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Lab-Report

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Objectives : The main objectives of this lab how to know Sockets provide the communication mechanism between two computers using TCP and java.net.ServerSocket class provides a mechanism for the server program to listen for clients and establish connections with them.

Theory : Sockets provide the communication mechanism between two computers using TCP. A client program creates a socket on its end of the communication and attempts to connect that socket to a server. The java.net.Socket class represents a socket, and the java.net.ServerSocket class provides a mechanism for the server program to listen for clients and establish connections with them.

Methodology :

1) Briefly explain the term IPC in terms of TCP/IP communication.

Ans :

The answer will be different each Operating System. Unix offers System V IPC, which gives you message queues, shared memory, and semaphores . InterProcess Communication is a term we use for interactions between two processes on the same host. William Westlake mentions TCP/IP, which is used for interactions with another host. It can be used locally as well, but it is relatively inefficient. Unix domain sockets are used in the same way, but are only for local use and a bit more efficient. Message queues are easy to use: processes and threads can send variable-sized messages by appending them to some queue and others can receive messages from them so that each message is received at most once. Those operations can be blocking or non-blocking. The difference with UDP is that messages are received in the same order as they are sent. Pipes are simpler, but pass a stream of data instead of distinct messages. Line feeds can be used as delimiters. Shared memory allows different processes to share fixed regions of memory in the same way that threads have access to the same memory. Unix allocates a certain amount of

memory after which it can be accessed like private memory. A similar mechanism is the memory-mapped file: the difference is that the memory segment is initialised from a disc file and changes can be permanent. The size of a file can change, which complicates shared files. Since two threads updating the same data can lead to inconsistencies, semaphores can be used to achieve mutual exclusion.

2) What is the maximum size of a UDP datagram? What are the implications of using a packet-based Protocol as opposed to a stream protocol for transfer of large files?

Ans :

This size is the theoretical maximum size of UDP Datagram, in practice though, this limit is further constrained by the MTU of data-link layer(which varies for each data-link layer technology, but cannot be less than 576 bytes), considering that, maximum size of UDP datagram can be further calculated as (for IPv4):

$576 \text{ bytes} - 20 \text{ bytes (IP header)} = 556$ (including 8 bytes UDP header)

It depends on the underlying protocol i.e., whether you are using IPv4 or IPv6

>> In IPv6, the maximum length of packet size allowed is 64 kB, so, you can have UDP datagram of size greater than that.

>> In IPv4, the maximum length of packet size is 65,536. So, for UDP datagram you have maximum data length as: $65,535 \text{ bytes} - 20 \text{ bytes (Size of IP header)} = 65,515$ bytes (including 8 bytes UDP header)

3) TCP is a reliable transport protocol, briefly explain what techniques are used to provide this reliability.

Ans:

A number of mechanisms help provide the reliability TCP . Each of these is described briefly below:

Duplicate data detection : It is possible for packets to be duplicated in packet switched network; therefore TCP keeps track of bytes received in order to discard duplicate copies of data that has already been received.

Retransmissions: In order to guarantee delivery of data, TCP must implement retransmission schemes for data that may be lost or damaged. The use of positive acknowledgements by the receiver to the sender confirms successful reception of data. The lack of positive acknowledgements, coupled with a timeout period calls for a retransmission.

Sequencing : In packet switched networks, it is possible for packets to be delivered out of order. It is TCP's job to properly sequence segments it receives so it can deliver the byte stream data to an application in order.

Timers: TCP maintains various static and dynamic timers on data sent. The sending TCP waits for the receiver to reply with an acknowledgement within a bounded length of time. If the timer expires before receiving an acknowledgement, the sender can retransmit the segment.

4) Why are the htons(), htonl(), ntohs(), ntohl() functions used?

Ans:

These are used for:

htons() host to network short

htonl() host to network long

ntohs() network to host short

ntohl() network to host long.

5) What is the difference between a datagram socket and a stream socket?

Ans:

The difference is given below :

Stream socket

A stream socket provides a bidirectional, reliable, sequenced, and unduplicated flow of data with no record boundaries. Stream sockets enable processes to communicate using TCP .After the connection has been established, data can be read from and written to these sockets as a byte stream. The socket type is SOCK_STREAM.

Datagram socket

A datagram socket supports a bidirectional flow of messages. Datagram sockets enable processes to use UDP to communicate. A process on a datagram socket might receive messages in a different order from the sending sequence. A process on a datagram socket might receive duplicate messages. Messages that are sent over a datagram socket might be dropped. The socket type is SOCK_DGRAM. Record boundaries in the data are preserved.

Daytime Protocol implementation:

Server Code :

''' Created on Feb 20, 2021

@author: Sabikun Nahar Piya '''

Server Program:

```
package je3.nio;

import java.nio;

import java.nio.channels;

import java.nio.charset;

import java.net;

import java.util.logging; import java.util;

/**
 * A more robust daytime service that handles TCP and UDP connections and *
 * provides exception handling and error logging.
 */

public class DaytimeServer {

    public static void main(String args[ ]){

        try {

            // Handle startup exceptions at the end of this block

            // Get an encoder for converting strings to bytes

            CharsetEncoder encoder = Charset.forName("US-ASCII").newEncoder( );

            // Allow an alternative port for testing with non-root accounts  int port = 13;

            // RFC867 specifies this port.  if (args.length > 0) port = Integer.parseInt(args[0]);

            // The port we'll listen on  SocketAddress localport = new
            InetSocketAddress(port);
```

```
// Create and bind a TCP channel to listen for connections on.
ServerSocketChannel tcpserver = ServerSocketChannel.open( ); tcpserver.socket(
).bind(localport);

// Also create and bind a DatagramChannel to listen on. DatagramChannel
udpserver = DatagramChannel.open( ); udpserver.socket( ).bind(localport);

// Specify non-blocking mode for both channels, since our

// Selector object will be doing the blocking for us.
tcpserver.configureBlocking(false); udpserver.configureBlocking(false);

// The Selector object is what allows us to block while waiting

// for activity on either of the two channels. Selector selector = Selector.open( );
// Register the channels with the selector, and specify what

// conditions (a connection ready to accept, a datagram ready

// to read) we'd like the Selector to wake up for.

// These methods return SelectionKey objects, which we don't

// need to retain in this example. tcpserver.register(selector,
SelectionKey.OP_ACCEPT); udpserver.register(selector,
SelectionKey.OP_READ);

// This is an empty byte buffer to receive empty datagrams with.

// If a datagram overflows the receive buffer size, the extra bytes

// are automatically discarded, so we don't have to worry about

// buffer overflow attacks here. ByteBuffer receiveBuffer =
ByteBuffer.allocate(0);
```

```

// Now loop forever, processing client connections for() {

try {

// Handle per-connection problems below

// Wait for a client to connect selector.select( );

// If we get here, a client has probably connected, so

// put our response into a ByteBuffer. String date = new java.util.Date( ).toString(
) + "\r\n"; ByteBuffer response=encoder.encode(CharBuffer.wrap(date));

// Get the SelectionKey objects for the channels that have // activity on them.
These are the keys returned by the

// register( ) methods above. They are returned in a

// java.util.Set. Set keys = selector.selectedKeys( );

// Iterate through the Set of keys. for(Iterator i = keys.iterator( ); i.hasNext( ); ) {
// Get a key from the set, and remove it from the set

SelectionKey key = (SelectionKey)i.next( ); i.remove( );

// Get the channel associated with the key Channel c = (Channel) key.channel( );
// Now test the key and the channel to find out

// whether something happened on the TCP or UDP channel if (key.isAcceptable(
) && c == tcpserver) {

// A client has attempted to connect via TCP.

// Accept the connection now. SocketChannel client = tcpserver.accept( );

// If we accepted the connection successfully,

```



```

// then send our response back to the client.  if (client != null) {
client.write(response);

// send response  client.close( );

// close connection  } } else if (key.isReadable( ) && c == udpserver) {

// A UDP datagram is waiting. Receive it now,

// noting the address it was sent from.  SocketAddress clientAddress =
udpserver.receive(receiveBuffer);

// If we got the datagram successfully, send

// the date and time in a response packet.  if (clientAddress != null)
udpserver.send(response, clientAddress); } } } catch(java.io.IOException e) {

// This is a (hopefully transient) problem with a single

// connection: we log the error, but continue running.

// We use our classname for the logger so that a sysadmin // can configure logging
for this server independently

// of other programs.  Logger l =
Logger.getLogger(DaytimeServer.class.getName( )); l.log(Level.WARNING,
"IOException in DaytimeServer", e); } catch(Throwable t) {

// If anything else goes wrong (out of memory, for example),

// then log the problem and exit.  Logger l =
Logger.getLogger(DaytimeServer.class.getName( ));

l.log(Level.SEVERE, "FATAL error in DaytimeServer", t); System.exit(1);

```

```

    }

    }

    }

    catch(Exception e) {

        // This is a startup error: there is no need to log it;

        // just print a message and exit  System.err.println(e);

        System.exit(1);

    }

}

}

}

```

Client code :

''' Created on Feb 20, 2021

@author: Sabikun Nahar Piya

'''

```
package je3.nio; import java.net;
```

```
/**
```

```
* Connect to a daytime server using the UDP protocol
```

```
. * We use java.net instead of java.nio because DatagramChannel doesn't honor
```

```
* the setSoTimeout( ) method on the underlying DatagramSocket
```

```
*/
```

```
public class DaytimeClient {
```

```
    public static void main(String args[ ]) throws java.io.IOException {
```

```
        // Figure out the host and port we're going to talk to  String host = args[0];
```

```
        int port = 13;
```

```
        if (args.length > 1) port = Integer.parseInt(args[1]);
```

```
        // Create a socket to use  DatagramSocket socket = new DatagramSocket( );
```

```
        // Specify a 1-second timeout so that receive( ) does not block forever.
```

```
        socket.setSoTimeout(1000);
```

```
        // This buffer will hold the response. On overflow, extra bytes are
```

```
        // discarded: there is no possibility of a buffer overflow attack here.  byte[ ] buffer
        = new byte[512];
```

```
        DatagramPacket packet = new DatagramPacket(buffer, buffer.length, new
        InetSocketAddress(host,port));
```

```
        // Try three times before giving up  for(int i = 0; i < 3; i++) {
```

```
        try {
```

```
            // Send an empty datagram to the specified host (and port)  packet.setLength(0);
```

```
            // make the packet empty  socket.send(packet); // send it out
```

```
            // Wait for a response (or time out after 1 second)
```

```
            packet.setLength(buffer.length);
```

```
// make room for the response socket.receive(packet); //

wait for the response // Decode and print the response System.out.print(new
String(buffer, 0, packet.getLength( ), "US-ASCII"));

// We were successful, so break out of the retry loop break; }
catch(SocketTimeoutException e) {

// If the receive call timed out, print error and retry System.out.println("No
response");

}

}

// We're done with the channel now

socket.close( );

}

}
```

Conclusion: From this lab , I have known that how to Sockets provide the communication mechanism between two computers using TCP. I have also known that how to communicate client and the server by writing to and reading from the socket. I learnt that how to java.net.ServerSocket class provides a mechanism for the server program to listen for clients and establish connections with them