# IT201 COMPUTER ORGANIZATION AND ARCHITECTURE

- 1)BASIC STRUCTURE OF COMPUTERS: Basic Computer Organization CPU Organization, Functional unit, Basic OPERATIONAL concepts, Bus structures, Performance, Data Representation. Fixed Point. Representation. Floating Point Representation. Error Detection codes.
- 2) REGISTER TRANSFER LANGUAGE AND MICRO OPERATIONS: Register Transfer language. Register Transfer Bus and memory transfers, Arithmetic Micro operations, logic micro operations, shift micro operations, Arithmetic logic shift unit.

- 3) Basic Computer Organization and Design: Instruction codes. Computer Registers Computer instructions, Timing and Control, Instruction cycle. Memory Reference Instructions, Input Output and Interrupt,
- 4) Micro Programmed Control: Control memory, Address sequencing, micro program example, design of control unit, micro Programmed control

STACK organization, Instruction Formats, Addressing modes, Data Transfer and Manipulation , Program Control. CISC and RISC.

5) Computer Processing Unit Organization: General Register Organization

6) Computer Arithmetic: Addition and subtraction, multiplication Algorithms,

Division Algorithms, Floating – point Arithmetic operations. BCD Adder

Asynchronous data transfer Modes of Transfer, Priority Interrupt Direct memory Access, Input –Output Processor (IOP) DMA controller

7) Input-Output Organization: Peripheral Devices, Input-Output Interface,

8) Pipeline And Vector Processing: Parallel Processing, Pipelining, Arithmetic Pipeline, Instruction Pipeline, Dependencies, Vector Processing.

9)Memory Organization: Memory Hierarchy, Main Memory –RAM And ROM Chips, Memory Address map, Auxiliary memory-magnetic Disks, Magnetic tapes, Associate Memory,-Hardware Organization, Match Logic, Cache Memory –Associative Mapping, Direct Mapping, Set associative mapping, Writing in to cache and cache Initialization, Cache Coherence, Virtual memory-Address Space and memory Space, Address mapping using pages, Associative memory page table, page Replacement.

#### **TEXTBOOKS**

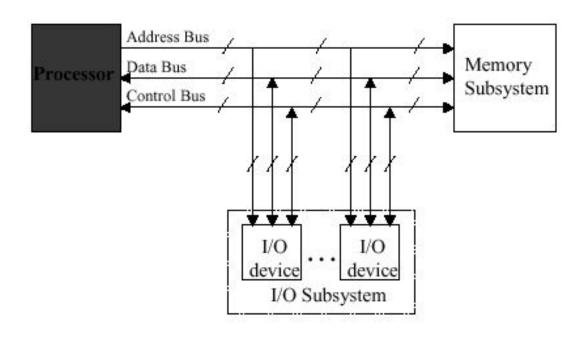
- Carl Hamacher et al., Computer Organization and Embedded Systems, Sixth Edition, McGraw-Hill, 2014.
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- Miles Murdocca and Vincent Heuring, Computer Architecture & Organization An Integrated Approach, Wiley, 2007.
- J. Hennesy and D. Patterson, Computer Architecture –A Quantitative Approach, 6th Ed., Morgan Kaufmann, 2017.
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#### **REFERENCES:**

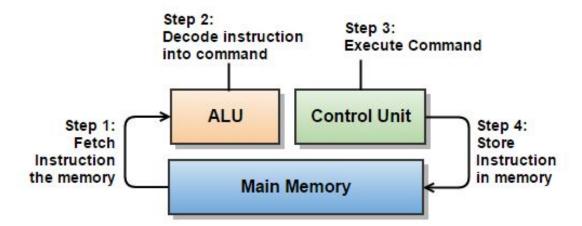
- Computer Organization and Architecture William Stallings Sixth Edition, Pearson/PHI
- Structured Computer Organization Andrew S. Tanenbaum, 4th Edition PHI/Pearson
- Fundamentals or Computer Organization and Design, Sivaraama Dandamudi Springer Int. Edition.

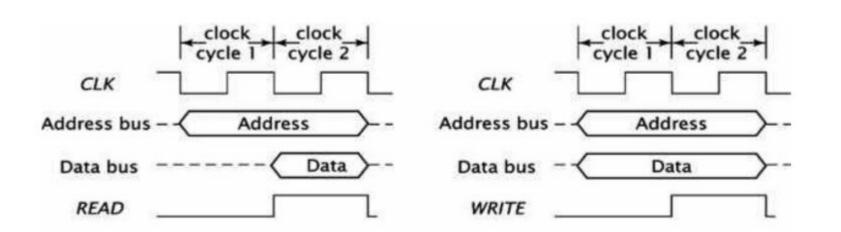
- 1) Quizzes : 20 Marks
- 2) Mid-sem: 30 Marks
- 3) End-sem: 50 Marks

## VON NEUMANN ARCHITECTURE



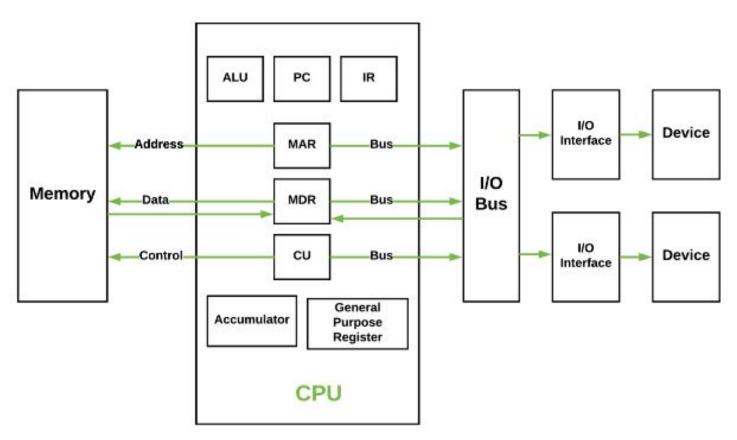
- 1. System bus: wires
- 2. Instruction cycles: fetch, decode and execute instructions





Timing diagram for read and write operations

# **CPU Organization**



- Main Memory Unit (Registers)
  - Accumulator
  - Program Counter (PC)
  - Memory Address Register (MAR)
  - Memory Data Register (MDR)
  - Instruction Buffer Register (IBR)

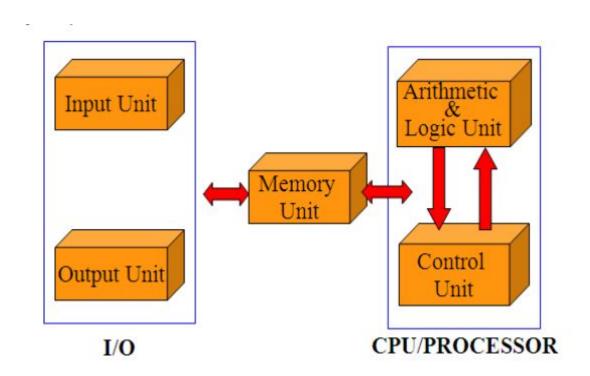
Current Instruction Register (CIR)

- inotituditon Buildi Regiotor (IBIR
- Input/Output Devices
- 1. Data Bus

Buses

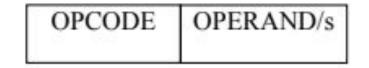
- 2. Address Bus
- 3. Control Bus

## **Functional Unit**



# **Basic Operational Concepts**

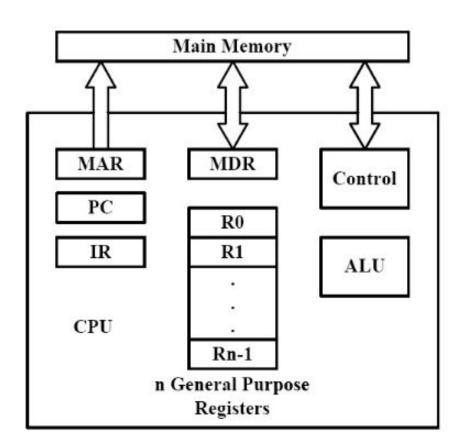
An Instruction consists of two parts, an Operation code and operand/s



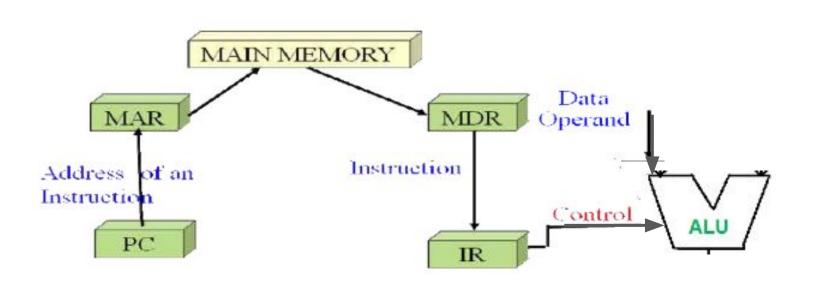
Eg 1) ADD LOCA, R0

Operation mnemonic	Name	Number of bits transferred	Description
L	Load	32	Transfer from memory in register
LH	Store half-word	16	Transler from memory to register
ST	store	32	Transfer from register to memory
STH	Store half-word	16	Transfer from register to memory

# Connections between the processor and the memory

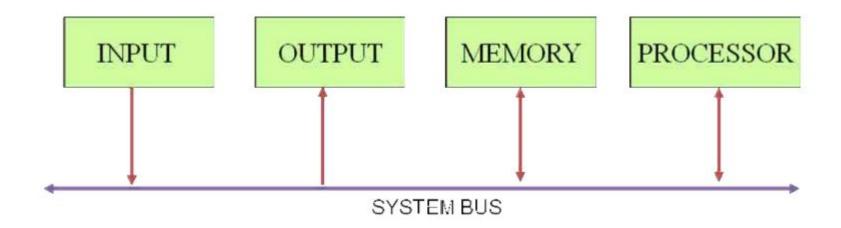


# Interaction between the memory and the ALU



#### **BUS STRUCTURES**

Group of lines that serve as connecting path for several devices is called a bus



## USER PROGRAM and OS ROUTINE INTERACTION

Assume computer with 1 processor, 1 disk and 1 printer and application program is in machine code on disk.

t0: the OS loads the program from the disk to

memory

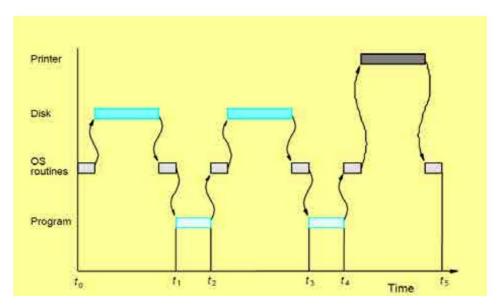
t1: program executes

t2: program accesses disk

t3: program executes

t4: program accesses printer

t5: program terminates



#### GENERATION OF COMPUTERS

**First generation**: 1946 to 1955: Computers of this generation used **Vacuum Tubes**. The computers were built using stored program concept. Ex: ENIAC, EDSAC, IBM 701. Computers of this age typically used about ten thousand vacuum tubes. They **were bulky in size had slow operating speed, short life time and limited programming facilities.** 

**Second generation**: 1955 to 1965: Computers of this generation used the **germanium transistors** as the active switching electronic device. Ex: IBM 7000, B5000, IBM 1401. Comparatively **smaller in size About ten times faster operating speed** as compared to first generation vacuum tube based computers. Consumed **less power, had fairly good reliability.** Availability of large memory was an added advantage.

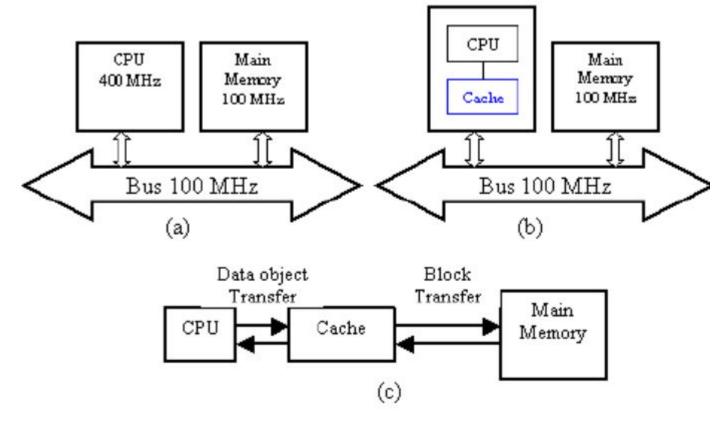
**Third generation**: 1965 to 1975: The computers of this generation used the **Integrated Circuits** as the active electronic components. Ex: IBM system 360, PDP minicomputer etc. They were still smaller in size. They had powerful CPUs with the capacity of executing 1 million instructions per second (MIPS). Used to consume very less power consumption.

**Fourth generation**: 1976 to 1990: The computers of this generation used the **LSI chips** like microprocessor as their active electronic element. HCL horizen III, and WIPRO"S Uniplus+ HCL"s Busybee PC etc. **They used high speed microprocessor as CPU.** They were more user friendly and highly reliable systems. They had large storage capacity disk memories.

**Beyond Fourth Generation**: 1990 onwards: Specialized and dedicated **VLSI chips** are used to control specific functions of these computers. Modern Desktop PC"s, Laptops or Notebook Computers.

#### PERFORMANCE

- The total time required to execute the program is elapsed time is a measure of the performance of the entire computer system.
- The time needed to execute a instruction is called the processor time.



a) CPU, main memory, and bus; (b) Processor cache; (c) Data transfer between CPU and cache and between cache and main memory

#### Processor clock

- Processor circuits are controlled by a timing signal called clock.
- The clock designs at the regular time intervals called clock cycles.

# Performance equation

T=N\*S/R

- T be the processor time required to execute a program
- N is the number of instruction
- S average number of basic steps executed in one clock cycle.
- R cycles per second

# Data Representation

Information that a Computer is dealing with is Data

Numeric Data: Numbers(Integer, real)

Non-numeric Data: Letters, Symbols

Relationship between data elements: Data Structures - Linear Lists, Trees, Rings, e

Program(Instruction)

# Numeric Data Representation

Fixed	
Point	

Decimal	Binary	Octal	Hexadecimal
00	0000	00	0
01	0001	01	1
02	0010	02	2
03	0011	03	3
04	0100	04	4
05	0101	05	5
06	0110	06	6
07	0111	07	7
08	1000	10	8
09	1001	11	9
10	1010	12	A
11	1011	13	В
12	1100	14	C
13	1101	15	D
14	1110	16	Е
15	1111	17	F

- Sign and Magnitude Representation: The most significant (leftmost) bit in the word as a sign bit. If the sign bit is 0, the number is positive; if the sign bit is 1, the number is negative.
- 2. One's Complement (1's) Representation: Forming the 1s complement of a given number is equivalent to subtracting that number from 2n -1 Eg.1s complement of 0101 is 1010
- **3.Two**"s Complement (2"s) Representation: Forming the 2s complement of a number is done by subtracting that number from 2n. So 2s complement of a number is obtained by adding 1 to 1s complement of that number. Eg: 2"s complement of 0101 is 1010 + 1 = 1011

In all systems, the leftmost bit is 0 for positive number and 1 for negative number.