

Name: Sabbir Ahmed

Section: 02

HW #: 2

Version: A

Username: sabbir1

**1. Prove  $2^n > n^2$  for every positive integer  $n$  when  $n \geq 5$ .**

Let  $P(n)$  represent the proposition  $2^n > n^2$ ,  $n \in \mathbb{N} \mid n \geq 5$ .

Let  $n = 5$ , so:

$$\rightarrow 2^5 > 5^2$$

$$\rightarrow 32 > 25$$

$\rightarrow P(n)$  to be true for  $n = 5$ .

Assume  $P(k)$  holds for all  $k \in \mathbb{N} \mid k \geq 5$ . So  $2^k > k^2$ .

Consider  $k+1$ . So:

$$\rightarrow 2^{k+1} > (k+1)^2$$

$$\equiv 2 \cdot 2^k > k^2 + 2k + 1$$

$$\equiv 2^k + 2^k > k^2 + 2k + 1$$

$$\text{But } 2^k > k^2$$

$$\therefore 2^{k+1} > (k+1)^2.$$

$2^n > n^2$  holds for all  $n \in \mathbb{N} \mid n \geq 5$ .

**2. Prove  $4^n - 1$  is divisible by 3 for all  $n \in \mathbb{N}$ .**

Let  $P(n)$  represent the proposition  $(4^n - 1) = 3k$  for all  $n \in \mathbb{N}$ .

Let  $n = 1$ . So:

$$\rightarrow 4^1 - 1$$

$$\rightarrow 3, \text{ which is divisible by 3.}$$

$$\rightarrow P(n) \text{ holds for } n = 1.$$

Assume  $P(a)$  holds for all  $a \in \mathbb{N}$ . So  $(4^a - 1) = 3k$ .

Consider  $a+1$ . So:

$$\rightarrow (4^{a+1} - 1)$$

$$\rightarrow (4 \cdot 4^a - 1)$$

$$\text{But } (4^a - 1) = 3k$$

$$\begin{aligned}
 \text{So, } (4^{a+1} - 1) &\equiv (4(3k + 1) - 1) \\
 &\rightarrow (4 \cdot 3k + 4 - 1) \equiv (12k + 3) \equiv 3(4k + 1) \\
 \therefore (4^{a+1} - 1) &\text{ is divisible by 3.}
 \end{aligned}$$

$$(4^n - 1) = 3k \text{ holds for all } n \in \mathbb{N}.$$

$$\text{3. Prove } \frac{1}{1 \cdot 2} + \dots + \frac{1}{n \cdot (n+1)} = \frac{n}{n+1}$$

$$\text{Let } P(n) \text{ represent the proposition } \frac{1}{1 \cdot 2} + \dots + \frac{1}{n \cdot (n+1)} = \frac{n}{n+1}$$

Let  $n = 1$ . So:

$$\begin{aligned}
 &\rightarrow \frac{1}{1 \cdot (1+1)} = \frac{1}{1+1} \\
 &\rightarrow \frac{1}{2} = \frac{1}{2} \\
 &\rightarrow P(n) \text{ holds for } n = 1.
 \end{aligned}$$

$$\text{Assume } P(k) \text{ holds for all } k \in \mathbb{N}. \text{ So } \frac{1}{1 \cdot 2} + \dots + \frac{1}{k \cdot (k+1)} = \frac{k}{k+1}$$

Consider  $k+1$ . Then:

$$\begin{aligned}
 &\rightarrow \frac{1}{1 \cdot 2} + \dots + \frac{1}{k \cdot (k+1)} + \frac{1}{(k+1) \cdot (k+2)} = \frac{k}{k+1} + \frac{1}{(k+1) \cdot (k+2)} \\
 &\qquad \qquad \qquad \rightarrow \frac{k}{k+1} + \frac{1}{(k^2 + 3k + 2)} \\
 &\qquad \qquad \qquad \rightarrow \frac{k+1}{k+2} \equiv \frac{(k+1)}{(k+1)+1} \\
 \therefore \frac{1}{1 \cdot 2} + \dots + \frac{1}{k \cdot (k+1)} + \frac{1}{(k+1) \cdot (k+2)} &= \frac{(k+1)}{(k+1)+1}
 \end{aligned}$$

$$\frac{1}{1 \cdot 2} + \dots + \frac{1}{n \cdot (n+1)} = \frac{n}{n+1} \text{ holds for all } n \in \mathbb{N}.$$

4.

Code	Cost	Times	Total
<pre>int sum = 0; for (int i = 0; i &lt; n; i++)     for (int j = 0; j &lt; n; j++)         ++sum;</pre>	1	1	1
	1 (int i) + n+1 (<) + n (++)	1	2n+2
	1 (int j) + n+1 (<) + n (++)	n	2n <sup>2</sup> +2n
	2	n*n	2n <sup>2</sup>
<b>FINAL</b>	<b>4n<sup>2</sup>+4n+3 = O(n<sup>2</sup>)</b>		

5.

Code	Cost	Times	Total
<pre>int sum = 0; for (int i = 0; i &lt; n; i += 2 )     for (int j = 0; j &lt; n; j++ )         ++sum;</pre>	1	1	1
	1 (int i) + n+1 (<) + n/2 (+= 2)	1	n+n/2+2
	1 (int j) + n+1 (<) + n (++)	n/2	n <sup>2</sup> +n
	2	n*n/2	n <sup>2</sup>
<b>FINAL</b>	<b>2n<sup>2</sup>+2n+n/2+3 = O(n<sup>2</sup>)</b>		

6.

Code	Cost	Times	Total
<pre>int sum = 0; for (int i = 1; i &lt; n; i *= 2 )     for (int j = 0; j &lt; n; j++ )         ++sum;</pre>	1	1	1
	1 (int i) + n+1 (<) + log <sub>2</sub> (n) (*= 2)	1	n+log <sub>2</sub> (n)+2
	1 (int j) + n+1 (<) + n (++)	log <sub>2</sub> (n)	2nlog <sub>2</sub> (n)+2log <sub>2</sub> (n)
	2	n*log <sub>2</sub> (n)	2nlog <sub>2</sub> (n)
<b>FINAL</b>	<b>4nlog<sub>2</sub>(n)+3log<sub>2</sub>(n)+n+3 = O(nlog<sub>2</sub>(n))</b>		

7.

Code	Cost	Times	Total
<pre>int sum = 0; for (int i = 0; i &lt; n; i++ )     for (int j = 0; j &lt; i * i; j++ )         for (int k = 0; k &lt; j; k++ )             ++sum;</pre>	1	1	1
	1 (int i) + n+1 (<) + n (++)	1	2n+2
	1 (int j) + n <sup>2</sup> +1 (< i*i) + n (++)	n	n <sup>3</sup> +n <sup>2</sup> +2n
	1 (int k) + n <sup>2</sup> +1 (< j) + n (++)	n <sup>3</sup>	n <sup>5</sup> +n <sup>3</sup> +2n
	2	n <sup>5</sup>	2n <sup>5</sup>
<b>FINAL</b>	<b>3n<sup>5</sup>+2n<sup>3</sup>+n<sup>2</sup>+6n+3 = O(n<sup>5</sup>)</b>		