- midtern review lecture 17 - Oct 14 ProctorTrack onboarding · did you get an email confirming success? · did you see how to enter password in UR Courses and see the exam question? · Midtern exam password will be shared prior to exam, but will also be shown in Proctor Track when starting the exam. Be sure to enter it in UR Courses! · Have blank paper and pencil ready on the side, But: All auswers entered in UR Courses via Keyboard (no mage upload) · be alone in your room! use UR Courses duat ONLY in order to ask me questions to clarify the meaning of an exam problem. · wnte Big-O, Omega, Thela, N^2, 2°k, etc., to ease typing of mathematical notation.

Chapter O. Introduction How did our two algorithms for Fibonacci numbers differ? To what effect? => algorithm design can greatly afect nutive note: space - time trade - of Chapter 1. Algorithan Analysis 1> suplified abstract model (why?) Tang (N), Tworst (N), Thest (N) Tang (N) is not related to amortized cost!  $T_2(N) = \theta(N)$ p recall 0,0,0,0  $/T_3(N) = \Theta(N^2)$   $/(T_3(N) = \Theta(N^2)$  $T_{\mu}(N)$ =0(1)  $T_2(N) = \Phi(T_2(N))$ ;  $T_2(N) = o(T_1(N))$ ;  $T_2(N) = \Omega(T_4(N))$ 

theorems and 2 about 0, 12, 0 are important to understand. to Running time calculations - consecutive fragments - if /else - loops, nested loop - recursion loops often yield standard sums:  $\sum_{k=1}^{1N} k = (+2+3+...+N) = \frac{N(N+1)}{2} = O(N^2)$  $\sum_{k=1}^{N} \frac{1}{2^{k}} = \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots + \frac{1}{2^{N}} < 1 = O(1)$ recursion: formulate reccurence relation, try it out for a few steps (small values of N), until you see a pattern, guess a general formula, and verify by induction D recall important basics on frees (BSTs o splay tree operations and the notion of amortized - th sotion into splay wee of size N has worst-care

- however, amortized cost for cusertion is O(log(N)) Chapter 2. Priority Queues D bihary heap (min-heap) array representation percolating up /down all operations, e.g., O(log (N)) user how deletallin G(N) build Heap D d-heaps 17 lefist heaps t skew heaps to binomial gueues Chapter 3. Sorhug always in increasing order D Insertion Sort Tworst (N) = Tang (N) = O(N2) to do better, one must sometimes exchange non-adjacent elements Dellsort: différent gap seguences => effect on nuntime?

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