## Maximum Absolute Difference

https://www.interviewbit.com/problems/maximum-absolute-difference/

https://www.geeksforgeeks.org/maximum-absolute-difference-value-index-sums/

keywords: Arrays, properties of absolute values.

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You are given an array of N integers, A_1, A_2,..., A_N. Return maximum value of f(i, j) for all 1 \le i, j \le N. f(i, j) is defined as |A[i] - A[j]| + |i - j|, where |x| denotes absolute value of x. For example,
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A=[1, 3, -1]

f(1, 1) = f(2, 2) = f(3, 3) = 0

f(1, 2) = f(2, 1) = |1 - 3| + |1 - 2| = 3

f(1, 3) = f(3, 1) = |1 - (-1)| + |1 - 3| = 4

f(2, 3) = f(3, 2) = |3 - (-1)| + |2 - 3| = 5

So, we return 5.
```

An **efficient** solution in O(n) time complexity can be worked out using the properties of absolute values.

f(i, j) = |A[i] - A[j]| + |i - j| can be written in 4 ways (Since we are looking at max value, we don't even care if the value becomes negative as long as we are also covering the max value in some way).

```
Case 1: A[i] > A[j] and i > j

|A[i] - A[j]| = A[i] - A[j]

|i -j| = i - j

hence, f(i, j) = (A[i] + i) - (A[j] + j)

Case 2: A[i] < A[j] and i < j

|A[i] - A[j]| = -(A[i]) + A[j]

|i -j| = -(i) + j

hence, f(i, j) = -(A[i] + i) + (A[j] + j)
```

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Case 3: A[i] > A[j] and i < j

|A[i] - A[j]| = A[i] - A[j]

|i -j| = -(i) + j

hence, f(i, j) = (A[i] - i) - (A[j] - j)

Case 4: A[i] < A[j] and i > j

|A[i] - A[j]| = -(A[i]) + A[j]

|i -j| = i - j

hence, f(i, j) = -(A[i] - i) + (A[j] - j)
```

Note that case 1 and 2 are equivalent and so are case 3 and 4 and hence we can design our algorithm only for two cases as it will cover all the possible cases.

- 1. Calculate the value of A[i] + i and A[i] i for every element of the array while traversing through the array.
- 2. Then for the two equivalent cases, we find the maximum possible value. For that, we have to store minimum and maximum values of expressions A[i] + i and A[i] i for all i.
- 3. Hence the required maximum absolute difference is maximum of two values i.e. max((A[i] + i) (A[j] + j)) and max((A[i] i) (A[j] j)). These values can be found easily in linear time.
- a. For max((A[i] + i) (A[j] + j)) Maintain two variables max1 and min1 which will store maximum and minimum values of A[i] + i respectively. max((A[i] + i) (A[j] + j)) = max1 min1
- b. For max((A[i] i) (A[j] j)). Maintain two variables max2 and min2 which will store maximum and minimum values of A[i] i respectively. max((A[i] i) (A[j] j)) = max2 min2

```
1. public class Solution {
     public int maxArr(ArrayList<Integer> A) {
3.
        int max1,max2,min1,min2;
4.
5.
        max1=max2=Integer.MIN VALUE;
6.
        min1=min2=Integer.MAX VALUE;
7.
8.
       int i;
9.
10.
       for(i=0;i<A.size();i++)
11.
12.
          max1=Math.max(max1,A.get(i)+i);
13.
          min1=Math.min(min1,A.get(i)+i);
14.
          max2=Math.max(max2,A.get(i)-i);
```