CSE 431/531: Algorithm Analysis and Design (Fall 2024) Dynamic Programming

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Outline

Optimum Binary Search Tree

2 Summary

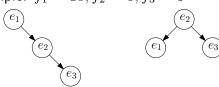
3 Summary of Studies Until Oct 30

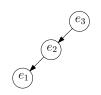
Def. Binary search tree (BST), also called an ordered or sorted binary tree, is a rooted binary tree data structure with the key of each internal node being greater than all the keys in the respective node's left subtree and less than the ones in its right subtree.

- n elements $e_1 < e_2 < e_3 < \cdots < e_n$
- e_i has frequency f_i
- goal: build a binary search tree for $\{e_1, e_2, \cdots, e_n\}$ with the minimum accessing cost:

$$\sum_{i=1}^{n} f_i \times (\text{depth of } e_i \text{ in the tree})$$

• Example: $f_1 = 10, f_2 = 5, f_3 = 3$





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- $10 \times 1 + 5 \times 2 + 3 \times 3 = 29$
- $10 \times 2 + 5 \times 1 + 3 \times 2 = 31$
- $10 \times 3 + 5 \times 2 + 3 \times 1 = 43$

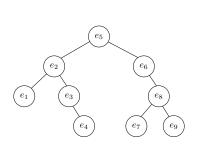
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- ullet suppose we decided to let e_k be the root
- $e_1, e_2, \cdots, e_{k-1}$ are on left sub-tree
- $e_{k+1}, e_{k+2}, \cdots, e_n$ are on right sub-tree
- d_i : depth of e_i in our tree
- C, C_L, C_R : cost of tree, left sub-tree and right sub-tree



•
$$d_1 = 3, d_2 = 2, d_3 = 3, d_4 = 4, d_5 = 1,$$

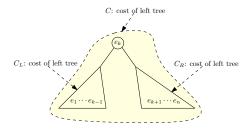
$$\bullet \ d_6 = 2, d_7 = 4, d_8 = 3, d_9 = 4,$$

•
$$C = 3f_1 + 2f_2 + 3f_3 + 4f_4 + f_5 + 2f_6 + 4f_7 + 3f_8 + 4f_9$$

$$C_L = 2f_1 + f_2 + 2f_3 + 3f_4$$

$$C_R = f_6 + 3f_7 + 2f_8 + 3f_9$$

•
$$C = C_L + C_R + \sum_{j=1}^{9} f_j$$



$$C = \sum_{\ell=1}^{n} f_{\ell} d_{\ell} = \sum_{\ell=1}^{n} f_{\ell} (d_{\ell} - 1) + \sum_{\ell=1}^{n} f_{\ell}$$

$$= \sum_{\ell=1}^{k-1} f_{\ell} (d_{\ell} - 1) + \sum_{\ell=k+1}^{n} f_{\ell} (d_{\ell} - 1) + \sum_{\ell=1}^{n} f_{\ell}$$

$$= C_{L} + C_{R} + \sum_{\ell=1}^{n} f_{\ell}$$

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$$opt[1, n] = (opt[1, k-1] + opt[k+1, n]) + \sum_{\ell=1}^{n} f_{\ell}$$

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• In general, opt[i, j] =

$$\begin{cases} 0 & \text{if } i = j + 1 \\ \min_{k:i \le k \le j} \left(opt[i, k - 1] + opt[k + 1, j] \right) + \sum_{\ell=i}^{j} f_{\ell} & \text{if } i \le j \end{cases}$$

```
1: fsum[0] \leftarrow 0
 2: for i \leftarrow 1 to n do fsum[i] \leftarrow fsum[i-1] + f_i
                                                                 \triangleright fsum[i] = \sum_{i=1}^{i} f_i
 3: for i \leftarrow 0 to n do opt[i+1,i] \leftarrow 0
 4: for \ell \leftarrow 1 to n do
         for i \leftarrow 1 to n - \ell + 1 do
 5:
              i \leftarrow i + \ell - 1, opt[i, j] \leftarrow \infty
 6:
               for k \leftarrow i to i do
 7:
                   if opt[i, k-1] + opt[k+1, j] < opt[i, j] then
 8:
                         opt[i, i] \leftarrow opt[i, k-1] + opt[k+1, i]
 9:
                         \pi[i, j] \leftarrow k
10:
               opt[i, j] \leftarrow opt[i, j] + fsum[j] - fsum[i-1]
11:
```

Printing the Tree

```
Print-Tree(i, j)

1: if i > j then

2: return

3: else

4: print('(')

5: Print-Tree(i, \pi[i, j] - 1)

6: print(\pi[i, j])

7: Print-Tree(\pi[i, j] + 1, j)

8: print(')')
```

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Dynamic Programming

- Break up a problem into many overlapping sub-problems
- Build solutions for larger and larger sub-problems
- Use a table to store solutions for sub-problems for reuse

Comparison with greedy algorithms

- Greedy algorithm: each step is making a small progress towards constructing the solution
- Dynamic programming: the whole solution is constructed in the last step

Comparison with divide and conquer

- Divide and conquer: an instance is broken into many independent sub-instances, which are solved separately.
- Dynamic programming: the sub-instances we constructed are overlapping.

Definition of Cells for Problems We Learnt

- Weighted interval scheduling: opt[i] = value of instance defined by jobs $\{1,2,\cdots,i\}$
- \bullet Subset sum, knapsack: opt[i,W']= value of instance with items $\{1,2,\cdots,i\}$ and budget W'
- Longest common subsequence: opt[i,j] =value of instance defined by A[1..i] and B[1..j]
- \bullet Shortest paths in DAG: $f[v] = \mbox{length}$ of shortest path from s to v
- Matrix chain multiplication, optimum binary search tree: opt[i,j] = value of instances defined by matrices i to j

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 - Quiz 7-8

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 - Quiz 9