

1 | INTERNAL CORE MODEL INDUCTION

1.1 OPERATORS AND HYBRID MICE

We'll need a generalisation of the concept of mice as we move up to the higher reaches of the core model induction, a generalisation usually known as either *hybrid mice* or *operator mice*. The basic concept is simple. When we're constructing "pure" mice we're traversing the \mathcal{J} -hierarchy, applying the $\mathcal{J}(x) := \text{rud}(x \cup \{x\})$ operator at every step, and taking unions at limit stages. In the hybrid case we're simply replacing \mathcal{J} with another *operator* \mathcal{F} , again applying it at every successor stage and taking unions at limits.

Figuring out what operators we're allowed to pick is the hard part, as we want to maintain all the fine structure that we get in the "pure" case. This has been done in great detail in [?], and we'll introduce a particularly simple case of their general definition here.

Definition 1.1 (Schlutzenberg-Trang). For a set x write $\rho_x : x \rightarrow \text{rk } x$ for the rank function of x and define the **rank closure** $\hat{x} := \text{trcl}(\{x, \rho_x\})$ of x and the **cone** $C_x := \{\hat{y} \in H_\kappa \mid x \in \mathcal{J}_1(\hat{y})\}$ over x . \circ

Definition 1.2 (Schlutzenberg-Trang). Let κ be an infinite cardinal, D a set of self-wellordered¹ sets and fix $b \in H_\kappa$. An **operator over H_κ on b with support D** is a partial function

$$\mathcal{F} : H_\kappa \dashrightarrow H_\kappa$$

such that $D \cap C_b \subseteq \text{dom } \mathcal{F}$ and $\text{dom } \mathcal{F}$ is closed under both unions and applications of \mathcal{F} . We also call C_b the **cone over b** and b the **base of C_b** . \circ

¹A set x is *self-wellordered* if there's a wellorder of x in $\mathcal{J}(x)$.

1.1. OPERATOR-INDUCED HYBRID INITIAL CORE MODEL INDUCTION

Before we move on, let's take a step back and have a look at a few examples of operators. As this is supposed to be a generalisation of the \mathcal{J} -function we'd want that to also be an operator, of course. Indeed, if $V = L$ then note that $\hat{x} = x$ for all x , so if we take \emptyset to be our base then $C_a = J_\kappa$, making $\text{dom } \mathcal{J} = J_\kappa$ trivially closed under both unions and applications of \mathcal{J} .

We could also let x be any set, assume that $V = L(\hat{x})$ and consider the operator \mathcal{J}_x , which is simply applying \mathcal{J} but with base \hat{x} instead of \emptyset . A similar argument as above would show that this is an operator on $J_\kappa(\hat{x})$ ($= H_\kappa^{L(\hat{x})}$) as well.

A slightly more sophisticated example would be $\mathcal{F} := (-)^\sharp$, where x^\sharp is the smallest initial segment of $\text{Lp}(x)$ with a measure. We can consider \mathcal{F} as an operator on \emptyset over HC, where $\mathcal{F}(\emptyset) = 0^\sharp$ and $\mathcal{F}^\omega(\emptyset) = \bigcup_{n < \omega} \mathcal{F}^n(\emptyset)$ is the smallest \sharp -closed structure.

Check the last example

We now, somewhat informally², define what hybrid mice are. Namely, an **\mathcal{F} -mouse on \hat{x}** is a structure \mathcal{M} built by successively applying \mathcal{F} to \hat{x} and taking unions at limits. We can stop the construction at any point and denote the amount of \mathcal{F} -applications by $l(\mathcal{M})$, the **length** of \mathcal{M} . The **initial segments** of \mathcal{M} are simply the intermediate models up to \mathcal{M} , being of the form $\mathcal{F}^\alpha(\hat{x})$ for some $\alpha < l(\mathcal{M})$.

As with regular mice, we require that they are *amenable*, *acceptable* and that proper initial segments \mathcal{F} -mice are *sound*. An extra property that we require in the hybrid context is that every proper initial segment $\mathcal{P} \triangleleft \mathcal{M}$ is *$<\omega$ -condensing*, roughly stating that if $\pi: \mathcal{N} \rightarrow \mathcal{P}$ is a sufficiently elementary embedding from a “nice” \mathcal{F} -structure then either \mathcal{N} is an initial segment of \mathcal{P} or an element of an ultrapower of \mathcal{P} ³.

Definition 1.3. Let κ be an infinite cardinal and \mathcal{F} an operator over H_κ on $b \in H_\kappa$. Then define the **\mathcal{F} -lower part model on b** as

$$\text{Lp}^\mathcal{F}(b) := \{\mathcal{M} \mid \mathcal{M} \text{ is a sound } \mathcal{F}\text{-mouse projecting to } b\} \quad \circ$$

²For a comprehensive definition, see [?, Section 2].

³For a proper definition of these concepts in the hybrid context, see [?, Section 2].

The fine-structural properties of operators from [?] can be used to show that, as in the pure case, $\text{Lp}^{\mathcal{F}}(b)$ is itself an \mathcal{F} -mouse. Further, $\text{Lp}^{\mathcal{F}}$ is itself an operator over H_κ on b and we write $\text{Lp}_\alpha^{\mathcal{F}}(b) := (\text{Lp}^{\mathcal{F}})^\alpha(b)$.

Check that the definition of $\text{Lp}^{\mathcal{F}}$ is correct and that it *does* in fact have the two properties above.

Definition 1.4 (Schindler-Steel). Let κ be an infinite cardinal and \mathcal{F} an operator over H_κ on $b \in H_\kappa$. Then \mathcal{F} **condenses well** if whenever $g \subseteq \text{Col}(\omega, \kappa)$ is V -generic and that there are models $\overline{\mathcal{M}}, \mathcal{M} \in H_\kappa$ and $\overline{\mathcal{M}}^+ \in V[g]$, all on b , with

- (i) $|\overline{\mathcal{M}}| = |b| \cdot \aleph_0$;
- (ii) $\overline{\mathcal{M}} \in \overline{\mathcal{M}}^+$;
- (iii) $\overline{\mathcal{M}}^+ = \text{Hull}_1^{\overline{\mathcal{M}}^+}(\overline{\mathcal{M}})$;
- (iv) Either
 - (a) There's a map $\pi: \overline{\mathcal{M}}^+ \rightarrow \mathcal{F}(\mathcal{M})$ in $V[g]$ with $\pi(\overline{\mathcal{M}}) = \mathcal{M}$ and $\pi \upharpoonright (b \cup \{b\}) = \text{id}$ which is Σ_0 -cofinal or Σ_2 -elementary; or
 - (b) There's a model $\mathcal{P} \in \text{dom } \mathcal{F}$ on b with $\mathcal{F}(\mathcal{P}) \in H_\kappa$ and maps $i: \mathcal{F}(\mathcal{P}) \rightarrow \overline{\mathcal{M}}^+$ and $\pi: \overline{\mathcal{M}}^+ \rightarrow \mathcal{F}(\mathcal{M})$, both in $V[g]$ but with their composition in V , with $i(\mathcal{P}) = \overline{\mathcal{M}}$, $\pi(\overline{\mathcal{M}}) = \mathcal{M}$,

$$i \upharpoonright (b \cup \{b\}) = \pi \upharpoonright (b \cup \{b\}) = \text{id},$$

i is Σ_0 -cofinal or Σ_2 elementary, and π is a weak Σ_1 -embedding.

Then $\overline{\mathcal{M}}^+ = \mathcal{F}(\overline{\mathcal{M}}) \in V$. ◦

It turns out that *condenses well* is a bit too strong to do proper core model theory, as shown in [?], and in that paper they propose a technical weakening of this concept which they call *condenses finely*, whose definition is of a similar spirit as the above. We therefore formally require that our desired operators only condense finely, but as all the operators that we will encounter “in the wild” in this thesis condense well, we will omit the definition of fine condensation here.

1.2. MOUSE WITNESS EQUIVALENCE

Definition 1.5. An operator \mathcal{F} **determines itself on generic extensions** if there exists a formula $\varphi(v_0, v_1)$ such that whenever \mathcal{M} is an \mathcal{F} -premouse with

$$\mathcal{M} \models \text{KP} + \ulcorner \text{there are arbitrarily large cardinals} \urcorner,$$

κ is an \mathcal{M} -cardinal and $g \subseteq \text{Col}(\omega, \kappa)$ is \mathcal{M} -generic, then $\text{HC}^{\mathcal{M}[g]}$ is closed under \mathcal{F} and $\mathcal{F} \restriction \text{HC}^{\mathcal{M}[g]} = (\tau_\kappa^{\mathcal{M}})^g$, with $\tau_\kappa^{\mathcal{M}}$ being the unique τ such that $\mathcal{M} \models \varphi[\kappa, \tau]$. \circ

Definition 1.6. Let κ be an infinite cardinal and \mathcal{F} an operator over H_κ on $b \in H_\kappa$. We then say that \mathcal{F} is **radiant**⁴ if \mathcal{F} condenses finely and determines itself on generic extensions. \circ

1.2 MOUSE WITNESS EQUIVALENCE

Definition 1.7. Define coarse (k, U, x) -Woodin pairs

\circ

Definition 1.8. Let \mathcal{F} be a total condensing operator and let α be an ordinal. Then the **coarse mouse witness condition at α with \mathcal{F}** , written $W_\alpha^*(\mathcal{F})$, states that given any scaled-co-scaled $U \subseteq \mathbb{R}$ whose associated sequences of prewellorderings are elements of $\text{Lp}_\alpha^{\mathcal{F}}(\mathbb{R})$, we have for every $k < \omega$ and $x \in \mathbb{R}$ a coarse (k, U, x) -Woodin pair (N, Σ) with $\Sigma \restriction \text{HC} \in \text{Lp}_\alpha^{\mathcal{F}}(\mathbb{R})$. \circ

Check if this is a reasonable definition.

Theorem 1.9 (Hybrid witness equivalence). *Let $\theta > 0$ be a cardinal, $g \subseteq \text{Col}(\omega, < \theta)$ V -generic, $\mathbb{R}^g := \bigcup_{\alpha < \theta} \mathbb{R}^{V[g \restriction \alpha]}$, \mathcal{F} a total radiant operator and α a critical ordinal of $\text{Lp}^{\mathcal{F}}(\mathbb{R}^g)$. Assume that*

$$\text{Lp}^{\mathcal{F}}(\mathbb{R}^g) \models \text{DC} + \ulcorner W_\beta^*(\mathcal{F}) \text{ holds for all } \beta \leq \alpha \urcorner.$$

⁴The terminology is meant to suggest that the operator is preserved when moving in “any direction”: down to smaller models or up to larger forcing extensions.

Then there is a hybrid mouse operator $\mathcal{N} \in V$ on $H_{\aleph_1^{V[g]}}$ such that

$$\text{Lp}^{\mathcal{F}}(\mathbb{R}^g) \models W_{\alpha+1}^*(\mathcal{F}) \quad \text{iff} \quad V \models \ulcorner M_n^{\mathcal{N}} \text{ is total on } H_{\aleph_1^{V[g]}} \text{ for all } n < \omega \urcorner$$

Furthermore, if $\theta < \aleph_1^V$ then we only need to assume that \mathcal{F} is total and condensing.

Be more explicit about what the given operator \mathcal{N} looks like.

1.3 CORE MODELS

Define $K(x)$, $K^{\mathcal{F}}(x)$, $\mathfrak{C}(X)$, \mathcal{Q} -structure

1.4 CORE MODEL DICHOTOMY

Lemma 1.10 (Mesken-N.). *Let θ be a regular uncountable cardinal or $\theta = \infty$ and let \mathcal{N} be a tame hybrid mouse operator on H_θ which relativises well. Then \mathcal{N} is countably iterable iff it's (θ, θ) -iterable, guided by \mathcal{N} . Furthermore, for every $x \in H_\theta$, if $M_1^{\mathcal{N}}(x)$ exists and is countably iterable, then it's also (θ, θ) -iterable, guided by \mathcal{N} .*

Change this to model operators; perhaps change parts of the proof and/or assumptions needed.

PROOF. Fix $x \in H_\theta$ and assume that $\mathcal{N}(x)$ is countable iterable. We first show that $\mathcal{N}(x)$ is (θ, θ) -iterable. Let $\mathcal{T} \in H_\theta$ be a normal tree of limit length on $\mathcal{N}(x)$. Let $\eta \gg \text{rk}(\mathcal{T})$ and let

$$\mathcal{H} := \text{cHull}^{H_\eta}(\{x, \mathcal{N}(x), \mathcal{T}\})$$

with uncollapse $\pi: \mathcal{H} \rightarrow H_\eta$. Set $\bar{a} := \pi^{-1}(a)$ for every $a \in \text{ran } \pi$. Note that $\overline{\mathcal{N}(x)} = \mathcal{N}(\bar{x})$ since \mathcal{N} relativises well. Now $\bar{\mathcal{T}}$ is a normal, countable iteration tree on $\mathcal{N}(\bar{x})$ and hence our iteration strategy yields a wellfounded cofinal branch $\bar{b} \in V$ for $\bar{\mathcal{T}}$. Note that $\bar{\mathcal{Q}} := \mathcal{Q}(\bar{b}, \bar{\mathcal{T}})$ exists, since if \bar{b} drops then there's nothing to do, and otherwise we have that

$$\rho_1(\mathcal{M}_{\bar{b}}^{\bar{\mathcal{T}}}) = \rho_1(\mathcal{N}(\bar{x})) = \text{rk } \bar{x} < \delta(\bar{\mathcal{T}}),$$

1.4. CORE MODEL INDUCTION

so $\delta(\bar{\mathcal{T}})$ is not definably Woodin over $\mathcal{M}_b^{\bar{\mathcal{T}}}$, as there is a definable surjection from $\rho_1(\mathcal{M}_b^{\bar{\mathcal{T}}})$ onto $\delta(\bar{\mathcal{T}})$.

Claim 1.11. $\bar{\mathcal{Q}} \trianglelefteq \mathcal{N}(\mathcal{M}(\bar{\mathcal{T}}))$

PROOF OF CLAIM. If $\bar{\mathcal{Q}} = \mathcal{M}(\bar{\mathcal{T}})$ then the claim is trivial, so assume that $\mathcal{M}(\bar{\mathcal{T}}) \triangleleft \bar{\mathcal{Q}}$. Note that $\bar{\mathcal{Q}} \trianglelefteq M_b^{\bar{\mathcal{T}}}$ by definition of \mathcal{Q} -structures, and that $M_b^{\bar{\mathcal{T}}}$ satisfies (2) of the definition of relativises well, meaning that

$$M_b^{\bar{\mathcal{T}}} \models \ulcorner \forall \eta \forall \zeta > \eta : \text{if } \eta \text{ is a cutpoint then } M_b^{\bar{\mathcal{T}}} \restriction \zeta \not\models \varphi_{\mathcal{N}}[\bar{x}, p_{\mathcal{N}}] \urcorner. \quad (1)$$

This statement is Π_2^1 and $\bar{\mathcal{Q}}$ is Π_2^1 -correct since it contains a Woodin cardinal, so that $\bar{\mathcal{Q}}$ satisfies the statement as well. Since \mathcal{N} is tame we get that $\delta(\bar{\mathcal{T}})$ is a cutpoint of $\bar{\mathcal{Q}}$, so that $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}})) = \mathcal{N}(\bar{\mathcal{Q}} \restriction \delta(\bar{\mathcal{T}}))$ is *not* a proper initial segment of $\bar{\mathcal{Q}}$. Further, as we're assuming that both $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}}))$ and $\mathcal{M}_b^{\bar{\mathcal{T}}}$ are (ω_1+1) -iterable above $\delta(\bar{\mathcal{T}})$ the same thing holds for $\bar{\mathcal{Q}} \trianglelefteq \mathcal{M}_b^{\bar{\mathcal{T}}}$, so that we can compare $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}}))$ with $\bar{\mathcal{Q}}$ (in V). Let

$$(\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}})), \bar{\mathcal{Q}}) \rightsquigarrow (\mathcal{P}, \mathcal{R})$$

be the result of the coiteration. We claim that $\mathcal{R} \trianglelefteq \mathcal{P}$. Suppose $\mathcal{P} \triangleleft \mathcal{R}$. Then there is no drop in $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}})) \rightsquigarrow \mathcal{P}$ and in fact $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}})) = \mathcal{P}$ since $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}}))$ projects to $\delta(\bar{\mathcal{T}})$. Furthermore, as we established that $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}})) = \mathcal{N}(\bar{\mathcal{Q}} \restriction \delta(\bar{\mathcal{T}}))$ isn't a proper initial segment of $\bar{\mathcal{Q}}$ it can't be a proper initial segment of \mathcal{R} either, as the coiteration is above $\delta(\bar{\mathcal{T}})$. But we're assuming that $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}})) = \mathcal{P} \triangleleft \mathcal{R}$, a contradiction. So $\mathcal{R} \trianglelefteq \mathcal{P}$.

Since $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}}))$ and $\bar{\mathcal{Q}}$ agree up to $\delta(\bar{\mathcal{T}})$ and there is no drop $\bar{\mathcal{Q}} \rightsquigarrow \mathcal{R}$ we have that $\bar{\mathcal{Q}} = \mathcal{R}$. If $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}})) \rightsquigarrow \mathcal{P}$ doesn't move either we're done, so assume not. Let F be the first exit extender of $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}}))$ in the coiteration. We have $\text{lh}(F) \leq o(\bar{\mathcal{Q}})$, $\bar{\mathcal{Q}} \trianglelefteq \mathcal{P}$ and $\text{lh}(F)$ is a cardinal in \mathcal{P} .

As $\bar{\mathcal{Q}}$ is $\delta(\bar{\mathcal{T}})$ -sound and projects to $\delta(\bar{\mathcal{T}})$ it follows that $J(\bar{\mathcal{Q}} \restriction \text{lh}(F))$ collapses $\text{lh}(F)$, so it has to be the case that $\bar{\mathcal{Q}} \restriction \text{lh}(F) = \mathcal{P}$ and thus $o(\mathcal{P}) = \text{lh}(F)$. But this means that $\mathcal{P} = \mathcal{N}(\mathcal{M}(\bar{\mathcal{T}}))$ even though we assumed that $\mathcal{N}(\mathcal{M}(\bar{\mathcal{T}})) \rightsquigarrow \mathcal{P}$ moved, a contradiction. \dashv

Now, in a sufficiently large collapsing extension extension of \mathcal{H} , \bar{b} is the unique cofinal, wellfounded branch of $\bar{\mathcal{T}}$ such that $\mathcal{Q}(\bar{b}, \bar{\mathcal{T}}) \trianglelefteq \mathcal{N}(\mathcal{M}(\bar{\mathcal{T}}))$ exists. Hence, by the homogeneity of $\text{Col}(\omega, \theta)$, $\bar{b} \in H$. By elementarity there is a unique cofinal, wellfounded branch b of \mathcal{T} such that $\mathcal{Q}(b, \mathcal{T}) \trianglelefteq \mathcal{N}(\mathcal{M}(\mathcal{T}))$. This proves that M is (uniquely) On-iterable and virtually the same argument yields the iterability of M via successor-many stacks of normal trees.

To show that M is fully iterable, it remains to be seen that the unique iteration strategy (guided by \mathcal{N}) of M outlined above leads to wellfounded direct limits for stacks of normal trees on M of limit length. Let λ be a limit ordinal and $\vec{\mathcal{T}} = (\mathcal{T}_i \mid i < \lambda)$ a stack according to our iteration strategy. Suppose $\lim_{i < \lambda} \mathcal{M}_\infty^{\mathcal{T}_i}$ is illfounded.

Redefine $\eta \gg \text{rk}(\vec{\mathcal{T}})$, $\mathcal{H} := \text{cHull}^{H_\eta}(\{x, M, \vec{\mathcal{T}}\})$ and $\pi : \mathcal{H} \rightarrow H_\eta$ the uncollapse, again with $\bar{a} := \pi^{-1}(a)$ for every $a \in \text{ran } \pi$. By elementarity we get that $\mathcal{H} \models \ulcorner \lim_{i < \bar{\lambda}} \mathcal{M}_\infty^{\bar{\mathcal{T}}_i} \text{ is illfounded} \urcorner$. But $\vec{\mathcal{T}}$ is countable and according to the iteration strategy guided by \mathcal{N} , so that

$$V \models \ulcorner \lim_{i < \bar{\lambda}} \mathcal{M}_\infty^{\bar{\mathcal{T}}_i} \text{ is wellfounded.} \urcorner$$

Now note that $(\lim_{i < \bar{\lambda}} \mathcal{M}_\infty^{\bar{\mathcal{T}}_i})^{\mathcal{H}} = (\lim_{i < \bar{\lambda}} \mathcal{M}_\infty^{\bar{\mathcal{T}}_i})^V$ and wellfoundedness is absolute between \mathcal{H} and V , a contradiction.

Now assume that $M_1^{\mathcal{N}}(x)$ exists for some $x \in H_\theta$, and that it's countably iterable. We then do exactly the same thing as with $\mathcal{N}(x)$ *except* that in the claim we replace (1) with

$$\bar{\mathcal{Q}} \models \forall \eta (\bar{\mathcal{Q}} \restriction \eta \not\models \ulcorner \delta(\bar{\mathcal{T}}) \text{ is not Woodin} \urcorner),$$

so that if $\mathcal{P} \triangleleft \mathcal{R}$ then $\delta(\bar{\mathcal{T}})$ is still Woodin in $\mathcal{P} = \mathcal{N}(\mathcal{M}(\bar{\mathcal{T}}))$, contradicting the defining property of $M_1^{\mathcal{N}}(x)$ (and thus also of \mathcal{R}). The rest of the proof is a copy of the above. \blacksquare

Theorem 1.12 (Hybrid core model dichotomy). *Let θ be a \beth -fixed point or $\theta = \infty$, and let \mathcal{F} be a model operator on H_θ that condenses well. Let $x \in H_\theta$. Then either:*

1.4. CORE MODEL INDUCTION

- (i) The core model $K^{\mathcal{F}}(x)|\theta$ exists and is (θ, θ) -iterable; or
- (ii) $M_1^{\mathcal{F}}(x)$ exists and is (θ, θ) -iterable.

PROOF. Assume first that $K^{c, \mathcal{F}}(x)|\theta$ reaches a premouse which isn't \mathcal{F} -small; let \mathcal{N}_ξ be the first part of the construction witnessing this. Then $\mathfrak{C}(\mathcal{N}_\xi) = M_1^{\mathcal{F}}(x)$, and by Lemma 1.10 it suffices to show that $M_1^{\mathcal{F}}(x)$ is countably iterable.

Insert argument?

Show that $M_1^{\mathcal{F}}(x)$ is countably iterable.

We can thus assume that $K^{c, \mathcal{F}}(x)|\theta$ is \mathcal{F} -small. Note that if $K^{c, \mathcal{F}}(x)|\theta$ has a Woodin cardinal then because the model is \mathcal{F} -closed we contradict \mathcal{F} -smallness, so the model has no Woodin cardinals either, making it (θ, θ) -iterable.

Let $\kappa < \theta$ be any uncountable cardinal and let $\Omega := \beth_\kappa(\kappa)^+$. Note that $\Omega < \theta$ since we assumed that θ is a \beth -fixed point and $\kappa < \theta$. If Ω is a limit cardinal in $K^{c, \mathcal{F}}(x)|\theta$ then let $\mathcal{S} := \text{Lp}(K^{c, \mathcal{F}}(x)|\Omega)$ and otherwise let $\mathcal{S} := K^{c, \mathcal{F}}(x)|\Omega$. Then by Lemma 3.3 of [?] we get that \mathcal{S} is countably iterable, with largest cardinal Ω in the “limit cardinal case”.

This also means that Ω isn't Woodin in $L[\mathcal{S}]$, as it's trivial in the case where Ω is a successor cardinal of $K^{c, \mathcal{F}}(x)|\theta$ by our case assumption, and in the “limit cardinal case” it also holds since

$$K^{c, \mathcal{F}}(x)|\Omega^{+K^{c, \mathcal{F}}(x)|\theta} \subseteq \mathcal{S}.$$

By [?] and [?] this means that we can build $K^{\mathcal{F}}(x)|\kappa$, as the only places they use that there's no inner model with a Woodin are to guarantee that $K^{c, \mathcal{F}}(x)|\Omega$ exists and has no Woodin cardinals, and in Lemma 4.27 of [?] in which they only require that Ω isn't Woodin in $L[\mathcal{S}]$.

As $\kappa < \theta$ was arbitrary we then get that $K^{\mathcal{F}}(x)|\theta$ exists. Note that $K^{\mathcal{F}}(x)|\theta$ has no Woodin cardinals either and is \mathcal{F} -small, so that \mathcal{Q} -structures trivially exist, making it (θ, θ) -iterable. ■