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Binary Search



- Binary search is used to search an element from a list of elements in sorted order
- The searching starts from middle element
- If searching value does not match with middle element, then compare it either it is smaller or greater than middle element
- If searching value is smaller than middle element, then drop upper half and consider lower half for new list for searching and vice versa.

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1 Data Structures & Algorithms (Binary Search & Its Complexity)

2 Binary Search

3 Binary Search Example

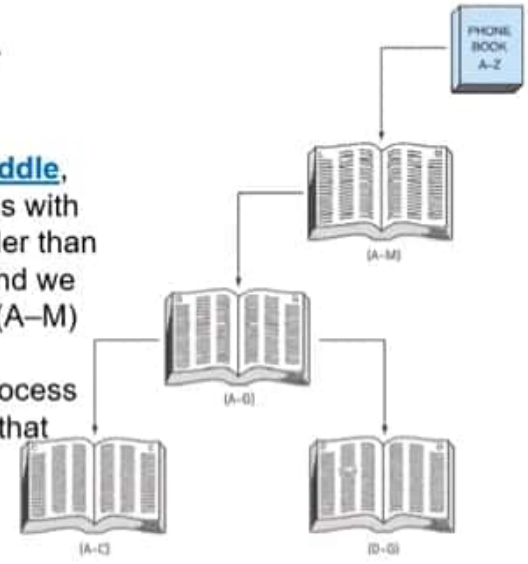
4 Binary Search Algorithm

5 Binary Search...

Binary Search Example



- Phone Book in sorted order
- Search "Daud" contact #
- We open the book from middle, found there are names begins with letter 'M' whereas 'D' is smaller than 'M', so we drop upper half, and we consider the lower half from (A-M)
- Again perform the above process until we found a single page that contains "Daud"



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Binary Search...



Working: Searching Value i.e. K = 40

BEG			MID			END			
1	2	3	4	5	6	7	8	9	10
5	10	15	20	25	30	35	40	45	50

MID = (1+10) / 2 = 5.5

1. Here K ≠ DATA [MID] i.e. 40 ≠ 25 But K > DATA[MID] so BEG = MID+1 i.e. BEG = 6

					BEG	MID	END		
1	2	3	4	5	6	7	8	9	10
5	10	15	20	25	30	35	40	45	50

MID = (6+10) / 2 = 8

2. Here K = DATA [MID] i.e. 40 = 40, So searching value found at MID location

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Binary Search...



Working: Searching Value i.e. K = 50

BEG				MID				END		
1	2	3	4	5	6	7	8	9	10	$MID = (1+10) / 2 = 5.5$
5	10	15	20	25	30	35	40	45	50	

1. Here $K \neq DATA[MID]$ i.e. $50 \neq 25$, But $K > DATA[MID]$ so $BEG = MID+1$ i.e. $BEG = 6$

					BEG	MID			END	
1	2	3	4	5	6	7	8	9	10	MID = (6+10) / 2 = 8
5	10	15	20	25	30	35	40	45	50	

2. Here $K \neq DATA[MID]$ i.e. $50 \neq 40$, Also $K > DATA[MID]$ so $BEG = MID+1$ i.e. $BEG = 9$

								MID	BEG	END	
1	2	3	4	5	6	7	8	9	10	MID = (9+10) / 2 = 9.5	
5	10	15	20	25	30	35	40	45	50		

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Binary Search...



Working: Searching Value i.e. K = 50

								MID			MID = (9+10) / 2 = 9.5
								BEG	END		
1	2	3	4	5	6	7	8	9	10		
5	10	15	20	25	30	35	40	45	50		

3. Here $K \neq \text{DATA}[\text{MID}]$ i.e. $50 \neq 45$, Also $K > \text{DATA}[\text{MID}]$ so $\text{BEG} = \text{MID} + 1$ i.e. $\text{BEG} = 10$

										MID BEG END	MID = (10+10) / 2 = 10
1	2	3	4	5	6	7	8	9	10		
5	10	15	20	25	30	35	40	45	50		

4. Here $K = \text{DATA}[\text{MID}]$ i.e. $50 = 50$, So, searching element is found at location MID

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Binary Search Algorithm



Study Point

(BINARY SEARCH) This algorithm will find an element K from a list DATA with N elements using Binary Search concept. The list DATA is already in sorted order.

Steps:

1. Set $BEG := 1$ and $END := N$
2. Repeat Step 3, and 4 While $BEG \leq END$
3. Set $MID := \text{INT}((BEG + END) / 2)$
4. IF $K = \text{DATA}[MID]$ THEN
 Write : 'ITEM FOUND'
ELSE IF $K < \text{DATA}[MID]$ THEN
 Set $END := MID - 1$
ELSE IF $K > \text{DATA}[MID]$ THEN
 Set $BEG := MID + 1$
ELSE
 Write : 'NOT FOUND'
[End of IF Structure]
[End of repeat loop]
5. Exit

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Complexity of Binary Search



- Let say the iteration in Binary Search terminates after k iterations. In the above example, it terminates after 3 iterations, so here $k = 3$
- At each iteration, the array is divided by half. So let's say the length of array at any iteration is n
- At Iteration 1,
Length of array = n
- At Iteration 2,
Length of array = $n/2$
- At Iteration 3,
Length of array = $n/2/2 = n/2^2$
- Therefore, after Iteration k ,
Length of array = $n/2^k$
- Also, we know that after
After k divisions, the length of array becomes 1
- Therefore
Length of array = $n/2^k = 1$
 $\Rightarrow n = 2^k$
- Applying log function on both sides:
 $\Rightarrow \log_2 (n) = \log_2 (2^k)$
 $\Rightarrow \log_2 (n) = k \log_2 (2)$
- As $(\log_2 (2) = 1)$
Therefore,
 $\Rightarrow k = \log_2 (n)$

fence, the time complexity of Binary Search is

$$\log_2 (n)$$

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Complexity of Binary Search



Study Point

- The complexity is measured by the number $f(n)$ of comparisons to locate the ITEM in DATA where DATA contains n elements.
- Observe that each comparison reduces the sample size in half.
- So,

$$f(n) = \log_2(n) = O(\log_2 n)$$

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