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CMPT 225 Fall 2024 Assignment 3
– 301 599 274 –

Question 1 : Cartesian Product

Cost Function:

Outside loop - 1 operation

Outer loop comparison - $n + 1$ operations

Outer loop - 3 operations (times n)

Inner loop comparison - $n + 1$ operations (times n)

Inner loop - 3 operations (times n^2)

$$= 3n^2 + n(n+1) + 3n + (n+1) + 1$$

$$= 3n^2 + n^2 + n + 3n + n + 1 + 1$$

$$= 4n^2 + 5n + 2$$

Barometer operations:

Comparisons: while ($i < n$) {
 while ($j < n$) {

Increments: $i++$;
 $j++$;

Time Complexity : $O(n^2)$

Question 2 : Triangle

Cost Function:

Outside loop - 1 operation

First outer loop - $5n + 1$

First inner loop - $n(3i)$.

Average- $n/2$

$$3n^2 / 2$$

Second outer loop - $5n + 1$

Second inner loop - $n(3i)$

Average - $n/2$

$$3n^2 / 2$$

$$= (3n^2 / 2) + (3n^2 / 2) + 5n + 1 + 5n + 1 + 1$$

$$= 3n^2 + 10n + 3$$

Barometer operations:

Comparisons: while ($i < x$) {

while ($j \leq i$) {

while ($i > 0$) {

while ($j \leq i$) {

Increments: $i++$;

$j++$;

$i--$;

$j++$;

Time Complexity : $O(n^2)$

Question 3 : Matrix Self Multiply

Cost Function:

Outside loop - 3 operations

Outer loop comparison - rows + 1

Outer Loop - 2 operations * rows

Middle loop comparison - (columns + 1) * rows

Middle loop - 5 columns * rows

Inner loop comparison - (rows + 1) * columns * rows

Inner loop - 4 * rows * columns * rows

Rows (n) = columns, throughout the code

$$= 4n^3 + n^3 + n^2 + 5n^2 + n^2 + n + 2n + n + 1 + 3$$

$$= 5n^3 + 7n^2 + 4n + 4$$

Barometer operations:

```
Comparisons: while ( r < rows ){  
    while ( c < columns ) {  
        while ( iNext = 0 ) {
```

```
Increments : iNext++;  
             c++;  
             r++;
```

Operation Executed the most:

```
next += m[rcIndex(r, iNext, columns)] * m[rcIndex(iNext, c, columns)];
```

Time Complexity : $O(n^3)$

Question 4 : selection sort

Cost Function:

Function is called n-1

First comparison 1 operation

Outside loop - 5 operations

While loop comparison - (n - i)

While loop - 3 operations * (n - i - 1)

Average n/2

Base case comparison - 1

Because of how the function is called at the end. It can be treated like a loop to count.

$$\begin{aligned} &= 6(n-1) + (n-1)(n/2) + 3(n-1)(n/2) + 1 \\ &= 6n - 6 + n^2/2 - n/2 + 3n^2/2 - 3n/2 + 1 \\ &= 2n^2 + 4n - 5 \end{aligned}$$

Barometer Operations:

Comparison: if (i < n - 1) {
 while (next < n) {

Increment: next++;

Function calls: ssort(arr, n, i + 1);

Time Complexity : $O(n^2)$

Question 5 : Pattern

Cost Function:

Function has 2 calls. Each one being called $\log_2(n)$ times.

Since each call has 2 other calls. This can be looked at as a tree.

$$2^{(\log_2(n))} \text{ calls: } n \text{ calls} \quad \text{or} \quad \log_2(n) + \log_2(n) = n$$

Each call:

Outside loop: 6 operations, including comparisons

Loop : $3 \cdot n$ operations

$$= 3n + 6 \text{ operations}$$

Total operations:

$$= n \cdot (3n + 6)$$

$$= 3n^2 + 6n$$

Barometer Operations:

Comparisons: while ($ast < n$) {

Function call: pattern($n/2$, i);

pattern($n/2$, i);

Time Complexity : $O(n^2)$

Question 6 : linear search

Cost Function:

If target not found, then n operations, however if the target is found (worst case to be at the last index), the return value of that recursive call being 0 will invoke another call (2 calls in total); causing the function to act as a binary recursive tree. In the worst case, the functions are called 2^n times.

Each call:

Returning Comparisons - 2 operations

If target not found - 1 operation

If target found - 1 operation

Operations other than recursive calls - 4 which can be written as 2^2

Each call is making two more calls

Cost function:

$$\begin{aligned}C(n) &= 2^2 + 2(C(n-1)) \\&= 2^2 + 2(2^2 + 2C(n-2)) \\&= 2^2 + 2^3 + 2^2 C(n-2) \\&= 2^2 + 2^3 + 2^2 (2^2 + 2C(n-3)) \\&= 2^2 + 2^3 + 2^4 + 2^3 C(n-3) \\&\dots \\&= 2^2 + 2^3 + \dots + 2^{n+1} + 2^n C(0)\end{aligned}$$

If len is 0, one operation. $C(0)$ is 1

$$= 2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^{n+1} + 2^n - 1 - 2$$

Geometric series + 2^n

$$\begin{aligned}&= ((1 - 2^{n+1}) / (1 - 2)) + 2^n - 1 - 2 \\&= 2^{n+1} - 1 + 2^n - 1 - 2 \\&= 2^{n+2} - 4\end{aligned}$$

Barometer Operations:

Comparisons: if (len == 0) return -1;

if (arr[0] == target) return 0;

Function calls: lsearch(arr+1, len-1, target)

Time Complexity : $O(2^n)$

Question 7 : pow

Cost Function:

Loop runs $\log_2(n)$ times

Outside loop - 1 operation

Loop comparison - $\log_2(n) + 1$

Inside loop - $4 * \log_2(n)$ operations

$$= 5 \log_2(n) + 2$$

Barometer Operations:

Comparison: while (exp > 0) {

Decrement: exp >>= 1;

Time Complexity : $O(\log_2(n))$