Assignment 10 – 29 - May - 2023

Implementation of LR parser

LR Parsing algorithm:

Input: An input string w and an LR parsing table with functions action and goto for grammar G.

Output: If w is in L(G), a bottom-up-parse for w; otherwise, an error indication.

Method: Initially, the parser has s_0 on its stack, where s_0 is the initial state, and w\$ in the input buffer. The parser then executes the following program :

```
set ip to point to the first input symbol of w$;
repeat forever begin
      let s be the state on top of the stack and
          a the symbol pointed to by ip;
      if action[s, a] = shift s' then begin
          push a then s' on top of the stack;
          advance ip to the next input symbol
      end
      else if action[s, a] = reduce A \rightarrow \beta then begin
           pop 2* | \beta | symbols off the stack;
           let s' be the state now on top of the stack;
           push A then goto[s', A] on top of the stack;
           output the production A \rightarrow \beta
      end
      else if action[s, a] = accept then
           return
      else error()
end
```

Write a Program to implement the LR parser. Consider 2 input Strings (to show accept and reject/error) for the Grammar given below & show the sequence of moves by input String using LR Parser.

[NOTE : Accept input via file]

1. Roll_No: 2110342211IT001 - 2110972211IT016

Grammar							
S →	E						
	E + T T						
T →	T * F F						
F →	id						

States	Action						
	id	+	*	\$	E	T	F
I ₀	S ₄				1	2	3
Iı		S ₅		Accept			
I_2		R_2	S ₆	R2			
I3		R ₄	R ₄	R4			
I ₄		R5	R5	R5			
I5	S4					7	3
I_6	S4						8
I ₇		R1	S6	R1			
I ₈		R3	R3	R3			

2. Roll_No:2110491211IT017 - 2110491211IT033

Grammar
$R \rightarrow S$
$S \rightarrow (L)L$
$S \rightarrow \epsilon$
$L \rightarrow int L$
$L \rightarrow S$

			GOTO			
State	()	int	\$	S	L
0	s2	r3		r3	1	
1				acc		
2	s2	r3	s4	r3	5	3
3		s6				
4	s2	r3	s4	r3	5	7
5		r5		r5		
6	s2	r3	s4	r3	5	8
7		r4		r4		
8		r2		r2		

3. Roll_No:2110491211IT034 - 2110491211IT051

Grammar

$$\mathbf{E} \rightarrow \mathbf{E} + \mathbf{T}$$

 $E \to T$

$$T \ \to \ T \ \hbox{\rm *} \ F$$

 $T \to F$

$$\mathbf{F} \rightarrow (\mathbf{E})$$

 $F \ \rightarrow \ id$

	Action							Goto			
State	id	+	*	()	\$	E	Т	F		
0	\$5			S4			1	2	3		
1		\$6				accept					
2		R2	S7		R2	R2					
3		R4	R4		R4	R4					
4	\$5			\$4			8	2	3		
5		R6	R6		R6	R6					
6	\$5			S4				9	3		
7	\$5			S4					10		
8		\$6			S11						
9		R1	S7		R1	R1					
10		R3	R3		R3	R3					
11		R5	R5		R5	R5					

4. Roll_No:2110491211IT052 - 2110491211IT068

Grammar

 $S' \rightarrow S$

 $S \rightarrow A S$

 $S \rightarrow b$

 $A \rightarrow S A$

 $A \rightarrow c$

ACTION GOTO								
State	c	b	\$	S	A			
0	s4	s3		1	2			
1	s4	s3	acc	6	5			
2	s4	s3		7	2			
3	r3	r3	r3					
4	r5	r5						
5	s4/r4	s3/r4		7	2			
6	s4	s3		6	5			
7	s4/r2	s3/r2	r2	6	5			

5. Roll_No:2110491211IT069 - 2110491211IT087

Grammar

 $S \to \boldsymbol{x}$

 $S \rightarrow (L)$

 $L \to S \,$

 $L \to L$, S

		GOTO					
	()	x	,	\$	S	L
1	s 3		s 2			4	
2	r 1	r 1	r 1	r 1	r 1		
3	s 3		s 2			6	5
4					Acc		
5		s 7		s 8			
6	r 3	r 3	r 3	r 3	r 3		
7	r 2	r 2	r 2	r 2	r 2		
8	s 3		s 2			- 9	
9	r 4	r 4	r 4	r 4	r 4		