

ECE 360P: Concurrent and Distributed Systems

Assignment 2

Instructor: Professor Vijay Garg

Deadline: 2:00 pm Feb. 11th, 2023

This homework contains 3 programming questions. Your submissions must be uploaded to Canvas by the deadline mentioned above. There will be a late penalty of 10% of the total grade for each day the assignment is late. You should use the homework 2 template available on the course GitHub. You should not change the file names and function signatures. In addition, you should not use package for encapsulation. The zip folder you submit should contain only java files with no sub-folders. The source files for both questions should be in the same folder. Please submit one zip file per team, with the name format [EID1_EID2].zip.

1. A `CyclicBarrier` is a synchronization aid that allows a set of threads to wait for all other threads to reach a common point. Whenever a thread calls the method `await`, it gets blocked until the number of parties specified in the constructor have also called `await`. When the specified number of threads have called `await`, all threads are released.

The implementation of a `CyclicBarrier` must allow for it to be reused after the waiting threads are released. In other words, a set of threads may use the same `CyclicBarrier` to synchronize across multiple barriers.

This implementation of a `CyclicBarrier` also has 2 states - active and inactive. An active `CyclicBarrier` exhibits the behavior specified above. In the inactive state, the `await` function in the `CyclicBarrier` is a no-op. Threads calling `await` are no longer blocked and proceed immediately.

```
public interface CyclicBarrier {  
  
    /*  
     * This method activates the cyclic barrier. If it is already in  
     * the active state, no change is made.  
     * If the barrier is in the inactive state, it is activated and  
     * the state of the barrier is reset to its initial value.  
     */  
    public void activate() throws InterruptedException;  
  
    /*  
     * This method deactivates the cyclic barrier.  
     * It also releases any waiting threads  
     */  
}
```

```

    */
    public void deactivate() throws InterruptedException;

    /*
     * An active CyclicBarrier waits until all parties have invoked
     * await on this CyclicBarrier. If the current thread is not
     * the last to arrive then it is disabled for thread scheduling
     * purposes and lies dormant until the last thread arrives.
     * An inactive CyclicBarrier does not block the calling thread. It
     * instead allows the thread to proceed by immediately returning.
     * Returns: the arrival index of the current thread, where index 0
     * indicates the first to arrive and (parties-1) indicates
     * the last to arrive.
     */
    public int await() throws InterruptedException;
}

```

- (a) **(15 points)** Implement a Java class `SemaphoreCyclicBarrier` that implements the `CyclicBarrier` interface using only semaphores.
- (b) **(15 points)** Implement a Java class `MonitorCyclicBarrier` that implements the `CyclicBarrier` interface using Java Monitor. (You are not allowed to use semaphores for its implementation)
2. **(35 points)** As you know, the Red River Showdown football game draws fans of UT and OU from everywhere to the Cotton Bowl. Naturally, some fans drink a lot of “apple juice,” and patrons constantly stumble to the bathroom throughout the game. However, there is only one bathroom shared between UT and OU fans. You have been tasked with designing a system conforming to the following rules:
- (a) The bathroom size is limited and can hold no more than 7 fans at any time.
 - (b) The bathroom can only hold fans from one team at any time, otherwise a fight might break out. More specifically, if there is a single UT fan in the bathroom there can be no OU fans in the bathroom and vice versa.
 - (c) We require some sort of fairness in that a fan requesting the bathroom is blocked until all other fans preceding them have entered the bathroom, regardless of the team (just as if fans were waiting in a line for the bathroom). Use a *ticketNumber* that ensures this fairness. Do not worry about the integer overflow for the ticket number.

Implement a Java class `FairUnifanBathroom` satisfying these rules using Java Monitors or ReentrantLock and Condition. It should support the following methods:

```

public class FairUnifanBathroom {
    public void enterBathroomUT() {...}
    public void enterBathroomOU() {...}
    public void leaveBathroomUT() {...}
    public void leaveBathroomOU() {...}
}

```

3. **(35 points)** Use Java ReentrantLock and Condition to implement a linked list based priority queue `PriorityQueue`. Each node in the queue has two fields: `name` (a `String`) and `priority` (an integer in the range 0..9). The order of the nodes is always kept sorted based on the priority with 9 as the highest priority and 0 as the least. Nodes with the same priority number are kept in the order of inserts. Furthermore, you should not use a global lock for the queue, but instead use the “hand-over-hand” locking technique discussed in class. Your `PriorityQueue` should support the following methods.

```
public class PriorityQueue {
    public PriorityQueue(int capacity) {
        // Creates a Priority queue with maximum allowed size as capacity
    }
    public int add(String name, int priority) {
        // Adds the name with its priority to this queue.
        // Returns the current position in the list where the name was inserted;
        // otherwise, returns -1 if the name is already present in the list.
        // This method blocks when the list is full.
    }
    public int search(String name) {
        // Returns the position of the name in the list;
        // otherwise, returns -1 if the name is not found.
    }
    public String getFirst() {
        // Retrieves and removes the name with the highest priority in the list,
        // or blocks the thread if the list is empty.
    }
}
```