CSE 431/531: Algorithm Analysis and Design (Fall 2024) Introduction II: Algorithm and Asymptotic Analysis

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Outline

- Introduction
 - What is an Algorithm?
 - More Computation Problems

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- Algorithm: Euclidean algorithm
- $gcd(270, 210) = gcd(210, 270 \mod 210) = gcd(210, 60)$
- $(270,210) \rightarrow (210,60) \rightarrow (60,30) \rightarrow (30,0)$

Algorithm = Computer Program?

- Algorithm: "abstract", can be specified using computer program, English, pseudo-codes or flow charts.
- Computer program: "concrete", implementation of algorithm, using a particular programming language

Pseudo-Code

Pseudo-Code:

Euclidean(a, b)

- 1: while b > 0 do
- 2: $(a,b) \leftarrow (b, a \mod b)$
- 3: return a

Python program:

- def euclidean(a: int, b: int):
- c = 0

•

- while b > 0:
 - c = b
 - $\mathsf{b} = \mathsf{a}~\%~\mathsf{b}$
 - $\mathsf{a}=\mathsf{c}$
- return a

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 - fundamental
 - it is fun!

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Sorting

Input: sequence of n numbers (a_1, a_2, \dots, a_n)

Output: a permutation $(a'_1, a'_2, \cdots, a'_n)$ of the input sequence such

that $a_1' \leq a_2' \leq \cdots \leq a_n'$

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Example:

• Input: 53, 12, 35, 21, 59, 15

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• Algorithms: insertion sort, merge sort, quicksort, ...

Insertion-Sort

ullet At the end of j-th iteration, the first j numbers are sorted.

```
iteration 1: 53, 12, 35, 21, 59, 15
iteration 2: 12, 53, 35, 21, 59, 15
iteration 3: 12, 35, 53, 21, 59, 15
iteration 4: 12, 21, 35, 53, 59, 15
iteration 5: 12, 21, 35, 53, 59, 15
iteration 6: 12, 15, 21, 35, 53, 59
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- $\bullet \ \, \mathsf{Input:} \ \, 53,12,35,21,59,15$
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insertion-sort(A, n)

- 1: **for** $j \leftarrow 2$ to n **do**
- 2: $key \leftarrow A[j]$
- 3: $i \leftarrow j-1$
- 4: while i > 0 and A[i] > key do
- 5: $A[i+1] \leftarrow A[i]$
- 6: $i \leftarrow i 1$
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- j = 6
- key = 15
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- **†**
 - i

Analysis of Insertion Sort

- Correctness
- Running time

Correctness of Insertion Sort

• Invariant: after iteration j of outer loop, A[1..j] is the sorted array for the original A[1..j].

```
after j=1:53,12,35,21,59,15

after j=2:12,53,35,21,59,15

after j=3:12,35,53,21,59,15

after j=4:12,21,35,53,59,15

after j=5:12,21,35,53,59,15

after j=6:12,15,21,35,53,59
```

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 - For the insertion sort algorithm: if input array is already sorted in ascending order, then algorithm runs much faster than when it is sorted in descending order.

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- Q2: Which input?
 - For the insertion sort algorithm: if input array is already sorted in ascending order, then algorithm runs much faster than when it is sorted in descending order.
- A2: Worst-case analysis:
 - \bullet Running time for size n= worst running time over all possible arrays of length n

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Important idea: asymptotic analysis

 Focus on growth of running-time as a function, not any particular value.

Exercise

- Ex 1: What is the worst-case runtime for the Insertion Sort algorithm when the input sequence A is already sorted?
 - A is sorted in ascending order
 - A is sorted in descending order
- Ex 2: What is the worst-case runtime for the Insertion Sort algorithm with any sequence input A?