

Explanation of "On Two Minimax Theorems in Graph"

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1 Theorems

THEOREM 1 (Edmonds). *The maximum number of edge-disjoint branchings (rooted at a) equals the minimum number of edges in a -cuts.*

2 Definitions

Branching and a-cut (from the paper)

branching In a directed graph $D = (V, A)$ a branching is a set of arcs not containing circuits s.t. each node of D is entered at most by one arc in A' . (So the arc in A' make up a forest).

a-cut A a -cut of G determined by a set $S \subset V(G)$ is the set of edges going from S to $V(G) - S$. It will be denoted by $\Delta_G(S)$. We also set that $\delta_G(S) = |\Delta_G(S)|$.

Branching and a-cut (from books)

branching A branching B is an arc set in a digraph D that is a forest such that every node of D is the head of at most one arc of B . A branching that is a tree is called an arborescence. A branching that is a spanning tree of D is called a spanning arborescence of D . Clearly, in a spanning arborescence B of D every node of D is the head of one arc of B except for one node. This node is called the root of B . If r is the root of arborescence B we also say that B is rooted at r or r -rooted.

r-arborescence An arborescence in a digraph $D = (V, A)$ is a set A' of arcs making up a spanning tree such that each node of D is entered by at most one arc in A' . It follows that there is exactly one node r that is not entered by any arc of A' . This node is called the root of A' , and A' is called rooted in r , or an r -arborescence.

r-cut An r -cut in a digraph is an edge set $\delta^+(X)$ for some $X \subset V(G)$ with $r \in X$.

edge-disjoint branchings Two branchings are edge-disjoint if they do not have any internal edge in common.

A note on orientation A rooted tree is a tree in which one vertex has been designated the root. The edges of a rooted tree can be assigned a natural orientation, either away from or towards the root, in which case the structure becomes a directed rooted tree. When a directed rooted tree has an orientation away from the root, it is called an arborescence, branching, or out-tree; when it has an orientation towards the root, it is called an anti-arborescence or in-tree.

3 Examples

3.1 Branching

3.1.1 Simple branching

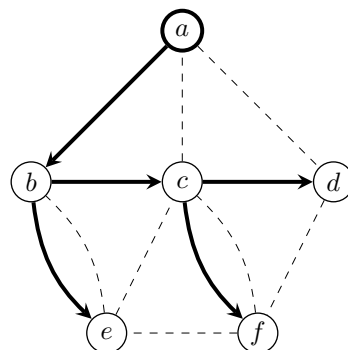
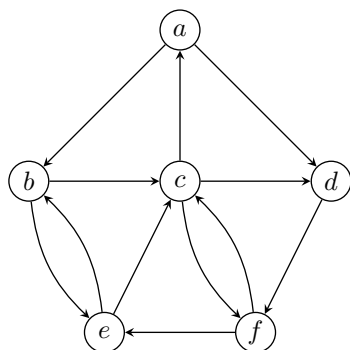


Figure 1: A digraph $D = (A, V)$ Figure 2: A branching of D routed at a .

3.1.2 Edge disjoint branching

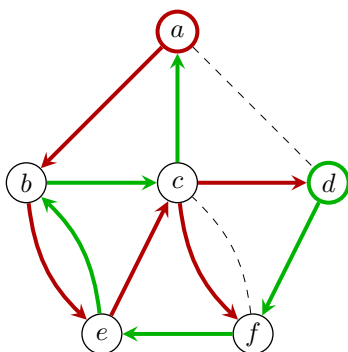


Figure 3: Two edge-disjoint branching routed respectively at a and d .

3.1.3 Edge disjoint a -routed branching

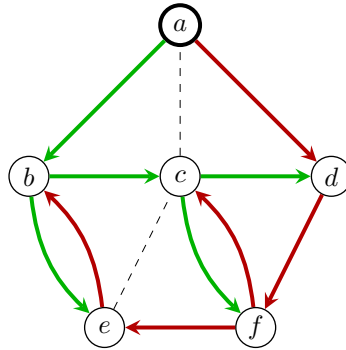


Figure 4: Two edge-disjoint a -routed branching.

3.2 a -cut exemples

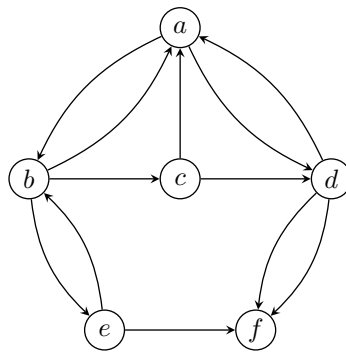


Figure 5: A digraph $D = (A, V)$.

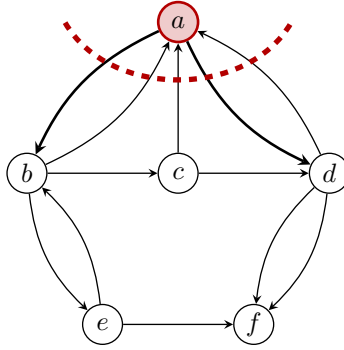


Figure 6: An a-cut of D ,
 $S = \{a\}$,
 $\Delta_D(S) = \{ab, ad\}$,
 $\delta_D(S) = 2$.

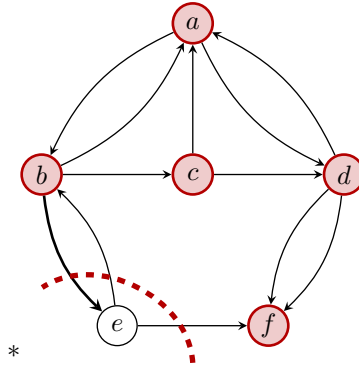


Figure 7: An minimal a-cut of D ,
 $S = \{a, b, c, d, f\}$,
 $\Delta_D(S) = \{be\}$,
 $\delta_D(S) = 1$.

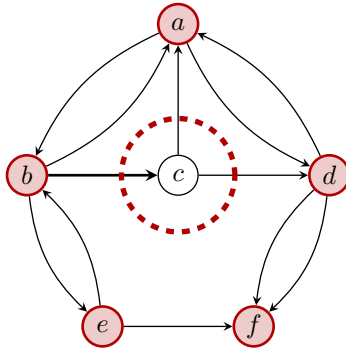


Figure 8: An minimal a-cut of D ,
 $S = \{a, b, d, e, f\}$,
 $\Delta_D(S) = \{bc\}$,
 $\delta_D(S) = 1$.