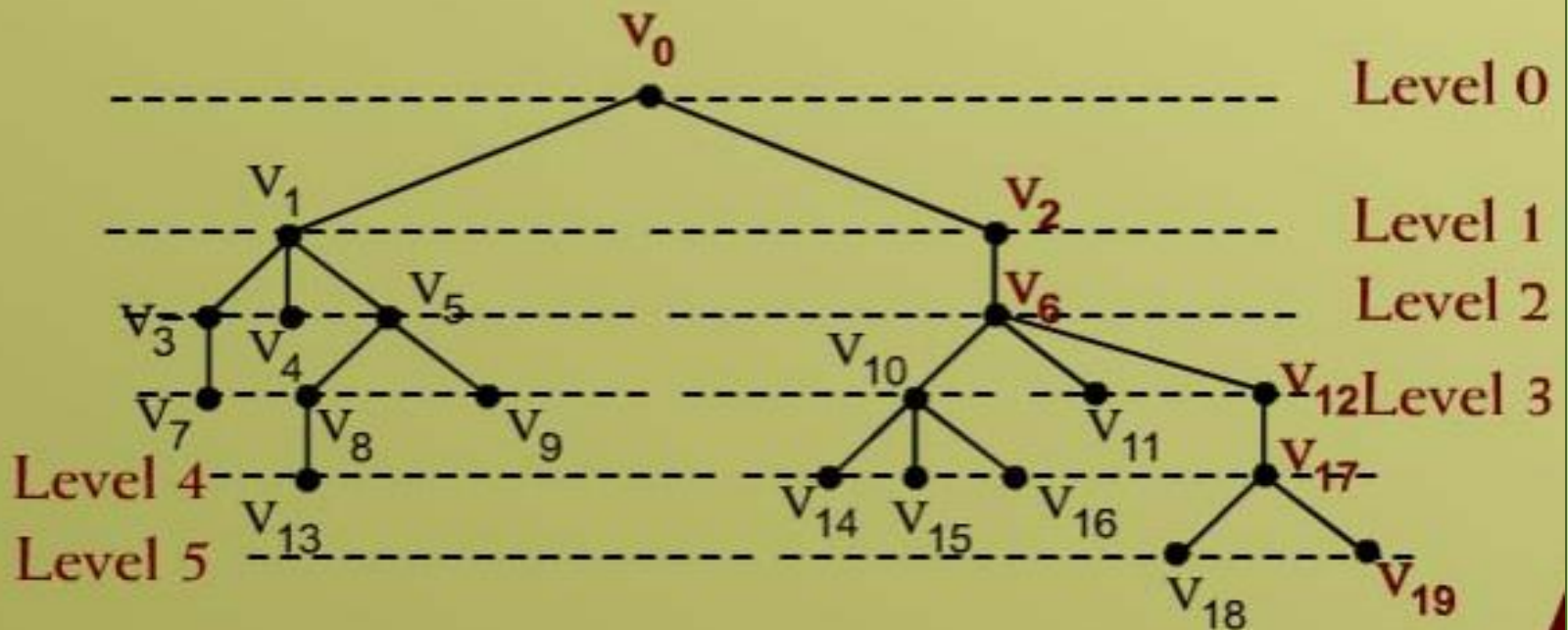
## EXERCISE

b. What is the **level** of  $v_0$ ?



## EXERCISE

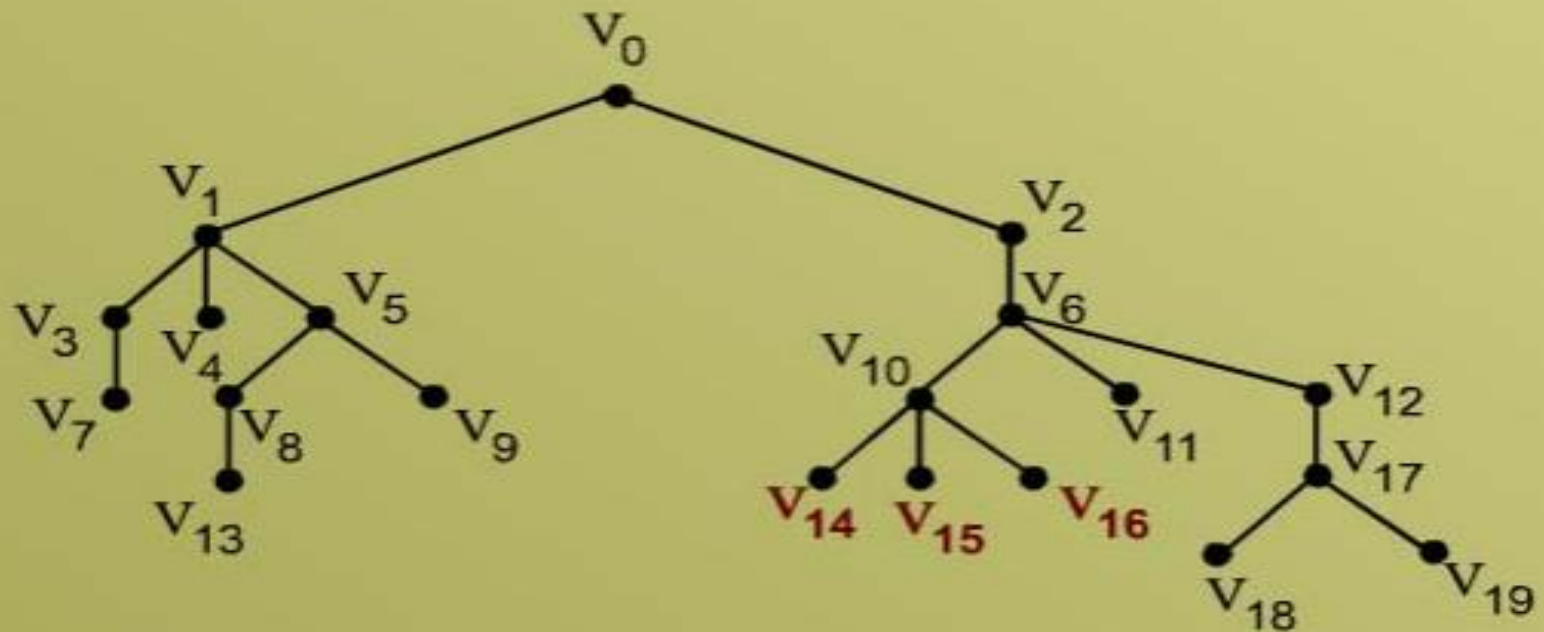
c. What is the **height** of this tree?





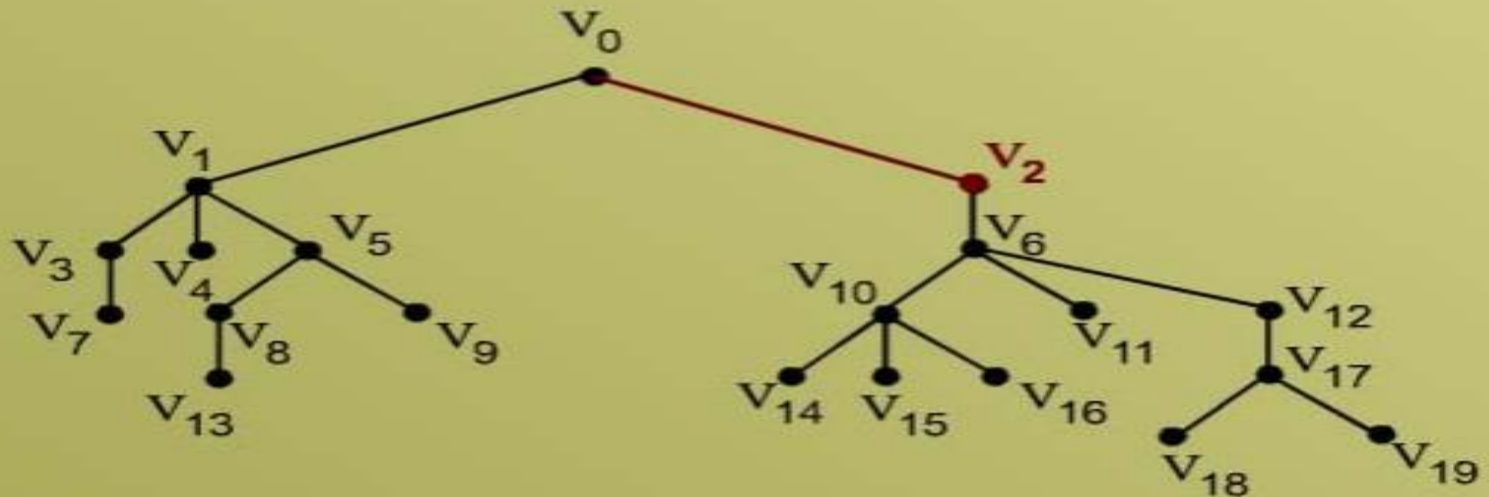
## EXERCISE

d. What are the **children** of  $v_{10}$ ?



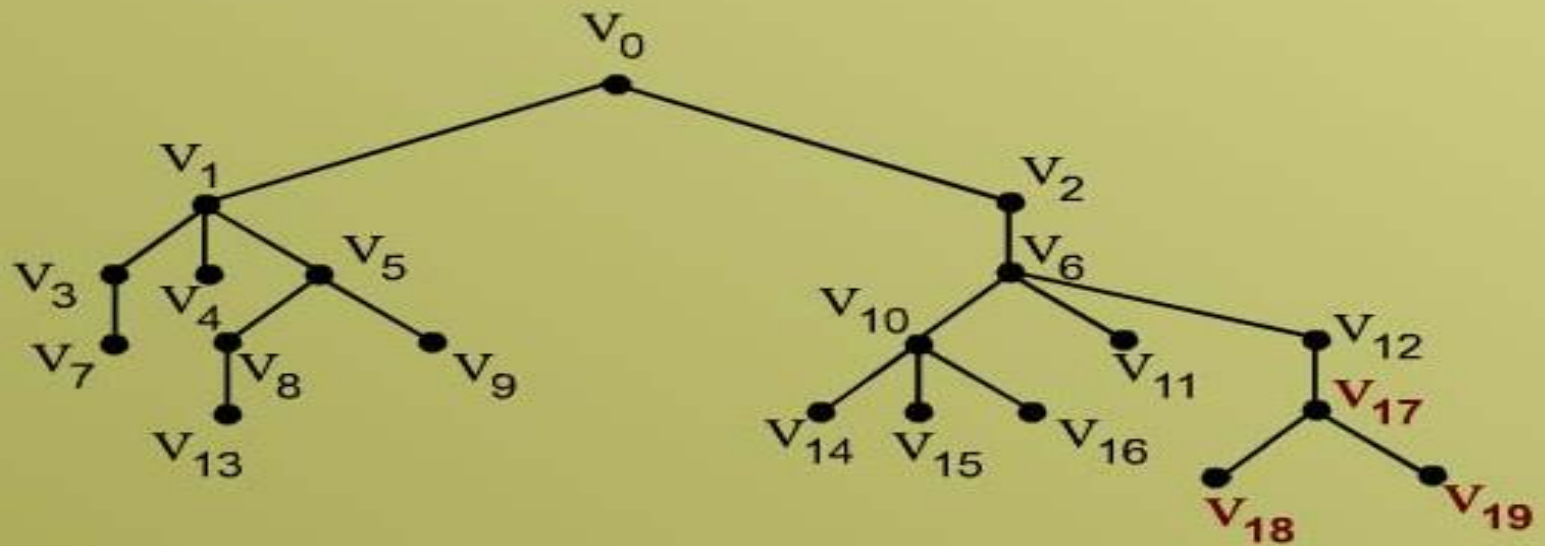
## EXERCISE

e. What are the **siblings** of  $v_1$ ?



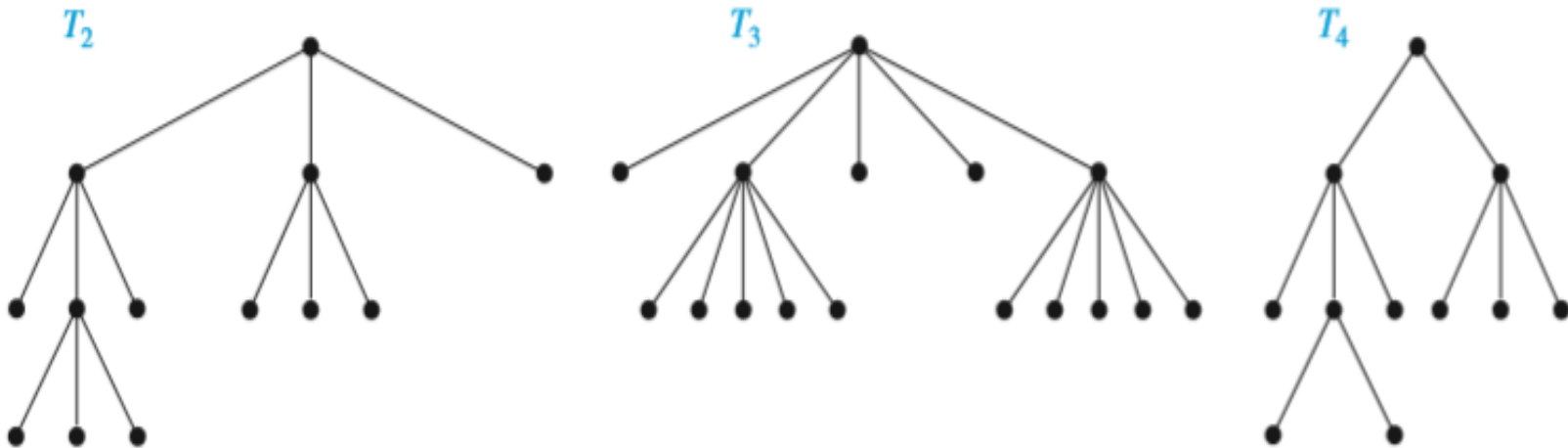
## EXERCISE

f. What are the **descendants** of  $v_{12}$ ?



## *m*-ary TREE

- A rooted tree is called *m*-ary tree if every internal vertex has no more than *m* children.
- The tree is called *full m*-ary tree if every internal vertex has exactly *m* children.
- An *m*-ary tree with  $m=2$  is called a *binary tree*.



# BINARY TREE

A **binary tree** is a **rooted tree** in which every **internal vertex** has at **most two children**.

Every child in a **binary tree** is designated either a **left child** or a **right child**.

# FULL BINARY TREE

A full binary tree is a binary tree in which each internal vertex has exactly two children.



## THEOREM

A full  $m$ -ary tree with  $k$  internal vertices contains  $n = mk + 1$  vertices.

1. If  $k$  is a positive integer and  $T$  is a full binary tree with  $k$  internal vertices, then  $T$  has a total of  $2k + 1$  vertices and has  $k + 1$  terminal vertices.



## THEOREM

A full  $m$ -ary tree with

1.  $n$  vertices has  $k = (n-1)/m$  internal vertices and  $l = [(m-1)n + 1]/m$  leaves.
2.  $k$  internal vertices has  $n = mk + 1$  vertices and  $l = (m-1)k + 1$  leaves.
3.  $l$  leaves has  $n = (ml-1)/(m-1)$  vertices and  $k = (l-1)/(m-1)$  internal vertices.

## THEOREM

There are at most  $m^h$  leaves (terminal vertices) in an  $m$ -ary tree of height  $h$ .

## THEOREM

2. If  $T$  is a **binary tree** that has  **$t$  terminal vertices** and **height  $h$** , then

$$t \leq 2^h$$

Equivalently,

$$\log_2 t \leq h$$

## EXERCISE

Explain why **graphs** with the given **specification** do not exist.

1. **Full binary tree** with **nine vertices** and **five internal vertices**.
2. **Binary tree** with **height 4** and **eighteen terminal vertices**.

## SOLUTION

1. Any **full binary tree** with **five internal vertices** has **six terminal vertices**, for a total of **eleven**, not **nine vertices** in all.

Thus there is **no full binary tree** with the given properties.

2. Any **binary tree** of **height 4** has at most  **$2^4 = 16$**  terminal vertices.

Hence, there is **no binary tree** that has **height 4** and **eighteen terminal vertices**.

## EXERCISE

Draw a **full binary tree** with **seven vertices**.

SOLUTION:

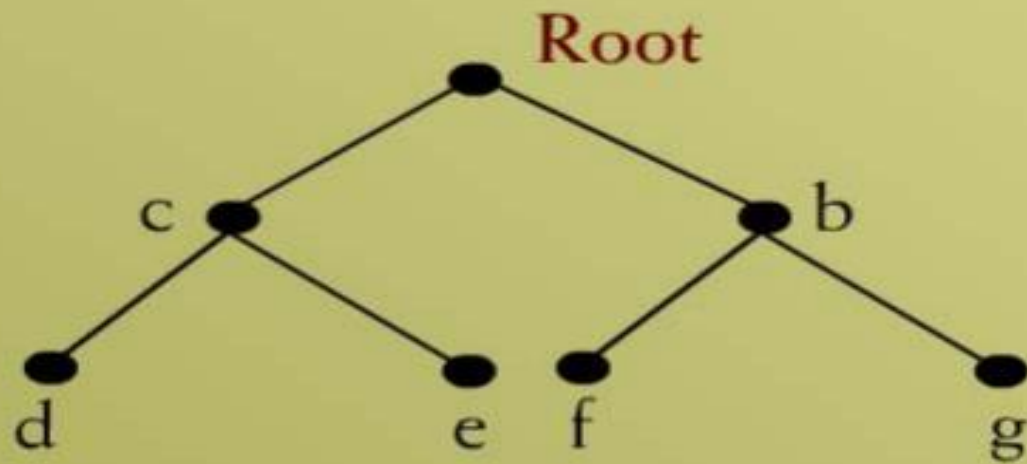
$$\text{Total vertices} = 2k + 1 = 7$$

$$\Rightarrow k = 3$$

Hence, number of **internal vertices** =  $k = 3$

Number of **terminal vertices** =  $k + 1 = 4$

## EXERCISE

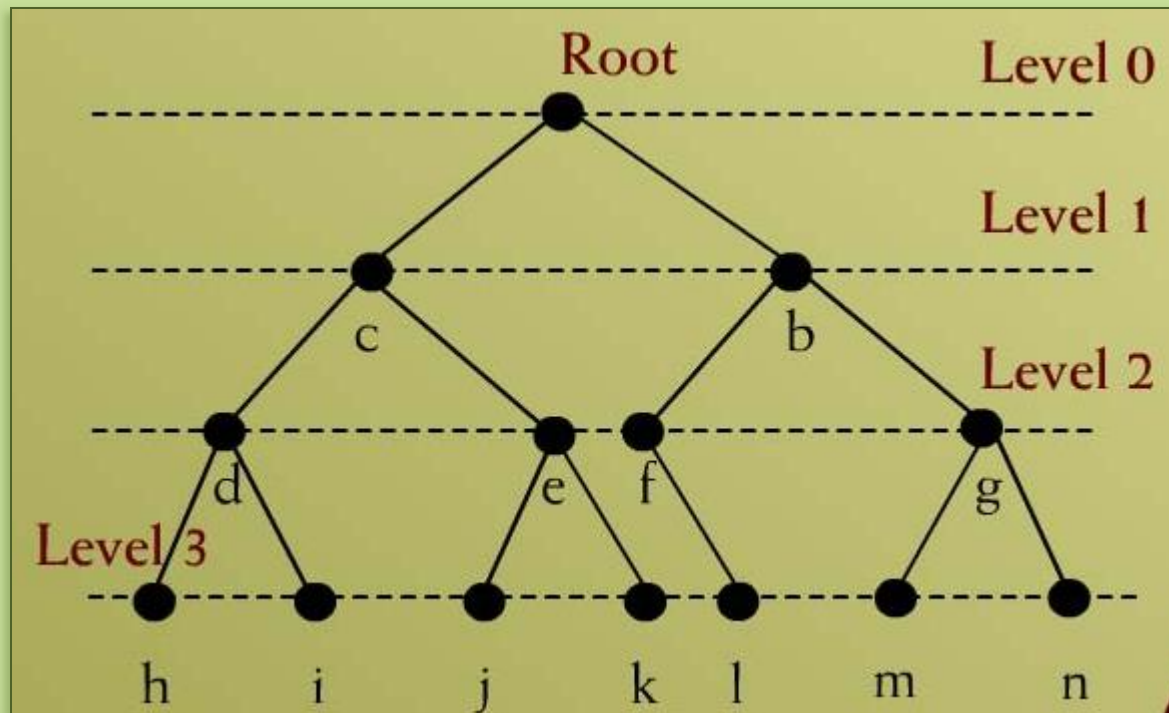


Which is required **full binary tree** with **seven vertices**.

## EXAMPLE

Draw a **binary tree** with **height 4 (level 3)**  
and having **seven terminal vertices**.

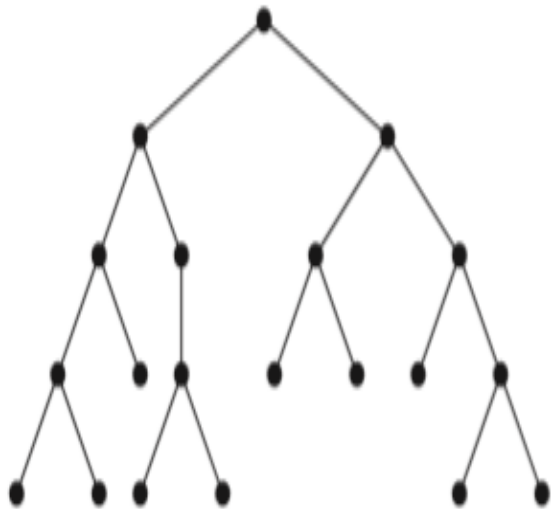
# SOLUTION





# Balanced Rooted Tree

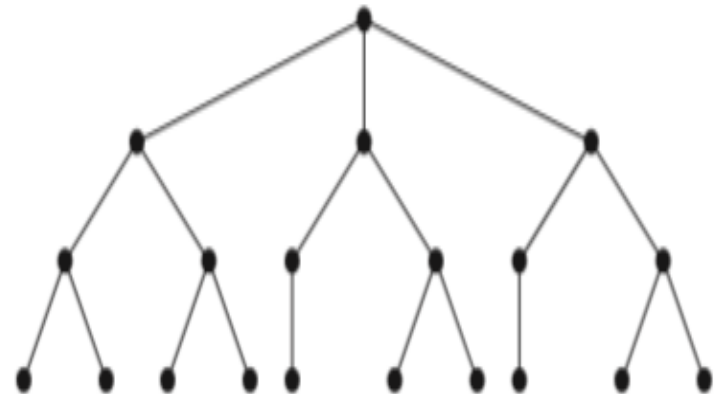
A rooted  $m$ -ary tree of height  $h$  is balanced if all leaves are at levels  $h$  or  $h-1$ .



$T_1$



$T_2$



$T_3$