

Authentication I

CSE 565: Fall 2024
Computer Security

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Announcement

- Please sign-up at course Piazza.
- Reminder of Quiz 0 (**Due 09/19**).
 - You must obtain full score of the Quiz.

Review of Last Lecture

- Key Exchange: how we establish the shared secret key
 - Diffie-Hellman Key Ex protocol
- Public Key Cryptography
 - Trapdoor function
 - RSA
 - ElGamal: PKC from Diffie-Hellman
 - Digital Signature & Certificates

Today's Topic

- Common techniques for authenticating users, locally and remotely
 - **Password-based authentication**
 - Token-based authentication
 - Biometric-based authentication
- Security challenges associated with different authentication methods
- Mitigations designed to address some of the above security challenges

Overview

Authentication

- Previously: Message Authentication
 - Message Authentication Code (Keyed Hash) to confirm that the message came from the stated sender (its authenticity) and has not been changed in transit (its *integrity*).
- Today: **User/entity Authentication**
 - Allow a user/computer to prove his/her/its identity to another entity (e.g., a system, a device, a remote server).

Entity Authentication

- How do we authenticate a human user to a system?
 - System is often remote server
 - Authenticate: ascertain who is interacting with the system
 - Necessary to apply appropriate security policy.
 - Only the intended subject should be able to authenticate to the system as that subject.



Entity Authentication

- How do we authenticate a human user to a machine?
 - Provide *identity* and *proof* of identity
 - Identity examples:
 - Name, username, student ID, others?
- User registration is required prior to an authentication protocol



Entity Authentication

- How can Alice prove that she's really Alice?
 - Three types of authentication factors
 - ▶ Password: **Something you know.**
 - ▶ Token: **Something you have.**
 - ▶ Biometrics: **Something you are.**
- Each factor can be used independently, or combined for multi-factor authentication.
 - ▶ Typically two-factor

Password-based Authentication: Basics

“Something you know”

- A **secret** that only the real Alice should know
 - ▶ A secret *passcode*.
 - Examples: PIN, password
 - PIN: Personal Identification Number (misnomer. Usually used for authentication, not identification.)
 - ▶ A secret about Alice
 - Examples: mother's maiden name, first pet, mortgage payment
- Technically, only proves knowledge of secret, not that it's really Alice
 - ▶ Secrets leak, can be shared, guessed.

“Something you know”



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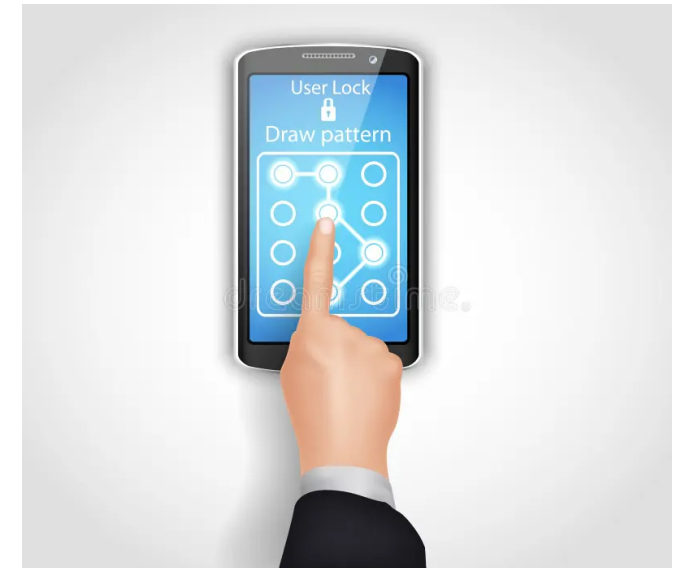
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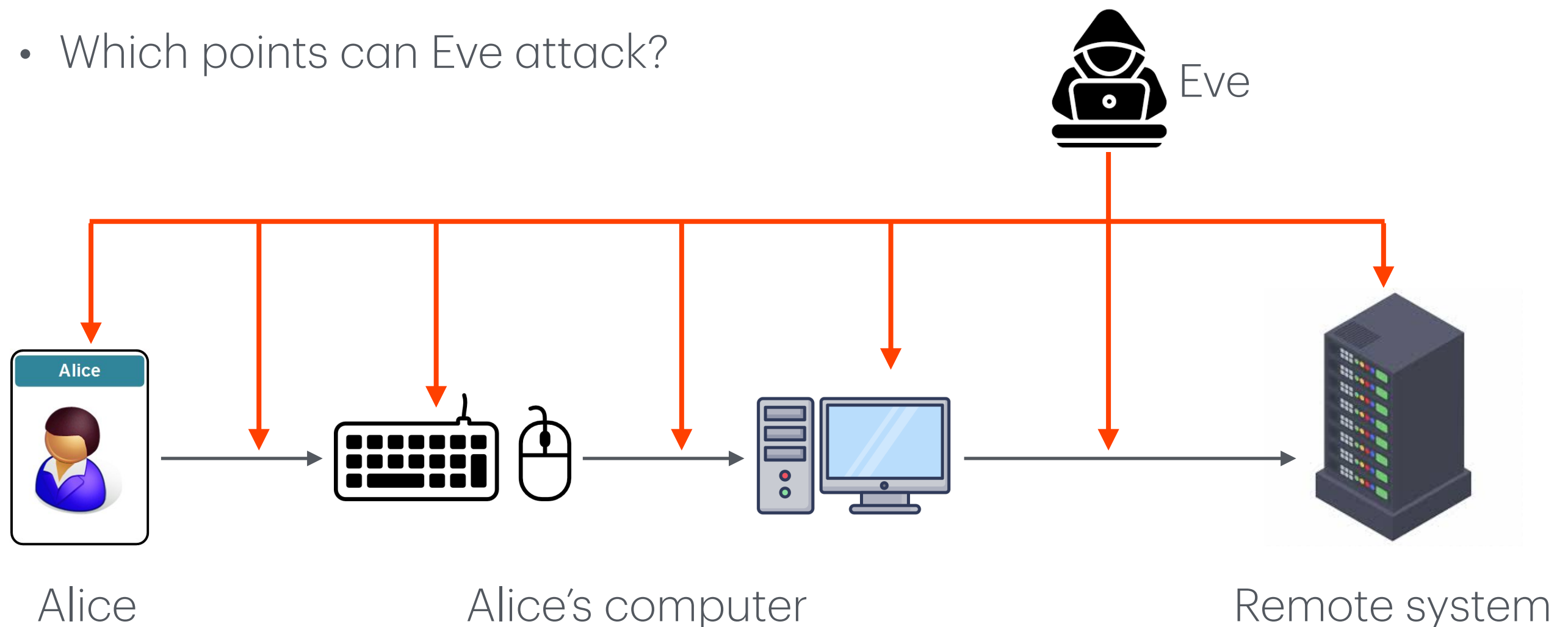
Password-Based Authentication

- How does Alice prove she knows the password?
 - Simplest: Alice provides the password to the system.
- Problems?
 - **Passive** adversary may observe password in transit
 - Need secure channel to protect confidentiality
 - **Active** adversary may impersonate the system
 - Alice needs a way of *authenticating the system*

Attacking Passwords

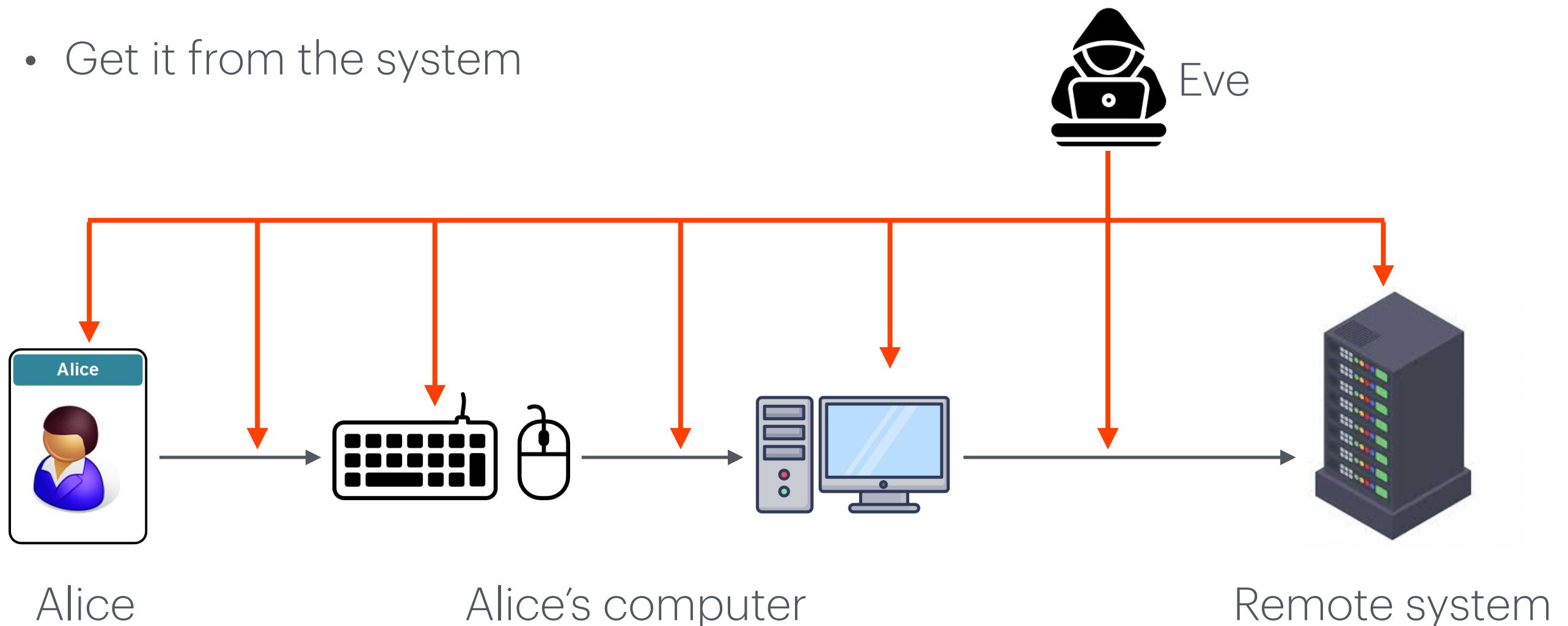
- Alice uses a keyboard to type her password into client software (e.g. browser) that sends it on to the remote system for authentication.

- Which points can Eve attack?



Attacking Passwords

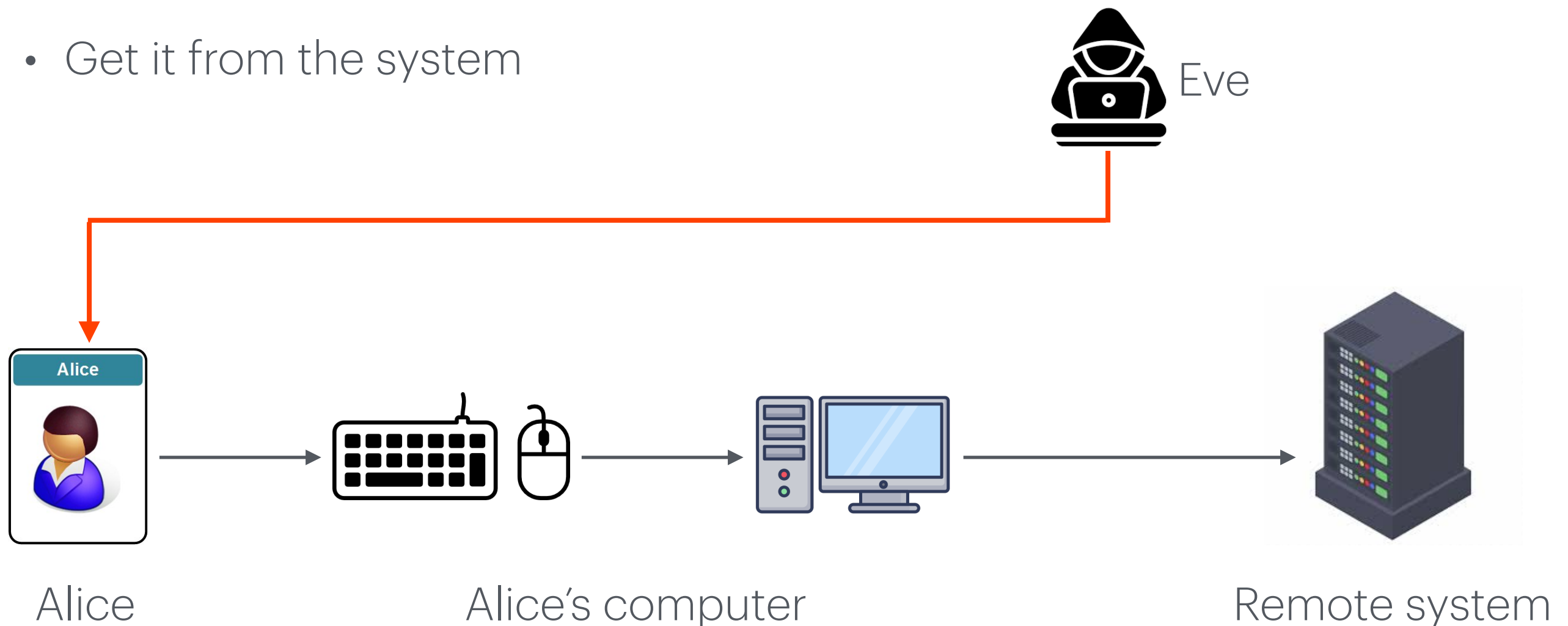
- Get it from Alice
- Intercept it
- Get it from the system



Attacking Passwords

- **Get it from Alice**

- Intercept it
- Get it from the system



Attacking Passwords

- Is Alice invested in keeping it a secret?
 - Debit card PIN number?
 - Personal email password?
 - Netflix password?
 - Corporate network password?
- Is it written down somewhere?
 - Good against remote attackers
 - Not good against targeted local attacks (co-workers, family, abusers)
 - Know your threat model!
- Can it be guessed based on available knowledge about Alice?

Attacking Passwords

A CRYPTO NERD'S IMAGINATION:

HIS LAPTOP'S ENCRYPTED.
LET'S BUILD A MILLION-DOLLAR
CLUSTER TO CRACK IT.

BLAST! OUR
EVIL PLAN
IS FOILED!

NO GOOD! IT'S
4096-BIT RSA!



WHAT WOULD ACTUALLY HAPPEN:

HIS LAPTOP'S ENCRYPTED.
DRUG HIM AND HIT HIM WITH
THIS \$5 WRENCH UNTIL
HE TELLS US THE PASSWORD.

GOT IT.



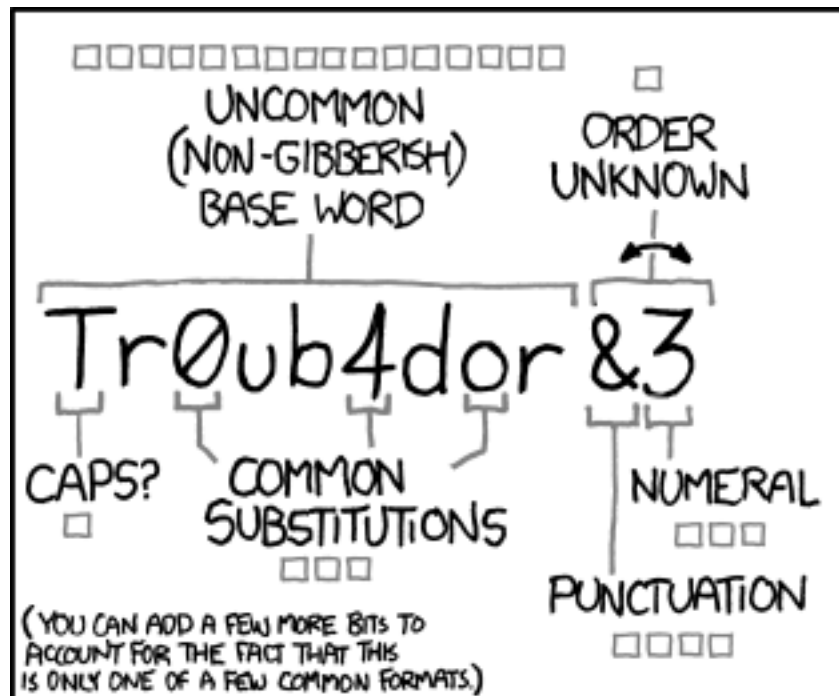
Strong passwords

- Challenge: come up with passwords that are hard to guess, but easy to remember.
- Common password rules:
 - Composition: Letters and numbers, mixed case, symbols, banned dictionary
 - Length
 - Lifetime
- Unintended consequences
 - Required letters/symbols \implies common substitution; Predictable pattern.
 - Monthly change requirement \implies Incremental; Recycling; Weaker pwd.

**Top 20 most
common passwords
according to
NordPass^[3]**

Rank	2021
1	123456
2	123456789
3	12345
4	qwerty
5	password
6	12345678
7	111111
8	123123
9	1234567890
10	1234567
11	qwerty123
12	000000
13	1q2w3e
14	aa12345678
15	abc123
16	password1
17	1234
18	qwertyuiop
19	123321
20	password123

Strong passwords



~28 BITS OF ENTROPY

$2^{28} = 3 \text{ DAYS AT } 1000 \text{ GUESSES/SEC}$

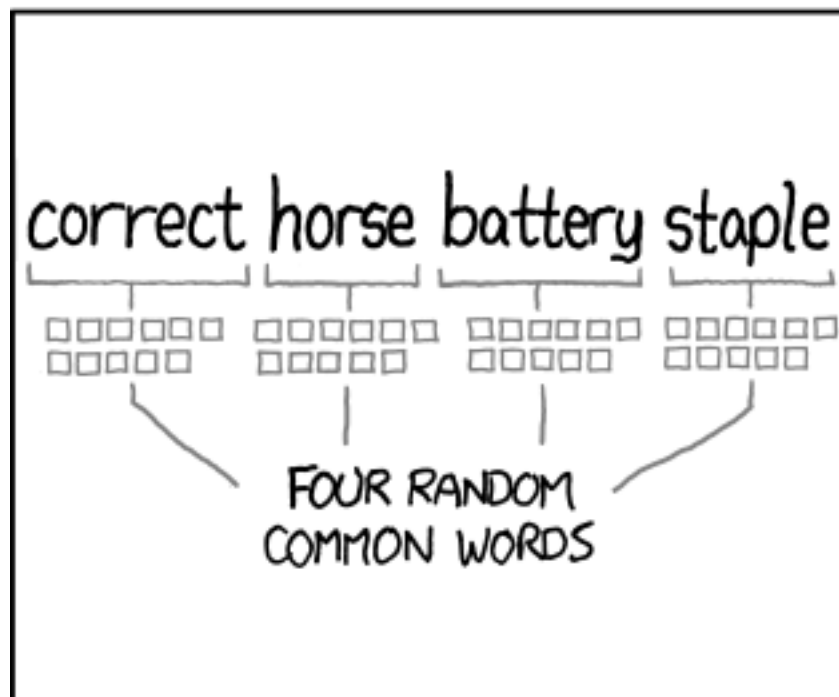
(PLAUSIBLE ATTACK ON A WEAK REMOTE WEB SERVICE. YES, CRACKING A STOLEN HASH IS FASTER, BUT IT'S NOT WHAT THE AVERAGE USER SHOULD WORRY ABOUT.)

DIFFICULTY TO GUESS: **EASY**

WAS IT TROMBONE? NO, TROUBADOR. AND ONE OF THE 0s WAS A ZERO?

AND THERE WAS SOME SYMBOL...

DIFFICULTY TO REMEMBER: **HARD**



~44 BITS OF ENTROPY

$2^{44} = 550 \text{ YEARS AT } 1000 \text{ GUESSES/SEC}$

DIFFICULTY TO GUESS: **HARD**

THAT'S A BATTERY STAPLE.

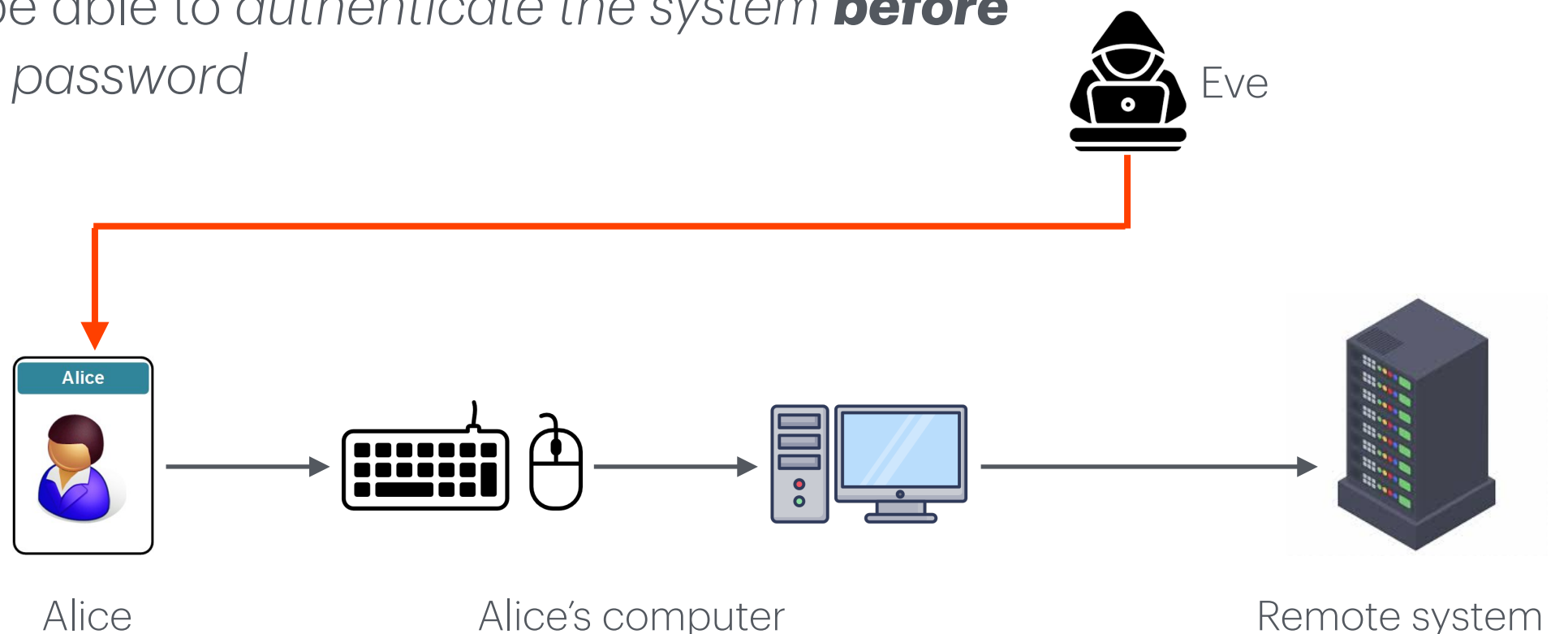
CORRECT!

DIFFICULTY TO REMEMBER: YOU'VE ALREADY MEMORIZED IT

THROUGH 20 YEARS OF EFFORT, WE'VE SUCCESSFULLY TRAINED EVERYONE TO USE PASSWORDS THAT ARE HARD FOR HUMANS TO REMEMBER, BUT EASY FOR COMPUTERS TO GUESS.

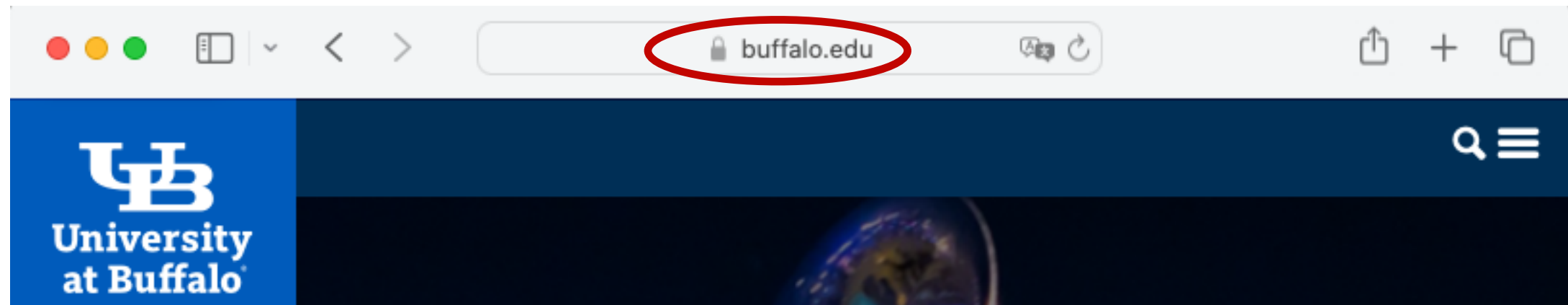
Attacking Passwords

- Can Eve trick Alice into revealing her password?
- How does Alice know she is logging into the real system?
- **Phishing!**
- Tricking Alice into revealing her password by impersonating the system she is trying to access
- Alice has to be able to *authenticate the system **before** providing her password*



Phishing

- How can Alice authenticate the system?
- HTTPS certificates validate the domain name in the URL.



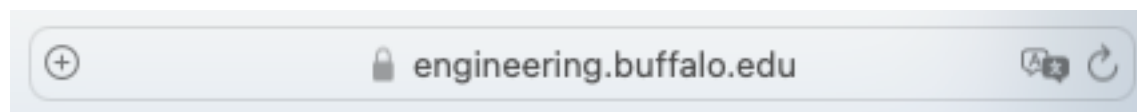
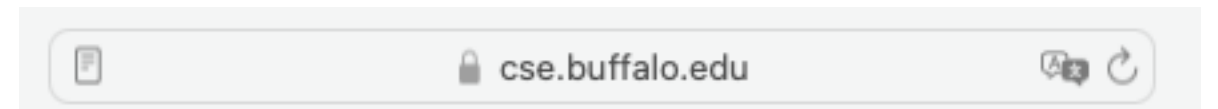
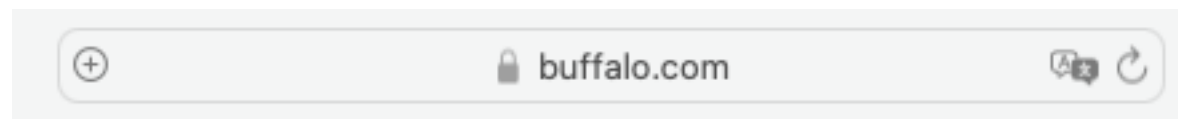
- What does it really tell you?
- That you are communicating with a server owned by UB?
- **No.** Only that you are communicating to www.buffalo.edu and your connection is secure (confidentiality and integrity are protected) against passive and active attackers on the link.

Phishing

- How do you know www.buffalo.edu is a legitimate UB web site?
- What about:
 - www.cse.buffalo.edu
 - www.buffalo.cse.edu
 - www.cse-buffalo.edu
 - www.buffalo.net
 -

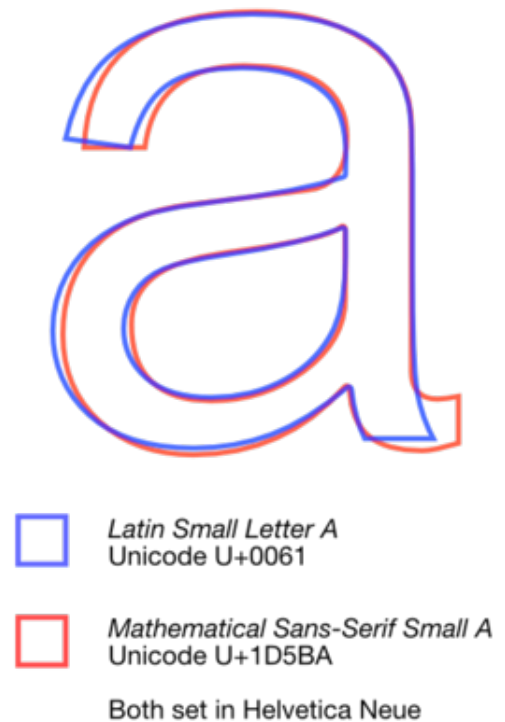
Phishing

- How do you know www.buffalo.edu is a legitimate UB web site?
- A user is *expected* to know which domains are associated with the entity they are trying to interact with.
- And how to properly parse the URL
- Some browsers now highlight the domain portion



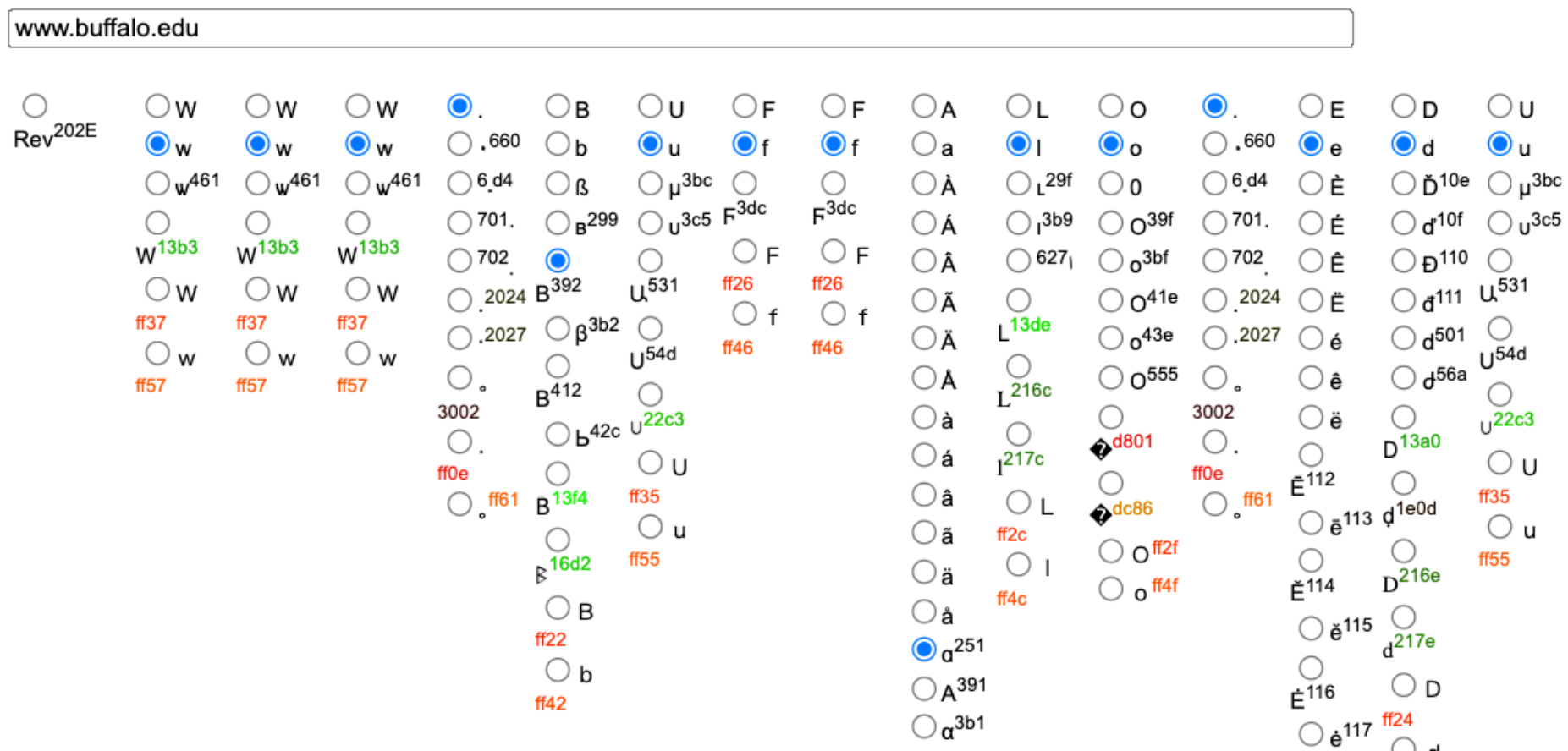
Phishing

- What if the user knows which domain to visit?
- **Homoglyphs**: symbols that appear identical or very similar
- Attacker: Register domain names that look just like the victim domain, but using a different character set.



Phishing

- <https://www.irongeek.com/homoglyph-attack-generator.php>

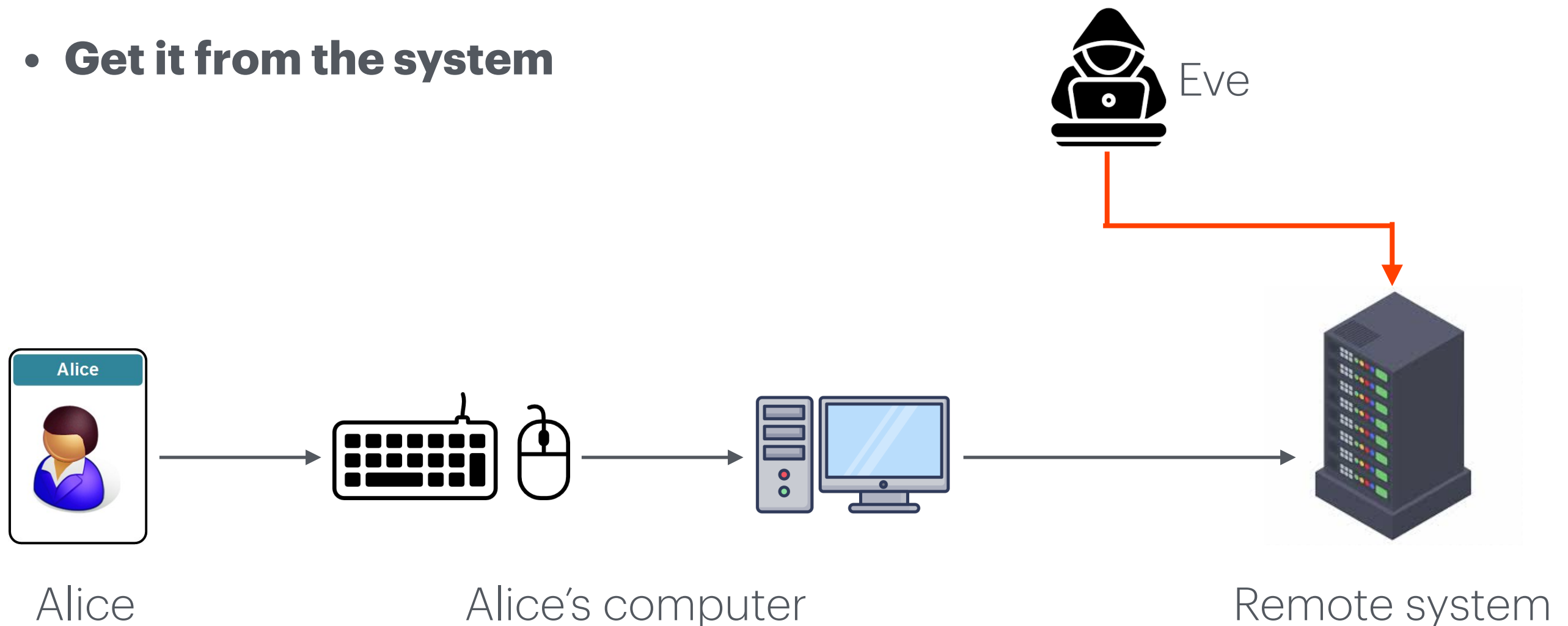


Unicode URL to give out: www.Buffalo.edu
Encoded label to set up in DNS: `www.xn--ufflo-2jc75t.edu`

Password-based
Authentication: Defend
against Direct Attack

Attacking Passwords

- Get it from Alice
- Intercept it
- **Get it from the system**



Attacking Passwords

- **Direct attack:** Use system as an *oracle*; try to log in with different passwords
 - Defense: Minimize error information
 - Defense: Slowing down password verification
 - Defense: Limit number of login attempts per user
- Attacker: Try different users for common passwords
- Compromise password database
 - Huge yield compared to user-side attacks
 - <https://haveibeenpwned.com/>
 - Password reuse issues

Protecting Passwords

- How can the system verify that the password Alice entered is correct?
- Naive solution:
 - Store a copy of the password and compare provided copy to the stored one.
- Problem?
 - If system is compromised, passwords are revealed
 - Same passwords may be used on other systems

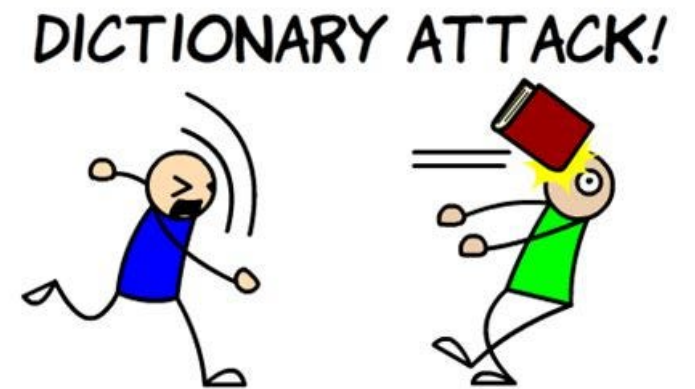
Protecting Passwords

- Other solutions?
 - Hint: System does not need to know the password, only be able to verify it is correct.
- What if the system stores a cryptographic hash of the password?
 - $H(\text{password})$
 - Hash must be *pre-image resistant*
- Better... but still problematic.

Attacking Passwords

- Given a hash of a password, Eve can use it to validate guesses
 - Also, obvious which users have identical passwords
- **Dictionary attacks**
 - Dictionary: collection of possible, or likely, password strings
 - Try every string in the dictionary until the correct entry is found.
- Pre-compute hashes of all strings in the dictionary, then perform reverse look-ups by hash to find corresponding password.

Attacking Passwords



- **Dictionary attacks** cost example
 - Assume passwords are composed of upper or lower case letters or digits
 - $26 + 26 + 10 = 62 \approx 64$ possible values per character
 - $64^n = 2^{6n}$ possible passwords of length n
 - For $n = 6$, 2^{36} possible password strings
 - ≈ 10 TB to store all possible 6-character passwords and respective SHA-1 hashes
- Can be reduced using techniques like *rainbow tables*.

Attacking Passwords

HACKERS RECENTLY LEAKED **153 MILLION** ADOBE USER EMAILS, ENCRYPTED PASSWORDS, AND PASSWORD HINTS. ADOBE ENCRYPTED THE PASSWORDS IMPROPERLY, MISUSING BLOCK-MODE 3DES. THE RESULT IS SOMETHING WONDERFUL:

USER	PASSWORD	HINT	
4e18acc1ab27a2d6		WEATHER VANE SWORD	<input type="text"/>
4e18acc1ab27a2d6			<input type="text"/>
4e18acc1ab27a2d6	a0a2876eb1ea1fca	NAME 1	<input type="text"/>
8babbb6279e06eb6d		DUH	
8babbb6279e06eb6d	a0a2876eb1ea1fca		<input type="text"/>
8babbb6279e06eb6d	85e9da81a8a78adc	57	
4e18acc1ab27a2d6		FAVORITE OF 12 APOSTLES	
1ab29ae86da6e5ca	7a2d6a0a2876eb1e	WITH YOUR OWN HAND YOU HAVE DONE ALL THIS	
a1f9b2b6299e7a2b	e0dec1e6ab797397	SEXY EARLOBES	<input type="text"/>
a1f9b2b6299e7a2b	617ab0277727ad85	BEST TOS EPISODE	<input type="text"/>
39738b7adb0b8af7	617ab0277727ad85	SUGARLAND	
1ab29ae86da6e5ca		NAME + JERSEY #	
877ab7889d3862b1		ALPHA	<input type="text"/>
877ab7889d3862b1			<input type="text"/>
877ab7889d3862b1			<input type="text"/>
877ab7889d3862b1		OBVIOUS	<input type="text"/>
877ab7889d3862b1		MICHAEL JACKSON	
38a7c9279codeb44	9dca1d79d4dec6d5		
38a7c9279codeb44	9dca1d79d4dec6d5	HE DID THE MASH, HE DID THE	<input type="text"/>
38a7c9279codeb44		PURLOINED	<input type="text"/>
a8ae5745a7b7af7a	9dca1d79d4dec6d5	FAV. LATER-3 POKEMON	

THE GREATEST CROSSWORD PUZZLE
IN THE HISTORY OF THE WORLD

Protecting Passwords

Make Dictionary Attack Harder

- **Salting**

- Note, the attacker only had to compute one dictionary of hashes that could then be used for any user's password hash from any system.
- We can parameterize, or "salt", password hashes with unique random numbers
 - *Public* salting: Instead of storing $H(p)$, store $(r, H(r||p))$, where r is random string (salt).
- Pre-computation is no longer possible. Attacker must compute unique hashes for every target.
 - What if the database is hacked & r is leaked?

Protecting Passwords

Make Dictionary Attack Harder

- **Salting**

- Note, the attacker only had to compute one dictionary of hashes that could then be used for any user's password hash from any system.
- We can parameterize, or "salt", password hashes with unique random numbers
 - *Secret* salting: Instead of storing $H(p)$, store $H(r||p)$, where r is a *short* random string (e.g. 12 bits).
 - ▶ Verification needs to enumerate all (4096 for 12 bits) possible r ; Tolerable for typical user login.
- Attacker's work is extended for 4096 times no matter what.

Protecting Passwords

Make Dictionary Attack Harder

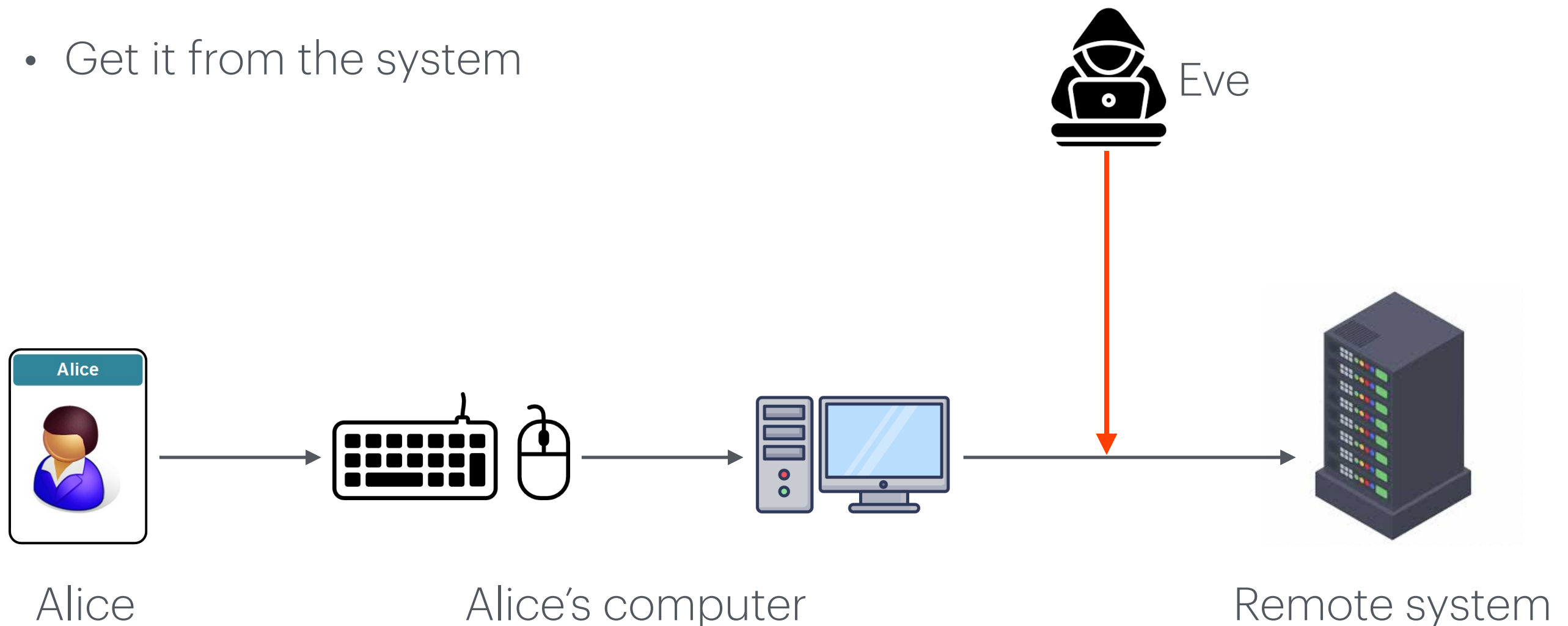
- **Slow hash function**

- The computation to verify a password for a given user on a legitimate system happens relatively infrequently, but an attacker attempting to crack a password hash must perform many, many attempts
- Conclusion: Use a deliberately slow and resource-consuming hashing function
 - PBKDF2: Time consuming but can be accelerated using parallel comp.
 - Scrypt: Time & space consuming. Provably no good time/space tradeoff exists.

Password-based
Authentication: Defend
against Evasdropper

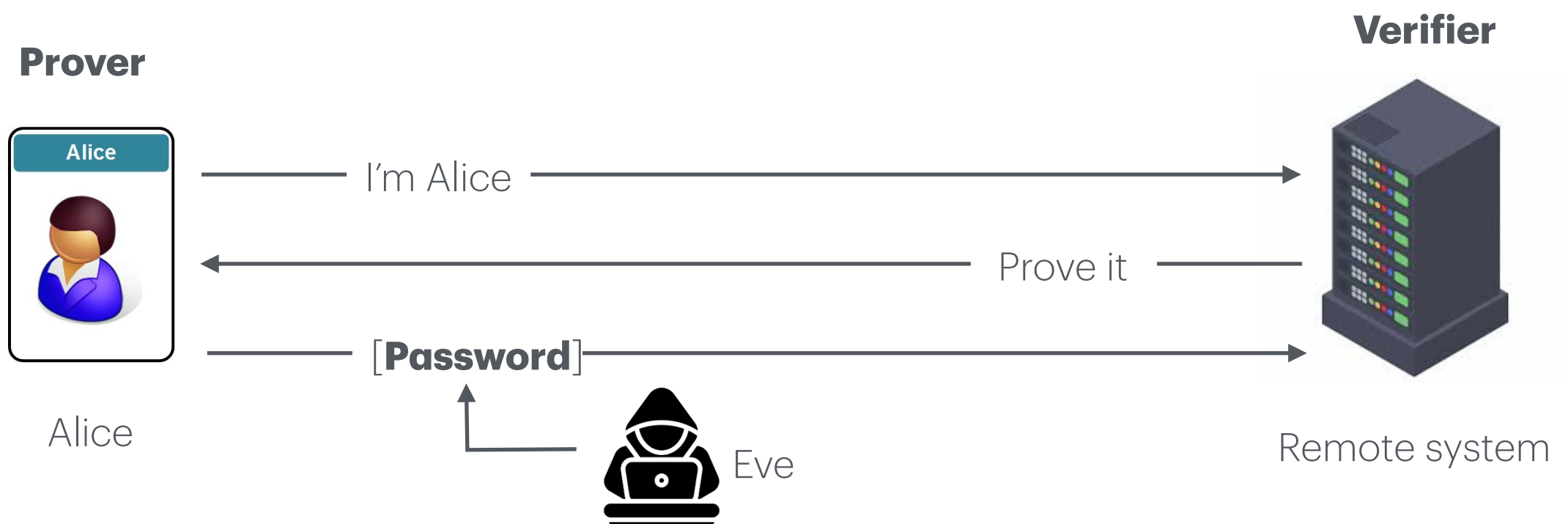
Attacking Passwords

- Get it from Alice
- **Intercept it**
- Get it from the system



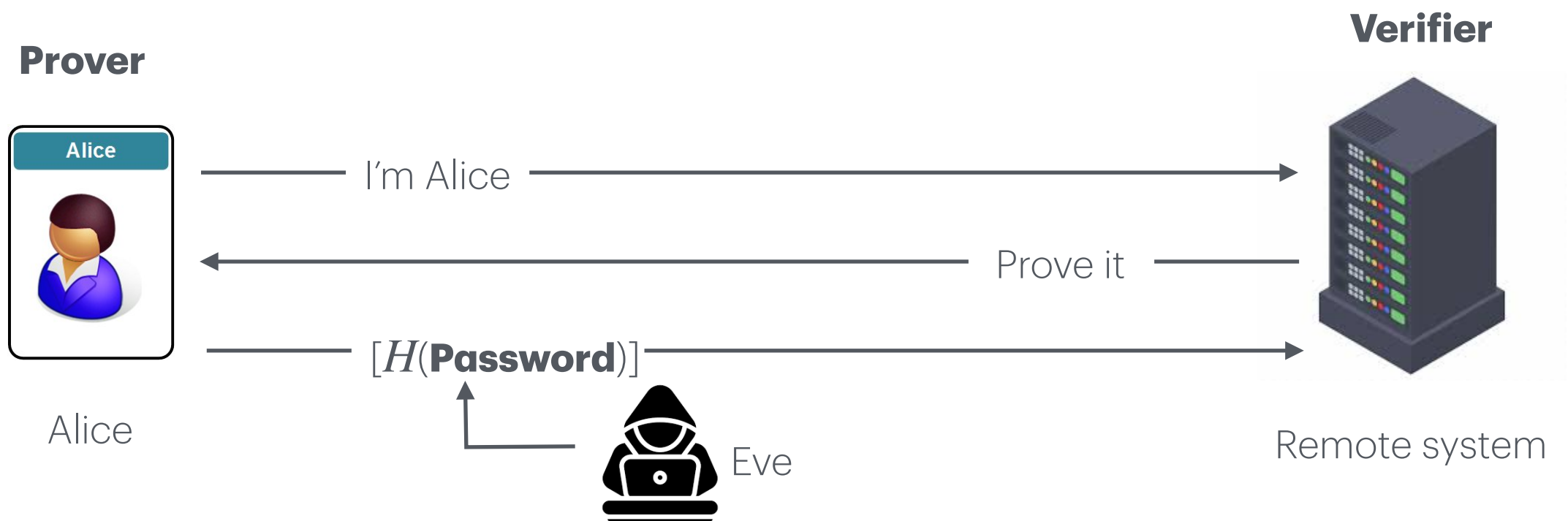
Attacking Passwords

- Alice sends password directly to remote system for authentication.
 - Sounds OK?: anyway the connection will be encrypted.
 - Be conservative:
 - We have seen so many attacks on encrypted comm. protocols (e.g. SSL)
 - Once leaked, the attacker can freely use it & deduce more information



Attacking Passwords

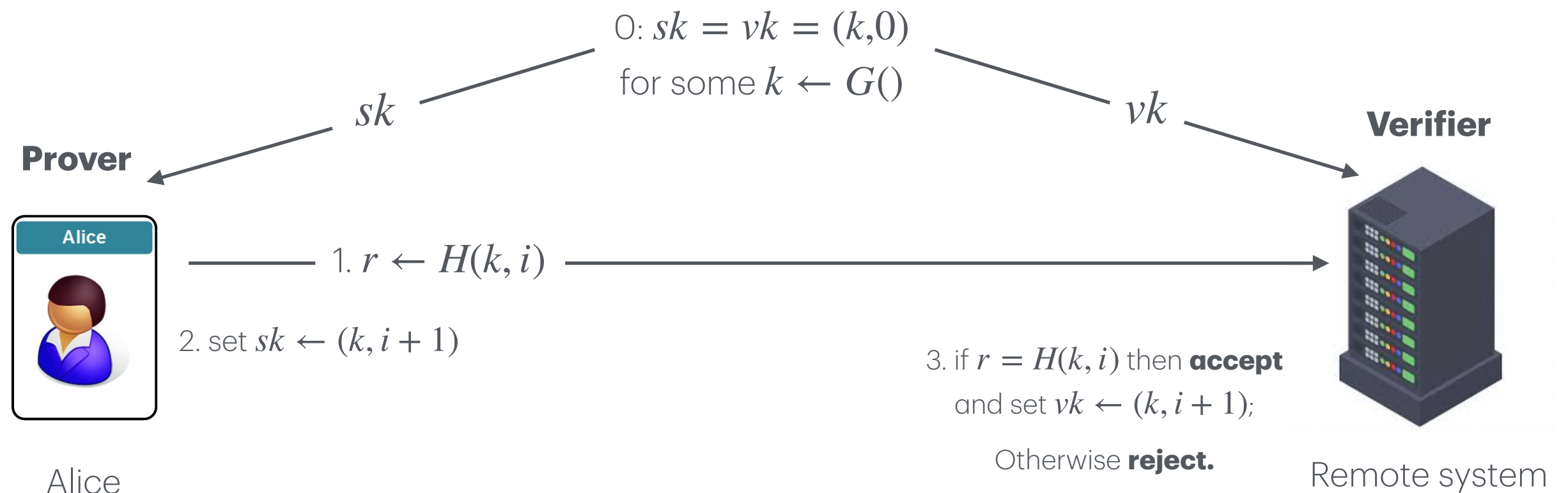
- Can Alice send password **hash** to remote system for authentication?
- Not feasible for most salting method (added at the server side)
- Still susceptible to *replay attack*.



Solution: One-time password

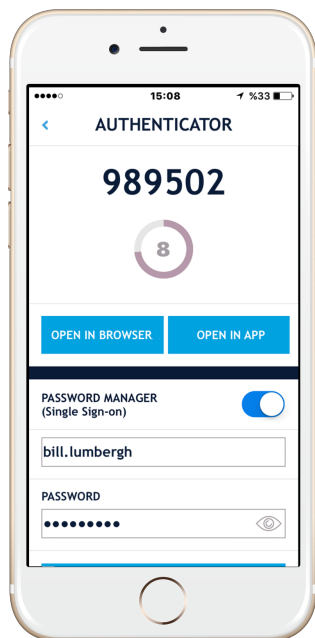
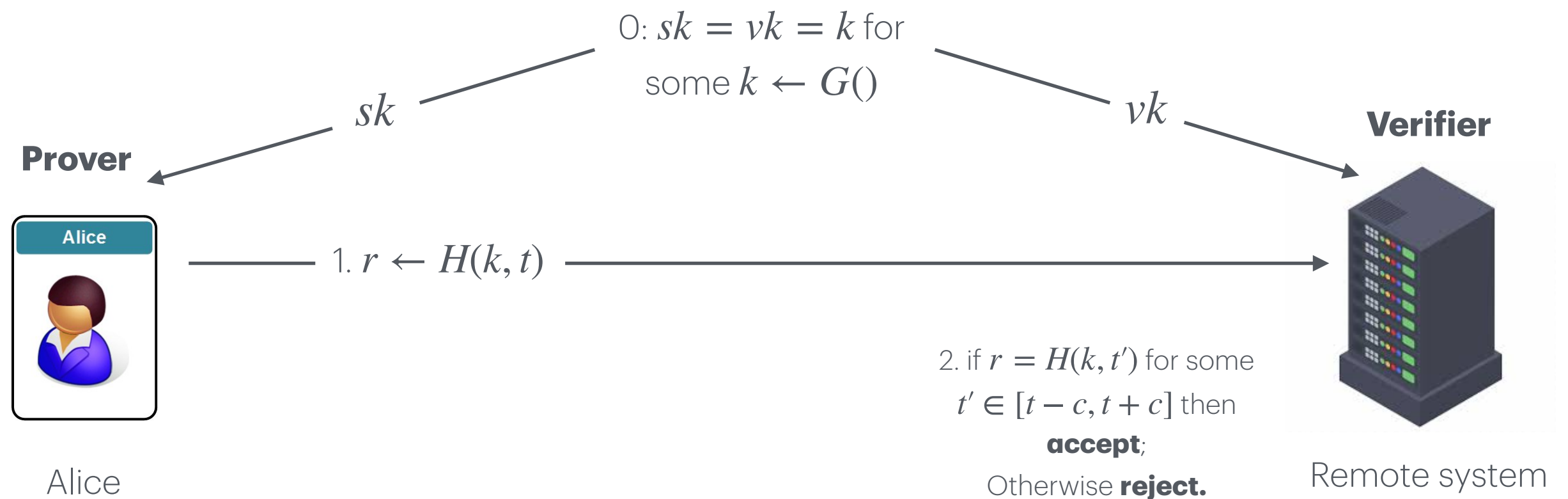
- Each password is used **only once**.
- Such authentication can be realized in the following ways:
 - The user and the system initially agree on a sequence of passwords
 - ▶ Simple solution but requires maintenance of the shared list
 - The user updates her password with each instance of the authentication protocol
 - ▶ E.g., the user sends the new password encrypted under a key derived from the current password
- Crucially relies on the correct communication of the new passwords to the system

Hash-based One-Time Password (HOTP)



- What is this (sk, vk) ?
 - Created at registration time. E.g. Hardware-embedded. Combined with token-based authentication.
- Problem: Needs explicit counter i . Can still be “almost static” if the protocol not executed frequently.

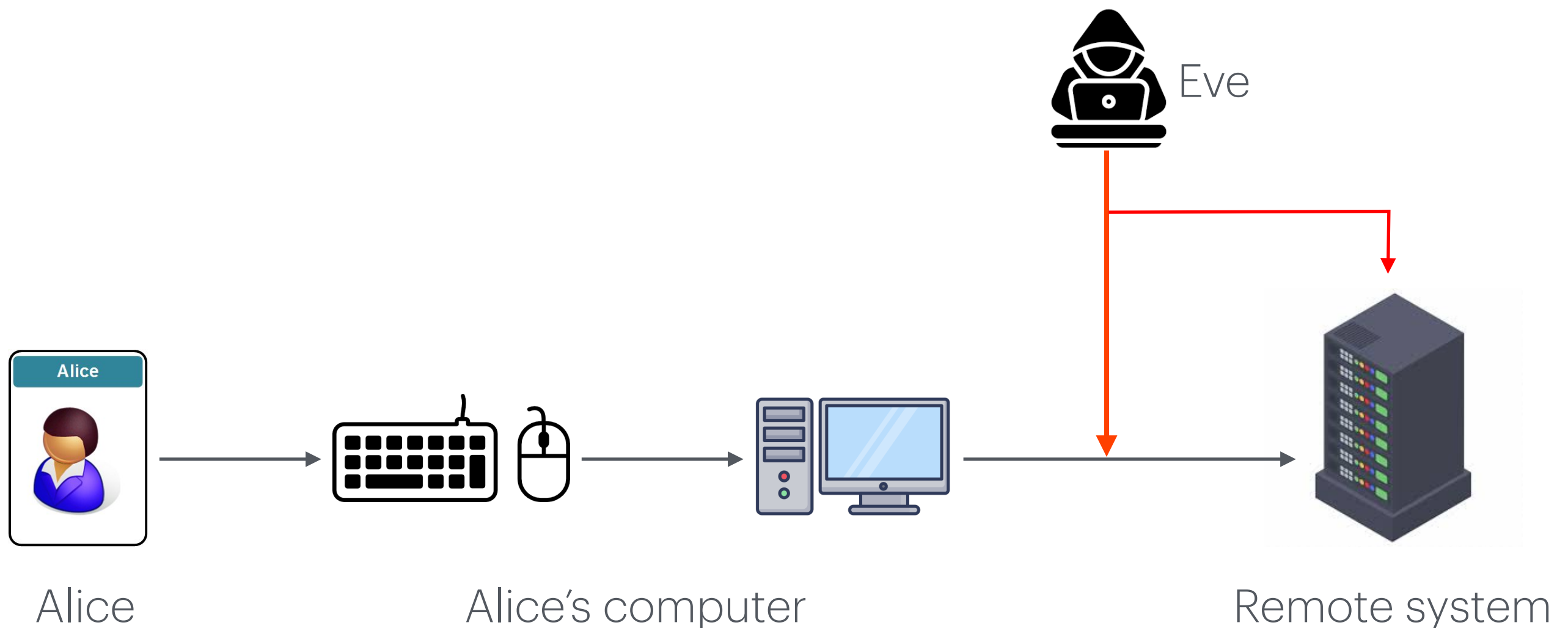
Time-based One-Time Password (TOTP)



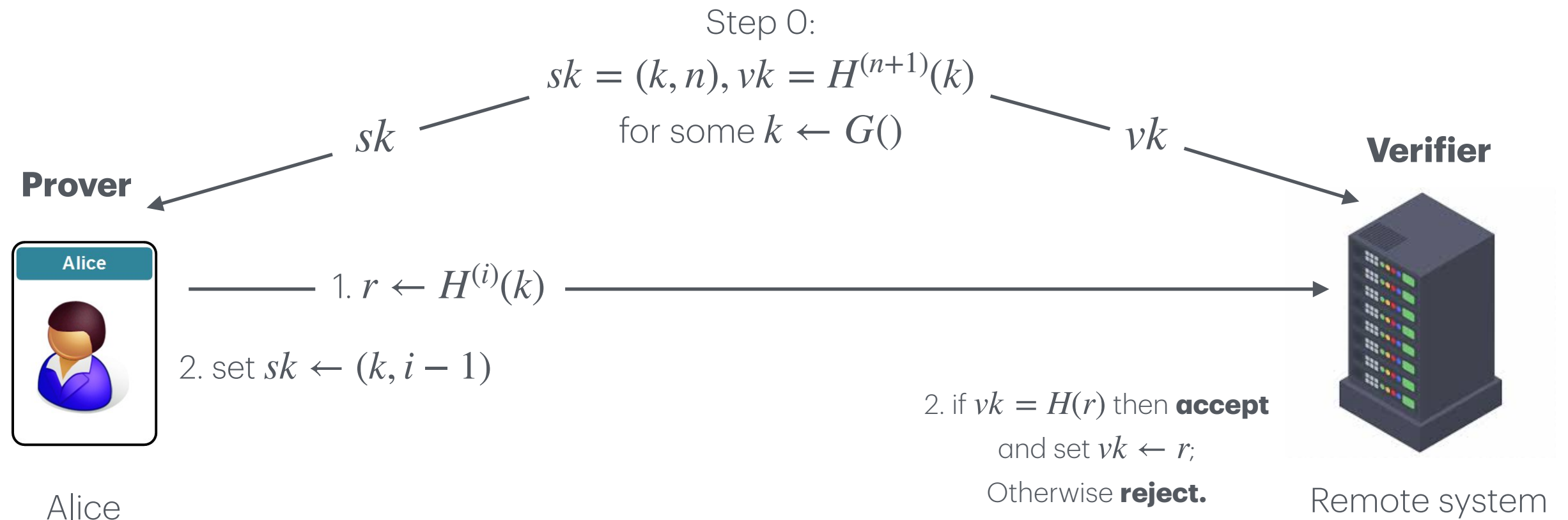
- t : a counter derived from current time.
- Server allow certain clock skew from the Prover: check all $t' \in [t \pm c]$
- The password is usually short (eg. 6 digits)

S/Key System

- A SPoF in HOTP & TOTP: Server stores vk for all clients.
 - If leaked, needs re-registration of all users.



S/Key System



- Secure even if vk is leaked (assuming $H^{(n)}$ is one-way, e.g. SHA256).

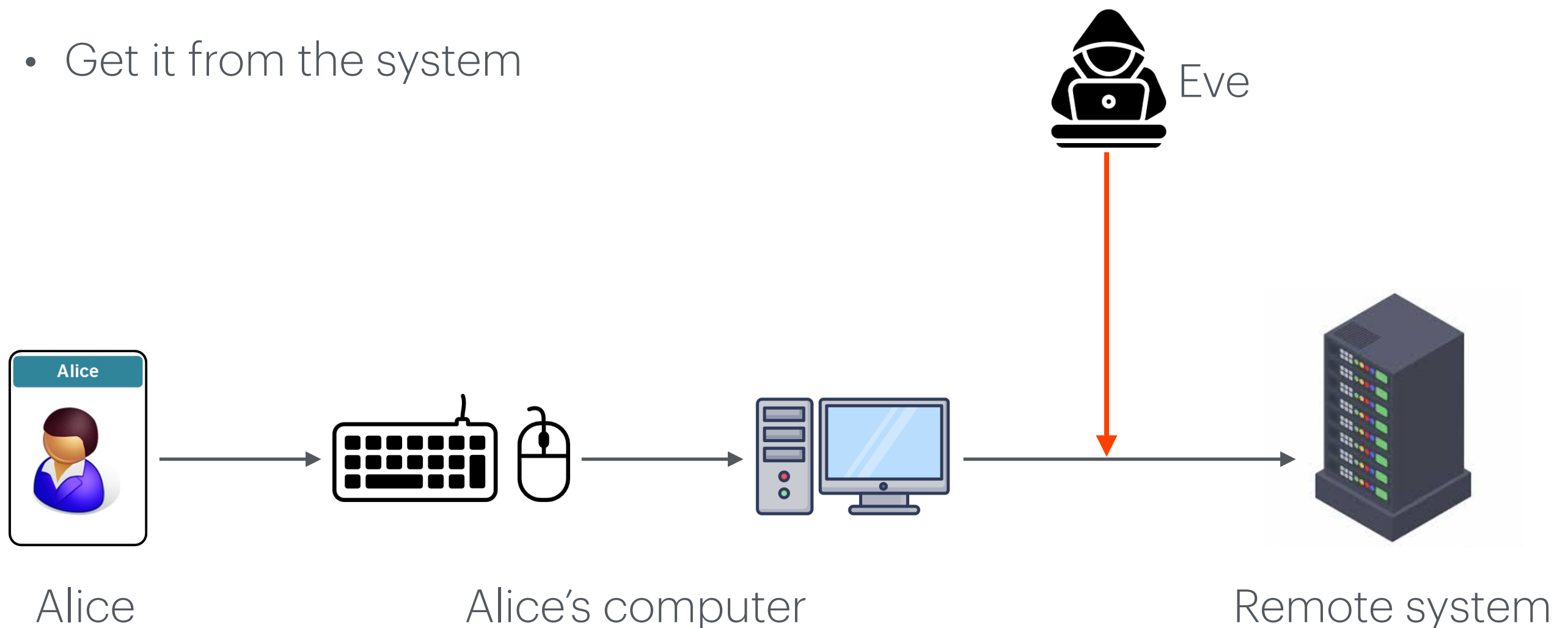
- Problem:

- Need re-registration after n authentications: usually $n < 10^6$
- r needs to be long enough (>128 bits) to ensure one-wayness

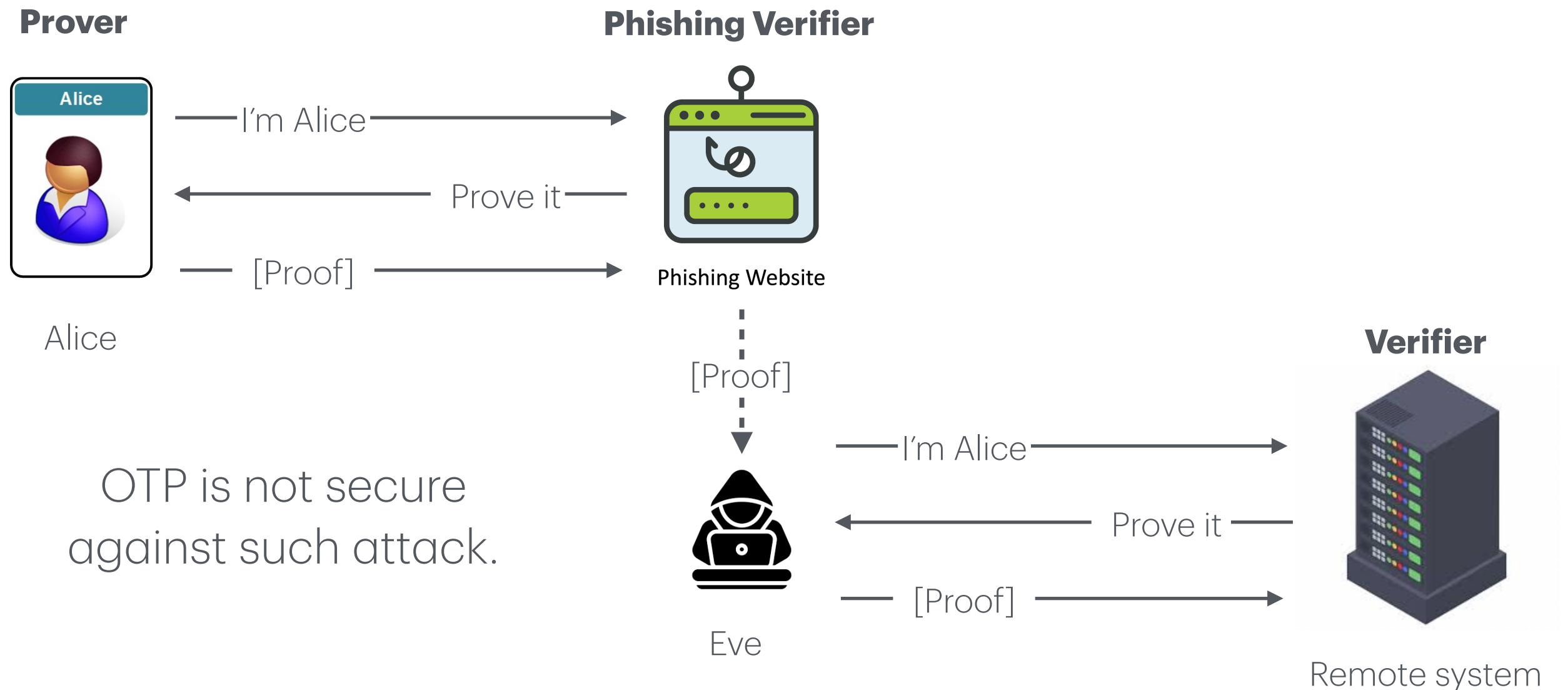
Password-based
Authentication: Defend
against Active Attack

Attacking Passwords

- Get it from Alice
- **Intercept it** Phishing-then-impersonate
- Get it from the system

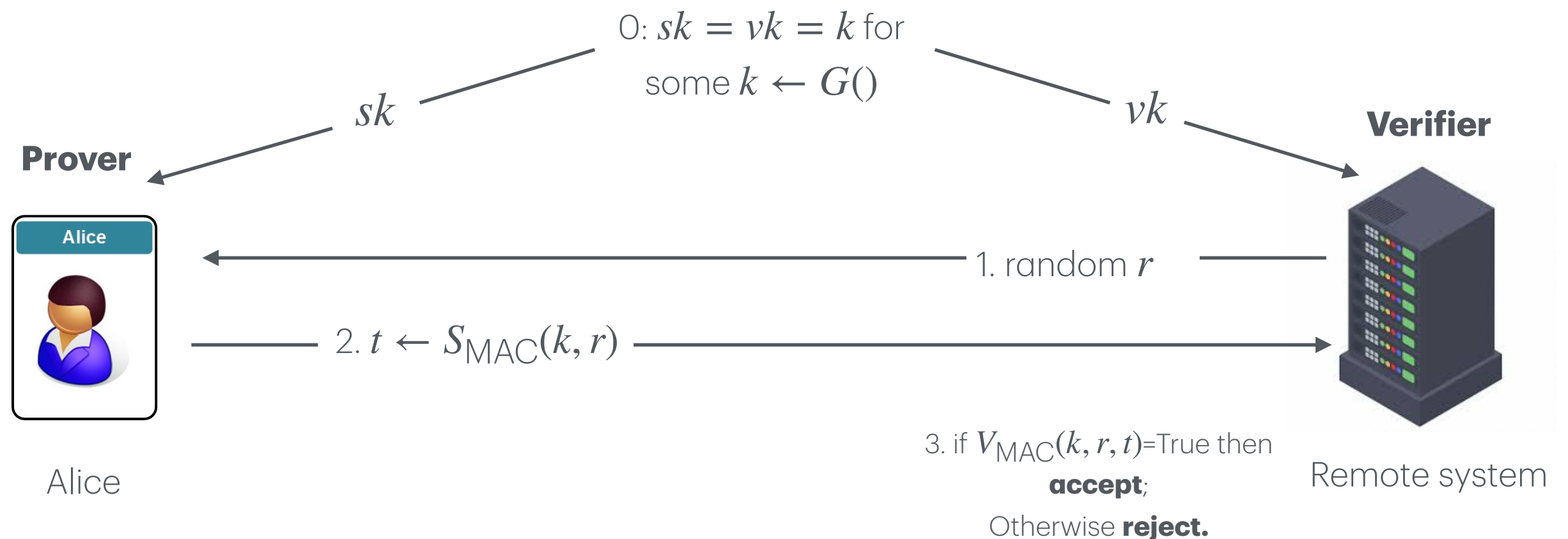


Active (Offline) Attack



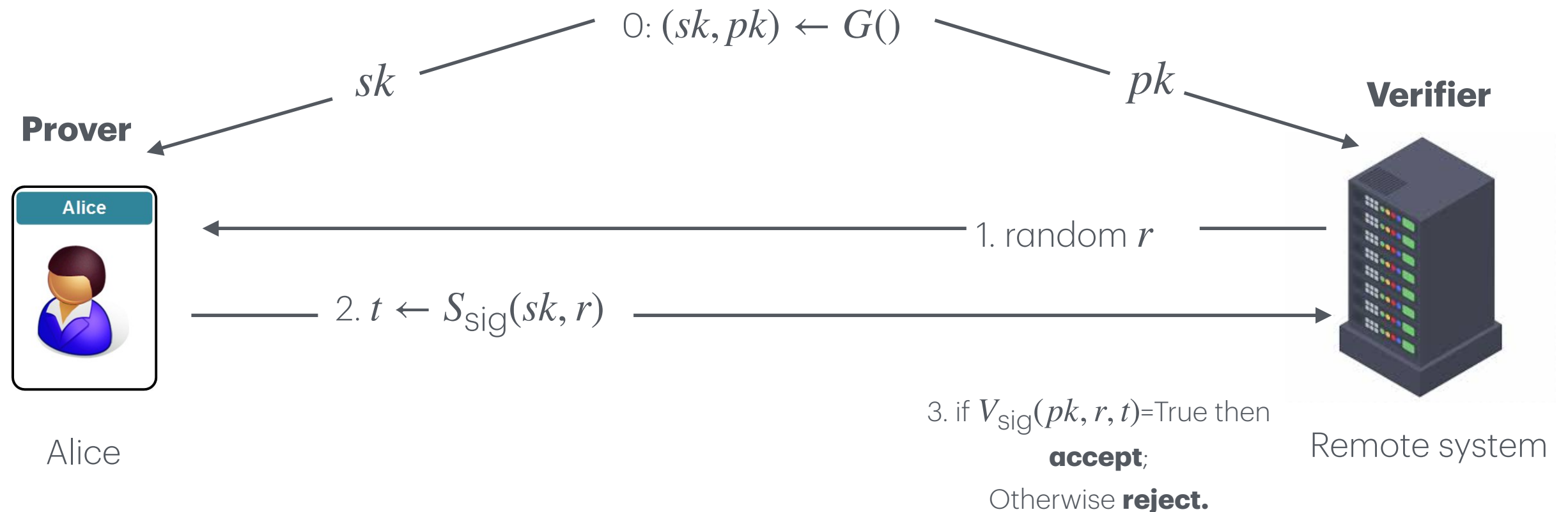
- **Offline** fake ATM:
 - interacts with user; *later* tries to impersonate user to real ATM
- **Offline** phishing:
 - phishing site interacts with user; *later* authenticates to real site

Challenge-response protocols



- (S_{MAC}, V_{MAC}) : a secure MAC
- Why is it secure against an active attacker?
 - What if vk is leaked?

Challenge-response protocols



- $(S_{\text{sig}}, V_{\text{sig}})$: a secure digital signature
- Secure against an active attacker even if pk is leaked.
 - What about Man-in-The-Middle?

Summary for Password-Based Authentication

- “Secret only you know”
- Issue: Can be leaked at various points
- Mitigations
 - Salting: defend against dictionary attack, leakage at remote system/database
 - One-time password: defend against communication intercepting
 - Idea: convince the verifier without leaking the secret
 - *Replace* static password: often used in token-based authentication. Generator (and secret key) embedded in hardware, e.g. car keyfob
 - *Combined* with static password: e.g. the Microsoft Authenticator, DUO

Acknowledgement

- The slides of this lecture is developed heavily based on
 - Slides from Prof Nadia Heninger's lecture on Computer Security (<https://cseweb.ucsd.edu/classes/wi23/cse127-a/slides/16-authentication.pdf>)
 - Slides from Prof Ziming Zhao's past offering of CSE565 (<https://zzm7000.github.io/teaching/2023springcse410565/index.html>)
 - Slides from Prof Marina Blanton's past offering of CSE565 (<https://www.acsu.buffalo.edu/~mblanton/cse565/>)
 - Slides from Prof Hongxin Hu's past offering of CSE565

Questions?