407 * Funcs and relations: A, B sets A × B = { (a,b): a ∈ A, b ∈ B} any R ⊆ A × B is a rel forom A → B. >P(X)=2= {S: \$CX} obinary rel from A to B if A=B

R is a binary rel on A > Fines: f: A → B is a function from A to B iff f is a rel. s.t. if a ∈ A at most one b ∈ B has (a,b) ⊆ f for any a EA at most one b EB

a call if f(a) f= > (a,f(a)); a ED} > Domain (f) = {a ∈ A: (a,b) ∈ f some b} = 3a ∈ A: Here is some b = f(a) in B} 3 Domain (R) = {aithere is b &B (a,b) CR}

> '=' = \$ (A,B): A S X, B S X, A=B3 Let Abeaset SCA define 1, on S by 1, (s) = S all s & S. Dom (1,) = S *Def! f: A > B is 1-1 (injective) iff a, az in A are equal if f(a) = f(az) * Indicator funcs: SSA

Is (a)= SI if a ES

O if a & S Ig ∈ {0,1} = {f:f:A → {0,1}}, BA domain (f)=A *Def: f: A > B is a bijection or 1-1 correspondence
iff f is 1-1 onto domain (f) = A

S-> Is is a bisection from P(A) to \subsection 13 A= 2A If f ∈ §0, 13 A s ∈ f, Sp= §a ∈ A: f(a)=13 * Comp. of funcs: A,B,C sets f:A→B,g:B→C funcs: gof(a)= g(f(a)) YaEA defines gof: A>C (a,b) in graph (f) (=f) (b,c) in graph (g) (=g) implies that (a,c) in graph (gof) C= (q of)(a) Relation from A to B S " B to C * (a,c) is in SOR iff FLEB s.t. (a,b) ER, (b,c) ES a(SOR)c arb bsc -> Let A, B, C be Z, R, S, be >. If a)b, b)c => a(>0>)c = a>>c asb bsc azb+1 b2c+1 i. a / c+7 => a (>0>)c

* Thm: Comp. of funcs. is associative. > Let f: A > B, g: B > C, L: C > D be funes. Then ho (gof) = (hog) of. Pf: Let d=ho(gof)(a), (a,d) & ho(gof).

Fc & C, d=h(c), Fc: (c,d) & ho (gof).

(gof)(a)=c.

Fb & b & ceg(b)

b & f(a)

Fb & (a,b) & f, (b,c) & g deh(c), c=g(b)

i. d=(hog)(b)

also f(a):b

i. d=(hog)(f(a))

=(hog)of)(a) (b, d) (hog) (a,d) E (hog) of (a,d) E ho (gof) If d= (hog) of)(a) => d=(ho(gof)(a) @ I dend ify funes: $\Rightarrow f: A \rightarrow B \Rightarrow f \circ 1_A = f, 1_B \circ f = f$ (f. 1A)(a) = f(1A(a)) = f(a) > if A=B: 1, of -fo1, =f any f: A > A *Thm: f: A > B, g: B > C i) f, g are inj = g of is inj ii) f, g are sur => g of is sur. iii) f, g are bij => g of is bij.

~

コート・トード・

iv) if (gof) is sur. => g is sw. v) if (gof) is ing. and range (f) Edomain(g) this is g(b) for b=f(a) DInv. funcs: f: A > B =) g of = 1 A or f og = 1 B gleft inv gright inv If f(x)=y => x=g(y) (y) (y) (y) Let RCA×B Las R'CBXA iff R'= {(b,a): (a,b) ER} ograph(f) = graph(g)

* gof = 1A, 1A is 1-1 so f is 1-1

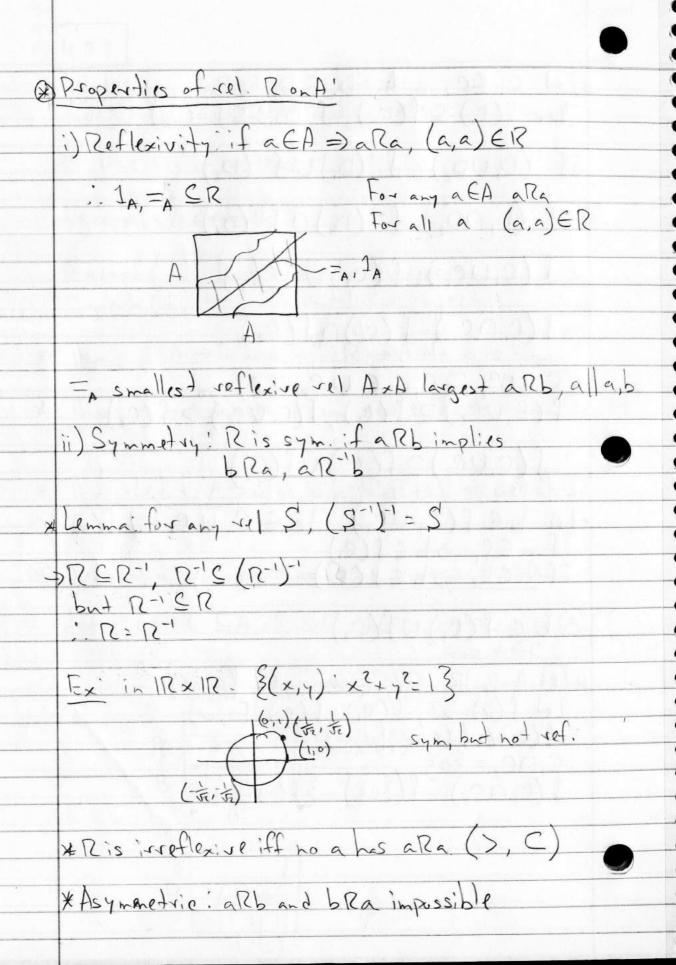
fog = 1B, 1B is onto so f in onto : fis big if f'exists -> graph (3) = draph (3)-1 => graph (f) has vertical line property
graph (f) has horizontal line property * Thmi f: A > B has f-1: B > A iff f is big. * Thm: f has inverse from range (f) to domain (f) iff * Thr: Let f: A > B be big. h: B > C as well (hof) Direct and inv. images under a fune f: A >B def: Let CCA, DCB. f(c)= {yEB: yef(x) Some x EC3 f-'(D) = {x EA : f(x) ED} Edual (f)

*Let C, EC, in A, D, ED, in B Then f(C,) Ef(C,), f'(D,) Ef-'(D,) 125-1(D, UD,) = 5-1(D,) U5-1(D,) >f-1(D, DD,)=f-1(D,) Df-1(D2) > f (C, UC2)= f(C,) Uf(C2) $\rightarrow f(C'\cup C^s) \cdot f(C')\cup f(C^s)$ $f(C' \cap C') = f(C') + f(C' \cap C') = f(C')$: f(c'ncs) 5 f(c') n f(cs) · Let bef(C, UC2). Fac C, UC2 w/f(a)=b. If a c () b c f (c)

If a c () b c f (c) : b e f (c,) u f (c2) * Let A=B=|R, C,= (-0,0], Cz=[0,0) Let f(x)= x2, f(c,): f(c,): [0, 0) : f(c,) \(f(c_2) = [0, \infty) \) $f(C, C_2 = {03}) = {f({03})} = {f({0})} = {03}$

1

カイナル



iii) Transituity: Ris transitive iff when aRb and bRC then a Rc If a (ROR) c then a RC. i. ROR SR @ Partial order: (Strong)
i) Transitive
+ ii) Irreflexive Ex. >, C Note if alb and blather a=b by trans. * A total complete linear partial order is a partial order which satisfies for an a 76 either alb or bla * Trichotomy: for a, b EA precisely one: aRb, a=b, bRa. * ex. A set P(A) subsets {0,13A funcs f: A > {0,13 SEP(A) => Is(a)= SI aES Is E {0,13} : define Ss = {a: f(a)=13 = support (f) 1: S → Is S: f → S; => S = I-1 (inverse)

-0

コココココココ

•
$$A = N \Rightarrow |P(N)| = |\{0, 13^N | = 2^{|N|} = 2^{N_0} \} S_0$$

*Weak partial order:

i) Transitive

ii) Reflexive

(May impose: aRb and bRa implies a=b)

Ex. Z, S

* Egynivalence rel:

i) Reflexive
ii) Symmetric
iii) Transitive

· Ex. = A XA

Ex. Congruence of triangles in plane

T, RTz, Tz R T3 then T, RT3

T, RTz iff area (T,) = area (Tz) olet f: A >B be a func w/domain(A)=A

a, ~g az iff f(a,)=f(az) Reflexive: a ~ a , f (a) cf(a) Symmetric: angb, f(a)=f(b) $\begin{array}{c}
\text{i. } f(b) = f(a) \\
\text{o. } b \sim f(a)
\end{array}$ Transitive: Let a ryb and bry c
then f(a)=f(b) and f(b)=f(c)
: f(a)=f(c) or ary c DEgnivalence Classes: Requivalence Relation on A a EA [a] = {b \in A: a \b} equivalence class of a modulo ~ · Let b ∈ [a]. ("a ∈ [a].)

a ∈ [b]. since and implies b ~a Assort [b] = [a] Let ce[b] sobre. But arb

-

73

-3

7

ついいい

-1

--

- 4

->

ナナナナナファ

[a] < [b], [a] = [b]Thm: Let ~ be an equivalence relion A

A/N= {[a]_n: a EA3 (factor set) forms a partition of A. i) A = U & [a] ~ : a ∈ A} ii) [a] ~ ∩ [b] ~ ≠ Ø => [a] ~ = [b] ~ Pf. let ce [a]~ U[p]~ then [a]~ [c]~ i. Using trans of equality: [a] = [b]~ *If m is an equivalence relation on A => \{b: anb} = [a]n \{[a]m: a\in A} Partitions A f:A >B, domain (f)=A => a, ~, az; ff f(a)=f(a) *[a] = f-1(b) if b= f(a) Sf-'(b): b (B) is factor set of all equivalence classes *P= §A: i E] be a partition of A. Let flace if a EAi, f: A > I, f'(i)=Ai if f(a)=i, [a] = Ai

@ A set theoretical model of Zt or N: Def: A is a set => S(A) = A+1 = AU & A } Def: $0 = \emptyset \Rightarrow S(0) = S(\emptyset)$ $\Rightarrow 1 = 0 + 1 = \sum \emptyset$ $\Rightarrow 2 = 1 + 1 = S(1) = \sum \emptyset$ $0 = \sum \emptyset$ If h= \{0,1,2,..,h-1} then h+1: S(x)= \{0,1,2,..,h-1} U 30, 1, 7, .., h-13 = {0,1,2,--, n-13 U {n}} = {0,1,2,..,n-1,n} Def: A collection C of sets is inductive iff AEC implies Zt is smallest inductive set containing O. N=Zt \ 203 is smallest inductive sed containing 1. @ Principle of induction: ACZ+ is Z+ iff i) OEA and ii) if nEA => N+1 EA Def : A sequence X is a function on Z+ (or N) to some * Proof by induction: Let (Pn) be a sequence of logical
propositions. If i) Po is true all kin ii) Ph is true inplies Phase drue then all Produce

しょしょしゅつ

-)

-)

* Definition by induction: If i) x, is defined and given definition of x then xn+1 is defined, then (xn) is well defined * Let of be a func. from A -> A (w/ domain A) f=1, f=f, f=fof Un∈ N let fn+1=fnof Thm: $f_1:A_1 \rightarrow A_2$, $f_2:A_3 \rightarrow A_3$, ..., $f_n:A_n \rightarrow A_n$, $f_n:A_n \rightarrow A_n$) $\Rightarrow ((f_n \circ f_{n-1}) \circ ((f_3 \circ f_e) \circ f_1)$ => (fro (fro (100 o (f30 (f20f,)))) · · · f3 o (f2 of1) = (f3 of2) of1 Pf: Suppose true for n funes.

Look at (n+1) funes.

(from ofjord) o (f. o...of,) =>(fn+1 o (fn o.. ofj+1)) o (f; o.. o (f, of,)) => fro (fro (fro (fro (fro fin)) o (fro (fro fin))) =>f, o (f, o. o (f, of, o. o (f, of,))) * Lemma: frol=frof = 10 (12-10t) = (20,10-10t) = (20,10-10t)

Thmifn+k = frofk
=fkofn Def: n= Sr(0) Def! htk= Sntk(o) (addition of two integers)

(associative, commutative) * m > h iff m= Sk(h), for some k EZ+ * m>n (translituity) * (m+1)-n = month (recursive def. of muldiplication) Def: m, n & Z+ m > n iff menth, for k \ Z+

m > n iff menth, for k \ N Thr: If A(# Ø) \(\int ZL+ then min (A) exists Pf: Let L= lowerbds of A st. if LEL, aEA, LEa OEL Assume min (A) does not exist. Thus no lelis in A

407 If a EA, I La. Thus I+1 & a any l & I Thus 1+1 ∈ 4 : 4 = Z+ (contradiction) Covi Every lower bold subset A (+0) of Z has a min Pf: If I lower bd A-l= Za-l: a EAZ = Z+ has a m+l is lover bd of A Def N= 21, ..., n3 initial segment of N tail or limb seg of Zt or Nif n)0 Def. A set A # Ø is finite iff In E N W/ Nn in 1-1 correspondence w/A. In this case, we set |A|= coud(A) = XA=n = \{0,1,2,...,n-1\} Thm: Let A be an infinite set. There is a non-repeating sequence is A. Pf Let A = A. Choose a, EA, Set Az=A, \ \ \(2a, \) \(\frac{1}{2} \) else A = }a,3

Suppose distinct a, az, ..., an have been chosen from A Ar # p, else A = {a,, ..., a, -13 = Nr. Pick an EAn Induction step Cos An infinites in 1-1 corresp. w/ a proper subset Lof f: A → A f(ai)= ain, i >1 f(A)= a, ta=ai, anyi f(A)= A\ {a} The There is no 1-1 corresp. bother No and a proper subset. Pf n=1 {13 => Ø (trivial base case) Assume tine n=k. Let f: Nk+1 > A = Nk+1 be 1-1 corresp. i) f(k+1)= k+1 some j Ek in Nx A Let g(i) = f(i) if 1 & i & k Then g: Nx > (Nx \ 3j3) (contradiction i. case i) is imp. ii) f (k+1) & k Let i: k+1 Let g (i)= k+1 q (h+1)=i g(m)= m, all others g: Nkm > Nkm

79999999

....

Look at gof: Nx > Nx.
LI corresp. both Nx and range (gof) = g(A) 3(5) & g.f(NK) : gof 1-1 corresp. from Nk+1 to g(A) # Nk+1
gof (k+1) = k+1 (back to case; contradiction) O Cor: If m <n there is no big. f: Nn > Nm @ Cor Als well defined unique n w/A, N in 1-1 musp.

N = A + N N 30 N There is no surjection from N to Nm if m < n

If |A| = |B| = k in N and f: A > B i) fishin ii) f is inj

m, n integers mis a moldiple of m)

iff $\exists k \in \mathbb{Z} \text{ s.t. } m \circ k = n$ Rephrased n Em Z = {mk: k ∈ Z} mln iff (-m) | n iff m | (-n) Case O: O.Z={Ok: k ∈ Z}={0} O·k if k #0 k·0=0 soklo +k Case 1: 1- Z= {1-k: k < Z3 1/k all k. But k | liff k= ±1 * Notes divisibility is weak partial order (autisymmetric) i) n/n all n
ii) n/m and m/n implies n=m If m | n then mk= n some k | m | | k | = | n If min EZt, so is k iii) llm, mln implies lln
m=lk', n=mk
n=l(k'k) * Note: Suppose mln, mlnz and that Ear, az} in Z then m (a, h, + azhe) 1 linear lintegral combination

mZ "subspace" of Z n=mk, n=mke (a,n,+azne)=a,(k,m)+ae(kem)=(a,k,+aeke)m Th 1.1.3 (Division Algorithm): Let a EZ, b EN. There are unique into ay (quotient) or (remainder) w/ a=bytr, O Ex 2b Pf: (Uniqueness) a=bgi+r', suppose r' &r pata = pata, Thus, v-v'=b(q-q') 0 = 1-1, Tp-0=p (contradiction) (contradiction) : q'-q=0=, q'=q=) x=x' Let R= 3 = a-by: + 30, q = 23 a-b (-lal) &G a+blal >B az-la/z-bla/ : (r=a-b(-lal)>0) ER(+Ø) SZ+ Let remin (R), a=bgy+r 17P3 It mot ~ 3p x = x, +p a=by+(x'+b) = b (q+1)+r

But 1' (1=1'+b, 1') OER (contradiction) : 03x 6p BACHURANUT WARRED