

# Public-Key Cryptography and RSA

An approach of asymmetric cipher

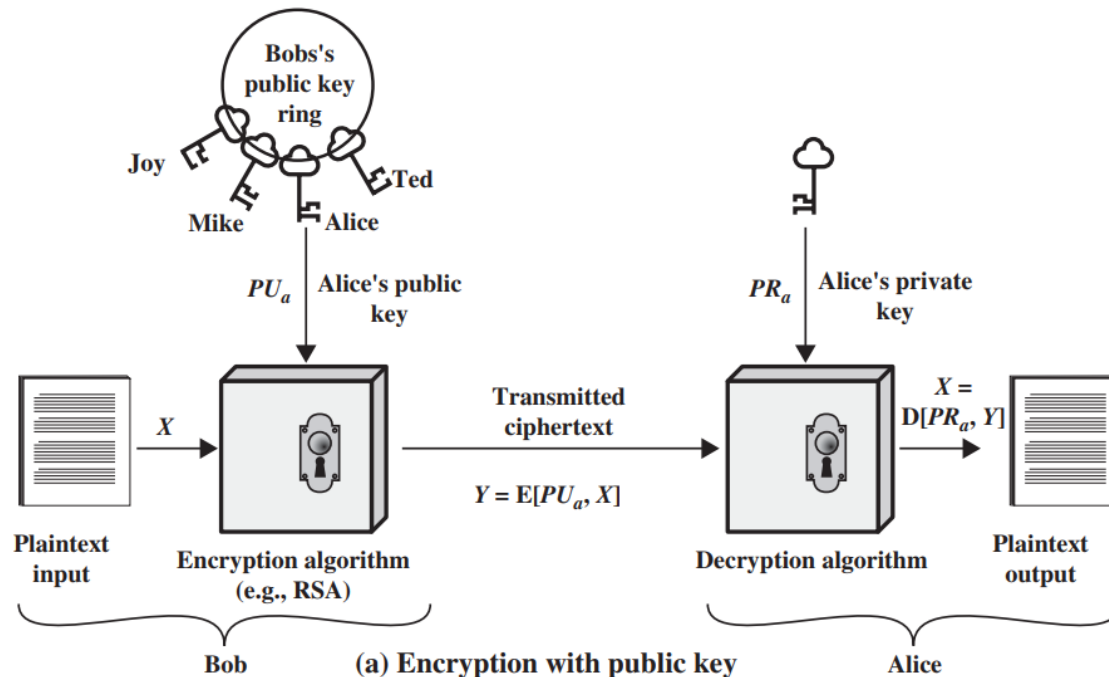
# Public-Key Cryptography

- Also known as Asymmetric Cryptography.
- Two related keys, a public key and a private key, that are used to perform encryption-decryption:
  - Public Key: shared openly; used for encryption.
  - Private Key: kept secret; used for decryption.

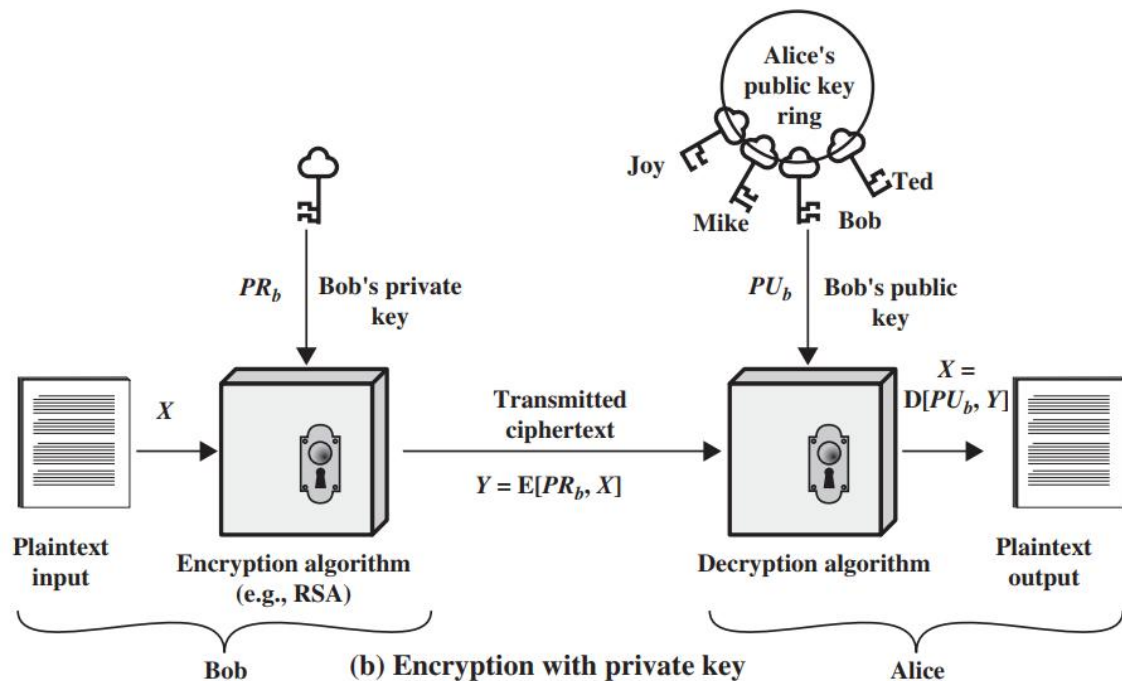
# A public-key encryption

- Asymmetric algorithms rely on one key for encryption and a different but related key for decryption.
- Either of the two related keys can be used for encryption, with the other used for decryption.

# public-key encryption with public key



# public-key encryption with private key



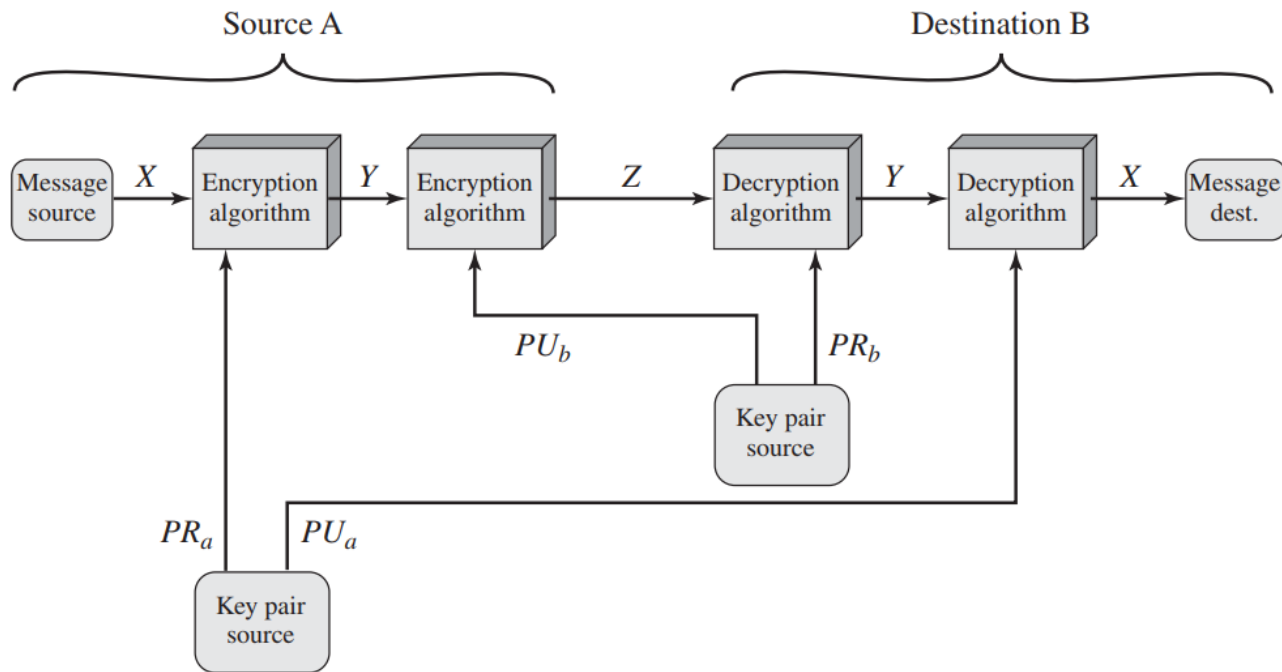
# The essential steps of public-key encryption

Each user generates a pair of keys to be used for the encryption and decryption of messages.

Each user places one of the two keys in a public register or other accessible file.

This is the public key. The companion key is kept private.  
Each user maintains a collection of public keys obtained from others.

# Public-Key Cryptosystem: Authentication and Secrecy



# RSA Algorithm

- Developed by Rivest, Shamir, and Adleman in 1977.
- Based on the mathematical difficulty of factoring large prime numbers.
- Most widely used public-key algorithm.
- RSA scheme is a block cipher in which the plaintext and ciphertext are integers between 0 and  $n - 1$  for some  $n$



# RSA Key Generation Steps

- Choose two large prime numbers,  $p$  and  $q$ .
- Compute  $n = p \times q$  (modulus).
- Compute  $\varphi(n) = (p - 1)(q - 1)$  (Euler's totient).
- Choose an encryption key  $e$ , such that  $1 < e < \varphi(n)$  and  $\gcd(e, \varphi(n)) = 1$ .
- Compute the decryption key  $d$ , such that  $d \times e \equiv 1 \pmod{\varphi(n)}$ .
- The public key =  $(e, n)$ ; the private key =  $(d, n)$ .

# Example

7 5 11

1. Select two prime numbers,  $p = 17$  and  $q = 11$ .
2. Calculate  $n = pq = 17 \times 11 = 187$ .
3. Calculate  $\phi(n) = (p - 1)(q - 1) = 16 \times 10 = 160$ . ✗ 0
4. Select  $e$  such that  $e$  is relatively prime to  $\phi(n) = 160$  and less than  $\phi(n)$ ; we choose  $e = 7$ .
5. Determine  $d$  such that  $de \equiv 1 \pmod{160}$  and  $d < 160$ . The correct value is  $d = 23$ , because  $23 \times 7 = 161 = (1 \times 160) + 1$ ;

# RSA Encryption Process

- To encrypt a message  $M$ :
- Convert the message into a numerical value  $m$ .
- Compute  $c = m^e \bmod n$ .
- The result  $c$  is the cipher text.

# RSA Decryption Process

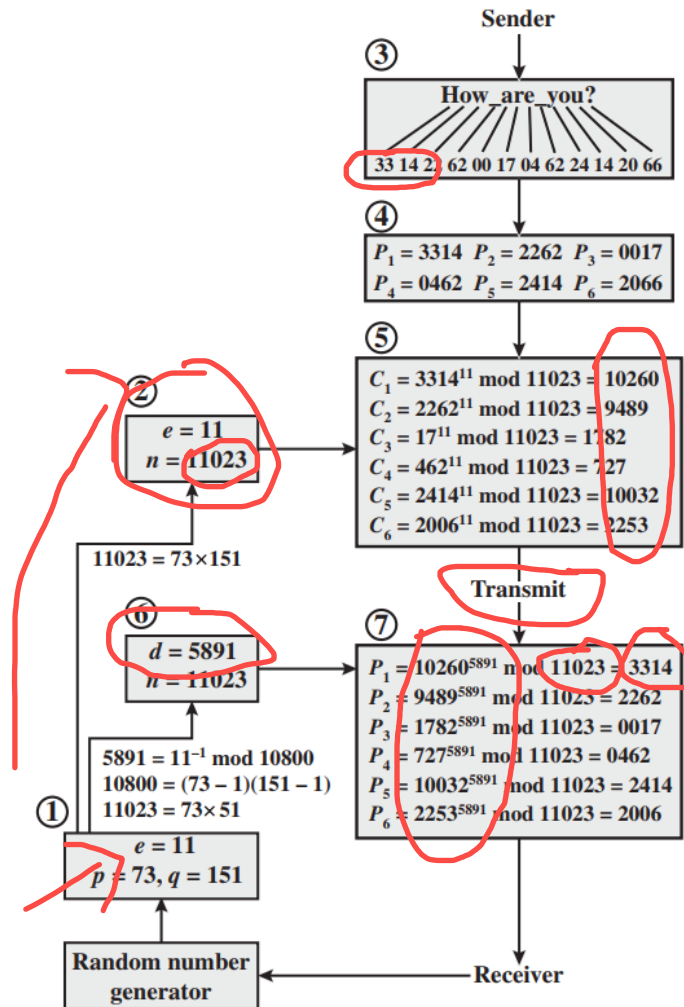
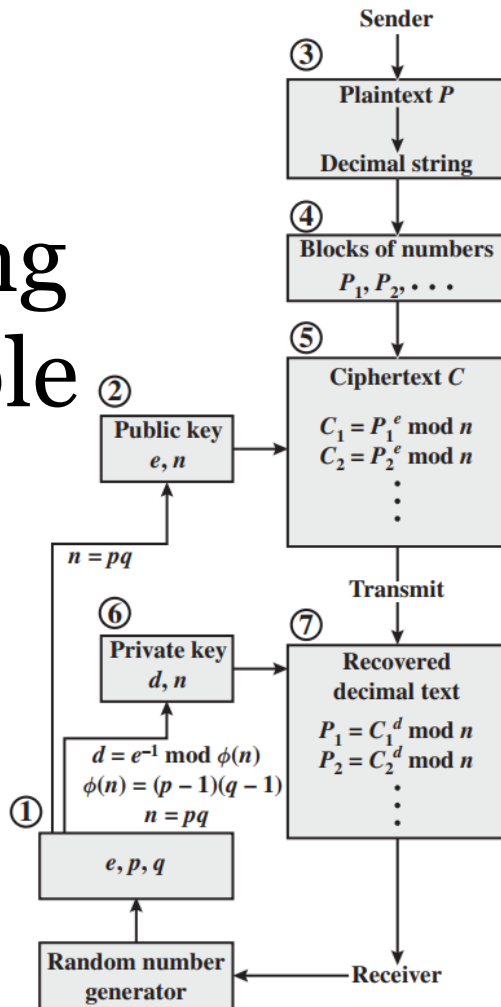
- To decrypt ciphertext  $c$ :
- Compute  $m = c^d \bmod n$ .
- Convert the numeric message  $m$  back to text.

# RSA Example (Small Numbers)

- Choose primes:  $\mathbf{p = 3, q = 11}$
- Compute  $\mathbf{n = 3 \times 11 = 33}$
- Compute  $\mathbf{\varphi(n) = (3-1)(11-1) = 20}$
- Choose  $\mathbf{e = 3}$  (since  $\gcd(3, 20) = 1$ )
- Compute  $\mathbf{d = 7}$  (since  $3 \times 7 \equiv 1 \pmod{20}$ )
- Public key =  $\mathbf{(3, 33)}$ , Private key =  $\mathbf{(7, 33)}$
- **Encryption:**  
Message  $m = 4 \rightarrow c = 4^3 \pmod{33} = 64 \pmod{33} = 31$
- **Decryption:**  
 $\rightarrow m = 31^7 \pmod{33} = 4$  (original message restored)

# RSA

## Processing of Multiple Blocks



# Strengths and Weaknesses of RSA

- **Strengths:**

- Proven mathematical foundation.
- Supports both encryption and authentication.
- Widely trusted and implemented.

- **Weaknesses:**

- Slower than symmetric ciphers.
- Vulnerable if key sizes are too small ( $< 2048$  bits).
- Quantum computing poses future risks.

Thank You