Hash functions from abelian threefolds

Leuven Isogeny Days

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The CGL hash function

The Charles-Goren-Lauter hash function '08

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Setup: supersingular elliptic curve E_o/\mathbb{F}_{p^2} Input: message (m_1 \dots m_n) Output: H(m) \in \mathbb{F}_{p^2} for i=1,\dots,n do E_i= \mathrm{radical2iso}(E_{i-1},m_i) return j(E_n)
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radical2iso computes a radical 2-isogeny

- · cost: essentially one sqrt
- m; decides the sign of the sqrt



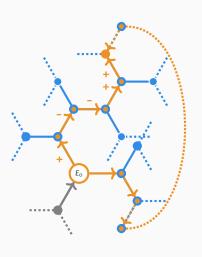
Security assumptions

Preimage resistance

- Given $j \in \mathbb{F}_{p^2}$, it is hard to find m with H(m) = j.
- ⇒ finding a 2*-isogeny between two supersingular elliptic curves

Collision resistance

- It is hard to find two messages $m \neq m'$ with $H(m) \neq H(m')$.
- \Rightarrow finding an endomorphism of degree 2* of E_0 .



CGL in dimension 2

The 2-dimensional setting

Principally polarized abelian surfaces

Product of elliptic curves,



 $E \times E'$

 $pprox p^2$

Jacobian of genus-2 curve C,



$$C: y^2 = f(x), \deg(f) \in \{5, 6\}.$$

 $\approx p^3$

number of superspecial varieties

(N, N) - Isogenies

- **Kernel**: maximal isotropic subgroup of A[N], isomorphic to $(\mathbb{Z}/N\mathbb{Z})^2$.
- $N=\ell$ prime: $\ell^3+\ell^2+\ell+1$ isotropic (ℓ,ℓ) -groups in total.
- Non-backtracking composition of two (ℓ,ℓ) -isogenies:

$$(\ell^2,\ell^2)$$
-isogeny , (ℓ^2,ℓ,ℓ) -isogeny



For an (ℓ, ℓ) -isogeny ϕ_1 , we are only interested in *good* extensions, i.e. ϕ_2 so that $\phi_2 \circ \phi_1$ is an (ℓ^2, ℓ^2) -isogeny.

 \Rightarrow Multiradical formulas: explicitly known for $\ell=2,3$ (Castryck-Decru '21)

CGL - dim2: Takashima '18, Castryck-Decru-Smith '20

```
Setup: superspecial abelian surface A_0/\mathbb{F}_{p^2}

Input: message (m_1 \dots m_{3n})

Output: H(m) \in \mathbb{F}_{p^2}

for i=1,\dots,n do

A_i = \text{radical2iso}(A_{i-1},m_{3i-2},m_{3i-1},m_{3i})

return \text{Igusa}(A_n)
```

radical2iso (Richelot)

Input:
$$A = Jac(C)$$
, $C : y^2 = l_1(x) \cdots l_6(x)$
Output: $A' = Jac(C')$, $C' : y^2 = l'_1(x) \cdots l'_6(x)$.
(1) $g_1 = l_1 \cdot l_4$, $g_2 = l_2 \cdot l_5$, $g_3 = l_3 \cdot l_6$
(2) $h_1 = g_2g'_3 - g_3g'_2$, $h_2 = g_3g'_1 - g_1g'_3$,
 $h_3 = g_1g'_2 - g_2g'_1$
(3) $l'_1 \cdot l'_2 = h_1$, $l'_3 \cdot l'_4 = h_2$, $l'_5 \cdot l'_6 = h_3$

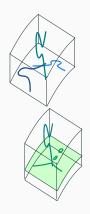
- cost: essentially 3 sqrts
- m_{3i-2}, m_{3i-1}, m_{3i}
 decide the
 signs of the
 sqrts

The 3-dimensional picture

Principally polarized abelian threefolds

reducible threefolds

product of elliptic curve *E* with abelian surface *A*



 $\approx p^3$

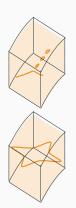
A reducible



A irreducible

irreducible threefolds

Jacobian of a genus-3 curve C



≈ p⁵ (?)

hyperelliptic



C plane quartic

number of superspecial varieties

Isogenies in dimension 3

(N, N, N)-isogenies

- kernel $G \subset A[N]$ maximal isotropic, and $G \cong \mathbb{Z}/N\mathbb{Z} \times \mathbb{Z}/N\mathbb{Z} \times \mathbb{Z}/N\mathbb{Z}$.
- $N=\ell$ prime: $(\ell^3+1)(\ell^2+1)(\ell+1)$ isotropic (ℓ,ℓ,ℓ) -groups of $A[\ell]$ in total.



Computational aspects

• Explicit radical formulas for $\ell=2$ using theta functions (Ohashi-Onuki-Yoshizumi-Kudo-Nuida '24, this work)

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Radical (2, ..., 2)-isogenies in theta coordinates

$(2, \ldots, 2)$ -isogenies in theta coordinates

Notation

- A: principally polarized abelian variety of dimension g.
- $\theta^{A}: A \to \mathbb{P}^{2^{g}-1}$: a level-2 theta structure.
- $\theta^A(O_A) = (a_0 : \cdots : a_{2^g-1})$: the level-2 theta null point of (A, θ^A) .

A maximal 2-isotropic group

$$G = \langle (a_0: -a_1: \cdots: a_{2^g-2}: -a_{2^g-1}), \ldots, (a_0: a_1: \cdots: -a_{2^g-2}: -a_{2^g-1}) \rangle$$

The $(2,\ldots,2)$ -isogeny 1 $\phi:A\to B=A/G$



Example: g = 1

\tilde{b} dual theta null point of (B, θ^B)



¹Building block for computing (2ⁿ, 2ⁿ)-isogenies (Dartois-Maino-Pope-Robert'24)

Radical formula for $\phi: A \rightarrow B$

Goal:

On input
$$\theta^A(O_A) = (a_0 : \cdots : a_{2^g-1})$$
, output $\theta^B(O_B) = (b_0 : \cdots : b_{2^g-1})$.

Idea: $o_B = \phi(o_A)$.

For
$$g=1$$

$$\begin{bmatrix}
a_0 \\
a_1
\end{bmatrix}
\xrightarrow{a_0^2}
\begin{bmatrix}
x_0 = a_0^2 + a_1^2 \\
x_1 = a_0^2 - a_1^2
\end{bmatrix}
\xrightarrow{\tilde{b}_0 = x_0/\tilde{b}_0}
\begin{bmatrix}
\tilde{b}_0 = x_0/\tilde{b}_0 \\
\tilde{b}_1 = x_1/\tilde{b}_1
\end{bmatrix}
\xrightarrow{b_0 = \tilde{b}_0 + \tilde{b}_1}
\begin{bmatrix}
b_0 = \tilde{b}_0 + \tilde{b}_1 \\
b_1 = \tilde{b}_0 - \tilde{b}_1
\end{bmatrix}
\xrightarrow{b_0}$$

In general dimension g:

$$\tilde{b}_i^2 = x_i \quad \text{for } i = 0, \dots, 2^g - 1.$$

Radical isogeny



Compatible square-roots

Essentially:
$$T_s: (X_0: \cdots: X_{2^g-1}) \to (\pm \sqrt{X_0}: \cdots: \pm \sqrt{X_{2^g-1}})$$

• Fact: A radical (2, ..., 2)-isogeny requires g(g+1)/2 square-root computations.

Trick: $(\sqrt{x_0}:\cdots:\sqrt{x_{2^g-1}})=(x_0:\sqrt{x_0x_1}:\cdots:\sqrt{x_0x_{2^g-1}}).$ \Rightarrow left with $2^g-1\geq g(g+1)/2$ square-root computations.

Lemma

Let $I=\{i_1,\ldots,i_{g(g+1)/2}\}$, a and assume $x_i\neq o$ for $i\in I\cap \{o\}$. Then fixing the square-roots $\sqrt{x_ox_i}$ with $i\in I$, uniquely determines the remaining square-roots.

^aThere are some restrictions on the index set I

Radical isogenies in dimensions 1 and 2

General description



dimension 1:

$$g(g+1)/2 = 2^g - 1 = 1$$

$$\mathcal{T}_{s}:(x_{0}:x_{1})\mapsto$$

$$(X_0: S \cdot \sqrt{X_0X_1})$$

dimension 2:

$$g(g+1)/2=2^g-1=3$$

$$\mathcal{T}_{s}:(X_{0}:X_{1}:X_{2}:X_{3})\mapsto$$

$$(X_0 : S_1 \cdot \sqrt{X_0 X_1} : S_2 \cdot \sqrt{X_0 X_2} : S_3 \cdot \sqrt{X_0 X_3})$$

Radical isogenies in dimension 3

$$2^g - 1 = 7 > 6 = g(g + 1)/2$$

Theorem

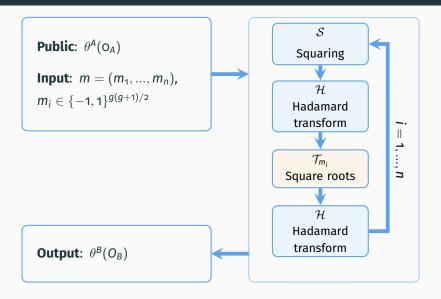
 $\mathcal{T}_s: (x_0: \cdots: x_7) \mapsto (x_0: s_1\sqrt{x_0x_1}: \cdots: s_6\sqrt{x_0x_6}: y_7)$, where $y_7 = F(s_1\sqrt{x_0x_1}, \ldots, s_6\sqrt{x_0x_6}, x_7)$ for an explicit algebraic expression F.

- **Evaluation of** *F***:** 22 multiplications (in the generic case).
- Proof idea: Not all points (y₀: · · · : y₇) correspond to a p.p. abelian threefold.

We use the *Riemann relations* to find a relation $G(y_0, ..., y_7) = 0$ among the coordinates of a valid theta null point.

Theta-CGL

Description of our hash function: Theta-CGL_a



Why go to dimension g > 1?

	<i>g</i> = 1			
superspecial a.v.s	$\approx p$	$pprox p^3$	$pprox p^6$	$p^{g(g+1)/2}$
solving the isogeny problem	$O(\sqrt{p})$	O(p)	$O(p^2)$	$O(p^{g-1})$

Theta-CGL $_q$ with security parameter $\lambda=$ 128

A: irreducible superspecial p.p.a.v. over \mathbb{F}_{p^2}

	g = 1	g = 2	g=3	<i>g</i> > 1
log p	256	128	64	$\lambda/(g-1)$
computations per bit	pprox 1 sqrt	pprox 1 sqrt	pprox 1 sqrt+	?

Going to higher dimensions allows us to significantly reduce the size of the prime, essentially at no additional cost!

Implementation in Rust

Running times of computing the hash of a 256-bit message in dimensions one, two and three.

	2-radical (μ s)	4-radical (μ s)	8-radical (μ s)
g = 1	3153	2037	1737
g = 2	989	742	_
g=3	432	_	_

Summary

This presentation

- Explicit formulas for radical (2, ..., 2)-isogenies in dimensions g = 1, 2, 3 using level-2 theta structure.
- Implementation of **Theta-CGL**g for g = 1, 2, 3 showing that dimension 3 is the best setting!

Our (upcoming) paper

- Radical isogeny formulas for 4- and 8-isogenies (g=1)
- Almost radical isogeny formulas for (4,4)-isogenies (g=2)
- Non-generic radical (2,2,2)-isogeny formulas (e.g. splitting and gluing isogenies).

Thanks for your attention!