

Chap 7.3

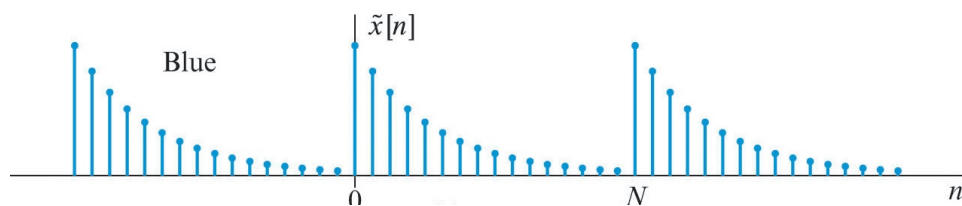
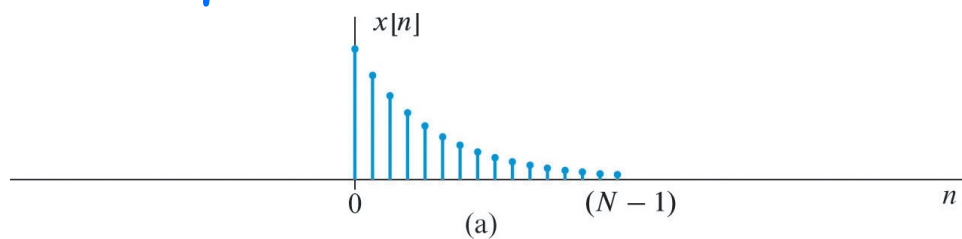
Discrete Fourier transform (DFT)

$$X[k] = \sum_{n=0}^{N-1} x[n] e^{-j \frac{2\pi}{N} kn}, \quad 0 \leq k < N$$

Inverse DFT

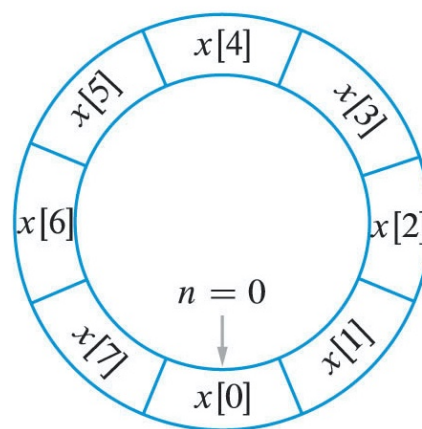
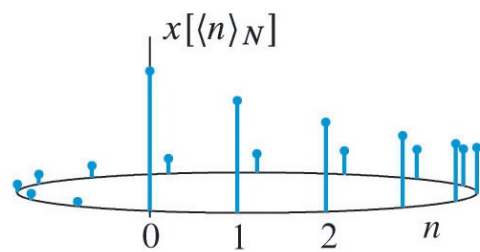
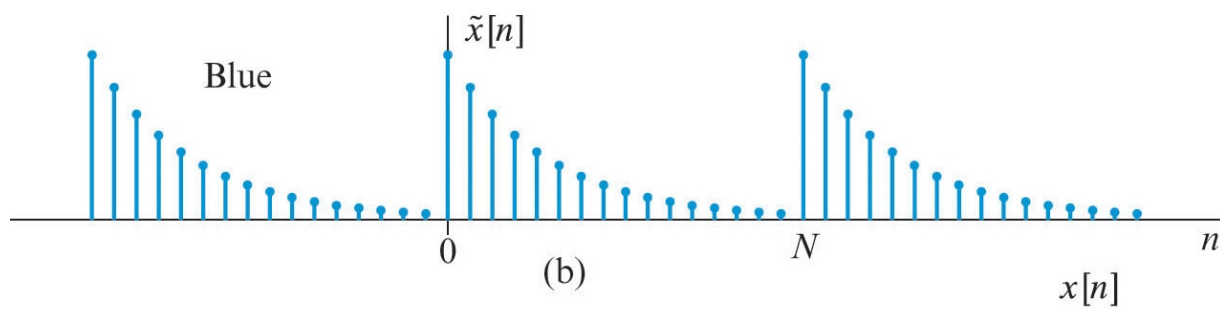
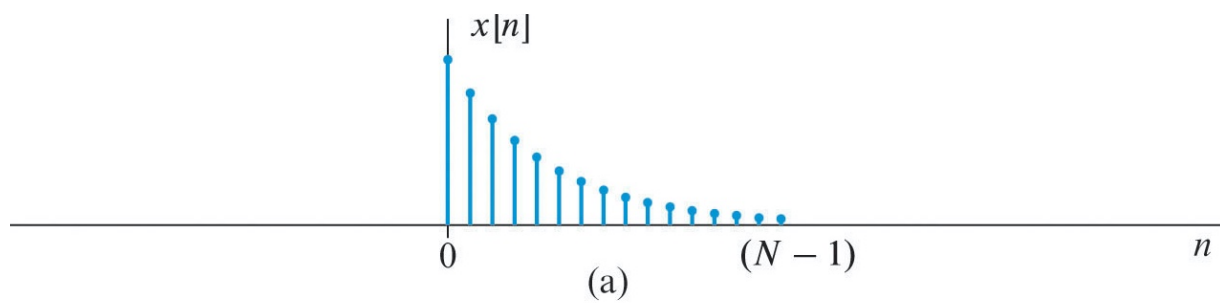
$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi}{N} kn}, \quad 0 \leq n < N$$

However, the reality is that the inverse is actually a periodic function with period N .



$$\tilde{x}[n] = \sum_{l=-\infty}^{\infty} x[n - lN]$$

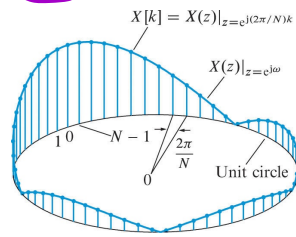
Of course because it is periodic, we only need 1 period $x[n]$, $0 \leq n \leq N-1$



(c)

(d)

Why is $\tilde{X}[n]$ a periodic function in time with period N ?



This is because the DFT are samples of the DTFT (i.e. N samples for 2π frequency). Since the DTFT is periodic, the samples are also periodic.

Time-domain aliasing happens if the length of $x[n]$ is $> N$

Ex 7.3

$$x[n] = a^n u[n], \quad 0 < a < 1 \xleftrightarrow{\text{DTFT}} X(e^{j\omega}) = \frac{1}{1 - a e^{-j\omega}}$$

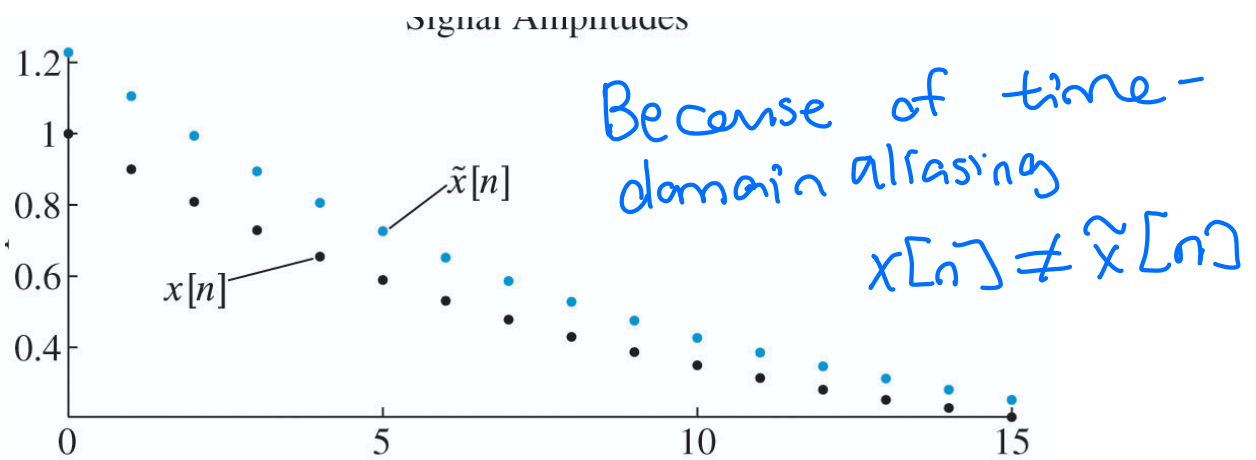
↓
note that $x[n]$ is infinite length, we will know that there will be time-domain aliasing

DFT gives us N samples over one period of the DTFT

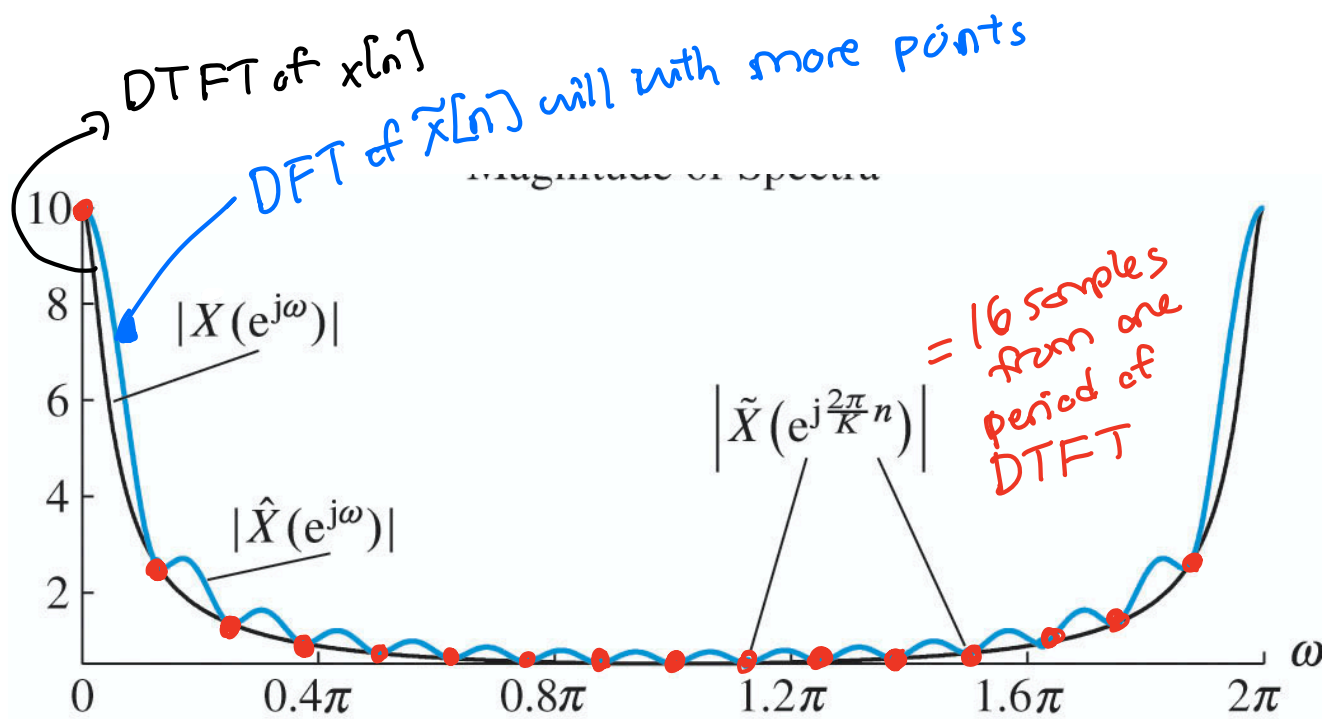
$$\tilde{X}(e^{j\frac{2\pi}{N}k}) \quad , \quad \text{Let } N=16$$

We can take the inverse DFT to get

$$\hat{x}[n] = \frac{1}{N} \sum_{k=0}^{N-1} \tilde{X}(e^{j\frac{2\pi}{N}k}) e^{j\frac{2\pi}{N}kn}$$



We can compute the DFT of $\tilde{x}[n]$



Aliasing in time is when N is not large enough. We should zero pad $x[n]$ to make sure N is large.

If $x[n]$ is infinite length, DFT will never exactly recover the true DTFT

Practical issues:

- 1) In many practical applications like speech processing and communication, the input sequence have indefinite length
- 2) Length of input sequence may be too large for storage and computation
- 3) Computation of output sequence can not be started until all input signal samples have been collected. This may cause unacceptable delays for many applications