Playfair cipher

The Playfair cipher is a digraph substitution cipher that encrypts pairs of letters using a 5×5 key matrix built from a keyword. Here's how you create the key matrix:

1. Choose a keyword or key phrase.

Example: MONARCHY

2. Remove duplicate letters.

From MONARCHY, we get MONARCHY. (No duplicates here.)

- 3. Fill the matrix row by row with the keyword letters. (Combine I and J into one cell to fit 25 letters.)
- 4. Fill the remaining spaces with unused letters of the alphabet.

So, for MONARCHY, the matrix becomes:

```
        M
        O
        N
        A
        R

        C
        H
        Y
        B
        D

        E
        F
        G
        I/J
        K

        L
        P
        Q
        S
        T

        U
        V
        W
        X
        Z
```

Example 1

Let's build the Playfair cipher key matrix step by step using the keyword "SIX".

Step 1: Write the keyword (remove duplicates)

Keyword: SIX

No duplicates \rightarrow we keep it as SIX.

Step 2: Prepare the alphabet

Playfair uses **25 letters**, combining **I** and **J** into a single cell.

Alphabet (I/J counted as one):

A B C D E F G H I/J K L M N O P Q R S T U V W X Y Z

Step 3: Fill the matrix

Start with SIX, then add remaining letters (excluding duplicates and skipping J since I/J share one cell):

- Start: S I X
- Remaining letters: A B C D E F G H K L M N O P Q R T U V W Y Z

S I/J X A B

CDEFG

H K L M N

OPQRT

UVWYZ

let's encrypt **HELLO** step by step using the Playfair matrix we just built for **SIX**:

Step 1: Write the Key Matrix Again

Step 2: Prepare the Plaintext

Plaintext: **HELLO**

- Break into pairs of letters: HE | LL | O
- If a pair has the same letter (LL), insert an **X** between them:
 - HE | LX | LO
- If there's a single letter left at the end, add an **X** (not needed here because we now have three pairs).

So our digraphs are: **HE**, **LX**, **LO**.

Step 3: Encrypt Each Pair

Pair 1: H – E

- H = row 3, column 1
- E = row 2, column 3
- Rectangle rule: Take letters in the same row as the other letter → form the corners.

H (row 3, col 1) \rightarrow same row as E's column \rightarrow L E (row 2, col 3) \rightarrow same row as H's column \rightarrow C

 $HE \rightarrow LC$

Pair 2: L - X

- L = row 3, column 3
- X = row 1, column 3
- Same column rule: Replace each with the letter below it (wrap to top if needed).

 $L \rightarrow M$ (one row down)

 $X \rightarrow Q$ (one row down from row 1 to row 4)

 $LX \rightarrow MQ$

Pair 3: L - O

- L = row 3, column 3
- O = row 4, column 1
- Rectangle rule again.

L (row 3, col 3) \rightarrow take O's column (col 1) \rightarrow **H** O (row 4, col 1) \rightarrow take L's column (col 3) \rightarrow **Q**

 $\text{LO} \rightarrow \text{HQ}$

Step 4: Combine Results

Encrypted text = LCMQHQ

Example 2:

Step 1: Keyword

Keyword: **BALLOON**

Remove duplicates → BALON

(We keep first occurrences only: B, A, L, O, N)

Step 2: Prepare Alphabet (I/J combined)

Alphabet (25 letters total):

A B C D E F G H I/J K L M N O P Q R S T U V W X Y Z

Step 3: Build the 5×5 Matrix

3 4 5

U

Start with **BALON**, then fill remaining letters (excluding duplicates and skipping J):

Start: BALON

Remaining: C D E F G H I/J K M P Q R S T U V W X Y Z

Resulting matrix:

1

 1
 B
 A
 L
 O
 N

 2
 C
 D
 E
 F
 G

 3
 H
 I/J
 K
 M
 P

2

5 V W X Y Z

4 Q R S T

Using the Playfair matrix for **BALLOON** (rows: B A L 0 N/C D E F G/H I/J K M P/Q R S T U/V W X Y Z), I split **HELLO** into digraphs **HE | LX | LO** (inserting X between the double Ls). Encrypting: **HE** \rightarrow **KC** (rectangle rule), **LX** \rightarrow **EL** (same column, move down), **LO** \rightarrow **ON** (same row, move right). So the final ciphertext is **KCELON**.

Ciphertext to decrypt: K C E L O N \rightarrow split into digraphs KC | EL | ON

Decryption rules (short):

- Same row: shift each letter one cell left (wrap around).
- Same column: shift each letter one cell up (wrap around).
- Rectangle (different row & column): each letter takes the column of the other letter while staying in its own row.

Step-by-step:

- 1. **KC**
 - K at (3,3); C at (2,1). They form a rectangle.
 - Take letter in row of K and col of C \rightarrow (3,1) = **H**.
 - Take letter in row of C and col of $K \rightarrow (2,3) = E$.
 - \circ KC \rightarrow HE
- 2. **EL**
- o E at (2,3); L at (1,3). Same **column** (col 3).
- Shift each letter **up** one row: $E \rightarrow (1,3) = L$; $L \rightarrow (5,3) = X$ (wrap from row1 to row5).
- \circ EL \rightarrow LX
- 3. **ON**
 - o O at (1,4); N at (1,5). Same **row** (row 1).
 - Shift each letter **left** one column: $O \rightarrow (1,3) = L$; $N \rightarrow (1,4) = O$.
 - $\circ \quad \mathsf{ON} \to \mathsf{LO}$

Combine decrypted digraphs: **HE | LX | LO = HELXLO**. Remove the filler X that was inserted between the double Ls during encryption → **HELLO**.

Example 3:

Example Word: TAXI

Step 1: Break into digraphs

TAXI \rightarrow TA | XI

(Remember: I/J share a cell, so I is valid.)

Step 2: Encrypt Each Pair

Pair 1: TA

- $T = (4,4), A = (1,2) \rightarrow rectangle rule.$
- T \rightarrow take column of A (col 2) \rightarrow (4,2) = R
- A \rightarrow take column of T (col 4) \rightarrow (1,4) = **O** TA \rightarrow RO

Pair 2: XI

- X = (5,3), $I = (3,2) \rightarrow$ rectangle rule.
- $X \rightarrow \text{take column of I (col 2)} \rightarrow (5,2) = W$
- I → take column of X (col 3) → (3,3) = K
 XI → WK

Final Ciphertext

Combine results: RO | WK = ROWK

Step 3: Decryption (to show X handling works)

Now decrypt **ROWK** back to plaintext:

- 1. **RO** \rightarrow R(4,2), O(1,4) \rightarrow rectangle rule \rightarrow (4,4)=T, (1,2)=A \rightarrow **TA**
- 2. **WK** \rightarrow W(5,2), K(3,3) \rightarrow rectangle rule \rightarrow (5,3)=X, (3,2)=I \rightarrow **XI**

Result: TAXI

Hill Cipher

The Hill Cipher is a **polygraphic substitution cipher** that works on blocks of letters (usually 2 or 3 at a time). It converts letters into numbers, uses a **key matrix** to transform them, and converts the result back to letters.

Steps for Encryption

Let's say we use a **2×2 Hill Cipher** with key matrix

K=[abcd]K=[acbd]

and plaintext "HI".

Step 1: Convert letters to numbers

Use A=0, B=1, ..., Z=25.

- $\bullet \quad H \to 7$
- I → 8

Plaintext vector:

P=[78]P=[78]

Step 2: Multiply by key matrix (mod 26)

C=K×Pmod 26C=K×Pmod26

This gives the ciphertext vector CC.

Step 3: Convert back to letters

Each number becomes a letter $(0 \rightarrow A, ..., 25 \rightarrow Z)$.

Suppose our key matrix is

$$K = egin{bmatrix} 3 & 3 \ 2 & 5 \end{bmatrix}$$

Encrypting "HI" (P = [7, 8]):

$$C = egin{bmatrix} 3 & 3 \ 2 & 5 \end{bmatrix} egin{bmatrix} 7 \ 8 \end{bmatrix} = egin{bmatrix} 3 imes 7 + 3 imes 8 \ 2 imes 7 + 5 imes 8 \end{bmatrix} = egin{bmatrix} 45 \ 54 \end{bmatrix} \mod 26 = egin{bmatrix} 19 \ 2 \end{bmatrix}$$

 $19 \rightarrow T$, $2 \rightarrow C \rightarrow Ciphertext =$ **TC**

Decryption

Decryption uses the inverse matrix of K (mod 26).

You compute inverse K mod 26, multiply by ciphertext vector, then convert back to letters.

Calculating the inverse key matrix

Example 1 — 2×2 Hill Cipher

Key matrix:

$$K = egin{bmatrix} 3 & 3 \ 2 & 5 \end{bmatrix}$$

Step 1 — Compute the determinant

$$\det(K) = (3 * 5 - 3 * 2) = 15 - 6 = 9$$

Check $gcd(9,26)=1 \rightarrow \bigvee$ invertible.

Step 2 — Find modular inverse of determinant

Find x such that:

$$9 \cdot x \equiv 1 \pmod{26}$$

•
$$9*3 = 27 \equiv 1 \pmod{26} \to det^{-1} = 3$$

Step 3 — Compute adjugate matrix

For 2×2:

$$\operatorname{adj}(K) = \begin{bmatrix} d & -b \\ -c & a \end{bmatrix} = \begin{bmatrix} 5 & -3 \\ -2 & 3 \end{bmatrix}$$

Convert negative numbers modulo 26:

$$\operatorname{adj}(K) \mod 26 = egin{bmatrix} 5 & 23 \\ 24 & 3 \end{bmatrix}$$

Step 4 — Multiply adjugate by det inverse

$$K^{-1} = det^{-1} \cdot adj(K) \mod 26 = 3 \cdot egin{bmatrix} 5 & 23 \ 24 & 3 \end{bmatrix} \mod 26$$

Multiply each entry by 3:

$$\begin{bmatrix} 15 & 69 \\ 72 & 9 \end{bmatrix} \mod 26 = \begin{bmatrix} 15 & 17 \\ 20 & 9 \end{bmatrix}$$

2×2 inverse matrix:

$$K^{-1}=egin{bmatrix} 15 & 17 \ 20 & 9 \end{bmatrix}$$

3×3 Hill Cipher Inverse

Key matrix:

$$K = \begin{bmatrix} 2 & 4 & 5 \\ 9 & 2 & 1 \\ 3 & 17 & 7 \end{bmatrix}$$

Step 1 — Compute determinant

$$\det(K) = 2(2*7-1*17) - 4(9*7-1*3) + 5(9*17-2*3)$$

Step by step:

1.
$$2*7-1*17=14-17=-3 \Rightarrow 2*(-3)=-6$$

2.
$$9*7-1*3=63-3=60 \rightarrow -4*60=-240$$

3.
$$9*17 - 2*3 = 153 - 6 = 147 \rightarrow 5*147 = 735$$

Sum: -6 - 240 + 735 = 489

Modulo 26: 489 mod 26 = 21 $\rightarrow \det(K) = 21$

Check $gcd(21,26)=1 \rightarrow \bigvee$ invertible.

Step 2 — Modular inverse of determinant

Find x: $21 * x \equiv 1 \mod 26$

- Try x=5 \rightarrow 21*5=105 \equiv 1 mod 26 $\sqrt{}$
- So $det^{-1} = 5$

Step 3 — Compute adjugate (cofactor matrix transposed)

- Compute each cofactor $C_{ij} = (-1)^{i+j} imes \det(\operatorname{minor})$
- I'll show a few:

1. C11: det of minor
$$\begin{bmatrix} 2 & 1 \\ 17 & 7 \end{bmatrix} = 2*7 - 1*17 = -3 \Rightarrow$$
 multiply by (+1) = -3

2. C12: det of minor $\begin{bmatrix} 9 & 1 \\ 3 & 7 \end{bmatrix} = 9*7 - 1*3 = 63 - 3 = 60 \Rightarrow$ multiply by (-1) = -60

3. C13: det of minor $\begin{bmatrix} 9 & 2 \\ 3 & 17 \end{bmatrix} = 9*17 - 2*3 = 153 - 6 = 147 \Rightarrow$ (+1)*147=147

... Continue for all 9 entries, then transpose to get adjugate matrix.

Step 4 — Multiply adjugate by det^{-1} modulo 26

$$K^{-1} = 5 * adj(K) \mod 26$$

- Multiply each entry by 5, then reduce mod26 (convert negatives to positive by adding 26).
- This gives the final 3x3 inverse matrix, which can then be used for decryption.

Example 2:

encryption and decryption HELP using the 2×2 Hill cipher with key matrix

$$K = egin{bmatrix} 3 & 3 \ 2 & 5 \end{bmatrix}.$$

Step 1 — Convert letters to numbers (A=0 ... Z=25)

H=7, E=4, L=11, P=15. I split the plaintext into 2-letter blocks: $\mathbf{HE} \mid \mathbf{LP} \rightarrow \text{vectors } P_1 = \begin{bmatrix} 7 \\ 4 \end{bmatrix}$, $P_2 = \begin{bmatrix} 11 \\ 15 \end{bmatrix}$.

Step 2 — Encrypt each block: compute $C = K \cdot P \mod 26$.

For HE:

$$\bullet \quad \text{Multiply:} \begin{bmatrix} 3 & 3 \\ 2 & 5 \end{bmatrix} \begin{bmatrix} 7 \\ 4 \end{bmatrix} = \begin{bmatrix} 3 \times 7 + 3 \times 4 \\ 2 \times 7 + 5 \times 4 \end{bmatrix} = \begin{bmatrix} 21 + 12 \\ 14 + 20 \end{bmatrix} = \begin{bmatrix} 33 \\ 34 \end{bmatrix}.$$

• Reduce mod 26: $33 \mod 26 = 7, \ 34 \mod 26 = 8$. So $C_1 = \begin{bmatrix} 7 \\ 8 \end{bmatrix} o$ numbers 7,8 map to letters **H** I.

For LP:

$$\bullet \quad \text{Multiply:} \begin{bmatrix} 3 & 3 \\ 2 & 5 \end{bmatrix} \begin{bmatrix} 11 \\ 15 \end{bmatrix} = \begin{bmatrix} 3 \times 11 + 3 \times 15 \\ 2 \times 11 + 5 \times 15 \end{bmatrix} = \begin{bmatrix} 33 + 45 \\ 22 + 75 \end{bmatrix} = \begin{bmatrix} 78 \\ 97 \end{bmatrix}.$$

• Reduce mod 26: $78 \mod 26 = 0, \ 97 \mod 26 = 19$. So $C_2 = \begin{bmatrix} 0 \\ 19 \end{bmatrix}$ \rightarrow numbers 0,19 map to letters A T.

Ciphertext: combine blocks → **HIAT**.

Step 3 — Decryption (recover plaintext from HIAT).

I need K^{-1} modulo 26. First compute $\det(K)=3\times 5-3\times 2=15-6=9$. The modular inverse of 9 modulo 26 is the number d^{-1} with $9\cdot d^{-1}\equiv 1\pmod{26}$. Since $9\cdot 3=27\equiv 1\pmod{26}$, we have $9^{-1}=3$.

Compute the adjugate of K: $\mathrm{adj}(K)=\begin{bmatrix} d & -b \\ -c & a \end{bmatrix}=\begin{bmatrix} 5 & -3 \\ -2 & 3 \end{bmatrix}$. Multiply adjugate by $9^{-1}=3$ and reduce mod26:

$$K^{-1}=3\cdot egin{bmatrix} 5 & -3 \ -2 & 3 \end{bmatrix} mod 26 = egin{bmatrix} 15 & -9 \ -6 & 9 \end{bmatrix} mod 26 = egin{bmatrix} 15 & 17 \ 20 & 9 \end{bmatrix},$$

because $-9\equiv 17$ and $-6\equiv 20$ (mod 26). So $K^{-1}=egin{bmatrix} 15 & 17 \ 20 & 9 \end{bmatrix}$.

Now decrypt each ciphertext block C with $P = K^{-1} \cdot C \mod 26$.

For first block $C_1 = egin{bmatrix} 7 \\ 8 \end{bmatrix}$ (HI):

$$\bullet \quad \text{Multiply:} \begin{bmatrix} 15 & 17 \\ 20 & 9 \end{bmatrix} \begin{bmatrix} 7 \\ 8 \end{bmatrix} = \begin{bmatrix} 15 \times 7 + 17 \times 8 \\ 20 \times 7 + 9 \times 8 \end{bmatrix} = \begin{bmatrix} 105 + 136 \\ 140 + 72 \end{bmatrix} = \begin{bmatrix} 241 \\ 212 \end{bmatrix}.$$

• Reduce mod26: $241 \bmod 26 = 7, \ 212 \bmod 26 = 4.$ That recovers $\begin{bmatrix} 7 \\ 4 \end{bmatrix} \to \mathbf{H} \, \mathbf{E}.$

For second block $C_2 = egin{bmatrix} 0 \\ 19 \end{bmatrix}$ (AT):

$$\bullet \quad \text{Multiply:} \begin{bmatrix} 15 & 17 \\ 20 & 9 \end{bmatrix} \begin{bmatrix} 0 \\ 19 \end{bmatrix} = \begin{bmatrix} 15 \times 0 + 17 \times 19 \\ 20 \times 0 + 9 \times 19 \end{bmatrix} = \begin{bmatrix} 323 \\ 171 \end{bmatrix}.$$

• Reduce mod26: $323 \bmod 26 = 11, \ 171 \bmod 26 = 15.$ That recovers $\begin{bmatrix} 11 \\ 15 \end{bmatrix} \rightarrow$ L P.

Combine recovered blocks → H E | L P → HELP. Decryption successful.

Example 3:

I split and number the plaintext: HELLO \rightarrow pad \rightarrow HELLOX \rightarrow blocks HE | LL | OX. Numerical vectors: HE = [7,4], LL = [11,11], OX = [14,23]. Encryption: $C = K \cdot P \mod 26$.

1. For HE:
$$K \begin{bmatrix} 7 \\ 4 \end{bmatrix} = \begin{bmatrix} 3 \times 7 + 3 \times 4 \\ 2 \times 7 + 5 \times 4 \end{bmatrix} = \begin{bmatrix} 21 + 12 \\ 14 + 20 \end{bmatrix} = \begin{bmatrix} 33 \\ 34 \end{bmatrix} \mod 26 = \begin{bmatrix} 7 \\ 8 \end{bmatrix} \Rightarrow \text{letters H I}.$$

2. For LL: $K \begin{bmatrix} 11 \\ 11 \end{bmatrix} = \begin{bmatrix} 3 \times 11 + 3 \times 11 \\ 2 \times 11 + 5 \times 11 \end{bmatrix} = \begin{bmatrix} 33 + 33 \\ 22 + 55 \end{bmatrix} = \begin{bmatrix} 66 \\ 77 \end{bmatrix} \mod 26 = \begin{bmatrix} 14 \\ 25 \end{bmatrix} \Rightarrow \text{letters O Z}.$

3. For OX: $K \begin{bmatrix} 14 \\ 23 \end{bmatrix} = \begin{bmatrix} 3 \times 14 + 3 \times 23 \\ 2 \times 14 + 5 \times 23 \end{bmatrix} = \begin{bmatrix} 42 + 69 \\ 28 + 115 \end{bmatrix} = \begin{bmatrix} 111 \\ 143 \end{bmatrix} \mod 26 = \begin{bmatrix} 7 \\ 13 \end{bmatrix} \Rightarrow \text{letters H N}.$

Combine blocks → ciphertext H I O Z H N → HIOZHN

Now decryption to verify: $\det(\mathsf{K}) = 3 \cdot 5 - 3 \cdot 2 = 9$, inverse of 9 mod 26 is 3, adjugate $= \begin{bmatrix} 5 & -3 \\ -2 & 3 \end{bmatrix}$, so $K^{-1} = 3 \cdot \mathrm{adj} \mod 26 = \begin{bmatrix} 15 & 17 \\ 20 & 9 \end{bmatrix}$. Apply $P = K^{-1} \cdot C \mod 26$ to each ciphertext block: $\mathsf{HI} \to [7,8] \to K^{-1}[7,8] = [7,4] \to \mathsf{HE}$; $\mathsf{OZ} \to [14,25] \to K^{-1}[14,25] = [11,11] \to \mathsf{LL}$; $\mathsf{HN} \to [7,13] \to K^{-1}[7,13] = [14,23] \to \mathsf{OX}$. Recovered plaintext vector sequence $\to \mathsf{HELLOX}$; remove the padding $\mathsf{X} \to \mathsf{HELLO}$.

Example 4: 3 x 3 Hill Cipher

Step 1 — Prepare plaintext

• Plaintext: **HELLO** (length 5)

• 3×3 Hill cipher requires blocks of 3 → pad with X to make it 6 letters: **HELLOX**

• Blocks: HEL | LOX

Step 2 — Convert letters to numbers (A=0 ... Z=25)

• H=7, E=4, L=11, L=11, O=14, X=23

Block vectors:

• HEL \rightarrow [7, 4, 11]

• LOX \rightarrow [11, 14, 23]

Step 3 — Choose a 3×3 key matrix

Example key:

$$K = egin{bmatrix} 2 & 4 & 5 \ 9 & 2 & 1 \ 3 & 17 & 7 \end{bmatrix}$$

Key must be invertible modulo 26 for decryption.

Step 4 — Encrypt each block

Ciphertext vector: $C = K \cdot P \mod 26$

1. HEL → [7,4,11]

$$C = \begin{bmatrix} 2 & 4 & 5 \\ 9 & 2 & 1 \\ 3 & 17 & 7 \end{bmatrix} \begin{bmatrix} 7 \\ 4 \\ 11 \end{bmatrix} = \begin{bmatrix} 2*7+4*4+5*11 \\ 9*7+2*4+1*11 \\ 3*7+17*4+7*11 \end{bmatrix} = \begin{bmatrix} 14+16+55 \\ 63+8+11 \\ 21+68+77 \end{bmatrix} = \begin{bmatrix} 85 \\ 82 \\ 166 \end{bmatrix} \mod 26 = \begin{bmatrix} 7 \\ 4 \\ 10 \end{bmatrix}$$

Numbers \rightarrow letters: 7=H, 4=E, 10=K \rightarrow HEK

2. LOX \rightarrow [11,14,23]

$$C = K \cdot P \bmod 26 = \begin{bmatrix} 2*11 + 4*14 + 5*23 \\ 9*11 + 2*14 + 1*23 \\ 3*11 + 17*14 + 7*23 \end{bmatrix} = \begin{bmatrix} 22 + 56 + 115 \\ 99 + 28 + 23 \\ 33 + 238 + 161 \end{bmatrix} = \begin{bmatrix} 193 \\ 150 \\ 432 \end{bmatrix} \mod 26 = \begin{bmatrix} 11 \\ 20 \\ 16 \end{bmatrix}$$

Numbers → letters: 11=L, 20=U, 16=Q → LUQ