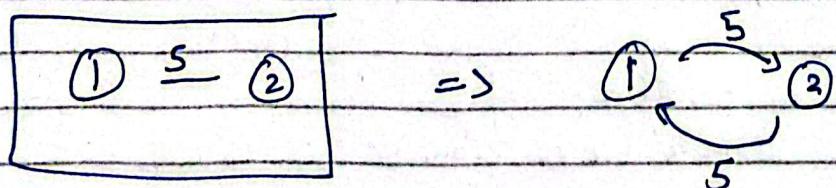


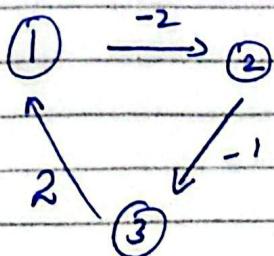
## BELLMAN FORD

### ALGORITHM

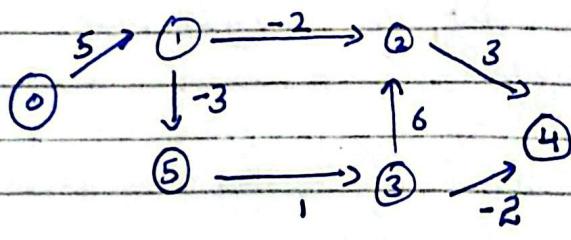
- It is also for shortest path
- Dijkstra's fail if the graph have negative weights.
- Dijkstra gives tle if negative cycles exist.
- It helps us to detect negative cycles as well.
- It is applicable only in DG directed graphs.

If it is undirected we should convert it to directed.





negative cycle if any weight is negative.



$(u, v, wt)$

- $(3, 2, 6)$
- $(5, 3, 1)$
- $(0, 1, 5)$
- $(1, 5, -3)$
- $(1, 2, -2)$
- $(3, 4, -2)$
- $(2, 4, 3)$

\* Relax all the edges  
N-1 times.

edges can  
be in any  
order

+ Relaxation:

if  $(\text{dist}[u] + \text{wt} < \text{dist}[v])$

Suppose if

$$\text{dist}[v] = \text{dist}[u] + \text{wt}$$

$$⑤ \xrightarrow{1} ③$$

$$\text{dist}[3] + 1$$

to reach 3

then

$$\text{if } N = 6$$

then 5 iterations of  
relaxation should be done.

$$\text{dist} [0, \infty, \infty, 0, \infty, \infty]$$

or 2 3 4 5

in every iteration we will go through all the edges.

1 iteration.

if ( $\text{dist}[3] + 6 < \text{dist}[2]$ )  
 $\infty < \infty$

if ( $\text{dist}[5] + 1 < \text{dist}[3]$ )  
 $\infty < \infty$

if ( $\text{dist}[0] + 5 < \text{dist}[1]$ ) ✓  
 $\infty < \infty$

So dist becomes

$\text{dist} = [0, 5, \infty, \infty, \infty, \infty]$ .

if ( $\text{dist}[1] + -3 < \text{dist}[5]$ )

So dist becomes

$\text{dist} = [0, 5, \infty, \infty, \infty, 2]$

if ( $\text{dist}[1] + (-2) < \text{dist}[2]$ )

$\text{dist} = [0, 5, 3, \infty, \infty, 2]$

if ( $\text{dist}[3] - 2 < \text{dist}[4]$ )  
 $\infty$

if ( $\text{dist}[2] + 3 < \text{dist}[4]$ )  
 $3+3 < \infty$  ✓

$\text{dist} = [0, 5, 3, \infty, 8, 2]$

At the end of first iteration we have

$$\text{dist} [0, 5, 3, \infty, 6, 2]$$

0 1 2 3 4 5

2nd iteration:

if  $\text{dist}[3] + 6 < \text{dist}[2]$

if  $(\text{dist}[5] + 1 < \text{dist}[3])$

2 + 1 < \infty

so dist becomes

$$\text{dist} = [0, 5, 3, 3, 6, 2]$$

:

if  $(\text{dist}[3] - 2 < \text{dist}[4])$

3 - 2 = 1 < 6

so

$$\text{dist} = [0, 5, 3, 3, 4, 2]$$

:

Similarly we will do 1<sup>st</sup>, 2<sup>nd</sup> ... 5<sup>th</sup>.

Q: why N-1 ?

Q: How to detect negative cycle.

Intuition why  $N-1$  iterations?

$$① \rightarrow ② \rightarrow ③ \rightarrow ④ \quad (u, v, wt)$$

$$(3, 4, 1)$$

$$(2, 3, 1)$$

$$(1, 2, 1)$$

$$(0, 1, 1)$$

So we have to do 4 iterations.

①

$$\text{dist} [0, \infty, \infty, \infty, \infty]$$

$$\text{dist}[3] + 1 < \text{dist}[4] \quad \times$$

$$\text{dist}[2] + 1 < \text{dist}[3] \quad \times$$

$$\text{dist}[1] + 1 < \text{dist}[2]$$

$$\text{dist}[0] + 1 < \text{dist}[1] \quad \checkmark$$

So

$$\text{dist} = [0, 1, \infty, \infty, \infty]$$

②

$$\text{dist} = [0, 1, 2, \infty, \infty]$$

③

$$\text{dist} = [0, 1, 2, 3, \infty]$$

④

$$\text{dist} = [0, 1, 2, 3, 4]$$

So that's why  $N-1$  iterations.

Q: How to detect Negative cycles?

if after the  $N-1$  iteration  
still distance reduces then  
the graph have negative  
cycle

