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# 1. C++ Macros and useful hearders

```
//#pragma warning (disable: 4786)
//#pragma comment (linker, "/STACK:0x800000")
//#define CRT SECURE NO WARNINGS 1
#include<bits/stdc++.h>
using namespace std;
template < class T > T _abs(T n) { return (n < 0 ? -n : n); }</pre>
template < class T > T _max(T a, T b) { return (!(a < b) ? a :</pre>
b);}
template < class T > T _min(T a, T b) { return (a < b ? a : b); }</pre>
template<class T > T sq(T x) { return x * x; }
#define ALL(p) p.begin(),p.end()
#define MP(x, y) make_pair(x, y)
#define SET(p) memset(p, -1, sizeof(p))
#define CLR(p) memset(p, 0, sizeof(p))
#define MEM(p, v) memset(p, v, sizeof(p))
#define CPY(d, s) memcpy(d, s, sizeof(s))
#define READ(f) freopen(f, "r", stdin)
#define WRITE(f) freopen(f, "w", stdout)
#define SZ(c) (int)c.size()
#define PB(x) push_back(x)
#define ff first
#define ss second
#define i64 long long
#define ld long double
#define pii pair< int, int >
#define psi pair< string, int >
#define cpp io() {ios base::sync with stdio(false);cin.tie(0);}
#define FOR(i,N) FORR(i,0,N)
#define FORR(i,a,b) FOTR(i,a,b,1)
#define FOTR(i,a,b,c) for(int i=(a);i<(b);i+=(c))
#define 11 long long
#define MAX 100000000
#define MOD 1000003
const double EPS = 1e-9;
const int INF = 0x7f7f7f7f;
const double PI = 3.141592653589793;
const double E = 2.718281828459045;
/** ----- **/
```

## 2. Data Structures

#### 2.1 SUFFIX ARRAY

```
/*sa:
```

The suffix array. Contains the n suffixes of s sorted in lexicographical order. Each suffix is represented as a single integer (the position in the string where it starts).

#### rank:

The inverse of the suffix array. rank[i] = the index of the suffix s[i..n) in the pos array. (In other words, sa[i] =  $k \le rank[k] = i$ ). With this array, you can compare two suffixes in O(1):

Suffix s[i..n) is smaller than s[j..n) if and only if rank[i] < rank[j].

#### 1cp:

The length of the longest common prefix between two consecutive suffixes:

```
lcp[i] = lcp(s + sa[i], s + sa[i-1]). lcp[0] = 0.
lcp(i,j) = min(lcp[i], lcp[i + 1], ..., lcp[j - 1]). */
int cmp(double x, double y = 0, double tol = EPS) {
     return (x \le y + tol)? (x + tol < y)? -1 : 0 : 1;}
const int MAXN = 100005;
char s[MAXN];
int t, n, rnk[MAXN], sa[MAXN], Lcp[MAXN];
// n = strlen(s) + 1
bool compare(int i, int j) {
     return rnk[i + t] < rnk[j + t];}</pre>
void setLCP() {
     int size = 0, i, j;
     for (i = 0; i < n; i++)
          if (rnk[i] > 0) {
               j = sa[rnk[i] - 1];
               while (s[i + size] == s[j + size]) ++ size;
               lcp[rnk[i]] = size;
               if (size > 0)--size;}
     Lcp[0] = 0;
int distinctSubstringsCount() {
     int ans = n - 1 - sa[1], i;
     for (i = 2; i < n; i++) {
          int ln = n - 1 - sa[i];
          ans += (ln - lcp[i]);
     return ans;}
```

```
void distinctSubstrings() {
     vector<pair<int, int>> substrs;
     int i, j, ln;
     for (i = 1; i < n; i++) {
          ln = n - 1 - sa[i];
          for (j = sa[i] + lcp[i]; j < n - 1; j++)
            substrs.push back(make pair(sa[i], j - sa[i] + 1);
     for (i = 0; i < substrs.size(); i++) {</pre>
          for (j = substrs[i].first; j < substrs[i].first +</pre>
                 substrs[i].second;j++)printf("%c", s[j]);
          printf("\n");}}
void build() {
     int bc[256];
     for (int i = 0; i < 256; ++i)
          bc[i] = 0;
     for (int i = 0; i < n; ++i)
          ++bc[s[i]];
     for (int i = 1; i < 256; ++i)
          bc[i] += bc[i - 1];
     for (int i = 0; i < n; ++i)
          sa[--bc[s[i]]] = i;
     for (int i = 0; i < n; ++i)
          rnk[i] = bc[s[i]];
     for (t = 1; t < n; t <<= 1)
          for (int i = 0, j = 1; j < n; i = j++) {
               while (j < n \&\& rnk[sa[j]] == rnk[sa[i]])
                     j++;
               if (j - i == 1)
                     continue;
               int *start = sa + i, *end = sa + j;
               sort(start, end, compare);
               int first = rnk[*start + t], num = i, k;
               for (; start < end; rnk[*start++] = num) {</pre>
                    k = rnk[*start + t];
                     if (k != first and (i > first or k >= j))
                          first = k, num = start - sa;
               }
          }
     setLCP();
}
void search(char *pat) {
     int m = strlen(pat);
     int l = 1, r = n - 1;
```

```
while (1 <= r) {
          int mid = 1 + (r - 1) / 2;
          int res = strncmp(pat, s + sa[mid], m);
          if (res == 0) {
                cout << s + sa[mid] << endl;</pre>
                return;
          if (res < 0)
                r = mid - 1;
          else
                l = mid + 1;
     cout << "Pattern not found\n";</pre>
}
void printSuffixStrings() {
     //prints suffixes in lexicographical order
     for (int i = 1; i < n; ++i) {
          printf("%s\n", (s + sa[i]));
     }
int main() {
     scanf("%s", s);
     n = strlen(s) + 1;
     build();
     //printSuffixStrings();
     distinctSubstrings();
     //cout<<distinctSubstringsCount();</pre>
     // for (int i = 0; i < 5; i++) {
     //
                char st[10];
               scanf("%s", st);
     //
     //
               search(st);
     //
          }
}
2.2A BIT / FENWICK TREE (LIGHTOJ Curious Robin Hood)
/// LIGHTOJ Curious Robin Hood
// SET N ** IMPORTANT
11 bit[MAX]; int N;
//returns sum from 1 to p
11 get(int p) {
     11 \text{ res} = 0;
     for (; p >= 0; p = (p & (p + 1)) - 1) res += bit[p];
     return res;
}
```

```
void update(int p, ll val) {
     for (; p < N + 1; p = (p | (p + 1)))
          bit[p] += val;
11 rangeSum(int low, int high) {
     return get(high) - (low == 0 ? 0 : get(low - 1));
int main() {
     //read();
     int T, m, x, y, z;
     scanf("%d", &T);
     FORR(cs, 1, T + 1) {
          MEM(bit, 0);
          scanf("%d%d", &N, &m);
          FOR(i, N) {
               scanf("%d", &x);
               update(i + 1, x);
          }
          printf("Case %d:\n", cs);
          FOR(LLL, m) {
               scanf("%d", &x);
               if (x == 1) {
                     scanf("%d", &y);
                     int v = rangeSum(y + 1, y + 1);
                     printf("%d\n", v);
                     update(y + 1, -v);
               else if (x == 2) {
                     scanf("%d %d", &y, &z);
                     update(y + 1, z);
                }
               else {
                     scanf("%d %d", &y, &z);
                     printf("%\frac{11d}{n}, rangeSum(y + 1, z + 1));
                }}}
2.2B BIT / FENWICK TREE (RANGE UPDATE - SPOJ HORRIBLE)
//range update, solution of spoj horrible
11 bit1[MAX], bit2[MAX];
int N;
11 get(11 *bit, int p) {
     11 \text{ res} = 0;
     for (; p \ge 0; p = (p & (p + 1)) - 1)
          res += bit[p];
     return res;
}
```

```
void update(ll *bit, int p, ll val) {
     for (; p < N + 1; p = (p | (p + 1)))
          bit[p] += val;
}
// adds val to each elements in [low, high]
void rangeUpdate(int low, int high, ll val) {
     update(bit1, low, val);
     update(bit1, high + 1, -val);
     update(bit2, low, val * (low - 1));
     update(bit2, high + 1, -val * high);
}
11 get_p(int pos) {
     return get(bit1, pos) * pos - get(bit2, pos);
}
// returns sum in range [low, high]
11 rangeSum(int low, int high) {
     return get_p(high) - (low == 0 ? 0 : get_p(low - 1));
}
int main() {
     cpp io();
     int t, c, m, p, q; 11 v;
     cin >> t:
     while (t--) {
          cin >> N >> c; // importent set N!
          memset(bit1, 0, sizeof(bit1));
          memset(bit2, 0, sizeof(bit2));
          while (c--) {
               cin >> m;
               if (m == 0) {
                    cin >> p >> q >> v;
                    rangeUpdate(p, q, v);
               }
               else {
                    cin >> p >> q;
                    cout << rangeSum(p, q) << "\n";}}}}</pre>
2.3 2D BIT ( SPOJ MATSUM )
//http://spoj-solutions.blogspot.com/2014/10/matsum-matrix-
summation.html
**** REMARKS ****
  1. update function adds given value with already set value.
  2. BIT can not start with index 0. That's why input is added
     with 1.
ll tree[1050][1050];
```

```
void update(int x, int y, ll val, int MAX) {
     while (x <= MAX) {</pre>
          int ty = y;
          while (ty <= MAX) {</pre>
                tree[x][ty] += val;
                ty += (ty \& -ty);
          x += (x \& -x);
     }
11 read(int x, int y) {
     11 \text{ sum } = 0;
     while (x) {
          int ty = y;
          while (ty) {
                sum += tree[x][ty];
                ty -= (ty & -ty);
          x -= (x \& -x);
     return sum;
11 readRegtangle(int x0, int y0, int x1, int y1) {
     return read(x1, y1) + read(x0 - 1, y0 - 1) - read(x0 - 1,
            y1)- read(x1, y0 - 1);}
int main() {
     int t; scanf("%d", &t);
     while (t--) {
          int n;
          scanf("%d", &n);memset(tree, 0, sizeof tree);
          while (1) {
                char s[10];scanf("%s", s);
                if (s[1] == 'E') {
                     int x, y, val;
                     scanf(" %d%d%d", &x, &y, &val);
                     11 p val = readRegtangle(x + 1, y + 1, x +
                     1, y + 1);
                     update(x + 1, y + 1, val - p_val, n + 9);
                } else if (s[1] == 'U') {
                     11 \text{ sum} = 0;
                     int x1, y1, x, y;
                     scanf(" %d%d%d%d", &x, &y, &x1, &y1);
                     sum = readRegtangle(x + 1, y + 1, x1 + 1,
                     y1 + 1);
                     printf("%11d\n", sum);
```

```
} else
                     break:
          }
          printf("\n");
     }
}
2.4 SEGMENT TREE (SPOJ GSS3)
const 11 MAX = 500000;
const 11 MIN = -500000000000000000LL;
struct NODE {
     11 rmax, 1max, tmax, sum;
} tree[MAX];
int inp[MAX];
int max3(ll a, ll b, ll c) {
     11 \text{ mx} = a;
     if (b > mx)
          mx = b;
     if (c > mx)
          mx = c;
     return mx;
}
void setNode(int node) {
     int ch1 = node \ll 1, ch2 = ch1 \mid 1;
     tree[node].sum = tree[ch1].sum + tree[ch2].sum;
     tree[node].lmax = max(tree[ch1].lmax, tree[ch1].sum +
tree[ch2].lmax);
     tree[node].rmax = max(tree[ch2].rmax, tree[ch1].rmax +
tree[ch2].sum);
     tree[node].tmax = max3(tree[ch1].tmax, tree[ch2].tmax,
               tree[ch1].rmax + tree[ch2].lmax);
NODE join(NODE n1, NODE n2) {
     NODE n;
     n.sum = n1.sum + n2.sum;
     n.lmax = max(n1.lmax, n1.sum + n2.lmax);
     n.rmax = max(n2.rmax, n1.rmax + n2.sum);
     n.tmax = max3(n1.tmax, n2.tmax, n1.rmax + n2.lmax);
     return n;
}
void build(int node, int first, int last) {
     if (first == last) {
          tree[node].rmax = tree[node].tmax = tree[node].sum =
          tree[node].lmax =inp[first];}
```

```
else {
          int mid = (first + last) >> 1, ch1 = node << 1, ch2 =</pre>
          ch1 1;
          build(ch1, first, mid);
          build(ch2, mid + 1, last);
          setNode(node);
     }
}
void update(int node, int first, int last, int p, int val) {
     if (first == last) {
          tree[node].rmax = tree[node].tmax = tree[node].sum =
          tree[node].lmax =val;
          return;
     }
     int mid = (first + last) >> 1;
     if (p <= mid)
          update(node << 1, first, mid, p, val);</pre>
     else
          update((node << 1) | 1, mid + 1, last, p, val);</pre>
     setNode(node);
}
NODE query(int node, int first, int last, int l, int r) {
     if (first > last | first > r | last < l)</pre>
          return {MIN,MIN,MIN,MIN};
     if (first >= 1 && last <= r)
          return tree[node];
     int mid = (first + last) >> 1;
     NODE n1 = query(node << 1, first, mid, 1, r);
     NODE n2 = query((node \langle\langle 1\rangle | 1, mid + 1, last, 1, r);
     if (n1.tmax == MIN)
          return n2;
     else if (n2.tmax == MIN)
          return n1;
     return join(n1, n2);
}
int main() {
     cpp_io();
     int n, m, x, y, z;
     cin >> n;
     FORR(i, 1, n + 1)cin >> inp[i];
     build(1, 1, n); cin >> m;
     FOR(i, m) {
          cin >> x >> y >> z;
```

```
if (x == 1) {
                 NODE nd = query(1, 1, n, y, z);
                 cout << nd.tmax << "\n";</pre>
           }
           else {
                 update(1, 1, n, y, z);
           }
     }
}
2.5 SEGMENT TREE WITH LAZY PROPAGATION (SPOJ MULTQ3)
const 11 MAX = 500000;
struct NODE {
     int zero, one, two;
     NODE() {
           zero = one = two = 0;
}tree[MAX];
int Lazy[MAX];
void setNode(int node) {
     int ch1 = node << 1, ch2 = ch1 ^ 1;</pre>
     tree[node].zero = tree[ch1].zero + tree[ch2].zero;
     tree[node].one = tree[ch1].one + tree[ch2].one;
     tree[node].two = tree[ch1].two + tree[ch2].two;
void lazyUpdate(int node, int lazyVal, bool push) {
     int v[3];
     v[0] = tree[node].zero, v[1] = tree[node].one,
     v[2] = tree[node].two;
     tree[node].zero = v[lazyVal];
     tree[node].one = v[(lazyVal + 1) % 3];
     tree[node].two = v[(lazyVal + 2) % 3];
     if (push) {
           lazy[node << 1] += lazyVal;</pre>
           lazy[(node << 1) | 1] += lazyVal;</pre>
     Lazy[node] = 0;
}
void build(int node, int first, int last) {
     if (first == last) tree[node].zero++;
     else {
           int mid = (first + last) >> 1;
           build(node << 1, first, mid);</pre>
           build((node << 1) | 1, mid + 1, last);</pre>
           setNode(node);
           }
}
```

```
void update(int node, int first, int last, int l, int r, int val) {
     if (Lazy[node] != 0) {
           lazyUpdate(node, Lazy[node] % 3, first != last);
     if (first > last | first > r | last < l) return;</pre>
     if (first >= 1 && last <= r) { //update</pre>
           lazyUpdate(node, val % 3, first != last);
           return;
     int mid = (first + last) >> 1;
     update(node << 1, first, mid, l, r, val);</pre>
     update((node << 1) | 1, mid + 1, last, l, r, val);</pre>
     setNode(node);
int query(int node, int first, int last, int l, int r) {
     if (first > last | | first > r | | last < 1) return 0;</pre>
     if (lazy[node] != 0) {
           lazyUpdate(node, Lazy[node] % 3, first != last);
     if (first >= 1 && last <= r) return tree[node].zero;</pre>
     int mid = (first + last) >> 1;
     int p1 = query(node << 1, first, mid, l, r);</pre>
     int p2 = query((node << 1) | 1, mid + 1, last, l, r);
     return (p1 + p2);
}
int main() {
     //read();
     cpp_io();
     int n, m, x, y, z;
     cin >> n >> m;
     build(1, 1, n);
     FOR(i, m) {
           cin >> x >> y >> z;
           if (x == 1) {
                 cout << query(1, 1, n, y + 1, z + 1) << "\n";</pre>
           }
           else {
                 update(1, 1, n, y + 1, z + 1, 1);
           }
     }
}
```

# 3. Dynamic Programming

```
3.1 1/0 Knapsack Top Down
int dp[MAX][MAX], P[]=\{60, 100, 120\}, W[]=\{10, 20, 30\}, N;
int knps(int i, int w) {
     if (i >= N \mid | w == 0)return 0;
     if (dp[i][w] != -1)return dp[i][w];
     int p1 = 0, p2 = 0;
     if (W[i] \le W)p1 = P[i] + knps(i + 1, W - W[i]);
     p2 = knps(i + 1, w);
     return dp[i][w] = max(p1, p2);
3.2 LIS (NON DP, N LOG N)
int n; int \nu[] = \{50, 3, 10, 7, 40, 80\};
int LISLength() {
     multiset<int> s; //use set if input doesn't have duplicates
     multiset<int>::iterator it;
     for (int i = 0; i < n; i++) {
           s.insert(v[i]);
           it = upper_bound(s.begin(), s.end(), v[i]);
          /* for strictly increasing:
          it = lower_bound(s.begin(), s.end(), x); it++; */
           if (it != s.end()) s.erase(it);
     return s.size();
}
3.3 LIS RETURN RESULT
vector<int> LISResult(vector<int> input) {
     vector<int>L(input.size());
     vector<int>L id(input.size());
     vector<int>P(input.size());
     vector<int>output;
     int lis = 0, lis_end = 0, i;
     for (int i = 0; i < input.size(); ++i) {</pre>
           int pos = lower_bound(L.begin(), L.begin() + lis,
        input[i]) - L.begin();
        //use upper bound for non decreasing
           L[pos] = input[i];
           L id[pos] = i;
           P[i] = pos ? L_id[pos - 1] : -1;
           if (pos + 1 > lis) {
                 lis = pos + 1; lis end = i;
           }}
     i = lis end;
     for (; P[i] >= 0; i = P[i]) output.push back(input[i]);
     output.push back(input[i]);
     reverse(output.begin(), output.end());
     return output;}
```

```
3.4 Longest Common Subsequence:
char str1[] = "AGGTAB", str2[] = "GXTXAYB";
int lcs(int n1, int n2) {
     int lcs[n1 + 1][n2 + 1];
     for (int i = 0; i <= n1; i++)</pre>
           for (int j = 0; j <= n2; j++) {</pre>
                 if (i == 0 || j == 0)
                       lcs[i][j] = 0;
                 else {
                       lcs[i][j] = max(lcs[i - 1][j], lcs[i][j - 1]);
                       if (str1[i - 1] == str2[j - 1])
                             lcs[i][j] = max(lcs[i][j],
                            lcs[i - 1][j - 1] + 1);
                 }
     return lcs[n1][n2];
3.5 Matrix Chain Multiplication:
int m[MAX][MAX], s[MAX][MAX];
int MatrixChainOrder(int p[], int n){
     n = n - 1; //length of p -1
     int L, i, j, k, q;
     for (int i = 1; i \le n; i++) m[i][i] = 0;
     for ( L = 2; L <= n; L++) {
           for (i = 1; i \le n - L + 1; i++) {
                 j = i + L - 1;
                 m[i][j] = INT MAX;
                 for (k = i; k \le j - 1; k++) {
                       q = m[i][k] + m[k + 1][j]
                       + p[i - 1] * p[k] * p[j];
                       if (q < m[i][j]) {
                            m[i][j] = q;
                            s[i][j] = k;
                       }
                 }
           }
     }
     return m[1][n];
}
void print(int p[], int i , int j) {
     if (i == j) cout << p[i]<< " ";</pre>
     else {
           cout << "(";</pre>
           print(p, i, s[i][j]);
           print(p, s[i][j]+1, j);
           cout << ")";</pre>
     }
}
```

# 3.6 Longest Common Increasing Subsequence (LCIS)

```
#define N 2000
int n, a[N], m, b[N], dp[N], pr[N], l;
void trace(int x, bool m) {
     if (!x)
           return;
     trace(pr[x], 1);
     printf("%d", b[x]);
     putchar(m ? ' ' : '\n');
}
void lcis() {
     int i, j;
     memset(dp, 0, sizeof(dp));
     memset(pr, 0, sizeof(pr));
     for (i = 1; i \le n; i++) {
           int c = 0, p = 0;
           for (j = 1; j \leftarrow m; j++) {
                 if (a[i] == b[j] && dp[j] < c + 1) {
                      dp[j] = c + 1;
                      pr[j] = p;
                 if (a[i] > b[j] \&\& c < dp[j]) {
                      c = dp[j];
                       p = j;
                 }
           }
     l = -1; // l = lcis length
     for (i = 1; i <= m; i++)
           L = \max(L, dp[i]);
}
int main() {
     scanf("%d", &n); /// let n = 4 int i, j;
     for (i = 1; i \le n; i++)
           scanf("%d", &a[i]); /// let a[] = {3, 4, 9, 1};
     scanf("%d", &m); /// let m = 7
     for (i = 1; i \le m; i++)
           scanf("%d", &b[i]); /// let b[] = \{5, 3, 8, 9, 10, 2, 1\};
     lcis();
     printf("Length of LCIS is %d\n\n", l); /// Length of LCIS is 2
     if (l) {
           for (i = 1; i <= m; i++)
                 if (dp[i] == L) {
                       cout << "The LCIS is : ";</pre>
                       trace(i, 0);
                       break;
                 }
     }
}
```

#### 4. Mathematics

```
4.1 Sieve of Eratosthenes
#define MAX 46656
#define LMT 216
#define LEN 4830
#define RNG 100032
#define chkC(x,n) (x[n>>6]&(1<<((n>>1)&31)))
#define setC(x,n) (x[n>>6]|=(1<<((n>>1)&31)))
unsigned base[MAX / 64], segment[RNG / 64], primes[LEN];
void sieve() {
     unsigned i, j, k;
     for (i = 3; i < LMT; i += 2)
           if (!chkC(base, i))
                for (j = i * i, k = i << 1; j < MAX; j += k)
                      setC(base, j);
     primes[0] = 2;
     j = 1;
     for (i = 3; i < MAX; i += 2)
           if (!chkC(base, i))
                primes[j++] = i;
}
bool isPrime(int n) {
     if (n == 2 | (n > 2 & n & 1 & ! chkC(base, n)))
           return true;
     return false;
}
4.2 Segmented Sieve
/// NOT OK
#define sq(x) ((x)*(x))
#define mset(x,v) memset(x,v,sizeof(x))
int primes[100], segment[100]; /// ADDED by SunW
void segmented_sieve(int a, int b) {
     unsigned i, j, k, cnt = (a \le 2 \&\& 2 \le b) ? 1 : 0;
     if (b<2) return;
     if (a<3) a = 3;
     if (a \% 2 == 0) a++;
     mset(segment, 0);
     for (i = 1; sq(primes[i]) <= b; i++) {
           j = primes[i] * ((a + primes[i] - 1) / primes[i]);
           if (j % 2 == 0) j += primes[i];
           for (k = primes[i] << 1; j <= b; j += k)
                if (j != primes[i]) setC(segment, (j - a));
     }
bool isPrimeSeg(int n, int a) {
     if (n == 2)return true;
     else if (!(n & 1) || n<2) return false;
```

```
if (!chkC(segment, n - a))return true;
     return false;
}
4.3 Prime Factors, Divisors & Numbder of Divisors
vector<int> v;
int primes[] = { 2, 3, 5, 7, 11, 13, 17, 19 }; /// ADDED by SunW
void primeFactors(int n) { //Ex. When n = 12, v[] = {2,2,3}
     for (int i = 0; primes[i] * primes[i] <= n; i++) {</pre>
           int x = primes[i];
           while (n \% x == 0) \{
                 v.push back(x);
                 n = n / x;
           }
     if (n \ge 2)
           v.push_back(n);
int noOfDivisors(int n) { //When n = 12, returns 6
     int c = 0, ans = 1;
     for (int i = 0; primes[i] <= sqrt(n); i++) {</pre>
           c = 0;
           int x = primes[i];
           while (n \% x == 0) \{
                 C++;
                n = n / x;
           if (c > 0)
                 ans *= (c + 1);
     if (n \ge 2)
           ans *= 2;
     return ans;
}
vector<int> divisors;
void allDivisors(int n) { //all divisors for small n
     divisors.clear();
     int i;
     for (i = 1; i < sqrt(n); i++) {
           if (n % i == 0) {
                divisors.push_back(i);
                 divisors.push back(n / i);
           }
     if (i * i == n)
           divisors.push back(i);
}
vector<pii> factors;
```

```
// When n = 12, factors[] = {(2,2), (3,1)}
void primeFactorsPair(int n) {
     int c = 0;
     for (int i = 0; primes[i] <= sqrt(n); i++) {</pre>
           c = 0;
           int x = primes[i];
           while (n \% x == 0) \{
                 C++;
                 n = n / x;
           }
           if (c > 0)
                 factors.push back(make pair(x, c));
     if (n \ge 2)
           factors.push back(make pair(n, 1));
}
//call primeFactorsPair with n(12)
//call allDivisors with (1, 0), divisors[] = {2,6,4,12,3}
//also needs sieve()
void allDivisors(int n, int i) {
     int j, x, k;
     for (j = i; j < factors.size(); j++) {</pre>
           x = factors[j].first * n;
           for (k = 0; k < factors[j].second; k++) {</pre>
                 divisors.push_back(x);
                 allDivisors(x, j + 1);
                 x *= factors[j].first;
           }
     }
}
///Divisors of numbers from 1 to N
/* When MAX = 500007 it takes about 600ms/51MB in codeforces
MAX = 10^6 \text{ takes } 1.4s \text{ in uva*/}
#define MAX 500007
vector<int> divisorz[MAX + 1];
void divisors 1 to n(int n) {
     for (int i = 1; i <= n; i++) {</pre>
           for (int j = i; j <= n; j += i)
                 divisorz[j].push back(i);
           // divisors[j]+=i for Sum of divisors
           // divisors[j]++ for Number of divisors
     }
}
4.4 GCD, LCM, EGCD
11 gcd(ll a, ll b) {
     return b ? gcd(b, a % b) : a;
}
```

```
11 lcm(ll a, ll b) {
     return a * (b / gcd(a, b));
//egcd solves a.x + b.y = gcd(a,b)
//if used to solve for ax+by=1, where a & b are coprime,
//check if egdc returns -1. If it does, swap x, y
11 egcd(11 a, 11 b, 11 & x, 11 & y) {
     if (a == 0) {
           x = 0;
           y = 1;
           return b;
     ll x1, y1;
     ll d = egcd(b \% a, a, x1, y1);
     x = y1 - (b / a) * x1;
     y = x1;
     return d;
}
4.5 Euler Phi/Totient Function
long long phi[MAX + 2], mark[MAX + 2];
//Phi from 1 to n
void ETF(int n) {
     int i, j;
     for (i = 0; i <= n; i++)
           phi[i] = i;
     mark[1] = 1;
     for (i = 2; i \le n; i++) {
           if (!mark[i]) {
                for (j = i; j \le n; j += i) {
                      mark[j] = 1;
                      phi[j] = phi[j] / i * (i - 1);
                 }
           }
     }
}
int phif(int n) { //phi for n
     int ret = n;
//iterate through primes[] for faster result
     for (int i = 2; i * i <= n; i++) {
           if (n % i == 0) {
                while (n % i == 0)
                      n /= i;
                ret -= ret / i;
           }
     if (n > 1)
           ret -= ret / n;
     return ret;
}
```

# 4.6 Permutation nPr, Combination nCr

```
///NOT OK
long long F[MAX];
void initF() {
     F[1] = 1;
     F[0] = 1;
     for (int i = 2; i < MAX; i++)</pre>
           F[i] = (F[i - 1] * i) \% MOD;
}
long long POW(int a, int b) { // a^b)%MOD
     long long x = 1, y = a;
     while (b > 0) {
           if (b & 1)
                x = (x * y) % MOD;
           y = (y * y) % MOD;
           b >>= 1;
     }
     return x;
long long inversePow(int n) {
     return POW(n, MOD - 2);
long long (int n, int r) {
     return (F[n] * ((inversePow(F[r]) * inversePow(F[n - r])) %
MOD)) % MOD;
}
long long P(int n, int r) {
     return (F[n] * inversePow(F[n - r])) % MOD;
}
```

## Some Formulas:

\*\* 
$$C(n, k) = C(n-1, k-1) + C(n-1, k)$$

| 1 | Permutation:                             | n!                      |
|---|--|-------------------------|
|   | Order ✓   Repetition ⊁                   | $\overline{(n-r)!}$     |
| 2 | Order ✓   Repetition ✓                   | n <sup>r</sup>          |
| 3 | Combination:                             | n!                      |
|   | Order 🗴   Repetition 🗴                   | $\overline{(r!)(n-r)!}$ |
| 4 | Permutation when P1 are of same kind, P2 | n!                      |
|   | are of same kind and so on               | $p_1! p_2! \cdots p_r!$ |
| 5 | Circular permutation of n distinct       | (n-1)!                  |
|   | object                                   |                         |

| 6  | Circular permutation when clockwise and                          |   |
|----|--|---|
|    | anti-clockwise arrangements are not                              | .5*(n-1)!   |
|    | different  |   |
| 7  | Number of <b>permutations</b> of <b>n</b> distinct               |   |
|    | things taking ${f r}$ at a time, when ${f s}$                    | $^{(n-s)}C_{(r-s)} \times r!$   |
|    | particular things are always to be                               | (r-s) · · · · ·   |
|    | <pre>included in each arrangement</pre>                          |   |
| 8  | Number of <b>permutations</b> of <b>n</b> distinct               | ()  |
|    | things taking ${f r}$ at a time, when ${f s}$                    | $^{(n-s)}C_r \times r!$   |
|    | particular things are <b>never included</b>                      |   |
| 9A | Number of <b>combinations</b> of <b>n</b> distinct               |   |
|    | objects taking $oldsymbol{r}$ at a time when each                | $^{(\mathrm{n+r-1})}\mathrm{C_r}$                                     |
|    | object may be <b>repeated any number of</b>                      | •   |
|    | times  |   |
| 9B | Number of ways in which <b>n</b> identical                       |   |
|    | things can be distributed among $oldsymbol{r}$                   | $^{(n+r-1)}C_{(r-1)}$   |
|    | persons, each one of them can receive                            | (1 1)   |
|    | 0,1,2 or more items  |   |
| 9C | Number of ways in which <b>n</b> identical                       |   |
|    | things can be distributed among $oldsymbol{r}$                   | $^{(n-1)}C_{(r-1)}$   |
|    | persons, each one of them can receive                            | (1 1)   |
|    | 1,2 or more items  |   |
| 10 | Number of <b>combinations</b> of <b>n</b> distinct               |   |
|    | things taking ${f r}$ at a time, when ${f s}$                    | $^{(n-s)}C_{(r-s)}$   |
|    | particular things are <b>always included</b> in                  | ()  |
|    | each selection   |   |
| 11 | Number of <b>combinations</b> of <b>n</b> distinct               | (n, a)  |
|    | things taking ${f r}$ at a time, when ${f s}$                    | $^{ m (n-s)}{ m C_r}$   |
|    | particular things are <b>never included</b>                      |   |
| 12 | Number of <b>combinations</b> of <b>n</b> distinct               |   |
|    | things taking ${f r}$ at a time, when ${f m}$                    | $^{\mathrm{n}}\mathrm{C_{r}}$ - $^{\mathrm{(n-m)}}\mathrm{C_{(r-m)}}$ |
|    | particular things <b>never come together</b>                     | - (1 111)   |
| 13 | Number of ways in which <b>n</b> distinct                        |   |
|    | things can be divided into <b>r unequal</b>                      | $\frac{n!}{a_1!\ a_2!\ a_3!\ \cdots\ a_r!}$                           |
|    | non-distinct groups containing a <sub>1</sub> , a <sub>2</sub> , | $a_1! a_2! a_3! \cdots a_r!$  |
|    | a <sub>3</sub> ,, a <sub>r</sub> things                          |   |
| 14 | Number of ways in which <b>n</b> distinct                        | $n!\ r!$  |
|    | things can be divided into <b>r unequal</b>                      | $\overline{a_1!  a_2!  a_3! \dots  a_r!}$                             |
|    | <b>distinct groups</b> containing $a_1$ , $a_2$ , $a_3$ ,        |   |
|    | , a <sub>r</sub> things  |   |
| 15 | Number of ways in which $\mathbf{m} 	imes \mathbf{n}$ distinct   | ( )   |
|    | things can be divided equally                                    | (mn)!   |
|    | into n groups (each group will have m                            | $\frac{(mn)!}{(m!)^n \ n!}$   |
|    | things and the groups are unmarked)                              | ` '   |

Number of ways in which  $m \times n$  distinct things can be divided equally into n groups (each group will have m things and the groups are distinct)  $\frac{(mn)!}{(m!)^n}$ 

| Distribution of |                 | How many balls boxes can contain                  |  |                                       |  |
|-----------------|-----------------|---|--|---------------------------------------|--|
| k Balls         | into n<br>Boxes | No<br>Restrictions                                | ≤ 1<br>(At most<br>one)                  | ≥ 1<br>(At least<br>one)              | = 1<br>(Exactly one)   |
| Distinct        | Distinct        | n <sup>k</sup><br>(formula 1)                     | <sup>n</sup> P <sub>k</sub> (formula 2)  | S(k,n) × n!  (formula 3)  (more info) | $^{n}P_{n} = n!$ if $k = n$ $o$ if $k \neq n$ (formula 4)            |
| Identical       | Distinct        | (k+n-1)C <sub>(n-1)</sub>                         | <sup>n</sup> C <sub>k</sub> (formula 6)  | (k-1) <sub>C(n-1)</sub> (formula 7)   | 1 if k = n<br>0 if k ≠ n<br>(formula 8)                              |
| Distinct        | Identical       | $\sum_{i=1}^{n} S(k,i)$ (formula 9) (more info)   | 1 if k ≤ n 0 if k > n (formula 10)       | S(k,n) (formula 11) (more info)       | 1 if k = n<br>0 if k ≠ n<br>(formula 12)                             |
| Identical       | Identical       | $\sum_{i=1}^{n} P(k, i)$ (formula 13) (more info) | 1 if k ≤ n<br>o if k > n<br>(formula 14) | P(k, n)  (formula 15)  (more info)    | <ul><li>1 if k = n</li><li>0 if k ≠ n</li><li>(formula 16)</li></ul> |

# 4.7 Chinese Remainder Theorem

```
while (true) {
                 if ((k * temp) % mods[i] == 1)
                      break;
                 k++;
           }
           s.push_back(k);
     for (int i = 0; i < int(s.size()); i++) {</pre>
           ret += ((m[i] * s[i]) % M * r[i]) % M;
           if (ret >= M)
                 ret -= M;
     }return ret;}
4.8 POLLARD's RHO (integer factorization algorithm)
const long long LIM = LLONG MAX;
uint64 t gcd(uint64 t a, uint64 t b) {
     return b ? gcd(b, a % b) : a;
long long mul(long long a, long long b, long long m) {
     long long x, res;
     if (a < b)
           swap(a, b);
     if (!b)
           return 0;
     if (a < (LIM / b))
           return ((a * b) % m);
     res = 0, x = (a \% m);
     while (b) {
           if (b & 1) {
                 res = res + x;
                 if (res >= m)
                      res -= m;
           }
           b >>= 1;
           X <<= 1;
           if (x >= m)
                x -= m;
     return res;
long long expo(long long x, long long n, long long m) {
     long long res = 1;
     while (n) {
           if (n & 1)
                 res = mul(res, x, m);
           x = mul(x, x, m);
           n >>= 1;
     return (res % m);
}
```

```
const int small_primes[] = { 2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31,
37, 41, 43,47, 51, 53, 59, 61, 67, 71 };
bool miller_rabin(long long p, int lim) {
     long long a, s, m, x, y;
     s = p - 1, y = p - 1;
     while (!(s & 1))
           s \gg 1;
     while (lim--) {
           X = S;
           a = (rand() \% y) + 1;
           m = expo(a, x, p);
           while ((x != y) \&\& (m != 1) \&\& (m != y)) {
                m = mul(m, m, p);
                x <<= 1;
           if ((m != y) && !(x & 1))
                 return false;
     return true;
}
void brent_pollard_rho(uint64 t n, vector<uint64 t> &v) {
     if (miller_rabin(n, 10)) {
           v.push_back(n);
           return;
     }
     uint64_t a, g, x, y;
     y = 1;
     a = rand() % n;
     x = rand() % n;
     for (int i = 0; ((i * i) >> 1) < n; i++) {
           x = mul(x, x, n);
           x += a;
           if (x < a)
                x += (ULLONG MAX - n) + 1;
           x \% = n;
           g = gcd(n, y - x);
           if ((g != 1) && (g != n)) {
                 n /= g;
                brent_pollard_rho(g, v);
                 if (n != g)
                      brent_pollard_rho(n, v);
                 else if (miller_rabin(n, 10))
                      v.push_back(n);
                 return;
           }
```

```
if (!(i & (i - 1)))
                 y = x;
     brent_pollard_rho(n, v);
}
void factorize(uint64_t n, vector<uint64_t> &v) {
     srand(time(0));
     int i, j, x;
     for (i = 0; i < 21; i++) {
           x = small_primes[i];
           while ((n \% x) == 0) \{
                 n /= x;
                 v.push_back(x);
           }
     if (n > 1)
           brent_pollard_rho(n, v);
     sort(v.begin(), v.end());
}
vector<pair<uint64 t, int> > fc;
vector<uint64_t> dv;
void divs(int n, int i) {
     uint64 t j, x, k;
     for (j = i; j < fc.size(); j++) {</pre>
           x = fc[j].first * n;
           for (k = 0; k < fc[j].second; k++) {</pre>
                 dv.push back(x);
                 divs(x, j + 1);
                 x *= fc[j].first;
           }
     }
void divisors(vector<uint64_t> &v) {
     dv.clear();
     fc.clear();
     int c = 0;
     for (int i = 0; (i + 1) < v.size(); i++) {
           if (v[i] == v[i + 1])
                 C++;
           else {
                 if (c)
                      fc.push_back(make_pair(v[i], c + 1));
                 else
                      fc.push back(make_pair(v[i], 1));
                 c = 0;
           }
     divs(1, 0);
     sort(dv.begin(), dv.end()); }
```

```
int main() {
     uint64_t n, i, t, x;
     cin >> t;
     while (t--) {
          cin >> n;
          vector<uint64_t> v;
          factorize(n, v);
          sort(v.begin(), v.end());
          v.push back(-1); /***IMPORTANT***/
          int len = v.size(), c = 0, counter = 0;
          printf("%11u = ", n);
          for (i = 0; (i + 1) < len; i++) {
                if (v[i] == v[i + 1])
                     counter++;
                else {
                     if (c)
                           printf(" * ");
                     if (counter)
                           printf("%llu^%d", v[i], ++counter);
                     else
                           printf("%llu", v[i]);
                     c++, counter = 0;
                }
           }
          puts("");
          divisors(v);
          for (auto i : dv)
                cout << i << " ";</pre>
          cout << endl;</pre>
     }
}
4.9 Millar Rabin
const int small_primes[] = { 2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31,
37, 41, 43,
          47, 51, 53, 59, 61, 67, 71 };
long long mul(long long a, long long b, long long m) {
     long long x, res;
     if (a < b)
          swap(a, b);
     if (!b)
          return 0;
     if (a < (LIM / b))
          return ((a * b) % m);
     res = 0, x = (a \% m);
     while (b) {
          if (b & 1) {
                res = res + x;
```

```
if (res >= m)
                      res -= m;
           }
           b >>= 1;
           x <<= 1;
           if (x >= m)
                x -= m;
     return res;
long long expo(long long x, long long n, long long m) {
     long long res = 1;
     while (n) {
           if (n & 1)
                res = mul(res, x, m);
           x = mul(x, x, m);
           n >>= 1;
     }
     return (res % m);
bool miller_rabin(long long p, int lim) {
     if (p < 2)
           return false;
     if (p == 2)
           return true;
     if (!(p & 1))
           return false;
     int i, val;
     long long a, s, m, x, y;
     for (i = 0; i < 20; i++) {
           val = small_primes[i];
           if (p == val)
                 return true;
           if ((p % val) == 0) {
                 return false;
           }
     }
     srand(time(0));
     s = p - 1, y = p - 1;
     while (!(s & 1))
           s \gg 1;
     while (lim--) {
           X = S;
           a = (rand() \% y) + 1;
           m = expo(a, x, p);
```

```
while ((x != y) \&\& (m != 1) \&\& (m != y)) {
                m = mul(m, m, p);
                x <<= 1;
          if ((m != y) && !(x & 1))
                return false;
     }return true;}
4.10 Miscellaneous Topics
a. Java BigInteger count digits:
public static int getDigitCount(BigInteger number) {
    double factor = Math.log(2) / Math.log(10);
    int digitCount = (int) (factor * number.bitLength() + 1);
    if (BigInteger.TEN.pow(digitCount - 1).compareTo(number) >
0) {
        return digitCount - 1;
    return digitCount;
BigDecimal Division:
a.divide(b, 1000, RoundingMode.CEILING); //java.math.*
b. ax + by = c
if ( c \% \gcd(a,b) == 0 ), then equation has a solution.
С.
                    primes(x)
               \boldsymbol{x}
               10
                          4
```

#### 100 25 1,000 168 10,000 1,229 100,000 9,592 1,000,000 78,498 10,000,000 664,579 100,000,000 5,761,455 1,000,000,000 50,847,534 10,000,000,000 455,052,511

```
d.
Gray Code
int gray(int n) { return n ^ (n >> 1); }
int rev_g(int g) {
    int n = 0;
    for (; g; g >>= 1) n ^= g;
    return n;
}
```

```
e. //roots of eqn ax^2 + bx + c. a!=0
void quadroots (double a, double b, double c, double &r1,
double &r2) {
     if ((b*b - 4 * a*c) >= 0.0) {
          double d = sqrt(b*b - 4 * a*c);
          r1 = (-b - d) / (2 * a);
          r2 = (-b + d) / (2 * a);
          return;
     }
     //imaginary root
     r1 = -b / (2 * a);
     r2 = sqrt(-(b*b - 4 * a*c)) / (2 * a);
f. Log base change formula
Log_bx = Log_ax / Log_ab
5. Geometry
5.1 Header/define
#include <cmath>
#include <iomanip>
#include <cfloat> //DBL MAX, DBL MIN
#define EPS (1e-9)
#define PI 3.14159265359 // acos(-1);
5.2 Structures
struct point {
     double x, y;
     point() { x = y = 0.0; }
     point(double _x, double _y) : x(_x), y(_y) {}
     bool operator <(const point & p) const {</pre>
          return ((x < p.x) \mid | (abs(x - p.x) < EPS \&\&
               (y < p.y));
     point operator-(point v) { return point(x - v.x, y - v.y);
}
     point operator+(point v) { return point(x + v.x, y + v.y);
};
struct line{ double a, b, c; };
struct vec {
       double x, y;
       vec( double _x, double _y) : x( _x), y( _y) {} };
```

```
5.3 Function list
1.
   dist
                                              29. tarea1
                     15. pointAtDist
                     16. linePointDist1
2.
                                              30. tarea2
   DEGtoRAD
3.
   ROTATE
                     17. linePointDist2
                                              31. tarea3
4.
                                              32. rInCircle
  rotatep
                     18. linePointDist3
                     19. linePointDist4
                                              33. rInCircleP
5.
   toVec
                                              34. rCircumCircle
6.
   scale
                     20. toLine
7.
   translate
                     21. parallel
                                              35. rCircumCircleP
8.
   dot
                     22. sameLine
9.
   norm_sq
                     23. lineIntersect1
10. len
                     24. lineIntersect2
11. dot2
                     25. cross3
                     26. lineIntersect3
12. cross
13. onLine
                     27. angle
14. onLineL
                     28. s perm
5.3 Point
1. Distance of 2 points
double dist(point p1, point p2) {
        return hypot(p1.x - p2.x, p1.y - p2.y); }
2. Degree to radian
double DEGtoRAD(double theta){
        return (theta * ( PI / 180.0)); }
3. Rotate p by theta degrees CCW w.r.t origin (0, 0)
point ROTATE(point p, double theta) {
        double rad = DEGtoRAD(theta);
        return point( p.x * cos(rad) - p.y * sin(rad),
              p.x * sin(rad) + p.y * cos(rad));
}
4. Rotate a point angle radian about center, CCW
point rotatep(point a , point center , double angle ){
       point p;
       double x1 = a.x - center.x;
       double y1 = a.y - center.y;
       double x2 = x1 * cos( angle) - y1 * sin( angle);
       double y2 = x1 * sin(angle) + y1 * cos(angle);
       p.x = x2 + center.x;
       p.y = y2 + center.y;
       return p; }
5.4 Vectors
5. vec toVec(point a , point b ) { return vec( b.x - a.x, b.y -
a.y); }
```

```
6. vec scale(vec v, double s) { return vec(v.x * s, v.y * s); }
7. point translate(point p, vec v) { return point(p.x+v.x,
p.y+v.y); }
8. double dot(vec a, vec b) { return (a.x * b.x + a.y * b.y); }
9. double norm_sq(vec v) { return v.x * v.x + v.y * v.y; }
10. double len(point a ) { return sqrt(a.x*a.x + a.y*a.y); }
11. double dot2(point a , point b ) { return a.x*b.x + a.y*b.y;
}
12. double cross(point a , point b ) { return a.x * b.y - a.y *
b.x; }
5.5 Point and Line
13. Is distance ab = ar + rb ?
bool onLine(point a , point b , point r ){
        if (fabs(dist( a, b) - (dist( a, r)
               + dist(b, r))) < EPS) return true;
        return false;
14. Is point b on line x?
bool onLineL(line x , point b ){
     if ( fabs(( x.a * b.x + x.b * b.y) + x.c) < EPS)
          return true;
     return false;
15. Find a point p which is on line ab and ap = d
point pointAtDist(point a , point b , double d ){
     point p;
     double d = dist( a, b);
     p.x = (a.x + (b.x - a.x)*(d/d));
     p.x = (a.y + (b.y - a.y)*(d/d));
     return p;
16. Distance of AB to C (When point on AB is not needed)
double linePointDist1
     (point A, point B, point C, bool isSegment) {
     double dis = cross(B - A, C - A) / len(B - A);
     if (isSegment) {
          if (\underline{dot2}(C - B, B - A) > EPS)
```

```
return len(B - C);
          if (dot2(C - A, A - B) > EPS)
               return len(A - C);
     return fabs(dis);
}
17. Distance ab to p. Closest point saved in c
double linePointDist2(point p, point a, point b, point &c) {
     vec ap = toVec(a, p), ab = toVec(a, b);
     double u = dot(ap, ab) / norm_sq(ab);
     c = translate(a, scale(ab, u));
     return dist(p, c);
18. Distance of p to ab(Segment). Closest point saved in c
double linePointDist3
     (point p, point a, point b, point &c) {
     vec ap = toVec(a, p), ab = toVec(a, b);
     double u = dot(ap, ab) / norm_sq(ab);
     if (u < 0.0) {
          c = point(a.x, a.y);
          return dist(p, a);
     if (u > 1.0) {
          c = point(b.x, b.y);
          return dist(p, b); }
     return distToLine(p, a, b, c);
}
19.
double linePointDist4(line 1, point p) {
     return fabs(p.x * 1.a + p.y * 1.b - 1.c) /
          sqrt(1.a* 1.a + 1.b* 1.b);
}
5.7 Lines
20.
line toLine(point p1 , point p2 ) {
       line x;
       x.a = p1.y - p2.y; x.b = p2.x - p1.x;
       x.c = -(x.a * p1.x + x.b* p1.y);
       return x;
21. 11 parallel of 12?
bool parallel(line 11, line 12) {
     return (fabs(11.a * 12.b - 12.a * 11.b) < EPS); }</pre>
22. line l1 == l2?
```

```
bool sameLine(line 11, line 12) {
     return (fabs(l1.a * l2.b - l1.b * l2.a) < EPS &&
          fabs(l1.b * 12.c - 11.c * 12.b) < EPS); }
23. Intersect of 2 lines(return false if there is no intersect)
bool lineIntersect1 (line 11 , line 12 , point &p ) {
        double a = 11.a * 12.b - 12.a * 11.b;
        if (fabs(a) < EPS) return false;</pre>
        double b = 11.b * 12.c - 12.b * 11.c;
        double c = 11.c * 12.a - 12.c * 11.a;
        p.x = b / a; p.y = c / a; return true;
24. Intersect of 2 line segments (return false if there is no
intersect)
bool lineIntersect2(point a, point b, point x, point y, point
&m) {
     line 11 = toLine(a, b), 12 = toLine(x, y);
     if (sameLine(11, 12)){
     //Check what's said about same lines
          if (onLine(a, b, x) || onLine(a, b, y) ||
               onLine(x, y, a) | | onLine(x, y, b))
                    return true;
          return false;
     if ( lineIntersect1 (l1, l2, m)){
          if (onLine(a, b, m) && onLine(x, y, m))
               return true;
     return false;
}
25.
double cross3(point p1, point p2, point p){
     return (p2.x - p1.x) * (p.y - p1.y) - (p2.y - p1.y) * (p.x
     - p1.x);
}
26. Checks if 2 line segment intersect, returns false if
segments touch
bool lineIntersect3(point p1, point p2, point q1, point q2){
     return cross3(p1, p2, q1) * cross3(p1, p2, q2) < 0 &&
          cross3(q1, q2, p1) * cross3(q1, q2, p2) < 0;
}
```

```
5.8 Triangle
27.
double angle(point a, point o, point b) {
     double ang = (fabs(atan2(a.y - o.y, a.x - o.x) -
           atan2(b.y - o.y, b.x - o.x)) * 180) / PI;
     //if (ang > 180.0)return(360.0 - ang);
     return ang; }
28. Semi perimeter
double s_perm(double a, double b, double c){
     return (a + b + c) / 2; }
29.
double tarea1(double h, double b){
     return h*b*0.5; }
30. Given length of all sides
double tarea2(double ab, double bc, double ca){
     double s = (ab + bc + ca) / 2;;
     return sqrt(s*(s-ab)*(s-bc)*(s-ca));
}
31.
double tarea3(point a , point b , point c ) {
      return fabs(.5 * ( a.x * ( b.y - c.y) + b.x * ( c.y -
      a.y) + c.x *(a.y - b.y)); }
Any triangle's area T can be expressed in terms of its
medians m_a, m_b, and m_c as follows. Denoting their semi-
sum (ma + mb + mc)/2 as \sigma, we have
T = \frac{4}{3}\sqrt{\sigma(\sigma - m_a)(\sigma - m_b)(\sigma - m_c)}.
5.9 Circle:
(x - h)^2 + (y - k)^2 = r^2
X^2 + y^2 + 2gx + 2fy + c = 0
h=-g, k=-f, c = h^2 + k^2 - r^2
h & k are the centre of the circle
Given 3 points(x1,y1; x2,y2; x3,y3)find circle's radius/centre:
double x1,x2,x3,y1, y2, y3, a1, a2, b1, b2, c1, c2, d, h, k, r;
a1 = 2*(x2-x1); a2 = 2*(x3-x2);
b1 = 2*(y2-y1); b2 = 2*(y3-y2);
c1 = x1*x1+y1*y1-x2*x2-y2*y2;
c2 = x2*x2+y2*y2-x3*x3-y3*y3;
d = a1*b2-a2*b1;
```

```
h=(b1*c2-b2*c1)/d;
k=(a2*c1-a1*c2)/d;
r=sqrt((h-x1)*(h-x1)+(k-y1)*(k-y1));
32. Radius of Incircle = Area / semi-perimeter
double rInCircle(double ab, double bc, double ca) {
             return tarea2(ab, bc, ca) / s perm(ab, bc, ca);
}
33. double rInCircleP(point a, point b, point c) {
             return rInCircle(dist(a, b), dist(b, c), dist(c, a));
}
34. Radius of circumscribed circle = R = a \times b \times c/(4 \times A).
double rCircumCircle(double ab, double bc, double ca) {
             return ab * bc * ca / (4.0 * tarea2(ab, bc, ca));
35. double rCircumCircleP(point a, point b, point c) {
             return rCircumCircle(dist(a, b), dist(b, c),
                          dist(c, a));
}
5.10 Monotone Chain Convex Hull + Points inside hull check
* Returns a list of convex hull in counter-clockwise order
* In case of collinear points, only end-points are taken
* The last point is the same as the first one
/// OK chkd by SunW
#define 1164 long long
struct point {
             1164 x, y;
             bool operator <(const point &p) const {</pre>
                          return (x < p.x | (x == p.x && y < p.y));
             }
};
vector<point> CH, P;
1164 cross(point p1, point p2, point p) {
             return (p2.x - p1.x) * (p.y - p1.y) - (p2.y - p1.y) * (p.x - p1.y) + (p.x - p1.
p1.x);
void convexHull() {
             sort(P.begin(), P.end());
             int n = P.size(), k = 0;
             vector<point> H(2 * n);
             for (int i = 0; i < n; i++) {
                          while (k \ge 2 \&\& cross(H[k - 2], H[k - 1], P[i]) <= 0)
                                       k--;
                          H[k++] = P[i];
             }
```

```
for (int i = n - 2, t = k + 1; i >= 0; i--) {
           while (k >= t && cross(H[k - 2], H[k - 1], P[i]) <= 0)
                k--;
           H[k++] = P[i];
     }
     H.resize(k);
     CH = H;
1164 ar(point a, point b, point c) {
     return abs(a.x * (b.y - c.y) + b.x * (c.y - a.y) + c.x * (a.y)
- b.y));
}
bool insideCH(point p) {
     point pi = CH[0];
     int low = 1, up = CH.size() - 2, mid;
     while (up - low > 1) {
           mid = (low + up) >> 1;
           1164 t = cross(pi, CH[mid], p);
           if (t < 0)
                up = mid;
           else
                low = mid;
     if (ar(pi, CH[low], CH[up])== (ar(pi, p, CH[low]) + ar(pi, p,
           CH[up]) + ar(CH[low], p, CH[up]))) return true;
     return false;
}
int main() {
     //rd();
     int x, y, 1, s;
     while (scanf("%d", &l) == 1) {
           P.clear();
           CH.clear();
           point pp;
           for (int i = 0; i < 1; i++) {
                scanf("%d %d", &x, &y);
                pp.x = x;
                pp.y = y;
                P.push_back(pp);
           }
           convexHull();
           int ans = 0;
           scanf("%d", &s);
           for (int i = 0; i < s; i++) {
                scanf("%d %d", &x, &y);
                pp.x = x;
                pp.y = y;
```

```
if (insideCH(pp)) {
                      ans++;
                 }
           }
           printf("%d\n", ans);
     }
}
6. Graph
6.1 Topological sort
/// Topological Sort with cycle check: (inside main)
// if ans.size()!= n TOPOSORT FAILS
//Memory efficient
int col[MAX], fl;
vector<int> G[MAX], ans;
void toposort(int u) {
     col[u] = 1;
     int i, y;
     //col, 0 = non visited, 1 = not finished, 2 = finished
     for (i = 0; i < G[u].size(); i++) {</pre>
           y = G[u][i];
           if (col[y] == 0)
                toposort(y);
           else if (col[y] == 1) { //cycle
                fl = 1;
                return;
           }
     if (fl == 1)
           return;
     col[u] = 2;
     ans.push_back(u);
// ans contains elements in reversed topo order
}
/// All possible Topo sort:
int vals[MAX], ANS[MAX], N; //vals = input
bool taken[MAX];
vector<int> GE[MAX];
void toposort(int u, int in[]) {
     if (u == N) {
     }
// Possible sort found print, return
     int i, j, x;
     vector<int> v;
     int in[MAX];
     for (i = 0; i < MAX; i++)
           _in[i] = in[i];
```

```
for (i = 0; i < N; i++) {
           x = vals[i];
           if (in[x] == 0 && !taken[x])
                 v.push_back(x);
     for (i = 0; i < v.size(); i++) {</pre>
           x = v[i];
           taken[x] = true;
           ANS[u] = x;
           for (j = 0; j < GE[x].size(); j++) {</pre>
                 if (_in[GE[x][j]] > 0)
                       _in[GE[x][j]]--;
           toposort(u + 1, _in);
           taken[v[i]] = false;
           for (j = 0; j \leftarrow N; j++)
                 _in[j] = in[j];
     }
int main() {
     int n = 5, x, y;
     priority_queue<int, vector<int>, greater<int> > q;
     int *indeg = new int[n + 2] { };
     for (int i = 1; i <= n; i++)</pre>
           if (indeg[i] == 0)
                 q.push(i);
     if (q.empty()) {
                        // TopoSort Fail
     else {
           while (!q.empty()) {
                 x = q.top();
                 q.pop();
                 ans.push_back(x);
                 for (int i = 0; i < G[x].size(); i++) {
                       y = G[x][i];
                       if (indeg[y] == 0)
                             break;
                       //cycle found ... break with flag
                       else if (indeg[y] == 1) {
                             indeg[y]--;
                             q.push(y);
                       } else
                             indeg[y]--;
                 }
           }
     }
}
```

## 6.2 Bipartite Graph

```
break because all the vertices in that graph needs to be set as
visited.
** Easier problems don't need number of colors
int colour[MAX]; bool visited[MAX]; vector <int> G[MAX];
bool bipartite(int &no, int u) {
     queue<int> q:
     q.push(u); colour[u] = 0;
     int x, i, tot = 0, zr = 1;
     bool isBip = true;
     visited[u] = true;
     while (!q.empty() /* && isBip */) {
           tot++;
           x = q.front(); q.pop();
           for (i = 0; i<G[x].size(); i++) {</pre>
                int y = G[x][i];
                if (!visited[y]) {
                      colour[y] = 1 - colour[x]; q.push(y);
            visited[y] = true;
                      if (colour[y] == 0)zr++;
                else if (colour[y] == colour[x]) { isBip = false; }
           }
if (!isBip) { no = 0; return false; }
     if (tot <= 2) { no = 1; return true; }</pre>
     else { no = min(zr, tot - zr); } return true;
// min or max
}
6.3 Articulation point
/// OK
int Low[MAX], num[MAX], parent[MAX], root, rootc, counter, n;
bool visited[MAX], artiVer[MAX];
vector<int> G[MAX];
void artculPoint(int u) {
     low[u] = num[u] = counter++;
     for (int i = 0; i < G[u].size(); i++) {</pre>
           int j = G[u][i];
           if (!visited[j]) {
                parent[j] = u;
                visited[j] = true;
                if (u == root)
                      rootc++;
                artculPoint(j);
                if (low[j] >= num[u])
                      if (!artiVer[u] && u != root)
                            artiVer[u] = true;
```

\*\* Sometimes even if isBip is false, the while loop shouldn't

```
//if (low[j] > num[u]) { Bridge
                      int k = u, l = j;
                //
                      if (k>1)swap(k, 1);
                //
                      edges.push_back(pii(k, 1));
                //
                //}
                Low[u] = min(Low[u], Low[j]);
           } else if (j != parent[u])
                Low[u] = min(Low[u], num[j]);
     }
** in main after calling articulPoint
if (rootc >= 2) artiVer[root] = true;
6.4 Strongly Connected Components (Tarjan's):
///NOT OK
vector<int> G[MAX], output[MAX];
int stck[MAX], num[MAX], component[MAX], low[MAX], onstack[MAX],
int idx, components, N;
void tarjan(int u) {
     num[u] = low[u] = idx++;
     stack[top++] = u;
     onstack[u] = 1;
     for (int i = 0; i < G[u].size(); i++) {</pre>
           int v = G[u][i];
           if (num[v] == -1) {
                tarjan(v);
                Low[u] = min(Low[u], Low[v]);
           } else if (onstack[v])
                low[u] = min(low[u], num[v]);
     }
     int v;
     if (Low[u] == num[u]) {
           do {
                v = stack[--top];
                onstack[v] = 0;
                component[v] = components;
                output[components].push_back(v);
// Storing SCCs in array of vector
           } while (u != v);
           components++;
     }
}
6.5 Prim's MST
struct node {
     int n, e;
     node() {
           n = 0;
           e = 0;
     }
```

```
node(int _a, int _b) :
                 n(_a), e(_b) {
     bool operator<(const node &B) const {</pre>
           return e > B.e;
};
vector<node> G[MAX];
bool visited[MAX];
int parent[MAX], C[MAX], root;
long long mst() {
     C[root] = 0;
     priority_queue<node> Q;
     Q.push(node(root, 0));
     long long cost = 0;
     while (!Q.empty()) {
           node u = Q.top();
           Q.pop();
           if (!visited[u.n]) {
                 cost += u.e;
                 visited[u.n] = true;
                 for (int i = 0; i < G[u.n].size(); i++) {</pre>
                       node v = G[u.n][i];
                       if (!visited[v.n] && C[v.n] > v.e) {
                             Q.push(v);
                             C[v.n] = v.e;
                             parent[v.n] = u.n;
                       }
                 }
           }
     return cost;
6.6 Kruskal's MST
struct edl {
     int a, b;
     long long w;
     bool operator <(const edl &a) const {</pre>
           return w < a.w;</pre>
} edlist[MAX * MAX];// Edgelist
int P[MAX];
int R[MAX], NE, NV; // R[]=Rank
void makeset(int n) {
     for (int i = 1; i <= n; i++) {</pre>
           P[i] = i;
      }
}
```

```
int findP(int n) {
     if (P[n] == n)
           return n;
     P[n] = findP(P[n]);
     return P[n];
void Union(int a, int b) {
     int x = findP(a);
     int y = findP(b);
     if (x == y)
           return;
     else {
           P[x] = y;
}
long long kruskal( /* bool flg */) {
//Should kruskal add MST edges to mark[]?
     makeset(NE);
     sort(edlist, edlist + NE);
     int i, j, k = 0;
     long long cost = 0;
     edl top;
     for (i = 0; i < NE; i++) {
           top = edlist[i];
           if (findP(top.a) != findP(top.b)) {
                 cost += top.w;
                 k++;
                 Union(top.a, top.b);
// if (flg)mark.push_back(i);
//Add MST edges to vector mark
                 if (k == NV - 1)
                       return cost;
           }
// return -1;
2<sup>nd</sup> best MST(Kruskal, UVa 10600):
///
struct edl {
     int a, b;
     long long w;
     bool operator <(const edl &a) const {</pre>
           return w < a.w;</pre>
} edlist[MAX * MAX], edlist2[MAX * MAX];
int P[MAX];
int R[MAX], NE, NV; // R[]=Rank
```

```
void makeset(int n) {
     for (int i = 1; i <= n; i++) {</pre>
           P[i] = i;
     }
}
int findP(int n) {
     if (P[n] == n)
           return n;
     P[n] = findP(P[n]);
     return P[n];
void Union(int a, int b) {
     int x = findP(a);
     int y = findP(b);
     if (x == y)
           return;
     else {
           P[x] = y;
     }
}
long long kruskal(bool flg) {
//Should kruskal add MST edges to mark[]?
     makeset(NE);
     sort(edlist, edlist + NE);
     int i, j, k = 0;
     long long cost = 0;
     edl top;
     for (i = 0; i < NE; i++) {
           top = edlist[i];
           if (findP(top.a) != findP(top.b)) {
                 cost += top.w;
                 k++;
                 Union(top.a, top.b);
// if (flg)mark.push back(i);
//Add MST edges to vector mark
                 if (k == NV - 1)
                      return cost;
           }
// return -1;
vector<int> mark;
int main() {
     int n, i, j, k;
     cin >> n;
     while (cin >> NV >> NE) {
           edl a;
           k = 0;
           for (i = 0; i < NE; i++) {
```

```
cin >> a.a >> a.b >> a.w;
                 edlist[k++] = a;
           cout << kruskal(true) << " ";</pre>
           int mi = 9999999;
           for (i = 0; i < NE; i++)
                 edlist2[i] = edlist[i];
           for (i = 0; i < mark.size(); i++) {</pre>
                 edlist[mark[i]].w = 100000;
                 j = kruskal(false);
                 if (j < mi)
                      mi = j;
                 for (j = 0; j < NE; j++)
                      edlist[j] = edlist2[j];
           }
           cout << mi << "\n"; //Cost of 2nd best MST</pre>
           mark.clear();
6.7 Dijkstra - All shortest paths
#define MEM(p, v) memset(p, v, sizeof(p))
struct Node {
     int node, cost, par;
     Node() {
           node = 0;
           cost = 0;
           par = 0;
     Node(int _a, int _b, int _c) :
                 node(_a), cost(_b), par(_c) {
     }
     bool operator<(const Node &A) const {</pre>
           return cost > A.cost;
     }
};
int dist[MAX], visited[MAX], inf;
vector<Node> G[MAX];
vector<int> P[MAX];
vector<pair<int, int> > spedges; //shortest path edges
void btrack(int n) {
     if (P[n].size() == 0 | visited[n])
           return;
     visited[n] = 1;
     for (int i = 0; i < P[n].size(); i++) {</pre>
           int pr = P[n][i];
           spedges.push back(make_pair(n, pr));
           btrack(pr);}}
```

```
void dijkstra(int s, int d) {
     MEM(visited, 0);
     MEM(dist, 127);
     dist[s] = 0;
     Node nd(s, 0, -1);
     priority queue<Node> Q;
     Q.push(nd);
     while (!Q.empty()) {
           int cn = Q.top().node;
           int dst = dist[cn];
           Node tp = Q.top();
           Q.pop();
           if (dst == tp.cost && cn != s) {
                 P[cn].push_back(tp.par);
           if (visited[cn])
                 continue;
           for (int i = 0; i < G[cn].size(); i++) {</pre>
                 int n = G[cn][i].node;
                 int c = G[cn][i].cost;
                 if (dst + c <= dist[n]) {
                       dist[n] = dst + c;
                       Q.push(Node(n, dist[n], cn));
                 }
           visited[cn] = 1;
     }
}
int main() {
     cpp_io();
     int n, e, a, b, c, s, d;
     cin >> n >> e;
     for (int i = 0; i<e; i++) {</pre>
           cin >> a >> b >> c;
           G[a].push_back(Node(b, c, -1));
           G[b].push back(Node(a, c, -1));
     cin >> s >> d;
     dijkstra(s, d);
     memset(visited, 0, sizeof(visited));
     btrack(d);
     for (int i = 0; i<spedges.size(); i++)</pre>
     cout << spedges[i].first << " " << spedges[i].second <<endl;</pre>
}
```

```
6.8 Bellman Ford
struct node {
     int n, e;
     node() \{n = 0; e = 0; \}
     node(int _a, int _b) :n(_a), e(_b) {}
};
int dist[MAX], NV;
vector<node> G[MAX];
void init(int n) {
     memset(dist, 127, sizeof(dist));
     for (int i = 0; i <= n; i++)</pre>
           G[i].clear();
bool bellman_ford(int s) {
     dist[s] = 0;
     for (int k = 0; k < NV; k++) {
           for (int i = 0; i < NV; i++) {
                for (int j = 0; j < G[i].size(); j++) {
                      int n = G[i][j].n;
                      int e = G[i][j].e;
                      if (dist[i] + e < dist[n]) {
                            dist[n] = dist[i] + e;
                            if (k == NV - 1)
                                 return false;
                      }}}
     return true;
}
6.9 Dinic's Algorithm
const int MAXN = 1000;
struct edge {
     int a, b, cap, flow;
};
int n, s, t, d[MAXN], ptr[MAXN], q[MAXN];
vector<edge> e;
vector<int> g[MAXN];
                     // Residual Graph
void add_edge(int a, int b, int cap) {
     edge e1 = { a, b, cap, 0 };
     edge e2 = { b, a, 0, 0 };
     g[a].push_back((int) e.size());
     e.push_back(e1);
     g[b].push_back((int) e.size());
     e.push back(e2);
                                       }
```

```
bool bfs() {
     int qh = 0, qt = 0;
     q[qt++] = s;
     memset(d, -1, sizeof(d));
     d[s] = 0;
     while (qh < qt && d[t] == -1) {
           int v = q[qh++];
           for (size_t i = 0; i < g[v].size(); ++i) {</pre>
                 int id = g[v][i], to = e[id].b;
                 if (d[to] == -1 && e[id].flow < e[id].cap) {
                      q[qt++] = to;
                      d[to] = d[v] + 1;
                 }
           }
     return d[t] != -1;
}
int dfs(int v, int flow) {
     if (!flow)
           return 0;
     if (v == t)
           return flow;
     for (; ptr[v] < (int) g[v].size(); ++ptr[v]) {</pre>
           int id = g[v][ptr[v]], to = e[id].b;
           if (d[to] != d[v] + 1)
                 continue;
           int pushed = dfs(to, min(flow, e[id].cap - e[id].flow));
           if (pushed) {
                 e[id].flow += pushed;
                 e[id ^ 1].flow -= pushed;
                 return pushed;
           }
     }
     return 0;
int dinic() {
     int flow = 0;
     for (;;) {
           if (!bfs())
                 break;
           memset(ptr, 0, sizeof(ptr));
           while (int pushed = dfs(s, INF))
                 flow += pushed;
     return flow;
}
```

## 6.10 Bipartite Matching

```
///
//Graph G strats from index 0
vector<int> G[MAX];
bool visited[MAX];
int Left[MAX], Right[MAX];
bool dfs(int u) {
     if (visited[u])
           return false;
     visited[u] = true;
     int len = G[u].size(), i, v;
     for (i = 0; i < len; i++) {
           v = G[u][i];
           if (Right[v] == -1) {
                Right[v] = u, Left[u] = v;
                return true;
           }
     for (i = 0; i < len; i++) {</pre>
           v = G[u][i];
           if (dfs(Right[v])) {
                Right[v] = u, Left[u] = v;
                return true;
           }
     return false;
                       // n = size of G
int match(int n) {
     int i, ret = 0;
     bool done;
     memset(Left, -1, sizeof Left);
     memset(Right, -1, sizeof Right);
     do {
           done = true;
           memset(visited, 0, sizeof visited);
           for (i = 0; i < n; i++) {
                if (Left[i] == -1 && dfs(i)) {
                      done = false;
                 }
     } while (!done);
     for (i = 0; i < n; i++)
           ret += (Left[i] != -1);
     return ret;}
```

```
7. String
/// OK chkd by SunW
7.1 Hasing
http://e-maxx-eng.appspot.com/string/string-hashing.html
typedef unsigned long long ull;
const int P = 31, SIZE = 100000;
// P = 53 if for both uppercase & lowercase
vector <ull> p pow(SIZE);
vector <ull> h(SIZE);
                           //hash of prefixes
void init_pow() {
     p_pow[0] = 1;
     for (int i = 1; i < SIZE; i++)</pre>
          p_pow[i] = p_pow[i - 1] * P;
void hash_pre(string s) { //hash of prefixes
     for (int i = 0; i < s.length(); i++){</pre>
          h[i] = (s[i] - 'a' + 1) * p_pow[i];
          if (i) h[i] += h[i - 1];
     }
}
ull hash_s(string s) { //hash of a string
     ull h s = 0;
     for (int i = 0; i < s.length(); i++)</pre>
          h s += (s[i] - 'a' + 1) * p pow[i];
     return h_s;
bool comp_substr(int 11, int 12, int r1, int r2) {
     //compares hashes of 2 substring in a string(0 based)
     ull h1 = h[r1];
     if (11) h1 -= h[11 - 1];
     ull h2 = h[r2];
     if (12) h2 -= h[12 - 1];
     // Get the two hashes multiplied by the same power of P
     // and then compare them
     if (11 <12 && h1 * p_pow[12 - 11] == h2 ||
          11> 12 && h1 == h2 * p_pow[11 - 12])
          return true; // substrings are equal
     else
          return false;
}
```

```
int count_substr(int n) {
     // Count number of differetn substrings in a string
     // n = length of string
     int result = 0;
     for (int l = 1; l <= n; l++) {
          // Need to find the number of distinct substrings of
          // length 1. Get the hashes for all substrs of length
     1
          vector \langle ull \rangle hs(n - 1 + 1);
          for (int i = 0; i < n - 1 + 1; i++){
               ull cur h = h[i + l - 1];
               if (i) cur_h -= h[i - 1];
               cur h *= p pow[n - i - 1];
               hs[i] = cur_h;
          }
          sort(hs.begin(), hs.end());
          hs.erase(unique(hs.begin(), hs.end()), hs.end());
          result += (int)hs.size();
     return result;
}
7.2 Rabin - Karp
/// OK chkd by SunW
vector<int> rabin_karp (string s, string t) {
     // s = pattern, t = text
     vector<int> res; //returns 0 based starting indexs of
matches
     init pow();
     ull h_s = hash_s(s);
     hash pre(t);
     for (int i = 0; i + s.length() - 1 < t.length(); i++){</pre>
          ull cur h = h[i + s.length() - 1];
          if (i) cur h -= h[i - 1];
          if (cur_h == h_s * p_pow[i])
               res.push back(i);
     return res;
}
```