

# Bit Manipulation

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## Introduction

Until now, we were using only integers, floating-point numbers, etc for our codes, but now we are going to look at the bits present in these data types for our calculations and problem-solving.

This is going to highly optimize our code, for example, an integer has 32 bits, now working on any of those 32 bits takes only  $O(1)$  time.

## Shift Operators

Shift operators are of two types:

1. **Left Shift:** This operator is represented by  $\ll$ . Example :  $8 \ll 1 = 16$

**This operator shifts every bit to its left position. For a clear understanding look at the image given below:**

Left Shift

$N = 8$

using left shift operator on  $N$

Now = 2 = 0 0 0 0 1 0 1 1

$\Rightarrow N \ll 1$

Binary Representation of 8 in 8 bits:-

$$\begin{array}{cccccccc} 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\ * & - & - & - & - & - & - & 0 \end{array} \Rightarrow (0 \ 0 \ 0 \ 1 \ 0 \ 0 \ 0 \ 0)_2 = (16)_{10}$$

The left most bit is lost & at the rightmost position 0 is added.

**General Formula:**  $N \ll i \rightarrow N * 2^i$

2. **Right Shift:** This operator is represented by  $\gg$ . Example :  $8 \gg 1 = 4$

This operator shifts every bit to its right position. For a clear understanding look at the image given below:

### Right Shift

$N = 8$

using right shift operator on  $N$

Now  $= 2 = 00001011$

$\Rightarrow N \gg 1$

Binary Representation of 8 in 8 bits:-

$00001000$   
↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓  
— — — — — — — —  $\Rightarrow 00000100$

If  $N$  is +ve , add 0

If  $N$  is -ve , add 1

**General Formula:**  $N \gg i \rightarrow \text{floor}(N/2^i)$

## Some other bitwise operators

1. **& (bitwise AND)** takes 2 numbers as operands and performs logical AND on every bit of the two numbers. The result of AND is 1 only if both bits are 1.
2. **| (bitwise OR)** takes 2 numbers as operands and performs logical OR on every bit of two numbers. The result of OR is 1 if any of the two bits is 1.
3. **~ (bitwise NOT)** takes 1 number and flips all the bits of it.
4. **^ (bitwise XOR)** takes 2 numbers as operands and performs XOR on every bit of 2 numbers. The result of XOR is 1 if the two bits are different.

| A | B | A   B | A & B | A ^ B | ~A |
|---|---|-------|-------|-------|----|
| 0 | 0 | 0     | 0     | 0     | 1  |
| 0 | 1 | 1     | 0     | 1     | 1  |
| 1 | 0 | 1     | 0     | 1     | 0  |
| 1 | 1 | 1     | 1     | 0     | 0  |

## Check $i^{\text{th}}$ bit

$i^{\text{th}}$  bit is set if its value is 1 and it is not set if its value is 0. We will be given a number **N** and a position **i** and we will have to check if the bit at  $i^{\text{th}}$  position is set or not. Now, let's take a look at an example to understand how we are going to proceed:

Given : **N = 5, i = 2**

The binary representation of 5 in 8 bits is **00000101**. We make a new number in 8 bit format that has only one bit set and that bit is at the  $i^{\text{th}}$  position. Then we take the **&** of these two numbers. If the answer is zero, then the  $i^{\text{th}}$  bit is set, else it is not set.

In the diagram, we have also shown an example where  $i^{\text{th}}$  is not set. The example has **N = 1, i = 2**.

$$N = (5)_{10} = (00000101)_2$$

$$New = (000000100)_2$$

$$\begin{array}{r} 00000101 \\ 00000100 \\ \hline 00000100 \end{array}$$

$$N = (1)_{10} = (00000001)_2$$

$$New = (000000100)_2$$

$$\begin{array}{r} 00000001 \\ 00000100 \\ \hline 00000000 \end{array}$$

## Flip $i^{\text{th}}$ bit

We know that :

$$1 \wedge 0 = 1$$

$$0 \wedge 0 = 0$$

Therefore,  $X \wedge 0 = X$  and

$$0 \wedge 1 = 1$$

$$1 \wedge 1 = 0$$

Therefore,  $X \wedge 1 = \sim X$

This means that to flip a bit we can take its XOR with  $2^i$ .

Example:

|  |  |
|--|--|
| $N = 11, i = 2$ <div style="text-align: center;"><math>\downarrow</math></div> $N = 00001011$ $N = 00001011$ $\text{Now} = 2^i = 00001011$ |  |
| <u>XOR</u>   | $\begin{array}{r} 00001011 \\ 00000101 \\ \hline 00001111 \end{array}$ |

$2^i$  can be calculated using the left shift operator, i.e.,  $2^i = 1 \ll i$ .

## Check N: odd or even

The basic way to do this is taking modulo by 2:

```
if ( N % 2 == 0 ) {  
    Even  
} else {  
    Odd  
}
```

**Optimized Approach:** Sometimes, in competitive programming our solution might not work by 0.1 seconds or a very small time, so to optimize this we can use bitwise operators.

Let's take a look at the binary representation of a few numbers:

```
11:  0 0 0 0 1 0 1 1  
8:   0 0 0 0 1 0 0 0  
9:   0 0 0 0 1 0 0 1  
10:  0 0 0 0 1 0 1 0
```

The last bit of every odd number is 1 and for every even number it is 0 because the last bit is represented by  $2^0$ . So we can just check the last bit of a number to check whether a number is odd or even. Try to implement this by using the concept of checking  $i^{\text{th}}$  bit taught before.

## Check N whether it is power of 2 or not

The basic way to do this is :

```
while ( ( N % 2 == 0 ) ) {  
    N = N / 2;  
}  
If ( N == 1 ) {  
    POWER OF 2  
} else {  
    NOT POWER OF 2  
}
```

The above approach runs in  $O(\log_2(N))$  time.

### Optimized Approach :

Let's take a look at the binary representation of some numbers and see if we can infer anything from them :

```
2:   0 0 0 0 0 0 0 1 0  
4:   0 0 0 0 0 0 1 0 0  
8:   0 0 0 0 0 1 0 0 0  
11:  0 0 0 0 0 1 0 1 1  
16:  0 0 0 0 1 0 0 0 0  
15:  0 0 0 0 0 1 1 1 1
```

There is one common pattern, that the numbers which are powers of 2 have only one bit set. Now take a look at the numbers 16 and 15 :

```
16:  0 0 0 0 1 0 0 0 0  
15:  0 0 0 0 0 1 1 1 1
```

If we take **16 & 15**, then we get 0. This happens for all numbers **N and N-1** if N is a power of 2.

**Generalized Formula:** If  $N \& (N-1) == 0$ , then N is a power of 2.

## Remove all set bits from LSB to i :

LSB or Least Significant Bit is the rightmost bit.

Let's take a binary number N and **i = 3**,

**N = 0 1 1 1 0 1 0 1 1 0**

Now, take another number **A = 1 << (i+1)**,

**A = 0 0 0 0 0 1 0 0 0 0**

**B = A-1 = 0 0 0 0 0 0 1 1 1 1**

**C = ~B = 1 1 1 1 1 1 0 0 0 0**

**Final answer = N & C = 0 1 1 1 0 1 0 0 0**

**Generalized Formula:**  $N \& M$  where  $M = \sim(1 \ll (i + 1)) - 1$