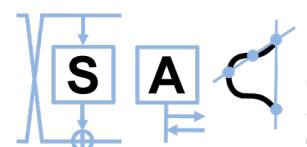


# Fast and Efficient Hardware Implementation of HQC

Sanjay Deshpande\*, Chuanqi Xu\*, Mamuri Nawan<sup>†</sup>, Kashif Nawaz<sup>†</sup>, and Jakub Szefer<sup>\*</sup>

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<sup>†</sup>Cryptography Research Centre, Technology Innovation Institute



**Selected Areas in Cryptography  
2023**

# Outline

- Motivation
- Introduction
- Goal and Existing work
- Components in HQC and their Hardware Implementation
- Comparison with Related State of the Art

# Motivation

- Quantum Computers with sufficient Qubits will be able to solve practical problems.
- Two algorithms are of special interest for cryptography community:
  - Shor's Algorithm
  - Grover's Search Algorithm
- For example, to break a 2048-bit RSA, a perfect quantum computer with 4099 ideal qubits can do it in 10 seconds.
- Most recently, IBM announced 433-qubits quantum computer "Osprey".
- There is a need for Quantum-safe Cryptography!



Image Source: MIT Technology Review

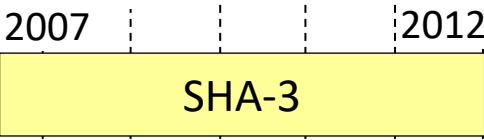
Source: <https://www.quintessencelabs.com/blog/breaking-rsa-encryption-update-state-art/>

Source: <https://newsroom.ibm.com/2022-11-09-IBM-Unveils-400-Qubit-Plus-Quantum-Processor-and-Next-Generation-IBM-Quantum-System-Two>

# Post Quantum Cryptography Competition

51 hash functions

→ 1 winner



Completed

In Progress

2013

2019

CAESAR

57 authenticated ciphers

→ multiple winners

69 Public Key Post Quantum  
Cryptography Schemes

56 Lightweight authenticated ciphers  
& hash functions

XII.2016

TBD

Post-Quantum

2018

2023

Lightweight

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Year

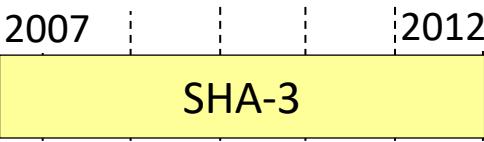


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# Post Quantum Cryptography Competition

51 hash functions

→ 1 winner

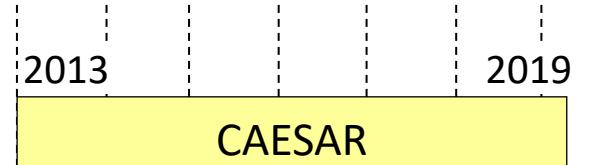


Completed

In Progress

57 authenticated ciphers

→ multiple winners



2019

69 Public Key Post Quantum  
Cryptography Schemes

XII.2016

TBD

Round 4 announced

1 candidate selected for standardization  
4 candidates will undergo further evaluation

2018

2023

56 Lightweight authenticated ciphers  
& hash functions



07

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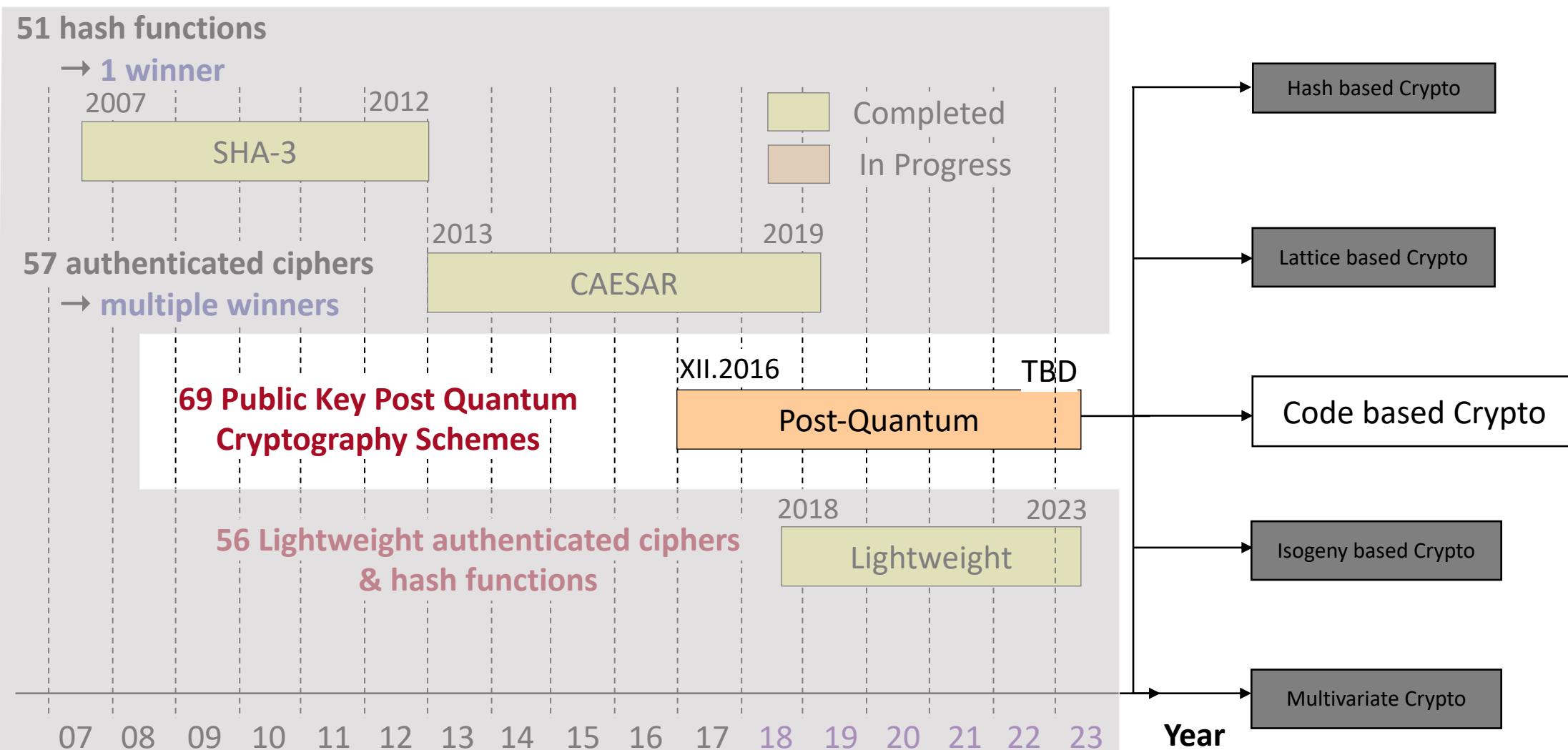
23

Year

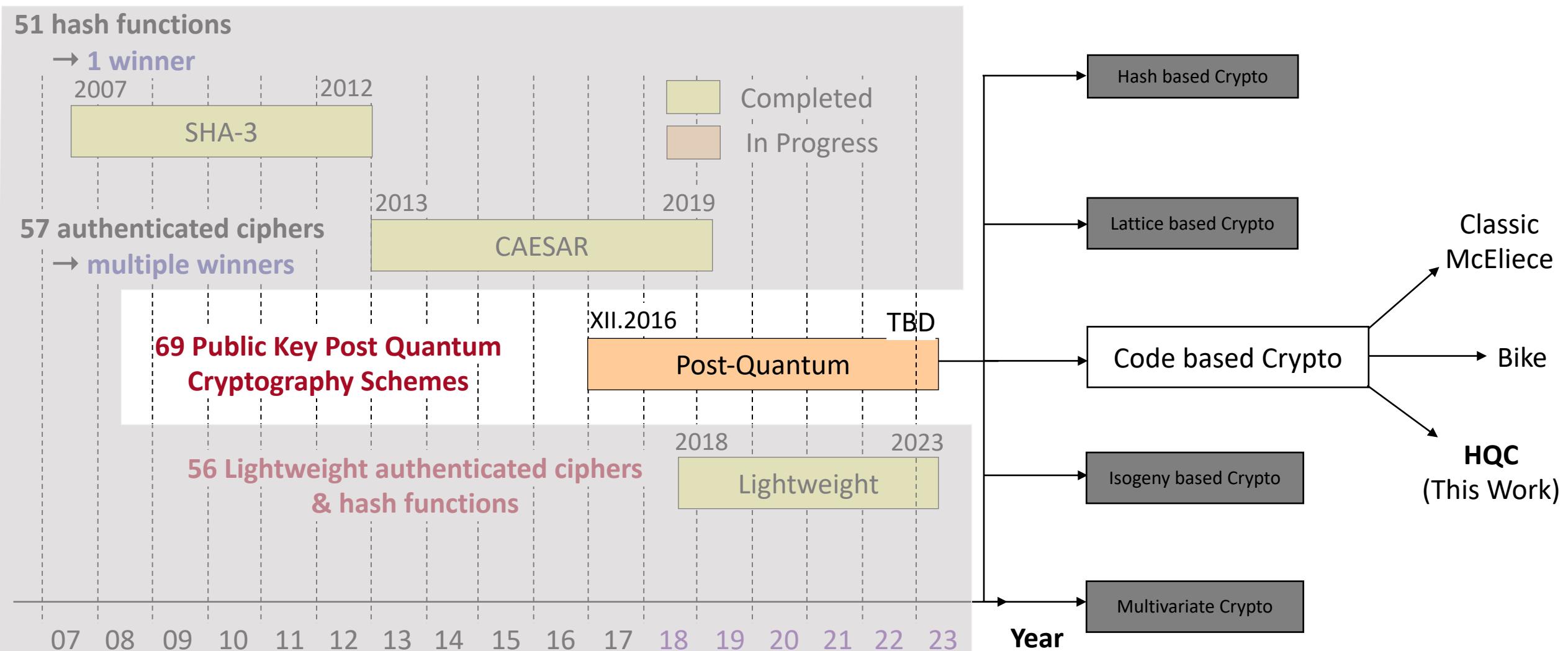


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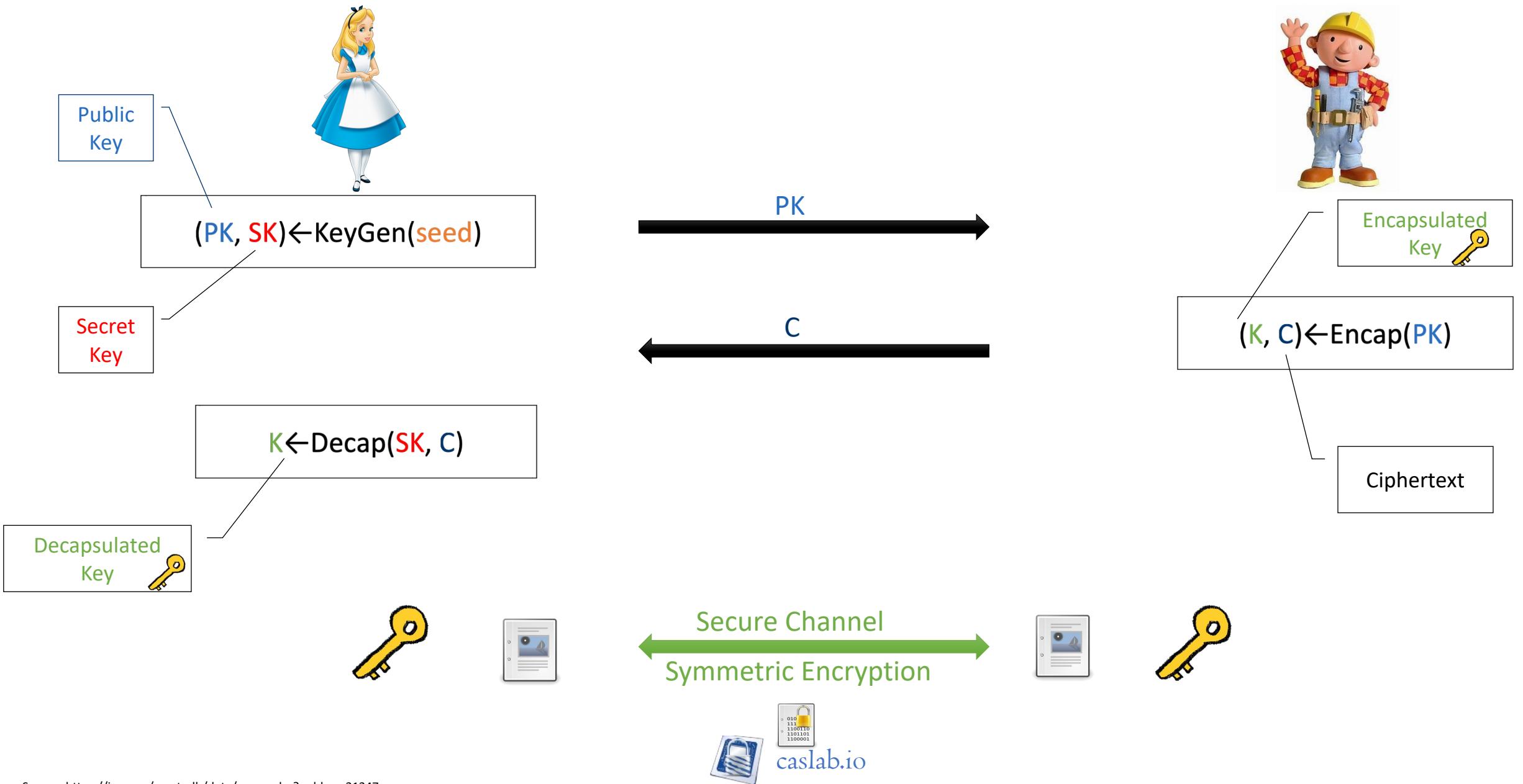
# Post Quantum Cryptography Competition



# Post Quantum Cryptography Competition



# Key Encapsulation Mechanism



## Design & Experimental Setup

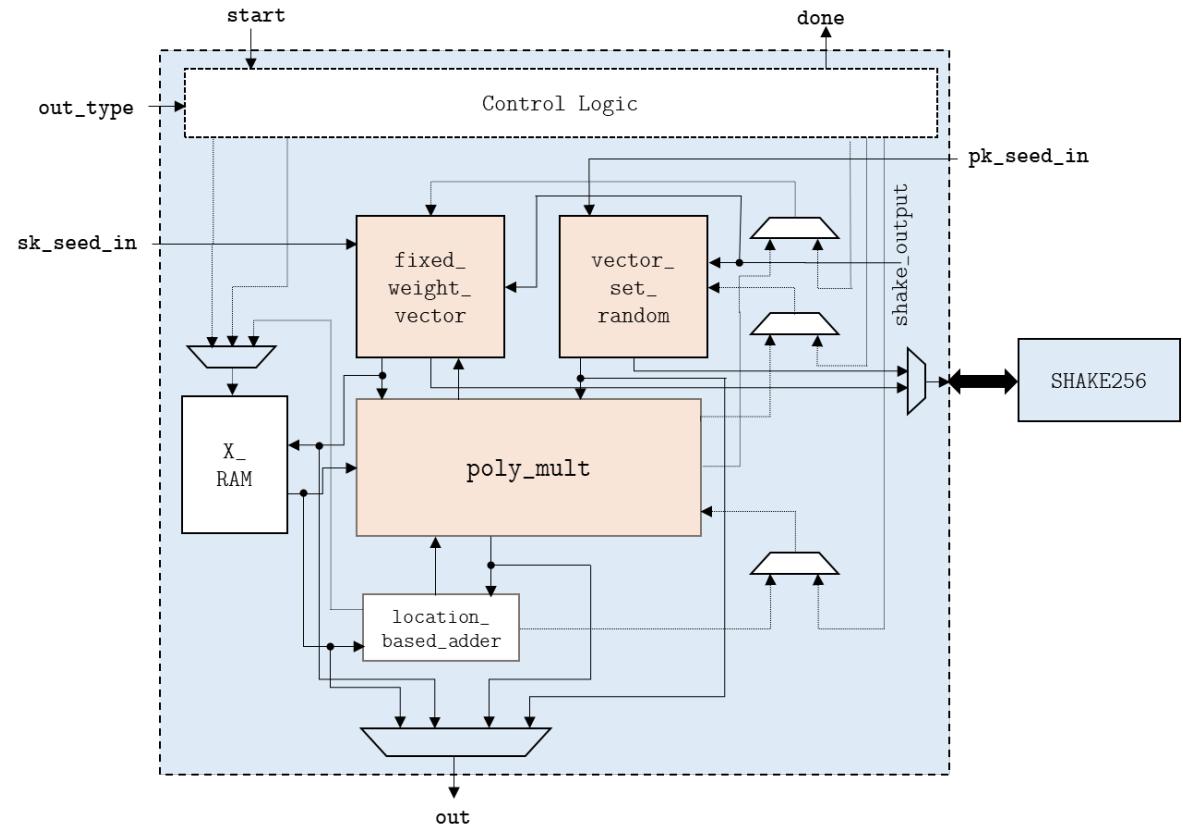
- HQC KEM parameter sets recommended as per the specification:
  - **hqc128**
  - **hqc192**
  - **hqc256**
- Target device: Xilinx - Artix 7 – xc7a200t-3 and xc7a100t-3
- Tool: Xilinx Vivado 2020.2.2 
- Verification: Using the Software Reference implementation by HQC team [AAB+20]
- Goal: Implement a constant-time hardware design which is parameterizable across
  - Security level
  - Performance parameters

Image Source: Xilinx

# Primitives

## 1. Key Generation

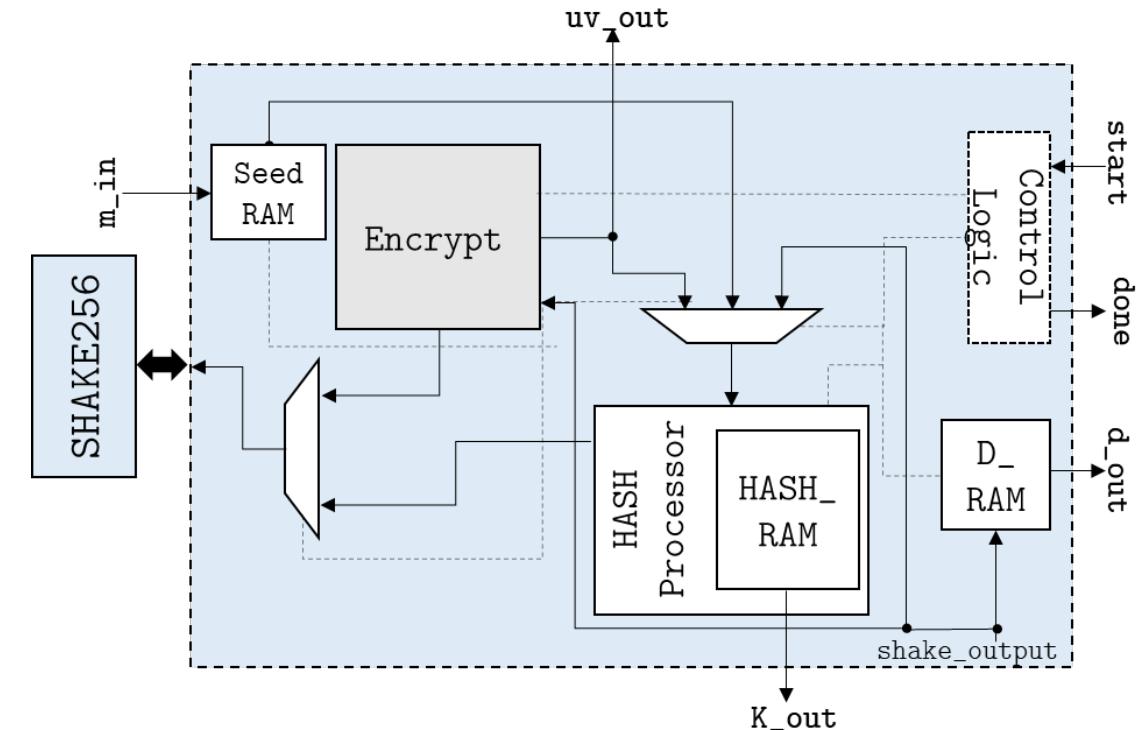
- Fixed weight error vector generation
  - SHAKE256
- Polynomial Multiplication
- Polynomial Addition



# Primitives

## 2. Encapsulation

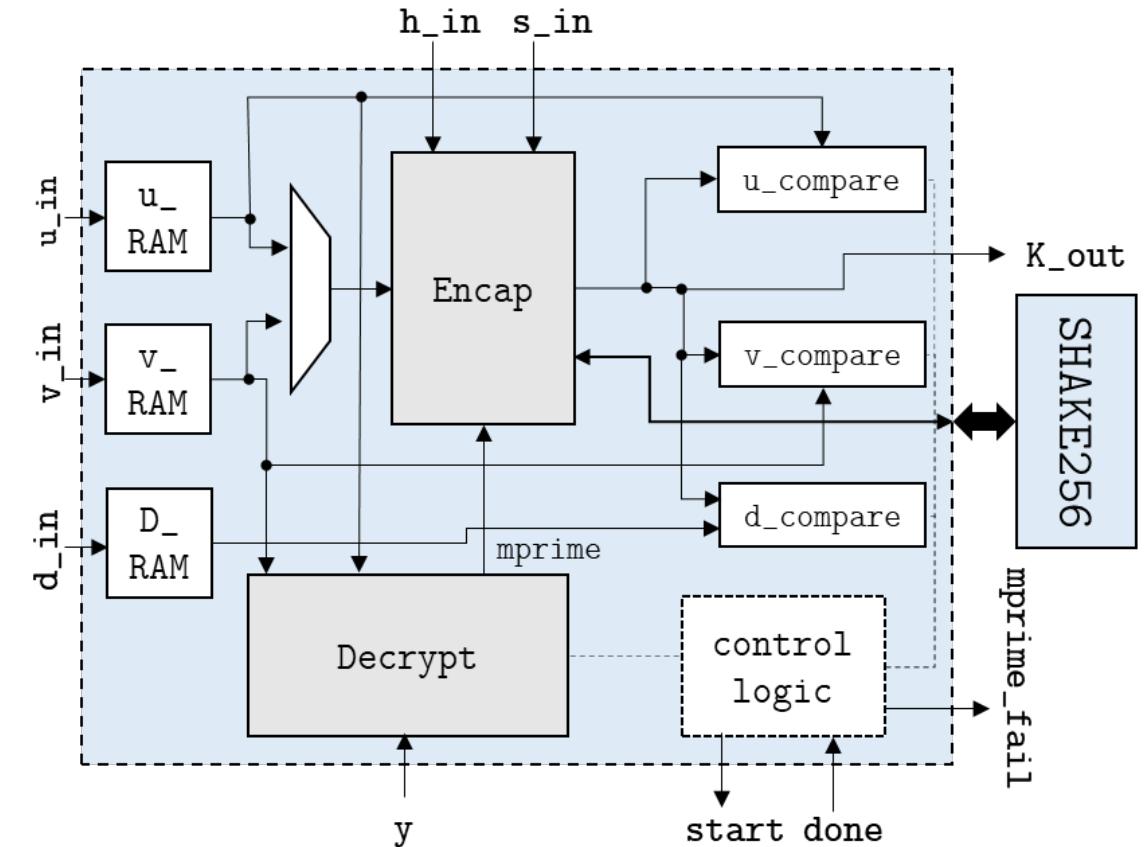
- Encrypt
  - Encode
  - Polynomial Multiplication
  - Polynomial Addition
  - Fixed weight error vector generation
    - SHAKE256
- Hash Computation
  - SHAKE256



# Primitives

## 3. Decapsulation

- Decrypt
    - Decode
    - Polynomial Subtraction
    - Polynomial Multiplication
  - Encrypt
  - Hash Computation
    - SHAKE256
- } Encapsulation



# Primitives – Common Modules

## 1. Key Generation

- Fixed weight error vector generation
  - SHAKE256
- Polynomial Multiplication
- Polynomial Addition

## 2. Encapsulation

- Encrypt
  - Encode
  - Polynomial Multiplication
  - Polynomial Addition
  - Fixed weight error vector generation
    - SHAKE256
- Hash Computation
  - SHAKE256

## 3. Decapsulation

- Decrypt
  - Decode
  - Polynomial Subtraction
  - Polynomial Multiplication
- Encrypt
- Hash Computation
  - SHAKE256

# Primitives - Existing Work

## 1. Key Generation

- Fixed weight error vector generation
  - SHAKE256 [CCD+22]
- Polynomial Multiplication
- Polynomial Addition

## 2. Encapsulation

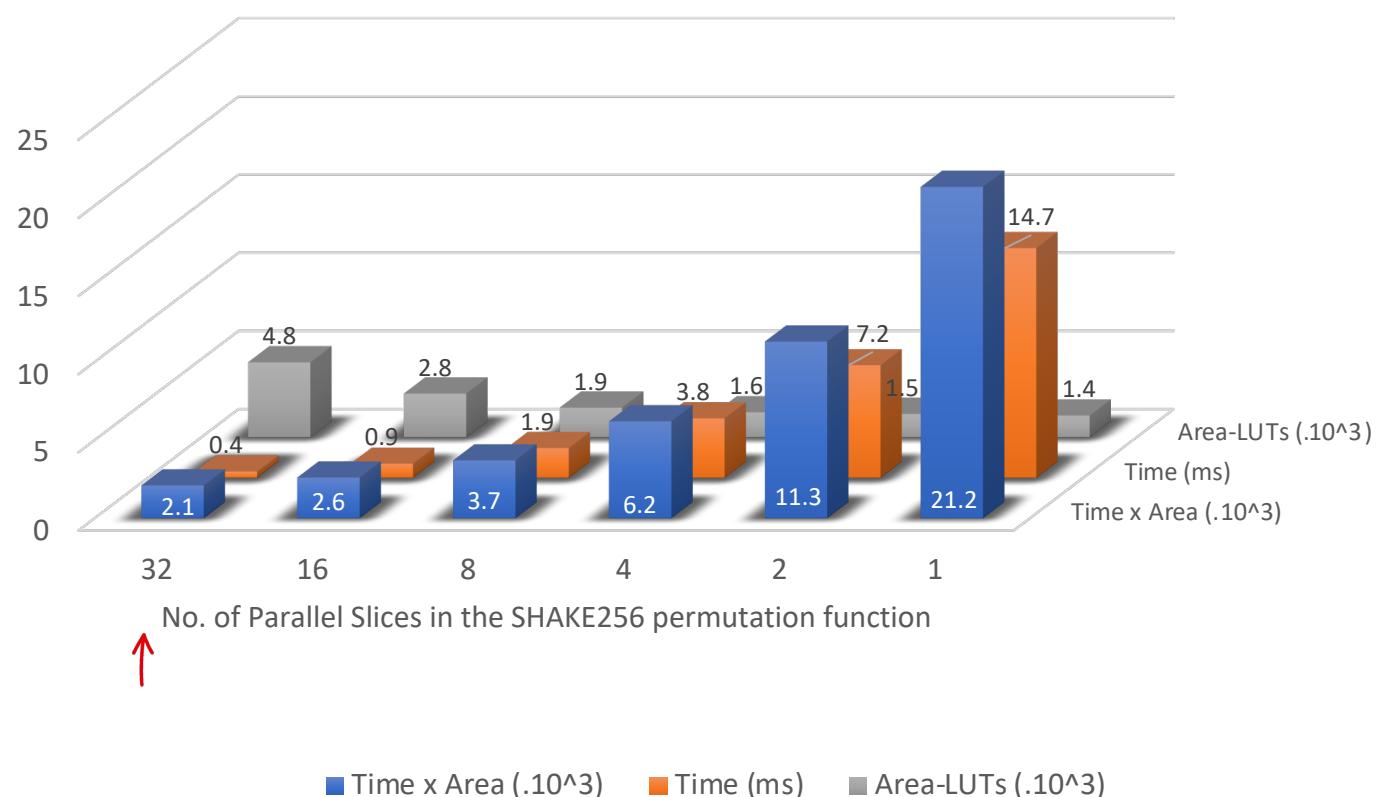
- Encrypt
  - Encode
  - Polynomial Multiplication
  - Polynomial Addition
  - Fixed weight error vector generation
    - SHAKE256 [CCD+22]
- Hash Computation
  - SHAKE256 [CCD+22]

## 3. Decapsulation

- Decrypt
  - Decode
  - Polynomial Subtraction
  - Polynomial Multiplication
- Encapsulation

# SHAKE256

- Improved existing design [CCD+22].
  - Added state preserving capability.
  - Added an extra mode in the performance parameter (Parallel Slices).



# Hardware Implementation of HQC KEM

## 1. Key Generation

- Fixed weight error vector generation
  - SHAKE256
- Polynomial Multiplication
- Polynomial Addition

## 2. Encapsulation

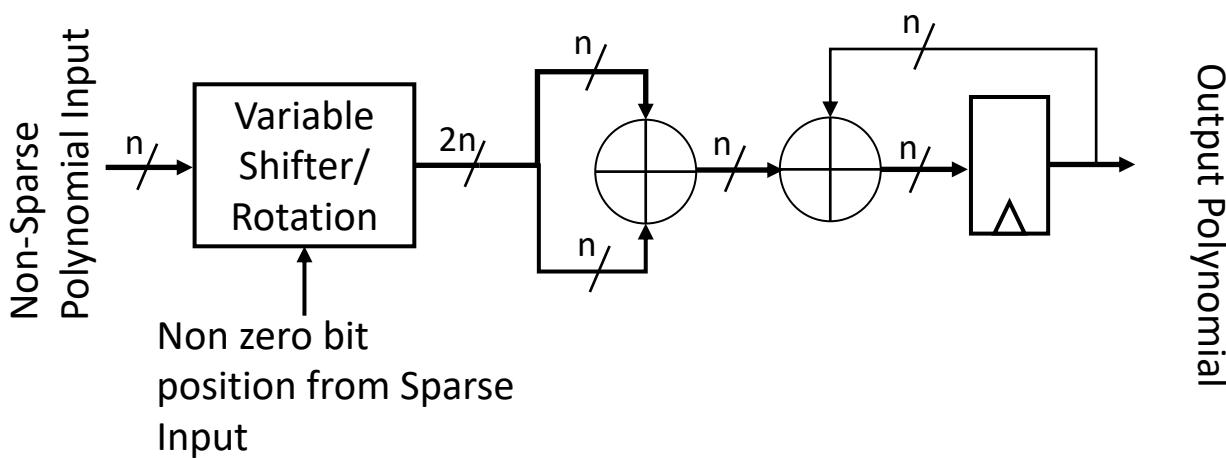
- Encrypt
  - Encode
  - Polynomial Multiplication
  - Polynomial Addition
  - Fixed weight error vector generation
    - SHAKE256
- Hash Computation
  - SHAKE256

## 3. Decapsulation

- Decrypt
  - Decode
  - Polynomial Subtraction
  - Polynomial Multiplication
- Encapsulation

# Polynomial Multiplication - Sparse Multiplication with Interleaved Reduction

- One of the inputs to polynomial multiplier is a sparse fixed weight vector.
- Indices of non-zero elements are used to shift non-sparse polynomial to imitate the multiplication.
- The multiplication and the modular reduction is interleaved.

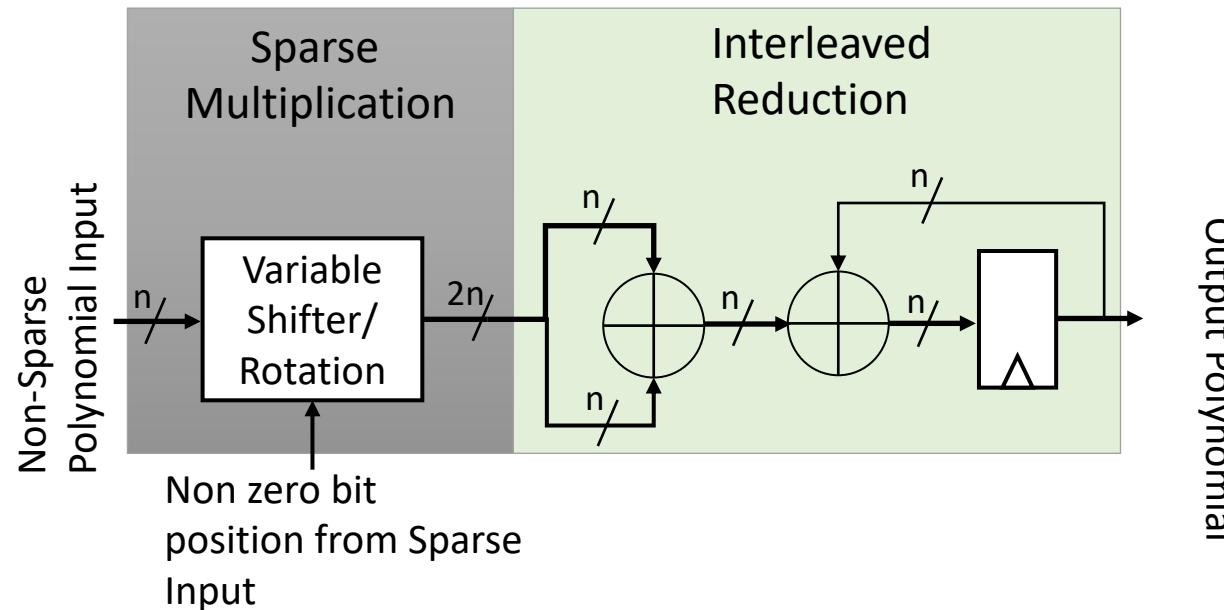


\*Conceptual Block Diagram



# Polynomial Multiplication - Sparse Multiplication with Interleaved Reduction

- One of the inputs to polynomial multiplier is a sparse fixed weight vector.
- Indices of non-zero elements are used to shift non-sparse polynomial to imitate the multiplication.
- The multiplication and the modular reduction is interleaved.



\*Conceptual Block Diagram

# Polynomial Multiplication - Results

BW (bits)	Resources						F (MHz)	Cycles (cyc.)	Time (us)	T x A
	Logic (SLICES)	Logic (LUT)	Memory (DSP)	Memory (FF)	Memory (BR)					
Our poly_mult module, Polynomial Length <sup>†</sup> = 12,323, $W_{SPARSE}^*$ = 71										
32	134	396	0	181	1	270	27,621	0.10	14	
64	202	599	0	205	2	277	13,918	0.05	10	
128	486	1,438	0	456	4	238	7,102	0.03	14	
General Sparse Multiplier, Polynomial Length <sup>†</sup> = 12,323, $W_{SPARSE}^*$ = 71 [19]										
32	132	319	0	127	2	234	27,691	0.12	16	
64	197	549	0	190	4	222	13,988	0.06	12	
128	378	1,136	0	381	8	185	7,172	0.04	15	
Sparse Multiplier, Polynomial Length <sup>†</sup> = 10,163, $W_{SPARSE}^*$ = 71 [15]										
32	100	—	—	2	240	158,614	0.66	66		
64	157	—	—	3	220	90,880	0.41	64		
128	292	—	—	5	210	51,688	0.24	70		

<sup>†</sup> = Slices (no info on LUTs), <sup>+</sup> Length of the non-sparse arbitrary polynomial, <sup>\*</sup> = Weight of the sparse polynomial input

Target Device: Artix 7 – xc7a200t



# Fixed Weight Vector Generation

- Generates a uniform random n-bit fixed-weight (w) vector.
- Algorithm given in [AAB+20].

Input: N, w, seed

Output: w distinct elements in range 0 to N – 1

```
1: pos = []*w
2. i,j,k,l = 0
3. prng_init(seed)
4: while (i < w) do
5:   rand_bits ← prng_draw(outsize = 24)
6:   if (rand_bits < THRESHOLD_VALUE)
7:     pos [i] = (rand_bits % N)
8:     i++
9:   end
10: end while
```

Guo et. al [GHJ+22] demonstrated a timing attack on the software reference implementation

Threshold Check

//Value for THRESHOLD is a constant based on the parameter set

```
11: while (j < w) do
12:   while (k < w) do
13:     if (pos[j] == pos[k] then
14:       while (l < 1) do
15:         rand_bits ← prng_draw(outsize=24)
16:         if (rand_bits < THRESHOLD) then
17:           pos [k] = (rand_bits % N)
18:           l++
19:         end if
20:       end while
21:     else then
22:       k++
23:     end
24:   end while
25:   j++
26: end while
27: end while
28: return pos
```

Duplicate Detection



# Fixed Weight Vector – Constant Weight Word (CWW) Method

- Constant time method proposed by Sendrier [SEN21].
- Recommended by HQC Team as a 4<sup>th</sup> round update.
- Barrett Reduction used in Index Generation.
- However, this method has small bias.

```
Input: N, w, seed
Output: w distinct elements in range 0 to N - 1
1: rand bits ← prng(input = seed, output size = 32 × w) RNG
2: for i ← w - 1 to 0 do
3:   pos[i] = i + (rand bits[32 + 32 * i - 1 : 32 * i])%(N - i) Index Generation
4: end for
5: for j ← w - 1 to 0 do
6:   duplicate found ← 0
7:   for k ← j + 1 to w - 1 do
8:     if pos[j] == pos[k] then
9:       duplicate found ← 1 Duplicate Detection and Index Replacement
10:    end if
11:   end for
12:   if duplicate found == 1 then
13:     pos[j] = j
14:   end
15: end for
16: return pos
```



# Fixed Weight Vector Generation – Fast and Non-Biased(FNB) Method

- Improvement on the original algorithm proposed given in [AAB+20]

Input: N, w, seed, ACC\_REJ

Output: w distinct elements in range 0 to N – 1

```
1: pos = []*(w + ACC_REJ)
2. i,j,k,m = 0
3. prng_init(seed)
4: while (i < (w+ ACC_REJ)) do
5:   rand_bits ← prng_draw(outsize = 24)          RNG
6:   if (rand_bits < THRESHOLD_VALUE) then
7:     pos [i] = (rand bits % N)
8:     i++
9:   end
10: end while

Threshold Check
```

//Value for THRESHOLD is a constant based on the parameter set  
// ACC\_REJ = ACCEPTABLE\_REJECTIONS  
// e.g., ACC\_REJ = 75

```
11: while (j < w + ACC_REJ) do
12:   while (k < (w+ ACC_REJ)) do
13:     if (j < w) then
14:       if pos[j] == pos[k] then
15:         if (m<ACC_REJ) then
16:           pos[k] = pos[m+w]
17:           m++
18:         end
19:       else
20:         USE PRNG TO DRAW MORE RAND BITS AND
REDO THRESHOLD CHECK and ASSIGN TO pos[k]
21:       end
22:     end
23:   end
24: else
25:   DUMMY OPERATIONS
26: else
27:   DUMMY OPERATIONS
28: end
29: k++
30: end while
31: j++
32: end while
33: return pos
```

Small prob non-constant time behavior

Duplicate Detection

## Fixed Weight Vector Generation - Evaluation

Design	Weight ( $w_r$ )	Resources						Time (us)	$T \times A$	Failure <sup>+</sup> Prob.
		Logic (LUT)	Memory (DSP)	Memory (FF)	F (MHz)	Cycles (cyc.)				
<b>Fast and Non-Biased Design (ACCEPTABLE REJECTIONS = <math>w_r</math>)</b>										
hqc128	75	316	0	124	2.0	223	1,479	6.63	2.10	$2.8 \times 2^{-199}$
hqc192	114	295	0	125	2.0	236	2,226	9.43	2.78	$1.1 \times 2^{-280}$
hqc256	149	314	0	192	2.5	242	3,248	13.42	4.21	$4.9 \times 2^{-355}$
<b>Constant Weight Word (CWW)</b>										
hqc128	75	201	4	229	1.0	201	3,062	15.23	3.06	0
hqc192	114	211	5	245	1.0	200	6,817	34.09	7.19	0
hqc256	149	216	5	248	1.0	204	11,487	56.31	1216	0

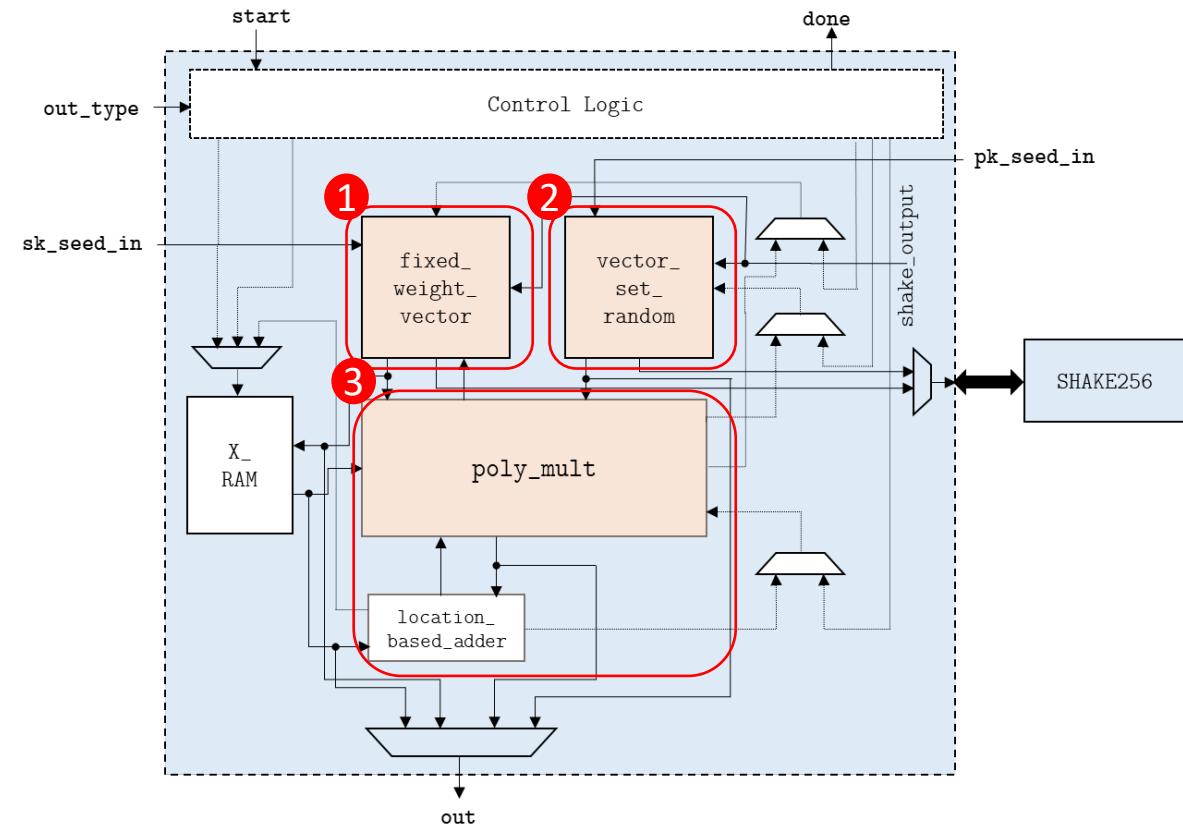
+ = Probability of our design failing to behave constant-time.

# Key Generation - Algorithm and Hardware Design

Inputs: pk\_seed, sk\_seed

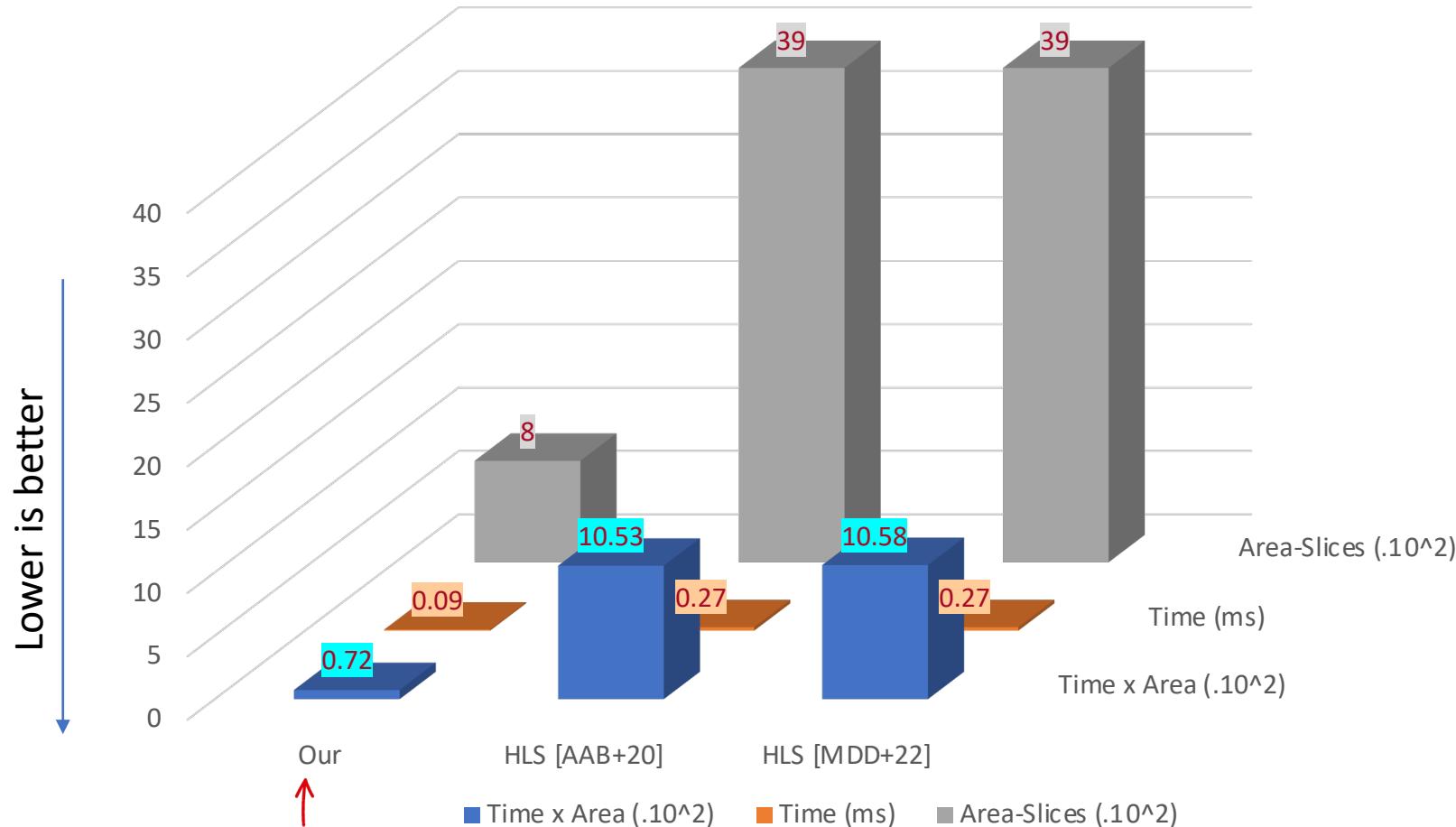
Outputs: public key( $h, s$ ), secret key ( $x, y$ )

1.  $(x, y) = \text{FixedWeight}(\text{sk\_seed})$
2.  $h = \text{VectorRandom}(\text{pk\_seed})$
3.  $s = x + h.y$



# Key Generation – Hardware Design Performance Comparison

HQC128



caslab.io

# Hardware Implementation of HQC KEM

## 1. Key Generation

- Fixed weight error vector generation
  - SHAKE256
- Polynomial Multiplication
- Polynomial Addition

## 2. Encapsulation

- Encrypt
  - Encode
  - Polynomial Multiplication
  - Polynomial Addition
  - Fixed weight error vector generation
    - SHAKE256
- Hash Computation
  - SHAKE256

## 3. Decapsulation

- Decrypt
  - Decode
  - Polynomial Subtraction
  - Polynomial Multiplication
- Encapsulation

# Encrypt - Algorithm and Hardware Design

Inputs: theta, Public Key ( $h_{in}$ ,  $s_{in}$ ), message ( $m_{in}$ )

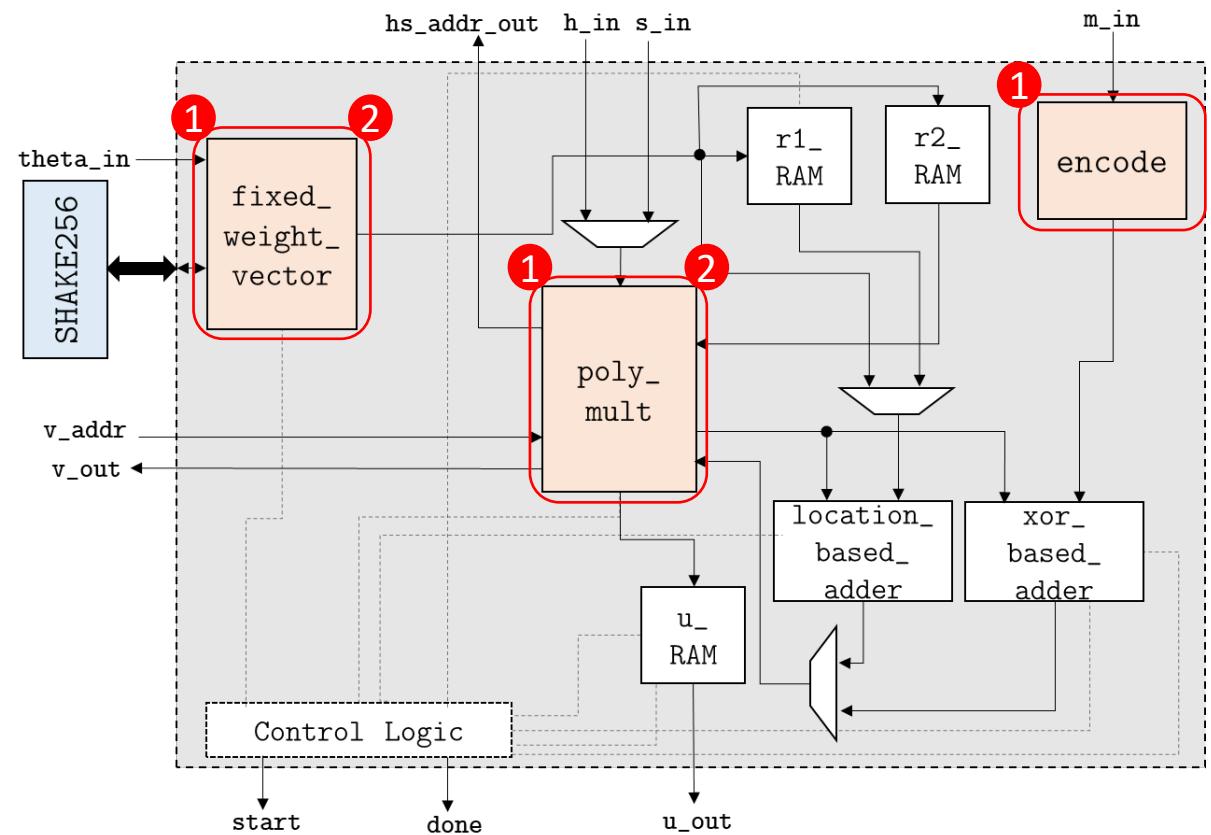
Outputs: Ciphertext ( $u_{out}$ ,  $v_{out}$ )

1.  $u_{out} = r_1 + h \cdot r_2$
2.  $v_{out} = m_{in}G + S \cdot r_2 + e$

Where  $m_{in}G$  = Encoded message

$r_1$ ,  $r_2$ ,  $e$  are fixed weight vectors

- Encode consists of Reed-Solomon and Reed-Muller Encoding



# Encrypt - Algorithm and Hardware Design

Inputs: theta, Public Key ( $h_{in}$ ,  $s_{in}$ ),  
message ( $m_{in}$ )

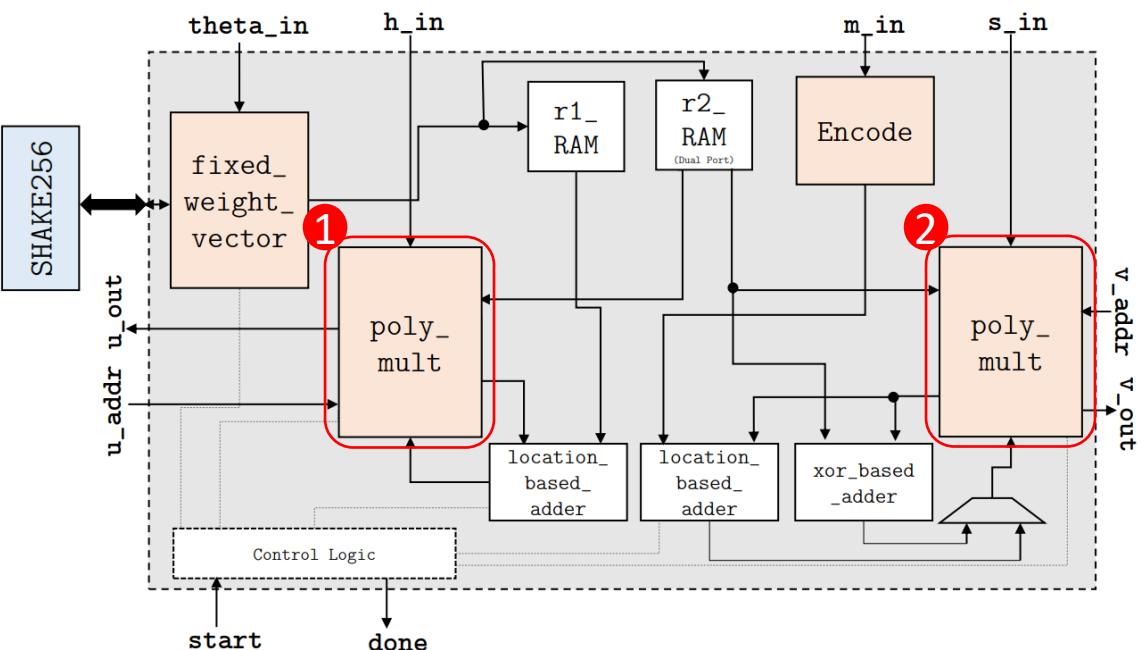
Outputs: Ciphertext ( $u_{out}$ ,  $v_{out}$ )

1.  $u_{out} = r_1 + h.r_2$
2.  $v_{out} = m_{in}.G + S.r_2 + e$

Where  $m_{in}G$  = Encoded message

$r_1, r_2, e$  are fixed weight vectors

- Two polynomial multiplications can be run in parallel.

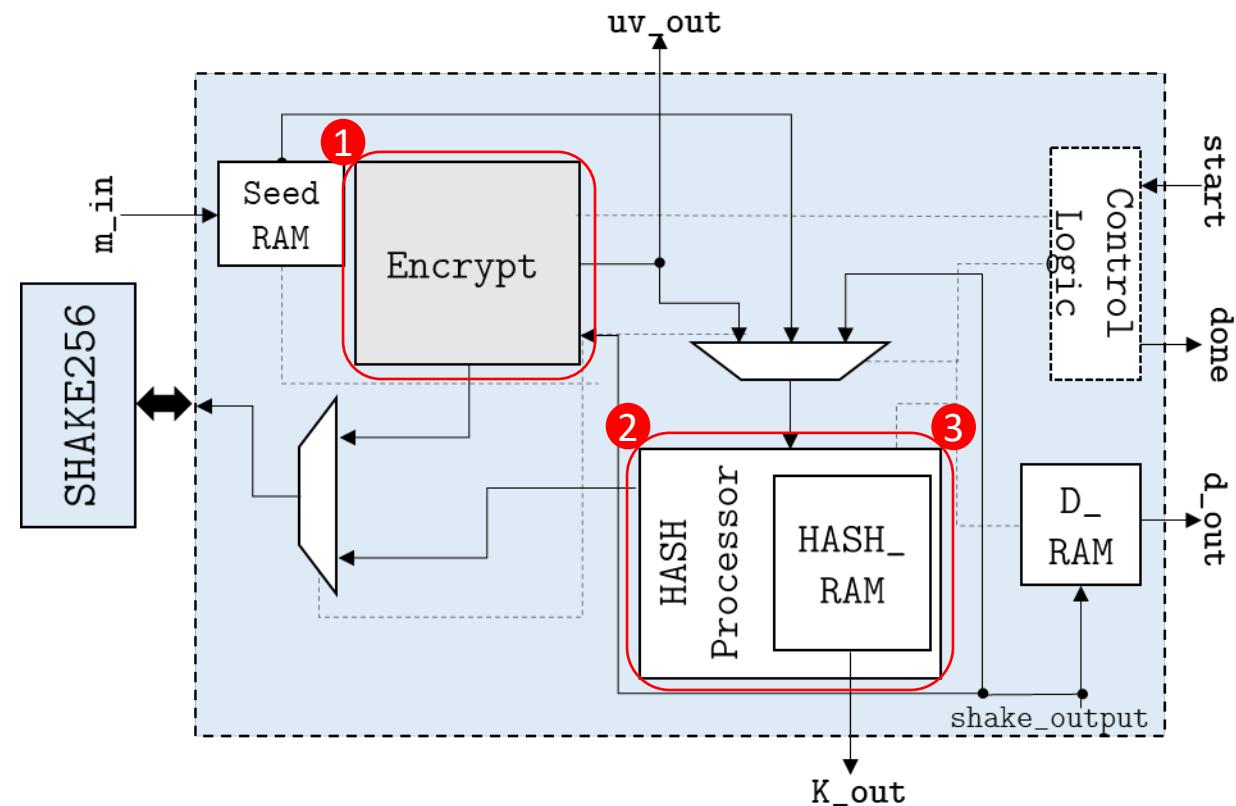


# Encapsulation - Algorithm and Hardware Design

Inputs: Public Key ( $h_{in}, s_{in}$ ), message ( $m_{in}$ )

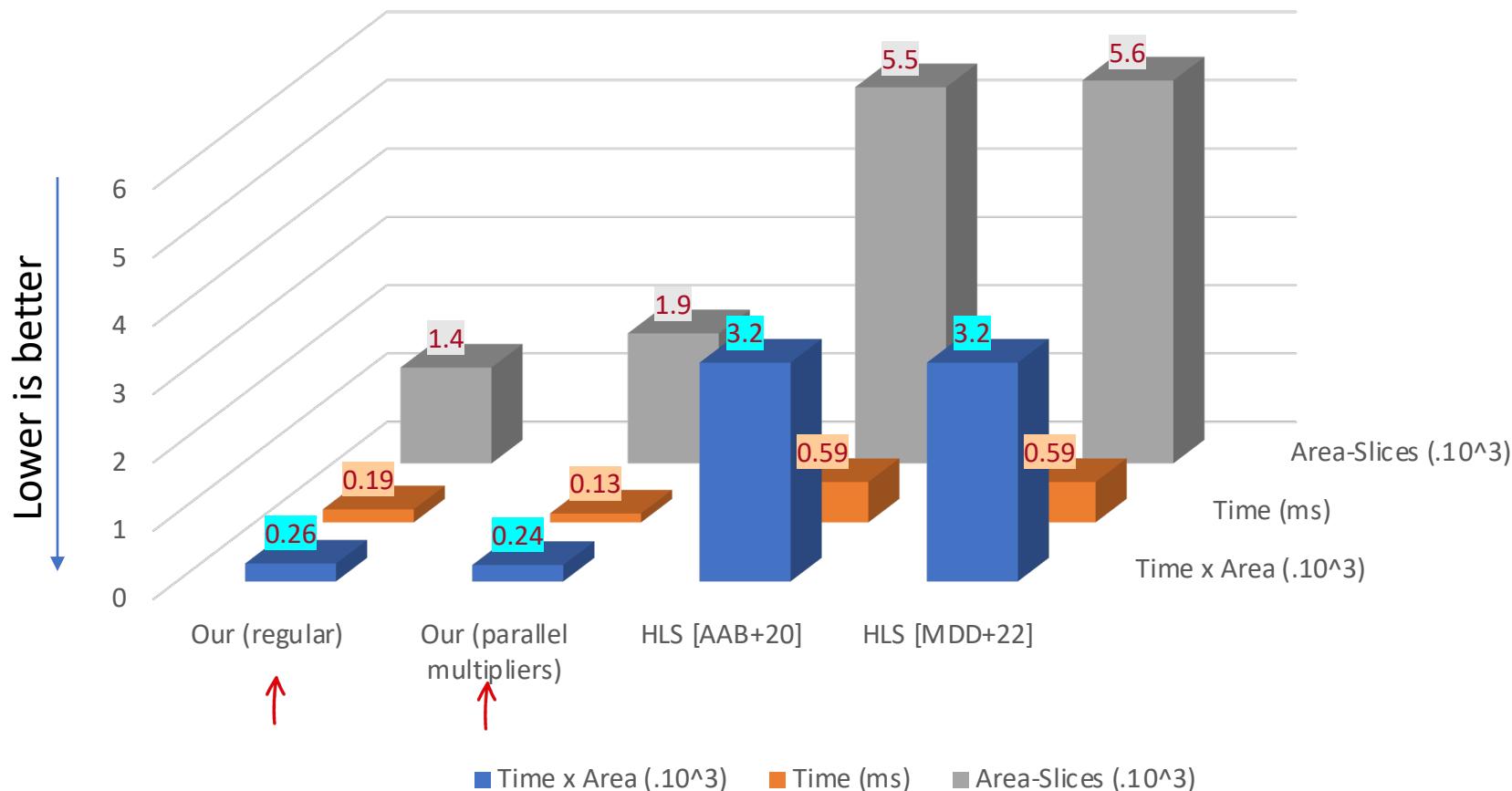
Outputs: Ciphertext ( $u_{out}, v_{out}$ ), Shared Secret ( $K_{out}$ ),  
Hashed message( $d_{out}$ )

1.  $(u_{out}, v_{out}) = \text{Encrypt}(m_{in})$
2.  $d = \text{Hash}(m_{in})$
3.  $\text{SS} = \text{Hash}(m_{in}, u_{out}, v_{out})$



# Encapsulation – Hardware Design Performance Comparison

HQC128



caslab.io

# Hardware Implementation of HQC KEM

## 1. Key Generation

- Fixed weight error vector generation
  - SHAKE256
- Polynomial Multiplication
- Polynomial Addition

## 2. Encapsulation

- Encrypt
  - Encode
  - Polynomial Multiplication
  - Polynomial Addition
  - Fixed weight error vector generation
    - SHAKE256
- Hash Processing
  - SHAKE256

## 3. Decapsulation

- Decrypt
  - Decode
  - Polynomial Subtraction
  - Polynomial Multiplication
- Encapsulation

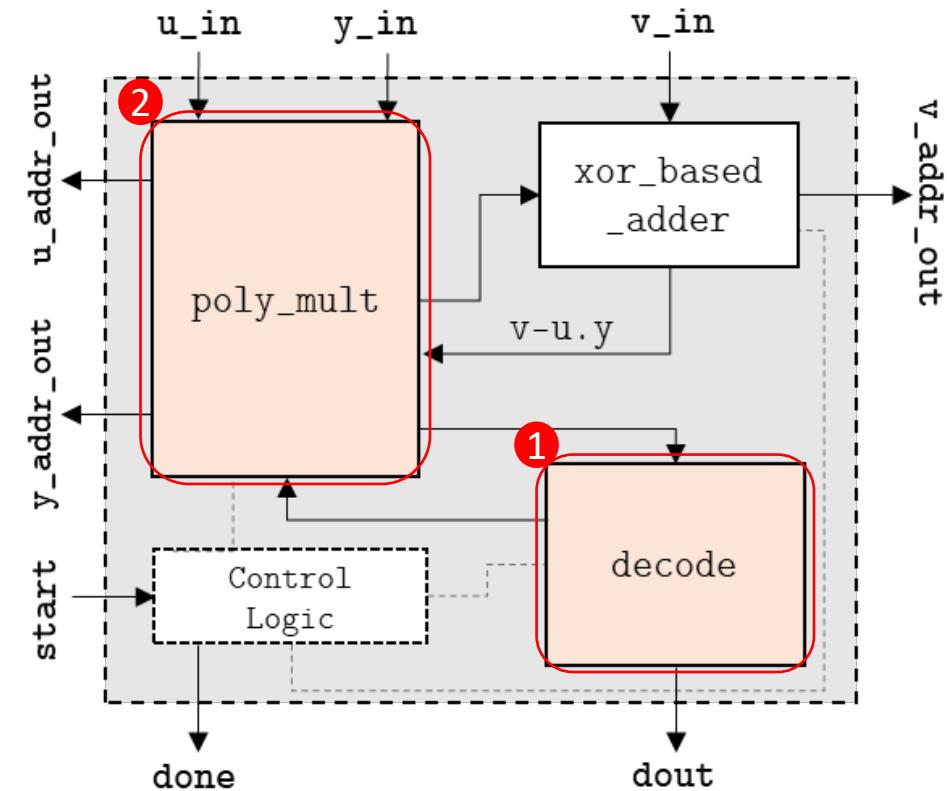
# Decrypt - Algorithm and Hardware Design

Inputs: Secret Key ( $y_{in}$ ), Ciphertext ( $u_{in}, v_{in}$ )

Outputs: Decoded Message ( $dout$ )

1.  $v_{minus\_uy} = v_{in} - u_{in} \cdot y_{in}$
2.  $dout = \text{Decode}(v_{minus\_uy})$

- Decode consists of Reed-Muller and Reed-Solomon decoding



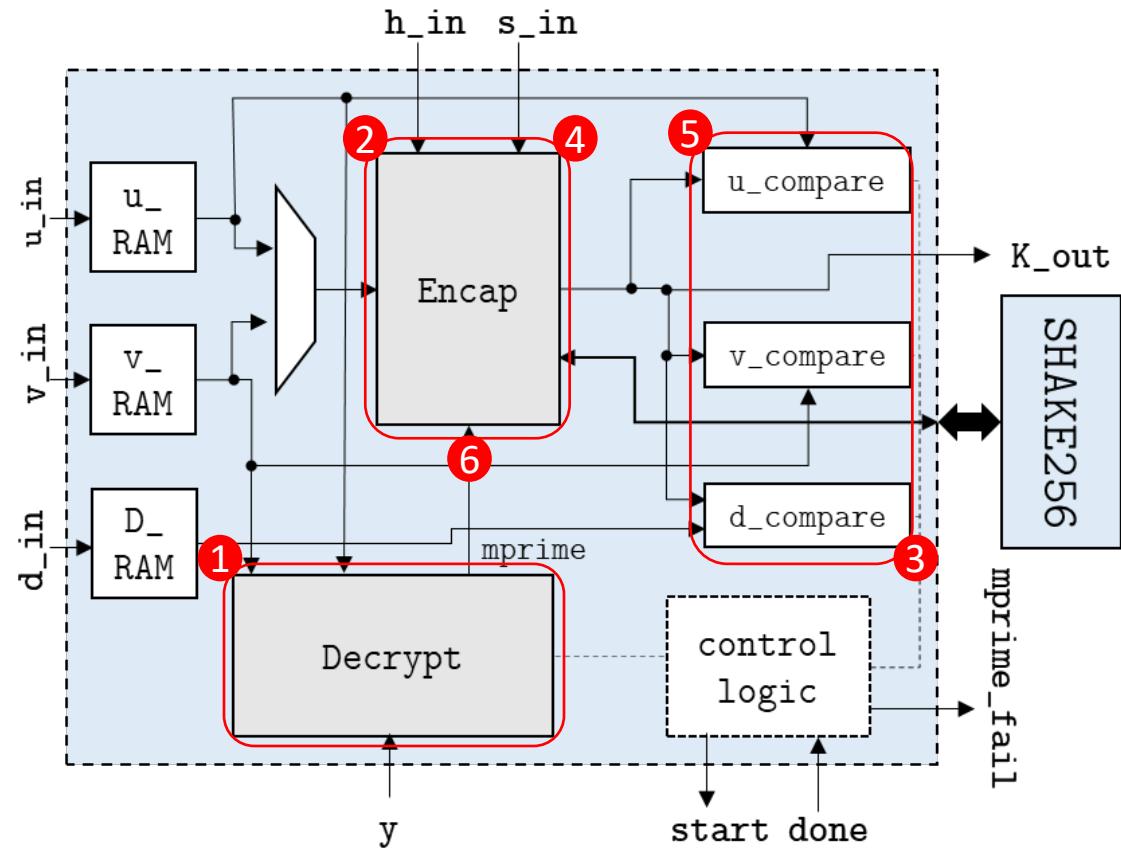
# Decapsulation

Inputs: Public Key( $h_{in}, s_{in}$ ), Secret Key ( $y_{in}$ ), Ciphertext ( $u_{in}, v_{in}$ ), Hashed message ( $d_{in}$ )

Outputs: Shared Secret ( $K_{out}$ )

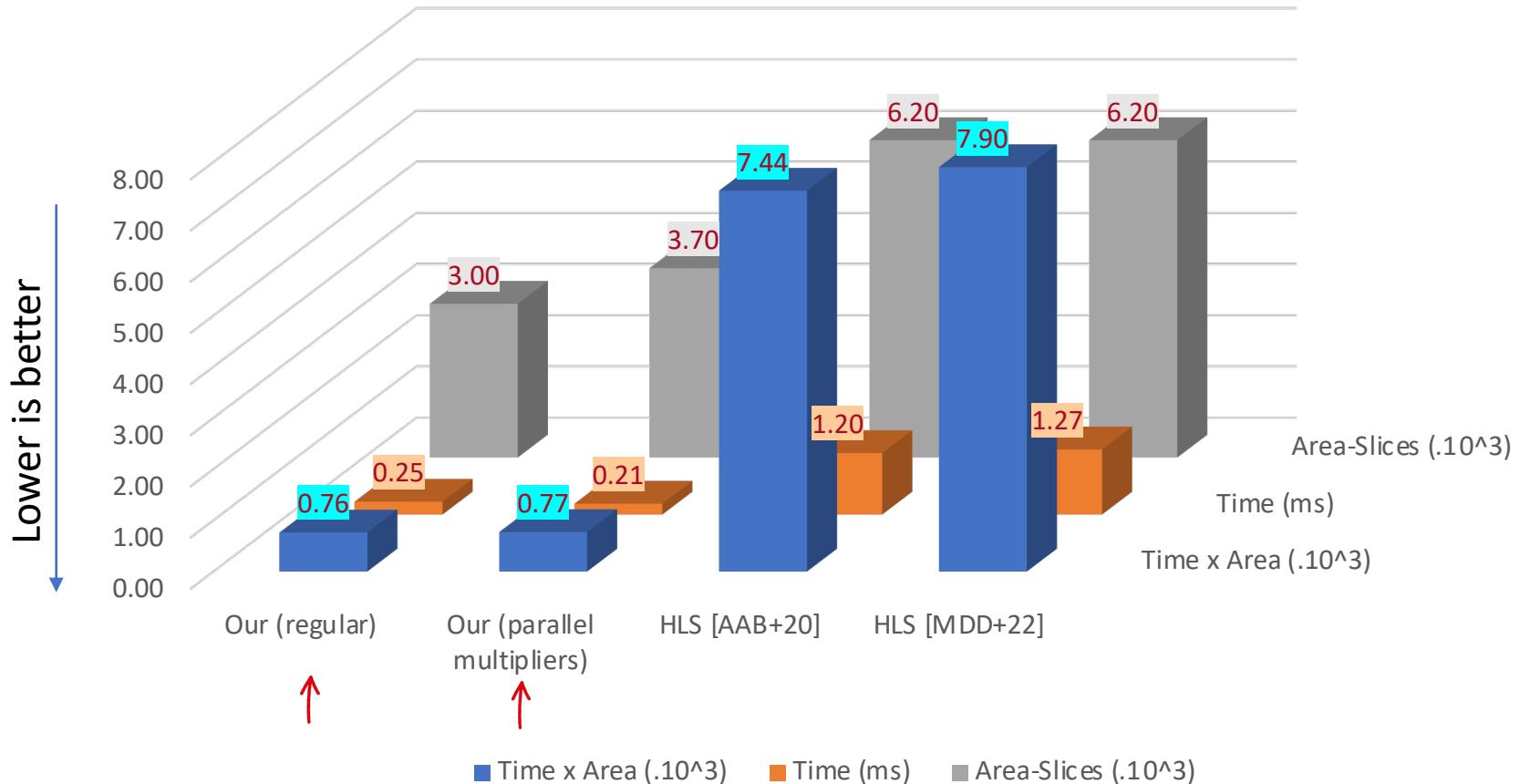
1.  $m' = \text{Decode}(v_{in} - u_{in} \cdot y_{in})$
2.  $d' = \text{Hash}(m')$
3. Verify  $d_{in} == d'$ ?
4.  $(u', v') = \text{Encrypt}(m')$
5. Verify  $(u', v') == (u_{in}, v_{in})$ ?
6.  $K_{out} = \text{Hash}(m', u_{in}, v_{in})$

- Decapsulation uses a tweaked version of Encapsulation module where  $u_{prime}$  and  $v_{prime}$  could be swapped with  $u_{in}$  and  $v_{in}$ .



# Decapsulation – Hardware Design Performance Comparison

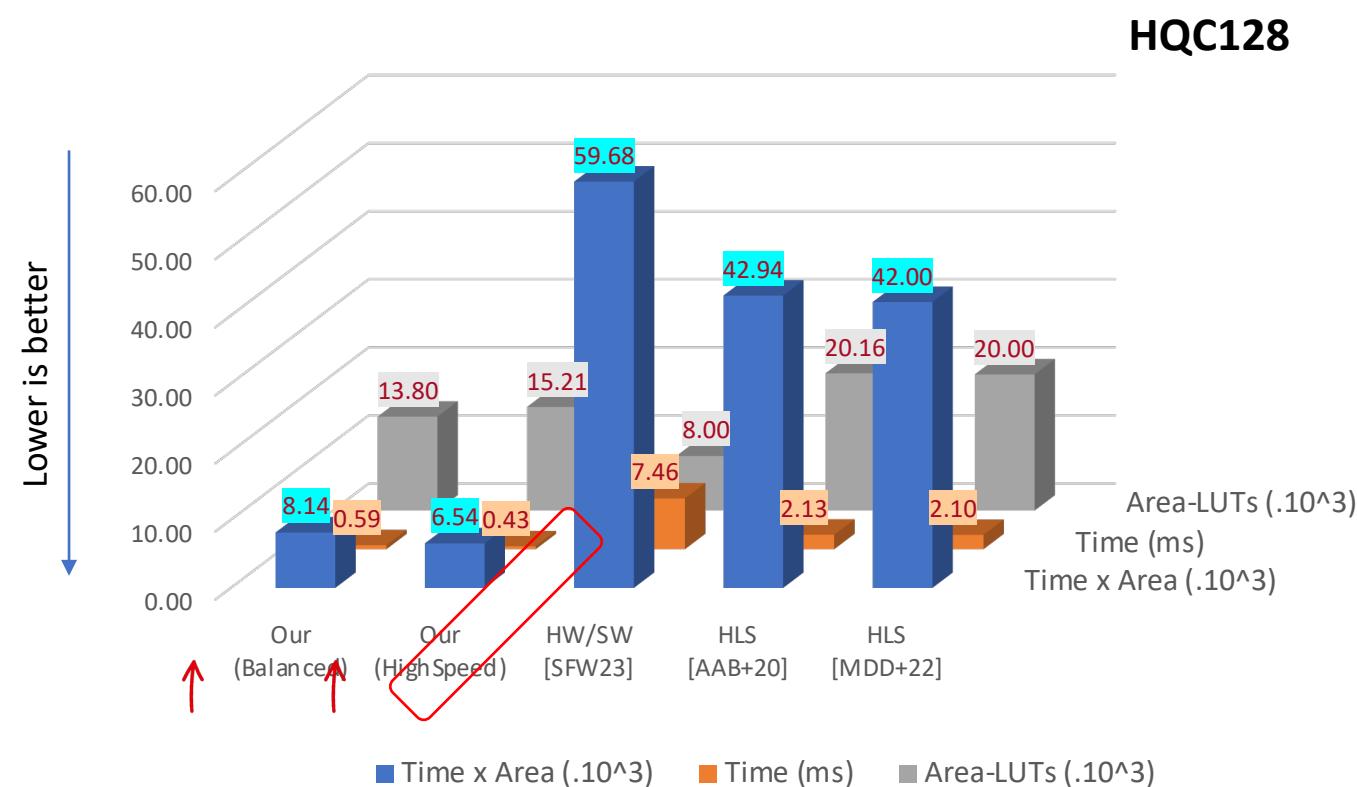
HQC128



caslab.io

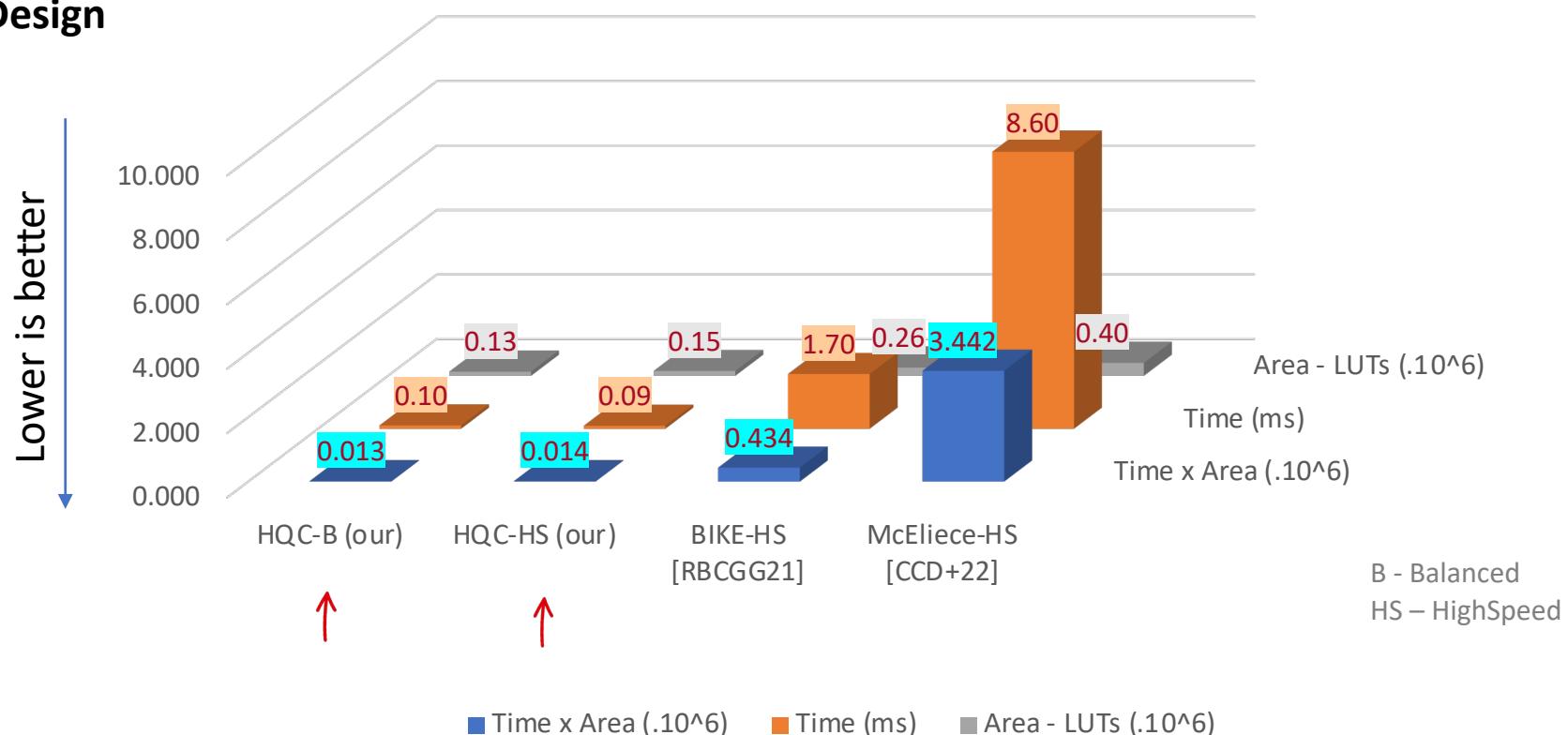
# Joint Design – Comparison with other HQC Designs

- Joint Design – combines KeyGen, Encap, and Decap in to one.
- Shared module among different primitives.
  - SHAKE256
  - Polynomial Multiplication
  - Encapsulation
  - Polynomial Addition
- Time = KeyGen + Encap + Decap.



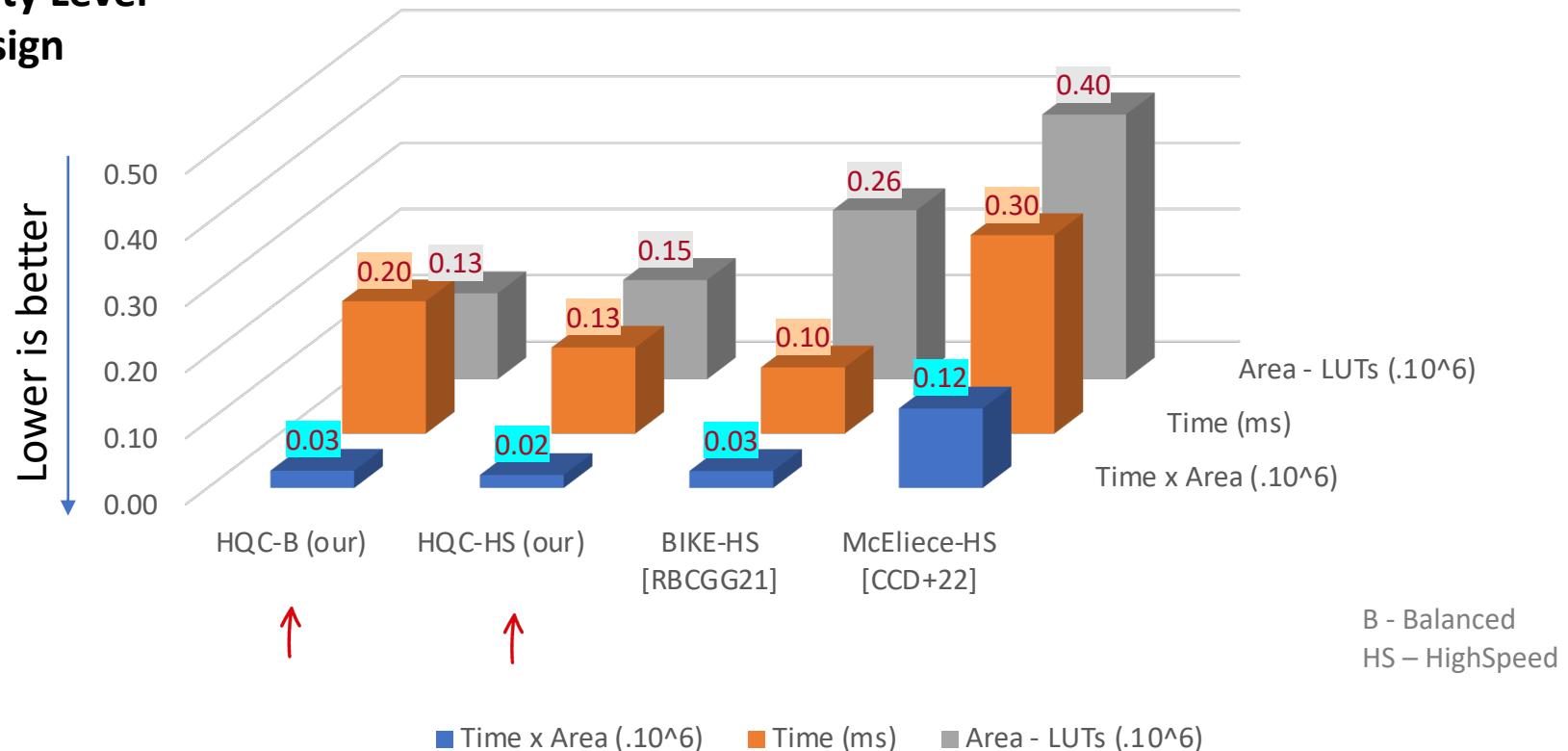
# Comparison with other Code Based Schemes - Key Generation

**128-bit Security Level**  
**Hardware Design**



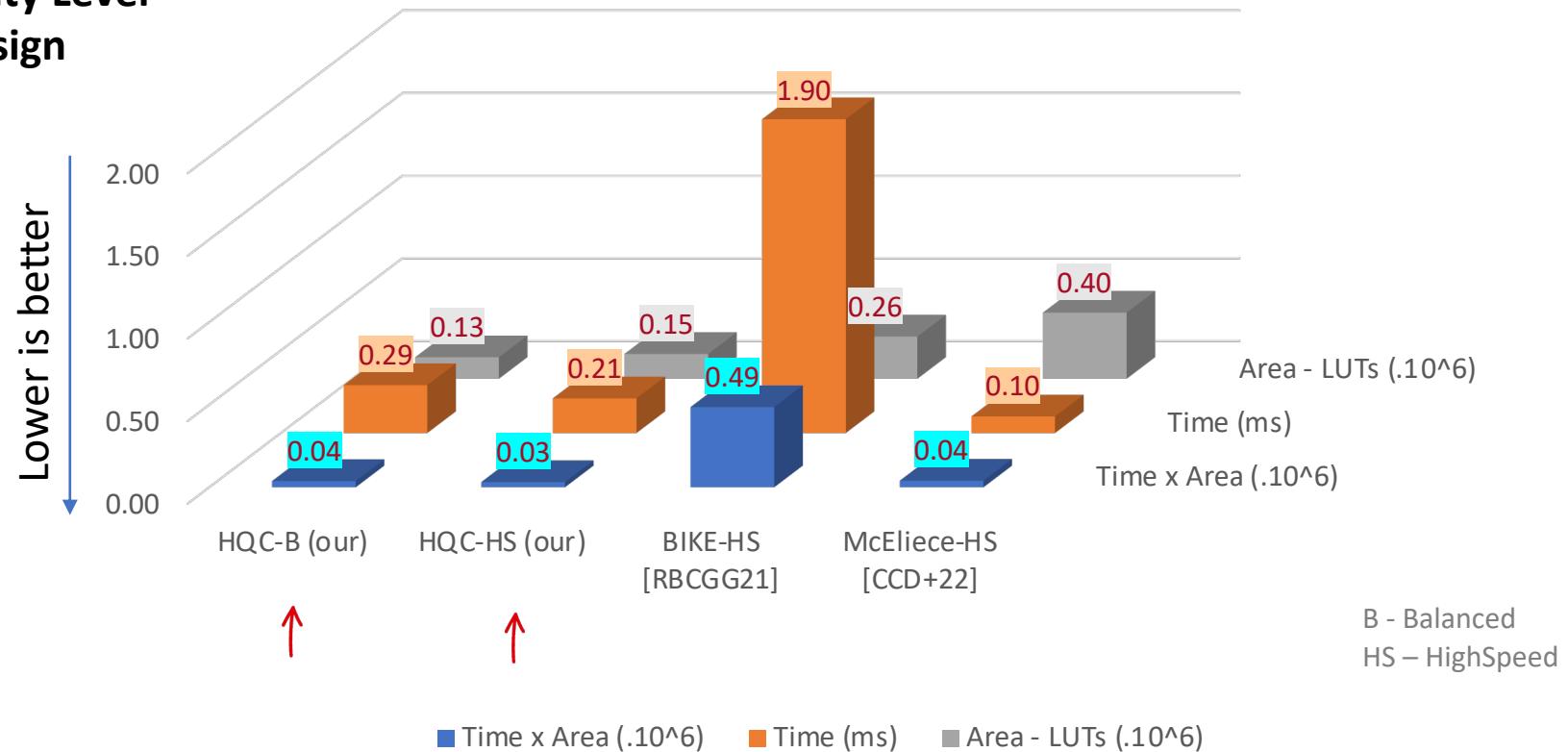
# Comparison with other Code Based Schemes - Encapsulation

## 128-bit Security Level Hardware Design



# Comparison with other Code Based Schemes - Decapsulation

## 128-bit Security Level Hardware Design



## Summary

- First (hand-tailored) hardware implementations of HQC Key Encapsulation Mechanism parameterizable at compile-time across all parameter sets.
- HQC can be a competitive candidate when optimized hardware is developed.

# References

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## Summary

- First (hand-tailored) hardware implementations of HQC Key Encapsulation Mechanism parameterizable at compile-time across all parameter sets.
- HQC can be a competitive candidate when optimized hardware is developed.

Thank you!  
Questions?



Link to the code base