

```

let rec concat l = match l with
| [] -> ""
| x::xs -> x ^ (concat xs)

let concat' l =
let rec conc l acc = match l with
| [] -> acc
| x::xs -> conc xs (acc ^ x)
in
conc l ""

```

Prove that `concat l` produces the same as `concat' l`.

Intuitively this is clear as both functions start off with the same initial state and keep joining elements together until the list is empty, the only difference being that `concat'` uses an inner function `conc` so as to store the result in an accumulator variable. `concat'` is better in the sense that there is no danger of a stack overflow since the string we want to eventually return is being stored in 'acc' rather than on the stack.

formal proof: (this is extremely boring)

NTS `concat l = concat' l`

(Step 1) Base Case: `l = []`

```

concat l = concat [] = "" = conc [] "" = concat' [] = concat' l

```

the first equality is due to the definition of `l` in the base case

the second equality is by the program `concat`

the third and fourth equalities are by the program `concat'`

(Step 2) Assume true for list `l`

(Step 3) Inductive Step: Prove that equality still holds if we add an extra item to list `l`.

```

let item = additional generic item in list l
let l1 = item :: l

```

```

concat l1 = concat (item :: l) = item ^ (concat l)
                                = item ^ (concat' l)
                                = item ^ (conc l "")
                                = conc l item
                                = concat' (item :: l)
                                = concat' l1

```

first equality -> by definition of `ll`  
second equality -> by `concat` function  
third equality -> by step 2  
fourth equality -> by `concat'` function  
fifth equality -> by `concat'` function  
sixth equality -> by `concat'` function  
seventh equality -> by definition of `ll`

Q.E.D

Again, `concat'` is better in the sense that there is no danger of a stack overflow since the string we want to eventually return is being stored in `'acc'` rather than on the stack.