

# EE5327 Optimization

## Presentation 2

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February 28, 2019

## Question:

**38)** The following table gives the cost matrix of a transportation problem

4	5	6
3	2	2
1	1	2

The basic feasible solution given by  $x_{11}=3$ ,  $x_{13}=1$ ,  $x_{23}=6$ ,  $x_{31}=2$ ,  $x_{32}=5$  is

1. degenerate and optimal
2. optimal but not degenerate
3. degenerate but not optimal
4. neither degenerate nor optimal

## Theoretical Solution:

### Check for degeneracy:

To check whether the solution is non-degenerate,

$m+n-1$  = number of allocated cells,

where  $m$  and  $n$  are the dimension of the cost matrix.

Given cost matrix is of dimension  $3 \times 3$ .

We calculate  $m + n - 1 = 3 + 3 - 1 = 5$ , which is equal to the number of allocated cells i.e. 5.

So, we conclude the problem is non degenerate.

## Check for optimality:

We use UV or MODI method to check for optimality.  
Arrange the cost matrix as follows:

	$v_1$	$v_2$	$v_3$
$u_1$	4	5	6
$u_2$	3	2	2
$u_3$	1	1	2

Table: 1

and  $x_{11} = 3$ ,  $x_{13} = 1$ ,  $x_{23} = 6$ ,  $x_{31} = 2$ ,  $x_{32} = 5$ ,  $x_{12}$ ,  $x_{21}$ ,  $x_{22}$ ,  $x_{33} = 0$ .

**To find the  $u_i$  and  $v_j$**

For allocated cells, we find the  $u_i$  and  $v_j$  using the formula

$$u_i + v_j = C_{ij} \quad (1)$$

Where  $C_{ij}$  are the coefficient of the cost matrix.

Assuming  $u_1 = 0$ , we find all the values of  $u$  and  $v$ , the cost matrix table becomes:

	$v_1=4$	$v_2=4$	$v_3=6$
$u_1=0$	4	5	6
$u_2=-4$	3	2	2
$u_3=-3$	1	1	2

Table: 2

## Computing penalties

Computing the penalties for the unallocated cells using the following equation.

$$P_{ij} = u_i + v_j - C_{ij} \quad (2)$$

Penalties are  $P_{12} = -1$ ,  $P_{13} = -3$ ,  $P_{22} = -1$ ,  $P_{33} = 1$ ,

For the feasible solution to be optimal the penalties should be zero or negative. But here we see  $P_{33}$  is positive. So the solution is not optimal.

## Optimal Solution

We select the unallocated cell with positive penalty and draw a loop in the table starting from that cell through allocated cells.

The loop path:  $C_{33} \rightarrow C_{13} \rightarrow C_{11} \rightarrow C_{31}$

We alternately assign (+) and (-) signs to the cells in the loop starting from the first loop being assigned (+).

We select the minimum value of the two negatively assigned cells (in our case  $C_{31} = 1$ ) and subtract this value from the (-) assigned cells and add to the (+) assigned cells.

## Optimal Solution

Thus, the new table with new allocated values is:

	$v_1$	$v_2$	$v_3$
$u_1$	4	5	6
$u_2$	3	2	2
$u_3$	1	1	2

Table: 1

with  $x_{11} = 4$ ,  $x_{23} = 6$ ,  $x_{31} = 1$ ,  $x_{32} = 5$  and  $x_{33} = 1$ .

We repeat the above process for finding new  $u_i$  and  $v_j$



## Optimal Solution

Again assuming  $u_1 = 0$ , we find all the values using equation (1).

	$v_1=4$	$v_2=4$	$v_3=5$
$u_1=0$	4	5	6
$u_2=-3$	3	2	2
$u_3=-3$	1	1	2

Table: 2

## Optimal Solution Found

We find the penalty using equation (2).

$$P_{12} = -1, P_{13} = -1, P_{21} = -2, P_{22} = -1,$$

Thus, we see all the penalties are negative, which satisfies the condition for optimality.

Thus, the optimal values are

$$x_{11} = 4, x_{23} = 6, x_{31} = 1, x_{32} = 5 \text{ and } x_{33} = 1.$$

Putting these values in the objective function

$$4x_{11} + 5x_{12} + 6x_{13} + 3x_{21} + 2x_{22} + 2x_{23} + x_{31} + x_{32} + 2x_{33}$$

yields 36.

# Conversion to LP

The cost matrix with demand and supply is:

	$D_1$	$D_2$	$D_3$	
$S_1$	4	5	6	4
$S_2$	3	2	2	6
$S_3$	1	1	2	7
	5	5	7	

Table: 3

## LP Form

We convert the given problem (Table 3) to LP.

Minimize

$$4x_{11} + 5x_{12} + 6x_{13} + 3x_{21} + 2x_{22} + 2x_{23} + x_{31} + x_{32} + 2x_{33}$$

s.t.

$$x_{11} + x_{12} + x_{13} \leq 4 \quad (3)$$

$$x_{21} + x_{22} + x_{23} \leq 6 \quad (4)$$

$$x_{31} + x_{32} + x_{33} \leq 7 \quad (5)$$

$$-x_{11} - x_{21} - x_{31} \leq -5 \quad (6)$$

$$-x_{12} - x_{22} - x_{32} \leq -5 \quad (7)$$

$$-x_{31} - x_{23} - x_{33} \leq -7 \quad (8)$$

$$-x_{11}, \dots, -x_{33} \leq 0 \quad (9)$$

# Code using CVXPY

```
from cvxpy import *
from numpy import matrix

A = matrix([

    [1,1,1,0,0,0,0,0,0],
    [0,0,0,1,1,1,0,0,0],
    [0,0,0,0,0,0,1,1,1],
    [-1,0,0,-1,0,0,-1,0,0],
    [0,-1,0,0,-1,0,0,-1,0],
    [0,0,-1,0,0,-1,0,0,-1],
    [-1,0,0,0,0,0,0,0,0],
    [0,-1,0,0,0,0,0,0,0],
    [0,0,-1,0,0,0,0,0,0],
    [0,0,0,-1,0,0,0,0,0],
    [0,0,0,0,-1,0,0,0,0],
    [0,0,0,0,0,-1,0,0,0],
    [0,0,0,0,0,0,-1,0,0],
    [0,0,0,0,0,0,0,-1,0],
    [0,0,0,0,0,0,0,0,-1]])

b = matrix([ 4,6,7,-5,-5,-7,0,0,0,0,0,0,0,0])

c=matrix([4.0 , 5.0 , 6.0 , 3.0 ,
2.0 , 2.0 , 1.0 , 1.0 , 2.0])

x = Variable(9,1)
#Cost function
f = c*x
obj = Minimize(f)
#Constraints
constraints = [A*x <= b.transpose()]

#solution
Problem(obj, constraints).solve()
```

## Result

f = 36.00

x1 = 4.00

x2 = 0.00

x3 = 0.00

x4 = 0.00

x5 = 0.00

x6 = 6.00

x7 = 1.00

x8 = 5.00

x9 = 1.00

Thus, we see these values matches our theoretical solution.