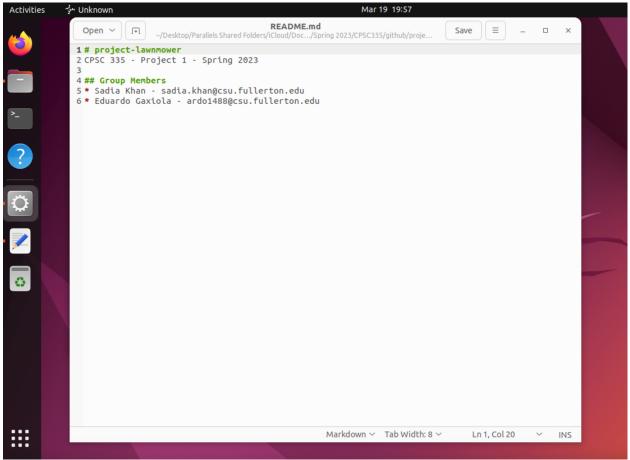
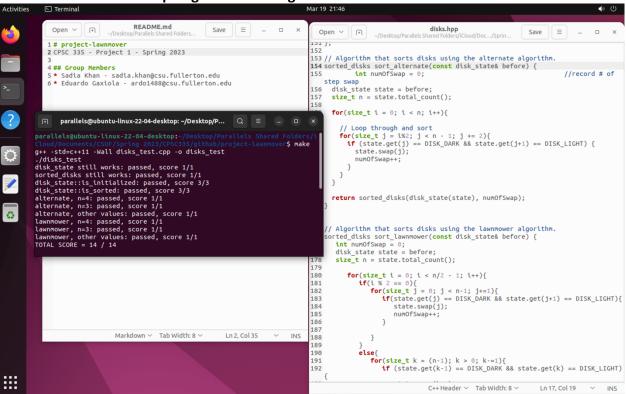
# Project Lawnmower (Project 1 for CPSC 335 course at CSUF – Spring 2023) Group Members: Sadia Khan (sadia.khan@csu.fullerton.edu) and Eduardo Gaxiola (ardo1488@csu.fullerton.edu)

#### Readme screenshot



**Screenshot of Code Compiling and Executing** 



Pseudocode listings, step count, and proof for time complexity:

#### **Alternate Algorithm:**

PseudoCode for Alternate Algorithm with step counts – Time Complexity: O(n^2):

```
Input: alternating disks of size 2n

numOfSwap = 0 // 1 tu

for (i = 0 to 2n, i++) do // 2n-0+1 = 2n+1 times

for (j = i%2 to 2n-1, j+=2) do // ((2n-1)-0)/2 + 1 = n + 1/2 or ((2n-1)-1))/2 + 1 = n

if(disks[j] == dark and disks[j+1] == light) // 4 tu

swap(disks[j], disks[j+1]) // 3 tu

numOfSwap++ // 1 tu

return state, numOfSwap

/* Step Count (SC) for the Alternate Algorithm */

s.c. for if inside inner for = 4 + 4 = 8

s.c. for inner for = max((n + 1/2)*8 , 8n) = 8n + 4
```

```
s.c. for outer for = (2n + 1)*(8n + 4) = 16n^2 + 16n + 4

Total for Alternate Algorithm:

s.c. = 1 + 16n^2 + 16n + 4 = 16n^2 + 16n + 5

===> This leads to O(n^2)
```

Proof:

```
Afternate algorithm:

\lim_{n\to\infty} f(n) = \lim_{n\to\infty} f(n)

= \lim_{n\to\infty} \frac{16n^{2} + 16n + 5}{n^{2}} = \lim_{n\to\infty} \frac{16 + \lim_{n\to\infty} 5}{n + \lim_{n\to\infty} n}

= 16 > 0 \text{ and a constant}

= 16 > 0 \text{ and a constant}
```

### **Lawnmower Algorithm:**

Input: alternating disks of size 2n

<u>PseudoCode for Lawnmower Algorithm with step counts – Time Complexity: O(n^2):</u>

return state, numOfSwap

## **Step Count for Lawnmower:**

s.c inside the outer for loop:  $2 + \max(\max(6n-3, 14n-7), \max(6n, 14n)) = (14n + 2)$  sc. Outer loop:  $n(14n+2) = 14n^2 + 2n // O(n^2)$ 

$$\lim_{n\to\infty} \frac{f(n)}{g(n)}$$
=  $\lim_{n\to\infty} \frac{14n^2 + 2n}{n^2} = \lim_{n\to\infty} \frac{14 + \lim_{n\to\infty} \frac{2}{n}}{n}$ 
=  $14 + 0 = 14 = 20$  and a constant
$$\frac{14n^2 + 2n}{n^2} \in O(n^2)$$