

Overview

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Contents

■ Part1: Course Overview

- Database Theory: Relational Model, Transaction, Indexing, and Optimization
- Database Practice: SQL, Oracle Practice, and database design

■ Part2: Introduction to Databases

Most parts are based on slides used in Stanford
(<http://web.stanford.edu/class/cs145>)

Introduction to Databases

Contents

1. Overview of the DBMS

1. Definition of DBMS
2. Data models & the relational data model
3. Schemas & data independence

What is a DBMS?

■ A large, integrated collection of data

■ Models a real-world enterprise

- *Entities* (e.g., Students, Courses)
- *Relationships* (e.g., Alice is enrolled in database course)

■ Consider building a course management system (CMS):

- Students
- Courses
- Professors

} *Entities*

- Who takes what
- Who teaches what

} *Relationships*

A Database Management System (DBMS) is a piece of software designed to store and manage databases

Data Models

■ A data model is a collection of concepts for describing data

- The relational model of data is the most widely used model today
 - Main Concept: the *relation*- essentially, a table

■ A schema is a description of a particular collection of data, using the given data model

- E.g. every *relation* in a relational data model has a *schema* describing types, etc.

		Columns			
Product	Rows	PName	Price	Category	Manufacturer
		Gizmo	\$19.99	Gadgets	GizmoWorks
		Powergizmo	\$29.99	Gadgets	GizmoWorks
		SingleTouch	\$149.99	Photography	Canon
		MultiTouch	\$203.99	Household	Hitachi

Modeling the CMS

Physical Schema: describes data layout

- Relations as unordered files
- Some data in sorted order (index)

Physical Schema : DB 물리적 구조 정의

-> Table / Index ... 의 물리적 저장소, Data file 위치 & 크기, Data 저장 방법 ... 포함
-> Logical Modeling 마친 후 물리적인 저장 방법으로 변환하는 과정을 거쳐야 함
=> 이때 물리적 구조 정의 역할

Logical Schema

- Students(sid: *string*, name: *string*, gpa: *float*)
- Courses(cid: *string*, cname: *string*, credits: *int*)
- Enrolled(sid: *string*, cid: *string*, grade: *string*)



Administrators



Applications

sid	Name	Gpa	Relations	cid	cname	credits						
101	Bob	3.2		564	564-2	4						
123	Mary	3.8		308	417	2						
Students			<table><tr><th>sid</th><th>cid</th><th>Grade</th></tr><tr><td>123</td><td>564</td><td>A</td></tr></table>	sid	cid	Grade	123	564	A	Courses		
sid	cid	Grade										
123	564	A										
Enrolled												

sid	Name	Gpa
101	Bob	3.2
123	Mary	3.8

Corresponding keys

cid	cname	credits
564	564-2	4
308	417	2

Students

sid	cid	Grade
123	564	A

Courses

Enrolled

Logical Schema : DB의 논리적 구조 정의

= DB에서 쓰이는 data 형식, 구조, 관계, .. 를 정의

-> 각 Entity & Attribute 간의 관계, 제약 조건, .. 정의

* 논리적 구조 : DB 사용하는 모든 사용자가 공유

=> data 일관성 & 무결성 보장 가능

-> Index, View, ... 의 구성 요소 정의 시에도 사용됨

External Schema: (Views)

- Course_info(cid: *string*, enrollment: *integer*)
- Derived from other tables

External Schema

: 특정 사용자 / 응용 프로그램이 DB에서 필요로 하는 data의 논리적 구조 정의

-> 필요한 data만 선택적으로 볼 수 있음

Data Independence

Concept: Applications do not need to worry about *how the data is structured and stored*

Logical data independence:

protection from changes in the *logical structure of the data*

I.e. should not need to ask:

can we add a new entity or attribute without rewriting the application?

Physical data independence:

protection from physical layout changes

I.e. should not need to ask:

which disks are the data stored on? Is the data indexed?

One of the most important reasons to use a DBMS

Data Independence

: 논리/물리적 구조에서의 변경이 다른 구조에 영향을 주지 않도록 하는 것

= 논리적 구조 변경해도 External Schema / App Code에 영향 X

= 물리적 구조 변경해도 논리적 구조 / 응용 프로그램에 영향 X

* 논리적 구조 : DB Schema

물리적 구조 : DB 저장 방식

2. Overview of DBMS Topics: Key Concepts & Challenges

Contents

- 1. Transactions**
- 2. Concurrency & locking**
- 3. Atomicity & logging**

Challenges with Many Users

■ Suppose that our CMS application serves 1000's of users or more- what are some challenges?

- Security: Different users, different roles

*We won't look at too much in this course,
but is extremely important*

사용자 인증, 접근 제어, 암호화 및 기타 보안 메커니즘을 제공
-> 데이터의 기밀성, 무결성 및 가용성을 보호

- Performance: Need to provide concurrent access

Disk/SSD access is slow,
DBMS hide the latency by doing more CPU work concurrently

많은 양의 데이터를 빠른 처리 & 많은 수의 동시 접속 처리 필요
-> DBMS는 쿼리 최적화, 인덱싱, 캐시 등의 기능을 제공

- Consistency: Concurrency can lead to update problems

DBMS allows user to write programs
as if they were the **only** user

여러 사용자가 동시에 DB 접근하면 data 일관성에 문제 발생 가능
⇒ Transaction & Lock Mechanism 제공
-> But, Lock mechanism은 Concurrency(동시성) 저해할 수 있음

Transactions

- A key concept is the transaction: an atomic sequence of DB actions (reads/writes)

Acct	Balance
a10	20,000
a20	15,000

Transfer \$3k from a10 to a20:

1. Debit \$3k from a10
2. Credit \$3k to a20

Acct	Balance
a10	17,000
a20	18,000

Atomicity: An action either completes *entirely* or *not at all*

Written naively, in which states is **atomicity** preserved?

- Crash before 1,
- After 1 but before 2,
- After 2.

DB Always preserves atomicity!

- 1) Transaction 실행 X -> DB에 영향 X
- 2) 1에서의 변경 사항은 commit X
 - > Transaction Rollback = DB 복원
- 3) Transaction 모든 작업 완료 -> commit

Transaction

1) Atomicity, 원자성

- > Transaction 은 전체 성공 / 실패인 원자적 작업 단위
- ⇒ 모든 작업 성공적 완료 시 DB 상태 변화
- ⇒ 작업 중 하나라도 실패 시 롤백 (일관성 유지에 중요한 역할)

2) Consistency, 일관성

- > Transaction 실행 전후의 DB 상태는 일관성 있어야 함
- = Transaction 실행되는 동안 DB는 일관된 상태 유지해야 함

3) Isolation, 격리성

- > 다른 Transaction에 영향 X = 독립적 실행 필요 (동시 작업이어도)

4) Durability, 지속성

- > 성공적으로 작업 완료 시 영구적으로 반영
- = system 장애 발생 시에도 해당 작업은 보존되어야 함

Concurrency: Scheduling Concurrent Transactions

- The DBMS ensures that the execution of $\{T_1, \dots, T_n\}$ is equivalent to some serial execution
- One way to accomplish this: Locking
 - Before reading or writing, transaction requires a lock from DBMS, holds until the end
- Key Idea: If T_i wants to write to an item x and T_j wants to read x , then T_i, T_j conflict.

Solution via locking:

- only one winner gets the lock
- loser is blocked (waits) until winner finishes

All concurrency issues handled by the DBMS...

A set of transactions is **isolated**
if their effect is as if all were executed serially

Locking

: Transaction이 Data 읽기/쓰기 전에 DBMS로부터 lock 요청 & 해당 lock 끝날 때까지 유지
-> 오직 하나의 transaction이 lock 얻을 수 있음
-> 다른 transaction은 잠금 해제 대기
⇒ 동시에 실행되는 여러 transaction 충돌 방지 & transaction 독립성 보장
⇒ DB 무결성 & 일관성 유지 가능

Isolated Transaction

: 모든 transaction이 일련의 serial 실행
& 동등한 결과를 만들도록 보장하는 것

Ensuring Atomicity

- DBMS ensures atomicity even if a transaction crashes!

- One way to accomplish this: Write-ahead logging (WAL)

Write-ahead Logging (WAL):

Before any action is finalized, a corresponding log entry is forced to disk

- Key Idea: Keep a log of all the writes done.

- After a crash, the partially executed transactions are undone using the log

All atomicity issues also handled by the DBMS...

WAL

:작업 수행 전 해당 작업에 대한 Log 기록을 disk에 저장

(Transaction 포함 & 작업 성공 여부 관계 X)

-> log 통해 system 장애 발생 / 새로운 요청 시 DB 복구 가능

⇒ System, 안정성 & 신뢰성 보장

A Well-Designed DBMS makes many people happy!

■ End users and DBMS vendors

- Reduces cost and makes money

Must still understand DB internals

■ DB application programmers

- Can handle more users, faster, for cheaper, and with better reliability / security guarantees!

■ Database administrators (DBA)

- Easier time of designing logical/physical schema, handling security/authorization, tuning, crash recovery, and more...

Summary

■ Key abstractions give data independence

■ DBMS are used to maintain, query, and manage large datasets.

- Provide concurrency, recovery from crashes, quick application development, integrity, and security