

Midterm 1



Hello everyone,

Midterm 1 grades are now being released from Crowdmark and the solution has been posted.

Check the solution. If you have any questions related to the marking of your Midterm 1, you can see me during my office hours.

Please, no email, no posting on Discussion forum.

Thank you! :)

Class scores distribution [Show](#)

My score

72.2% (32.5/45)

Q1

4



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1. [6 marks] Consider the following syntactically correct C main function:

```
int main(int argc, char *argv[]) {
```

```
    int x = 0xDECAF000;
```

```
    short y = 0xCAFE;
```

```
    char z = y;
```

```
    // z = 0xFE
```

```
    if ( x > y ) printf("Cafe\n");
```

```
    // Your modified condition is:
```

```
    // if ( (unsigned short) x > y )
```

2. uncorrect second cast -2

```
    if ( x < y ) printf("Decaf\n");
```

```
    // Your modified condition is:
```

```
    // if ( (unsigned int) x < y )
```

```
    return 0;
```

```
}
```

Using only two casts (data type casts), and these two casts must be different, modify the conditions of the two conditional statements in the main function above such that only "Cafe" is printed on the computer monitor screen when the program executes on our target machine.

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Q2

3

A=10
B=11
C=12
D=13
E=14

2. [3 marks] Consider the following syntactically correct C code fragment:

```
float aFloat = 13.75;
int sum = (int) aFloat + 0xFFFFFEE;
```

13 + FFFF

Which value does the variable sum contain when the above C code fragment has been executed on our target machine? Represent your answer as a hexadecimal number.

0x 0000 000B

0001101
1111110
1000001011

Correct 3

3. [3 marks] Consider the following C code fragment:

```
char char1 = 101;
char char2 = 27;
char sumOfChar = char1 + char2;
```

128 = 101 + X
X = 27

Q3

3

Which value must be assigned to char2 in order for the sum of char1 and char2 to create a positive overflow? 275 because 27+101=128

= 1111111 = -1

4. [3 marks] Consider the disassembled function below. It was obtained by disassembling an executable containing this function on our target machine.

```
7002 : f3 0f 1e fa +4      endbr64
7006 : 48 83 ec 08 +4      sub    $0x8,%rsp
700A : 48 8b 05 d9 2f 00 00 +7 mov    0x2fd9(%rip),%rax
700F : 48 85 c0 +3          test   %rax,%rax
7014 : 74 02 +2             je     7016 <_init+0x16>
7016 : ff d0 +2             callq  *%rax
7018 : 48 83 c4 08 +4      add    $0x8,%rsp
701C : c3 +1               retq
```

B.
C.
D.
E.
F.
G.
H.
I.
J.

As you can see, six (6) of the memory addresses are missing. Complete the left column of the above disassembled function by adding these six (6) missing memory addresses.

701C

7002 7003 7004 7005
7006 7007 7008 7009
700A 700B 700C 700D 700E 700F 7010
7011 7012 7013
7014 7015
7016 7017
7018 7019 701A 701B



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5. [Total marks: 9] Consider the following function `mystery` written in x86-64 assembly code, where the parameter `x` is in the register `%edi` and the parameter `y` is in the register `%esi`:

<code>.globl mystery</code>	Line 1
<code>mystery:</code>	Line 2
<code>endbr64</code>	Line 3
<code>testl %esi, %esi</code>	Line 4
<code>jle .L4</code>	Line 5
<code>if (3 > 0) movl \$0, %eax</code>	Line 6
<code>movl \$1, %edx</code>	Line 7
<code>.L3:</code>	Line 8
<code>imull %edi, %edx</code>	Line 9
<code>addl \$1, %eax</code>	Line 10
<code>← cmpl %esi, %eax</code>	Line 11
<code>jne .L3</code>	Line 12
<code>.L1:</code>	Line 13
<code>movl %edx, %eax</code>	Line 14
<code>ret</code>	Line 15
<code>.L4:</code>	Line 16
<code>movl \$1, %edx</code>	Line 17
<code>jmp .L1</code>	Line 18

esi != eax
3 != 1
3 != 2
3 != 3

eax = 0
edx = 1

edx = 4, 16, 64
eax = 1, 2, 3

eax = 64

2
16.
x4
64

- a. [3 marks] If we call this function as follows: `mystery(x, y)` where `x = 4` and `y = 3`, which value will it return?

it will return 64

Q5a. 3

Correct 3

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Q5b.

2

b. [2 marks] If you were to replace the two occurrences of the label .L3, in the mystery function, listed on the previous page, with a more descriptive label name, which label name would you choose? Write this more descriptive label name below:

I would choose the name ~~do~~while, or while. because the label does something, and the checks for a condition to continue to back to the label, just like a loop.

Correct 2

Q5c.

2

c. [2 marks] If you were to replace Line 6 in the mystery function, listed on the previous page, with another equivalent x86-64 instruction, i.e., an x86-64 instruction that would produce the same result, but is not a mov* instruction, which instruction would you choose? Write this entire x86-64 instruction below:

xorl %eax, %eax

Correct 2

I don't remember exact word but it $\llcorner \wedge \llcorner$ will make it 0.

Q5d.

0

d. [2 marks] If you were to replace Line 10 in the mystery function, listed on the previous page, with another equivalent x86-64 instruction, i.e., an x86-64 instruction that would produce the same result, but is not a add* instruction, which instruction would you choose? Write this entire x86-64 instruction below:

incl \$1, %eax

Incorrect -2



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2168420
43210

Q6a.

2

6. [Total marks: 12] Consider a floating point number encoding scheme based on the IEEE 754 standard for floating point representation and defined as follows:

- It uses 8 bits.
- There is one sign bit s .
- The exp can be any number in the range $[0 \dots 31]$. 5 bits

a. [2 marks] Compute the bias of this IEEE-like encoding scheme described above and show your work:

$$\text{bias} = 2^{k-1} - 1 = 2^{5-1} - 1 = 2^4 - 1 = 16 - 1 = 15$$

b. [1 mark per cell for a total of 7 marks] Complete the table found on the next page with bit patterns representing the real numbers listed in the left column using the IEEE-like encoding scheme described above.

c. [3 marks] Can 1.875×2^{-15} be encoded exactly (not approximated) as a normalized number using the IEEE-like encoding scheme described above? Explain and show why or why not.

Q6c.

0.5

no because -15 would
make $\text{exp} = 11111$ which
is $\frac{31}{2}$, meaning cannot
be represented.

Your conclusion is correct but
your work is faulty.

0.5

Q6b.

1

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Real Number	IEEE-like floating point number represented using its bit pattern	Show your work
2	010000 000	Show your work for 2: $2 \Rightarrow 1.0$ $\rightarrow 1.0 \times 2^1$ $= 000$ $\text{exp} - \text{bias}$ $1 = \text{exp} - 15$ $\text{exp} = 16$ > 0 0 10000 000
zero	0.00000.00	
Largest positive denormalized	0 00000 100	
$\frac{1}{4}$	0 01101 000	Show your work for $\frac{1}{4}$: $\frac{1}{4} = 0.25_{10}$ $= 0.012$ $\times 10^{-2}$ $= 000$ $E = \text{exp} - \text{bias}$ $-2 = \text{exp} - 15$ $\text{exp} = 13$ $S = 0$ $E = 0$ $E = \text{exp} - \text{bias}$ $-2 = \text{exp} - 15$
Smallest positive normalized	0 11110 001	
Largest positive normalized	0 10001 001	

Incorrect cell. -1

Partially correct. -0.5

Incorrect cell. -1

Incorrect cell. -1

Partially correct. -0.5

Incorrect cell. -1

Incorrect cell. -1



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7. [9 marks] Consider the following C function declaration:

```
long average( rdilong w, rsilong x, rdxlong y, rcxlong z);
```

which, as its name implies, calculates and returns the average of its four (4) parameters w, x, y and z. Below, implement this function in x86-64 assembly language. Make sure your x86-64 assembly code is as syntactically correct as possible, i.e., it would execute on our *target machine*. You can assume that your x86-64 assembly code will not create an overflow of any kinds.

Requirement: When answering this question, you cannot use any of the `div*` or `idiv*` x86-64 instructions.

```
# Description: Calculates and returns the average of its
#               four parameters.
#
```

```
.globl average
average:
```

```
movq $0, %rax
addq %rdi, %rax
addq %rsi, %rax
addq %rdx, %rax
addq %rcx, %rax
shrq $2, %rax
```

add all
store in
rax
shift right 2
= divide by 2^2
= 4

```
ret
```

Q7 9

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**Table of x86-64 Registers**

64-bit (quad)	32-bit (double)	16-bit (word)	8-bit (byte)		
63..0	31..0	15..0	15..8	7..0	
rax	eax	ax	ah	al	Return value
rbx	ebx	bx	bh	bl	Callee saved
rcx	ecx	cx	ch	cl	4th arg
rdx	edx	dx	dh	dl	3rd arg
rsi	esi	si		sil	2nd arg
rdi	edi	di		dil	1st arg
rbp	ebp	bp		bpl	Callee saved
rsp	esp	sp		spl	Stack pointer
r8	r8d	r8w		r8b	5th arg
r9	r9d	r9w		r9b	6th arg
r10	r10d	r10w		r10b	Caller saved
r11	r11d	r11w		r11b	Caller saved
r12	r12d	r12w		r12b	Callee saved
r13	r13d	r13w		r13b	Callee saved
r14	r14d	r14w		r14b	Callee saved
r15	r15d	r15w		r15b	Callee saved



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Table of Powers of 2

Power of 2 ^x	Value	Power of 2 ^x	Value	
0	1			
1	2	-1	1/2	0.5
2	4	-2	1/4	0.25
3	8	-3	1/8	0.125
4	16	-4	1/16	0.0625
5	32	-5	1/32	0.03125
6	64	-6	1/64	0.015625
7	128	-7	1/128	0.0078125
8	256	-8	1/256	0.00390625
9	512	-9	1/512	0.001953125
10	1024	-10	1/1024	0.0009765625
11	2048			
12	4096			
13	8192			
14	16384			
15	32768			
16	65536			

Table of x86-64 Jumps

Instruction	Condition	Description
jmp	always	Unconditional jump
jz / jnz	ZF	Jump if equal / zero
jnz / jnz	~ZF	Jump if not equal / not zero
js	SF	Jump if negative
jns	~SF	Jump if nonnegative
jo	OF	Jump if overflow
jno	~OF	Jump if not overflow
jg / jnle	~(SF ^ OF) & ~ZF	Jump if greater (signed >)
jge / jnl	~(SF ^ OF)	Jump if greater or equal (signed ≥)
jl / jnge	SF ^ OF	Jump if less (signed <)
jle / jng	(SF ^ OF) ZF	Jump if less or equal (signed ≤)
ja / jnbe	~CF & ~ZF	Jump if greater (unsigned >)
jae / jnb	~CF	Jump if greater or equal (unsigned ≥)
jnb / jnae / jc	CF	Jump if less (unsigned <)
jbe / jna / jnc	CF ZF	Jump if less or equal (unsigned ≤)

