

## APS 105 Lecture 35 Notes

Last time: selection sort, bubble sort and introduced quicksort

Today: Recap quicksort and develop its source code

Recap:

Eg.      2    1    3    5    4

↑  
sorted as it is in its correct position, also called as "pivot" where all #s before are  $< 3$  and all #s after are  $> 3$ .

Quicksort works by ordering a pivot in its location, then work on left subarray then on right subarray

Example:

left right

0 1 2 3 4 5 6 7 8

pivot = 10

10	14	8	13	20	3	6	9	4
----	----	---	----	----	---	---	---	---

$A[\text{right}] < 10$  From right, look at number that belongs to the left of pivot  
 $\text{left}++$  From left, look at number that belongs to the right of pivot

left right

0 1 2 3 4 5 6 7 8

swap 14 & 4

10	4	8	13	20	3	6	9	14
----	---	---	----	----	---	---	---	----

Increment left till you find  $A[\text{left}] > \text{pivot}$

Decrement right till you find  $A[\text{right}] < \text{pivot}$

left right

0 1 2 3 4 5 6 7 8

swap 13 & 9

10	4	8	13	20	3	6	9	14
----	---	---	----	----	---	---	---	----

left right

0 1 2 3 4 5 6 7 8

10	4	8	9	20	3	6	13	14
----	---	---	---	----	---	---	----	----

Increment left till you find  $A[\text{left}] > \text{pivot}$

Decrement right till you find  $A[\text{right}] < \text{pivot}$

left right

0 1 2 3 4 5 6 7 8

swap 20 & 6

10	4	8	9	20	3	6	13	14
----	---	---	---	----	---	---	----	----

Increment left till you find  $A[\text{left}] > \text{pivot}$

Decrement right till you find  $A[\text{right}] < \text{pivot}$

right left

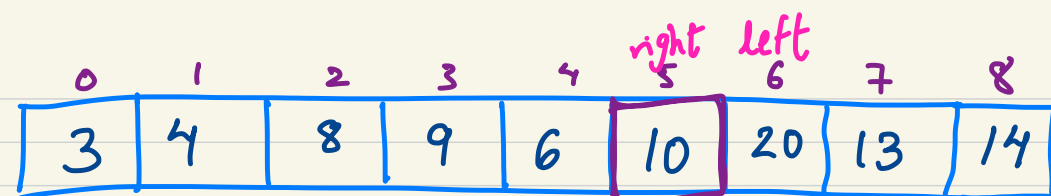
0 1 2 3 4 5 6 7 8

Now (left > right)

10	4	8	9	6	3	20	13	14
----	---	---	---	---	---	----	----	----

Exchange pivot with  $A[\text{right}]$

(3)



Continue doing QuickSort  
on left sub-array

all elements on left is  $< 10$   
all elements on right  $> 10$

Continue doing  
QuickSort on  
right sub-array

Step 1: implement a function that works on an array from index low to index high. Finds a pivot and makes all #s before it  $<$  pivot and all # after it  $>$  pivot. The function returns the final index of the pivot.

```
int partition (int list[], int low, int high) {
    bool done = false;
    int pivotInd = low, left = low + 1, right = high;
    while ( ! done ) {
```

I want  
left & right  
to cross one  
another to  
stop!

```
        while (left < right && list[pivotInd] > list[left])
            left ++;
```

```
        while (left <= right && list[pivotInd] < list[right])
```

```
            right --;
        if (left < right) swap (&list[left], &list[right]);
        else { swap (&list[right], &list[pivotInd]);
              done = true; }
```

```
    return right;
}
```

(4)

Step 2: Call partition but on the subarray on the left of pivot and right of pivot

```
void quickSort(int list[], int length) {
    quickSortHelper(list, 0, length-1);
```

```
}
```

```
void quickSortHelper(int list, int low, int high) {
```

```
    if (low < high) {
```

```
        int pivotInd = partition(list, low, high);
```

```
        quickSortHelper(list, low, pivotInd-1);
```

```
        quickSortHelper(list, pivotInd+1, high);
```

```
    }
```

```
}
```

Base Case is to do nothing!

sorts the  
left subarray

sorts the  
right subarray

# Searching Algorithms

Search for an element in an array

```
int sequentialSearch ( int list[], int length, int data){
    int index = -1;
```

At most we do  $\left\{ \begin{array}{l} \text{for (int } i=0; i < \text{length} \ \&\& \text{index} == -1; i++) \\ \text{length comparisons} \left\{ \begin{array}{l} \text{if (list[i] == data)} \end{array} \right. \end{array} \right.$

Best case : 1 comp.  $\text{index} = i;$

Average:  $n/2$  comp.  $\text{return } i;$

Is there a better way?

If my array was sorted, e.g.  $\begin{matrix} 0 & 1 & 2 & 3 & 4 \\ 1 & 3 & 5 & 10 & 13 \end{matrix}$   
and I want to look for 10?

1) look at  $\text{arr}[n/2] = \text{arr}[5/2] = \text{arr}[2] = 5$

if  $10 > 5 \rightarrow$  look right subarray

$10 < 5 \rightarrow$  look left subarray

2) Repeat (1)

$\text{arr}[\frac{n/2 + n}{2}] = \text{arr}[3] = 10$

if  $10 == 10 \rightarrow$  found

We eliminate half of the array everytime.

We do  $\log_2(\text{length})$  comparisons.

We call the method "binary search"

# Iterative Solution

⑥

```
int binarySearch (int list[], int length, int data){
```

```
    int low = 0, high = length - 1;
```

```
    while ( low <= high ) {
```

```
        int middle = (low + high) / 2;
```

```
        if ( list[middle] == data)
```

```
            return middle;
```

```
        else if ( list[middle] > data)
```

```
            high = middle - 1;
```

```
        else
```

```
            low = middle + 1;
```

```
    }
```

↘

0 1 3 5 8 13

look for 1

① low = 0      middle = 2      high = 5

② low = 0      middle = 0      high = 1

③ low = 1      middle = 1      high = 1

found, so we  
need to enter  
loop when low == high

Recursive

⑦

```
int binarySearchHelper(int list[], int length,  
                        int data, int low, int high){
```

```
    if (low > high)  
        return -1;
```

```
    int middle = (low + high) / 2;  
    if (list[middle] == data)  
        return middle;
```

```
    if (list[middle] > data) go left  
        return binarySearchHelper(list, length, data,  
                                    low, middle - 1);
```

```
    else go right  
        return binarySearchHelper(list, length, data,  
                                    middle + 1, high);
```

```
}
```

```
int binarySearch(int list[], int length, int data){
```

```
    return binarySearchHelper(list, length, data, 0,  
                               length - 1);
```